

Machine Words in Isabelle/HOL

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Abstract

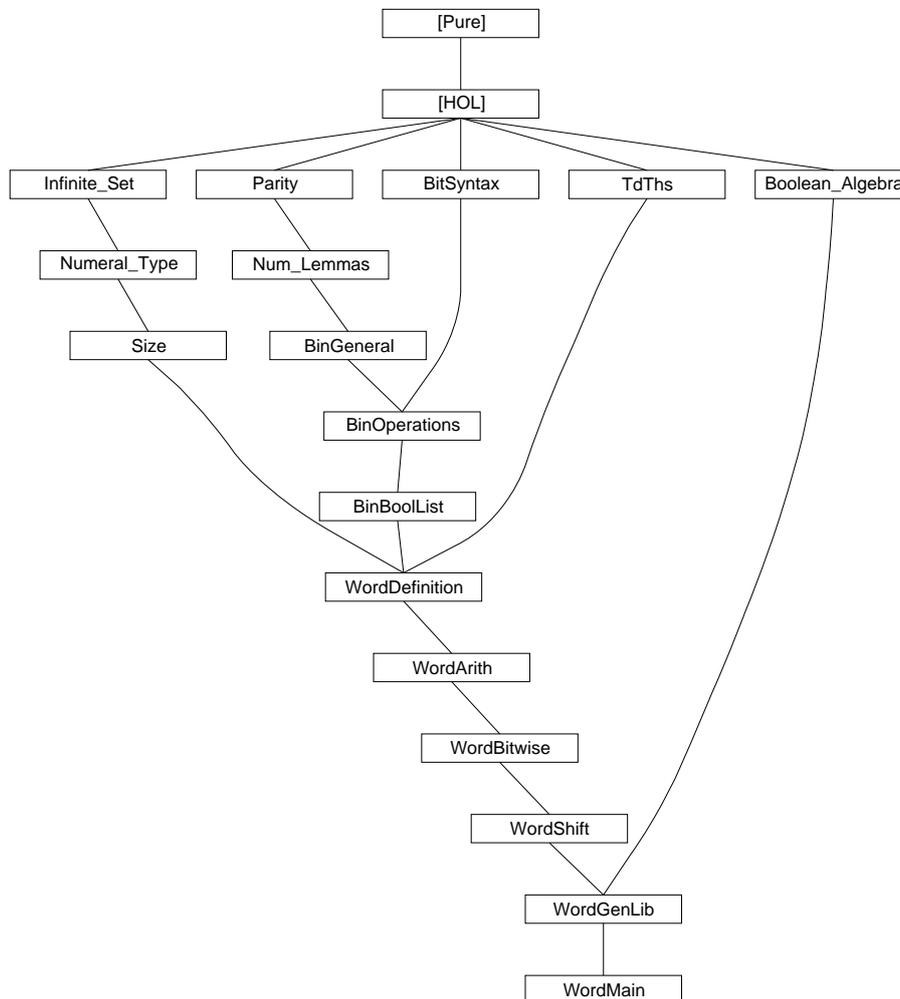
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
  imports Infinite-Set
begin
```

1.1 Preliminary lemmas

```
lemma inj-Inl [simp]: inj-on Inl A
  by (rule inj-onI, simp)
```

```
lemma inj-Inr [simp]: inj-on Inr A
  by (rule inj-onI, simp)
```

```
lemma inj-Some [simp]: inj-on Some A
  by (rule inj-onI, simp)
```

```
lemma card-Plus:
  [| finite A; finite B |] ==> card (A <+> B) = card A + card B
  unfolding Plus-def
  apply (subgoal-tac Inl ' A ∩ Inr ' B = {})
  apply (simp add: card-Un-disjoint card-image)
  apply fast
done
```

```
lemma (in type-definition) univ:
  UNIV = Abs ' A
```

proof

```
show Abs ' A ⊆ UNIV by (rule subset-UNIV)
```

```
show UNIV ⊆ Abs ' A
```

proof

```
fix x :: 'b
```

```
have x = Abs (Rep x) by (rule Rep-inverse [symmetric])
```

```
moreover have Rep x ∈ A by (rule Rep)
```

```
ultimately show x ∈ Abs ' A by (rule image-eqI)
```

qed

qed

```
lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  by (simp add: univ card-image inj-on-def Abs-inject)
```

1.2 Cardinalities of types

```
syntax -type-card :: type => nat ((1CARD/(1'(-))))
```

```
translations CARD(t) => card (UNIV::t set)
```

typed-print-translation \ll

let

```
fun card-univ-tr' show-sorts - [Const (@{const-name UNIV}, Type(-,[T]))] =
```

```

    Syntax.const -type-card $ Syntax.term-of-typ show-sorts T;
  in [(card, card-univ-tr')]
  end
  >>

```

```

lemma card-unit:  $CARD(unit) = 1$ 
  unfolding univ-unit by simp

```

```

lemma card-bool:  $CARD(bool) = 2$ 
  unfolding univ-bool by simp

```

```

lemma card-prod:  $CARD('a::finite \times 'b::finite) = CARD('a) * CARD('b)$ 
  unfolding univ-prod by (simp only: card-cartesian-product)

```

```

lemma card-sum:  $CARD('a::finite + 'b::finite) = CARD('a) + CARD('b)$ 
  unfolding univ-sum by (simp only: finite card-Plus)

```

```

lemma card-option:  $CARD('a::finite option) = Suc CARD('a)$ 
  unfolding univ-option
  apply (subgoal-tac (None::'a option) \notin range Some)
  apply (simp add: finite card-image)
  apply fast
  done

```

```

lemma card-set:  $CARD('a::finite set) = 2 ^ CARD('a)$ 
  unfolding univ-set
  by (simp only: card-Pow finite numeral-2-eq-2)

```

1.3 Numeral Types

```

typedef (open) num0 = UNIV :: nat set ..
typedef (open) num1 = UNIV :: unit set ..
typedef (open) 'a bit0 = UNIV :: (bool * 'a) set ..
typedef (open) 'a bit1 = UNIV :: (bool * 'a) option set ..

```

```

instance num1 :: finite

```

```

proof
  show finite (UNIV::num1 set)
    unfolding type-definition.univ [OF type-definition-num1]
    using finite by (rule finite-imageI)
qed

```

```

instance bit0 :: (finite) finite

```

```

proof
  show finite (UNIV::'a bit0 set)
    unfolding type-definition.univ [OF type-definition-bit0]
    using finite by (rule finite-imageI)
qed

```

```

instance bit1 :: (finite) finite
proof
  show finite (UNIV::'a bit1 set)
    unfolding type-definition.univ [OF type-definition-bit1]
    using finite by (rule finite-imageI)
qed

lemma card-num1: CARD(num1) = 1
  unfolding type-definition.card [OF type-definition-num1]
  by (simp only: card-unit)

lemma card-bit0: CARD('a::finite bit0) = 2 * CARD('a)
  unfolding type-definition.card [OF type-definition-bit0]
  by (simp only: card-prod card-bool)

lemma card-bit1: CARD('a::finite bit1) = Suc (2 * CARD('a))
  unfolding type-definition.card [OF type-definition-bit1]
  by (simp only: card-prod card-option card-bool)

lemma card-num0: CARD (num0) = 0
  by (simp add: type-definition.card [OF type-definition-num0])

lemmas card-univ-simps [simp] =
  card-unit
  card-bool
  card-prod
  card-sum
  card-option
  card-set
  card-num1
  card-bit0
  card-bit1
  card-num0

```

1.4 Syntax

```

syntax
  -NumeralType :: num-const => type (-)
  -NumeralType0 :: type (0)
  -NumeralType1 :: type (1)

```

```

translations
  -NumeralType1 == (type) num1
  -NumeralType0 == (type) num0

```

```

parse-translation <<
  let

```

```

  val num1-const = Syntax.const Numeral-Type.num1;

```

```

val num0-const = Syntax.const Numeral-Type.num0;
val B0-const = Syntax.const Numeral-Type.bit0;
val B1-const = Syntax.const Numeral-Type.bit1;

fun mk-bintype n =
  let
    fun mk-bit n = if n = 0 then B0-const else B1-const;
    fun bin-of n =
      if n = 1 then num1-const
      else if n = 0 then num0-const
      else if n = ~1 then raise TERM (negative type numeral, [])
      else
        let val (q, r) = Integer.div-mod n 2;
            in mk-bit r $ bin-of q end;
    in bin-of n end;

fun numeral-tr (*-NumeralType*) [Const (str, -)] =
  mk-bintype (valOf (Int.fromString str))
| numeral-tr (*-NumeralType*) ts = raise TERM (numeral-tr, ts);

in [(-NumeralType, numeral-tr)] end;
>>

print-translation <<
let
  fun int-of [] = 0
  | int-of (b :: bs) = b + 2 * int-of bs;

  fun bin-of (Const (num0, -)) = []
  | bin-of (Const (num1, -)) = [1]
  | bin-of (Const (bit0, -) $ bs) = 0 :: bin-of bs
  | bin-of (Const (bit1, -) $ bs) = 1 :: bin-of bs
  | bin-of t = raise TERM(bin-of, [t]);

  fun bit-tr' b [t] =
    let
      val rev-digs = b :: bin-of t handle TERM - => raise Match
      val i = int-of rev-digs;
      val num = string-of-int (abs i);
    in
      Syntax.const -NumeralType $ Syntax.free num
    end
  | bit-tr' b - = raise Match;

in [(bit0, bit-tr' 0), (bit1, bit-tr' 1)] end;
>>

```

1.5 Classes with at least 1 and 2

Class `finite` already captures “at least 1”

```
lemma zero-less-card-finite [simp]:
  0 < CARD('a::finite)
proof (cases CARD('a::finite) = 0)
  case False thus ?thesis by (simp del: card-0-eq)
next
  case True
  thus ?thesis by (simp add: finite)
qed
```

```
lemma one-le-card-finite [simp]:
  Suc 0 <= CARD('a::finite)
  by (simp add: less-Suc-eq-le [symmetric] zero-less-card-finite)
```

Class for cardinality “at least 2”

```
class card2 = finite +
  assumes two-le-card: 2 <= CARD('a)
```

```
lemma one-less-card: Suc 0 < CARD('a::card2)
  using two-le-card [where 'a='a] by simp
```

```
instance bit0 :: (finite) card2
  by intro-classes (simp add: one-le-card-finite)
```

```
instance bit1 :: (finite) card2
  by intro-classes (simp add: one-le-card-finite)
```

1.6 Examples

```
term TYPE(10)
```

```
lemma CARD(0) = 0 by simp
lemma CARD(17) = 17 by simp
```

```
end
```

2 Size: The size class

```
theory Size
imports Numeral-Type
begin
```

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in `Numeral_Type`.

axclass *len0* < *type*

consts

len-of :: ('a :: *len0* itself) => *nat*

Some theorems are only true on words with length greater 0.

axclass *len* < *len0*

len-gt-0 [*iff*]: 0 < *len-of TYPE* ('a :: *len0*)

instance *num0* :: *len0* ..

instance *num1* :: *len0* ..

instance *bit0* :: (*len0*) *len0* ..

instance *bit1* :: (*len0*) *len0* ..

defs (**overloaded**)

len-num0: *len-of* (*x*::*num0* itself) == 0

len-num1: *len-of* (*x*::*num1* itself) == 1

len-bit0: *len-of* (*x*::'a::*len0* *bit0* itself) == 2 * *len-of TYPE* ('a)

len-bit1: *len-of* (*x*::'a::*len0* *bit1* itself) == 2 * *len-of TYPE* ('a) + 1

lemmas *len-of-numeral-defs* [*simp*] = *len-num0 len-num1 len-bit0 len-bit1*

instance *num1* :: *len* **by** (*intro-classes*) *simp*

instance *bit0* :: (*len*) *len* **by** (*intro-classes*) *simp*

instance *bit1* :: (*len0*) *len* **by** (*intro-classes*) *simp*

— Examples:

lemma *len-of TYPE*(17) = 17 **by** *simp*

lemma *len-of TYPE*(0) = 0 **by** *simp*

— not simplified:

lemma *len-of TYPE*('a::*len0*) = *x*

oops

end

3 Num-Lemmas: Useful Numerical Lemmas

theory *Num-Lemmas* **imports** *Parity* **begin**

lemma *contentsI*: $y = \{x\} \implies \text{contents } y = x$

unfolding *contents-def* **by** *auto*

lemma *prod-case-split*: *prod-case* = *split*

by (*rule ext*)⁺ *auto*

lemmas *split-split* = *prod.split* [*unfolded prod-case-split*]

lemmas *split-split-asm* = *prod.split-asm* [*unfolded prod-case-split*]
lemmas *split.splits* = *split-split split-split-asm*

lemmas *funpow-0* = *funpow.simps(1)*
lemmas *funpow-Suc* = *funpow.simps(2)*

lemma *nonemptyE*: $S \sim = \{\}$ $\implies (!x. x : S \implies R) \implies R$
apply (*erule contrapos-np*)
apply (*rule equals0I*)
apply *auto*
done

lemma *gt-or-eq-0*: $0 < y \vee 0 = (y::nat)$ **by** *auto*

constdefs

mod-alt :: $'a \implies 'a \implies 'a :: \text{Divides.div}$
mod-alt $n\ m == n \bmod m$

— alternative way of defining *bin-last*, *bin-rest*

bin-rl :: $int \implies int * bit$
bin-rl $w == \text{SOME } (r, l). w = r \text{ BIT } l$

declare *iszero-0* [*iff*]

lemmas *xtr1* = *xtrans(1)*
lemmas *xtr2* = *xtrans(2)*
lemmas *xtr3* = *xtrans(3)*
lemmas *xtr4* = *xtrans(4)*
lemmas *xtr5* = *xtrans(5)*
lemmas *xtr6* = *xtrans(6)*
lemmas *xtr7* = *xtrans(7)*
lemmas *xtr8* = *xtrans(8)*

lemma *Min-ne-Pls* [*iff*]:
Numeral.Min $\sim = \text{Numeral.Pl}$
unfolding *Min-def Pls-def* **by** *auto*

lemmas *Pls-ne-Min* [*iff*] = *Min-ne-Pls* [*symmetric*]

lemmas *PlsMin-defs* [*intro!*] =
Pls-def Min-def Pls-def [*symmetric*] *Min-def* [*symmetric*]

lemmas *PlsMin-simps* [*simp*] = *PlsMin-defs* [*THEN Eq-TrueI*]

lemma *number-of-False-cong*:
 $\text{False} \implies \text{number-of } x = \text{number-of } y$
by *auto*

lemmas *nat-simps* = *diff-add-inverse2 diff-add-inverse*

lemmas *nat-iffs* = *le-add1 le-add2*

lemma *sum-imp-diff*: $j = k + i \implies j - i = (k :: \text{nat})$
by (*clarsimp simp add: nat-simps*)

lemma *nobm1*:

$0 < (\text{number-of } w :: \text{nat}) \implies$
 $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Numeral.pred } w)$
apply (*unfold nat-number-of-def One-nat-def nat-1 [symmetric] pred-def*)
apply (*simp add: number-of-eq nat-diff-distrib [symmetric]*)
done

lemma *of-int-power*:

of-int ($a \wedge n$) = (*of-int* $a \wedge n :: 'a :: \{\text{recpower, comm-ring-1}\}$)
by (*induct n*) (*auto simp add: power-Suc*)

lemma *zless2*: $0 < (2 :: \text{int})$

by *auto*

lemmas *zless2p* [*simp*] = *zless2* [*THEN zero-less-power*]

lemmas *zle2p* [*simp*] = *zless2p* [*THEN order-less-imp-le*]

lemmas *pos-mod-sign2* = *zless2* [*THEN pos-mod-sign* [**where** $b = 2 :: \text{int}$]]

lemmas *pos-mod-bound2* = *zless2* [*THEN pos-mod-bound* [**where** $b = 2 :: \text{int}$]]

— the inverse(s) of *number-of*

lemma *nmod2*: $n \bmod (2 :: \text{int}) = 0 \mid n \bmod 2 = 1$

using *pos-mod-sign2* [*of n*] *pos-mod-bound2* [*of n*]
unfolding *mod-alt-def* [*symmetric*] **by** *auto*

lemma *emep1*:

even n \implies *even d* \implies $0 \leq d \implies (n + 1) \bmod (d :: \text{int}) = (n \bmod d) + 1$

apply (*simp add: add-commute*)
apply (*safe dest!: even-equiv-def [THEN iffD1]*)
apply (*subst pos-zmod-mult-2*)
apply *arith*
apply (*simp add: zmod-zmult-zmult1*)
done

lemmas *eme1p* = *emep1* [*simplified add-commute*]

lemma *le-diff-eq'*: $(a \leq c - b) = (b + a \leq (c :: \text{int}))$

by (*simp add: le-diff-eq add-commute*)

lemma *less-diff-eq'*: $(a < c - b) = (b + a < (c :: \text{int}))$

by (*simp add: less-diff-eq add-commute*)

lemma *diff-le-eq'*: $(a - b \leq c) = (a \leq b + (c :: \text{int}))$

by (*simp add: diff-le-eq add-commute*)

lemma *diff-less-eq'*: $(a - b < c) = (a < b + (c::int))$
by (*simp add: diff-less-eq add-commute*)

lemmas *m1mod2k* = *zless2p* [*THEN zmod-minus1*]
lemmas *m1mod22k* = *mult-pos-pos* [*OF zless2 zless2p, THEN zmod-minus1*]
lemmas *p1mod22k'* = *zless2p* [*THEN order-less-imp-le, THEN pos-zmod-mult-2*]
lemmas *z1pmod2'* = *zero-le-one* [*THEN pos-zmod-mult-2, simplified*]
lemmas *z1pdiv2'* = *zero-le-one* [*THEN pos-zdiv-mult-2, simplified*]

lemma *p1mod22k*:
 $(2 * b + 1) \bmod (2 * 2 ^ n) = 2 * (b \bmod 2 ^ n) + (1::int)$
by (*simp add: p1mod22k' add-commute*)

lemma *z1pmod2*:
 $(2 * b + 1) \bmod 2 = (1::int)$
by (*simp add: z1pmod2' add-commute*)

lemma *z1pdiv2*:
 $(2 * b + 1) \bmod 2 = (b::int)$
by (*simp add: z1pdiv2' add-commute*)

lemmas *zdiv-le-dividend* = *xtr3* [*OF zdiv-1 [symmetric] zdiv-mono2, simplified int-one-le-iff-zero-less, simplified, standard*]

lemma *BIT-eq*: $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$
apply (*unfold Bit-def*)
apply (*simp (no-asm-use) split: bit.split-asm*)
apply *simp-all*
apply (*drule-tac f=even in arg-cong, clarsimp*)+
done

lemmas *BIT-eqE* [*elim!*] = *BIT-eq* [*THEN conjE, standard*]

lemma *BIT-eq-iff* [*simp*]:
 $(u \text{ BIT } b = v \text{ BIT } c) = (u = v \ \wedge \ b = c)$
by (*rule iffI*) *auto*

lemmas *BIT-eqI* [*intro!*] = *conjI* [*THEN BIT-eq-iff [THEN iffD2]*]

lemma *less-Bits*:
 $(v \text{ BIT } b < w \text{ BIT } c) = (v < w \ | \ v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$
unfolding *Bit-def* **by** (*auto split: bit.split*)

lemma *le-Bits*:

$(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ (b \sim = \text{bit.B1} \mid c \sim = \text{bit.B0}))$
unfolding *Bit-def* **by** (*auto split: bit.split*)

lemma *neB1E* [*elim!*]:
assumes *ne*: $y \neq \text{bit.B1}$
assumes *y*: $y = \text{bit.B0} \implies P$
shows *P*
apply (*rule y*)
apply (*cases y rule: bit.exhaust, simp*)
apply (*simp add: ne*)
done

lemma *bin-ex-rl*: $EX \ w \ b. \ w \ \text{BIT} \ b = \text{bin}$
apply (*unfold Bit-def*)
apply (*cases even bin*)
apply (*clarsimp simp: even-equiv-def*)
apply (*auto simp: odd-equiv-def split: bit.split*)
done

lemma *bin-exhaust*:
assumes *Q*: $\bigwedge x \ b. \ \text{bin} = x \ \text{BIT} \ b \implies Q$
shows *Q*
apply (*insert bin-ex-rl [of bin]*)
apply (*erule exE*)
apply (*rule Q*)
apply *force*
done

lemma *bin-rl-char*: $(\text{bin-rl } w = (r, l)) = (r \ \text{BIT} \ l = w)$
apply (*unfold bin-rl-def*)
apply *safe*
apply (*cases w rule: bin-exhaust*)
apply *auto*
done

lemmas *bin-rl-simps* [*THEN bin-rl-char [THEN iffD2], standard, simp*] =
Pls-0-eq Min-1-eq refl

lemma *bin-abs-lem*:
 $\text{bin} = (w \ \text{BIT} \ b) \implies \sim \text{bin} = \text{Numeral.Min} \dashrightarrow \sim \text{bin} = \text{Numeral.Pls} \dashrightarrow$
 $\text{nat } (\text{abs } w) < \text{nat } (\text{abs } \text{bin})$
apply (*clarsimp simp add: bin-rl-char*)
apply (*unfold Pls-def Min-def Bit-def*)
apply (*cases b*)
apply (*clarsimp, arith*)
apply (*clarsimp, arith*)
done

lemma *bin-induct*:

```

assumes PPls: P Numeral.Pls
and PMin: P Numeral.Min
and PBit: !!bin bit. P bin ==> P (bin BIT bit)
shows P bin
apply (rule-tac P=P and a=bin and f1=nat o abs
in wf-measure [THEN wf-induct])
apply (simp add: measure-def inv-image-def)
apply (case-tac x rule: bin-exhaust)
apply (frule bin-abs-lem)
apply (auto simp add : PPls PMin PBit)
done

```

```

lemma no-no [simp]: number-of (number-of i) = i
unfolding number-of-eq by simp

```

```

lemma Bit-B0:
  k BIT bit.B0 = k + k
by (unfold Bit-def) simp

```

```

lemma Bit-B1:
  k BIT bit.B1 = k + k + 1
by (unfold Bit-def) simp

```

```

lemma Bit-B0-2t: k BIT bit.B0 = 2 * k
by (rule trans, rule Bit-B0) simp

```

```

lemma Bit-B1-2t: k BIT bit.B1 = 2 * k + 1
by (rule trans, rule Bit-B1) simp

```

```

lemma B-mod-2':
  X = 2 ==> (w BIT bit.B1) mod X = 1 & (w BIT bit.B0) mod X = 0
apply (simp (no-asm) only: Bit-B0 Bit-B1)
apply (simp add: z1pmod2)
done

```

```

lemmas B1-mod-2 [simp] = B-mod-2' [OF refl, THEN conjunct1, standard]

```

```

lemmas B0-mod-2 [simp] = B-mod-2' [OF refl, THEN conjunct2, standard]

```

```

lemma axbby:
  a + m + m = b + n + n ==> (a = 0 | a = 1) ==> (b = 0 | b = 1) ==>
  a = b & m = (n :: int)
apply auto
apply (drule-tac f=%n. n mod 2 in arg-cong)
apply (clarsimp simp: z1pmod2)
apply (drule-tac f=%n. n mod 2 in arg-cong)
apply (clarsimp simp: z1pmod2)
done

```

```

lemma axxmod2:

```

$(1 + x + x) \text{ mod } 2 = (1 :: \text{int}) \ \& \ (0 + x + x) \text{ mod } 2 = (0 :: \text{int})$
by *simp* (rule *z1pmod2*)

lemma *axxdiv2*:

$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int})$
by *simp* (rule *z1pdiv2*)

lemmas *iszero-minus* = *trans* [*THEN trans*,

OF iszero-def neg-equal-0-iff-equal iszero-def [*symmetric*], *standard*]

lemmas *zadd-diff-inverse* = *trans* [*OF diff-add-cancel* [*symmetric*] *add-commute*,
standard]

lemmas *add-diff-cancel2* = *add-commute* [*THEN diff-eq-eq* [*THEN iffD2*], *standard*]

lemma *zmod-uminus*: $-(a :: \text{int}) \text{ mod } b \text{ mod } b = -a \text{ mod } b$
by (*simp add* : *zmod-zminus1-eq-if*)

lemma *zmod-zsub-distrib*: $((a :: \text{int}) - b) \text{ mod } c = (a \text{ mod } c - b \text{ mod } c) \text{ mod } c$
apply (*unfold diff-int-def*)
apply (*rule trans* [*OF - zmod-zadd1-eq* [*symmetric*]])
apply (*simp add*: *zmod-uminus zmod-zadd1-eq* [*symmetric*])
done

lemma *zmod-zsub-right-eq*: $((a :: \text{int}) - b) \text{ mod } c = (a - b \text{ mod } c) \text{ mod } c$
apply (*unfold diff-int-def*)
apply (*rule trans* [*OF - zmod-zadd-right-eq* [*symmetric*]])
apply (*simp add* : *zmod-uminus zmod-zadd-right-eq* [*symmetric*])
done

lemma *zmod-zsub-left-eq*: $((a :: \text{int}) - b) \text{ mod } c = (a \text{ mod } c - b) \text{ mod } c$
by (*rule zmod-zadd-left-eq* [**where** $b = -b$, *simplified diff-int-def* [*symmetric*]])

lemma *zmod-zsub-self* [*simp*]:

$((b :: \text{int}) - a) \text{ mod } a = b \text{ mod } a$
by (*simp add*: *zmod-zsub-right-eq*)

lemma *zmod-zmult1-eq-rev*:

$b * a \text{ mod } c = b \text{ mod } c * a \text{ mod } (c :: \text{int})$
apply (*simp add*: *mult-commute*)
apply (*subst zmod-zmult1-eq*)
apply *simp*
done

lemmas *rdmods* [*symmetric*] = *zmod-uminus* [*symmetric*]
zmod-zsub-left-eq zmod-zsub-right-eq zmod-zadd-left-eq
zmod-zadd-right-eq zmod-zmult1-eq zmod-zmult1-eq-rev

lemma *mod-plus-right*:

```
((a + x) mod m = (b + x) mod m) = (a mod m = b mod (m :: nat))
apply (induct x)
apply (simp-all add: mod-Suc)
apply arith
done
```

lemma *nat-minus-mod*: $(n - n \text{ mod } m) \text{ mod } m = (0 :: \text{nat})$

```
by (induct n) (simp-all add : mod-Suc)
```

lemmas *nat-minus-mod-plus-right* = trans [OF *nat-minus-mod mod-0* [symmetric],
THEN *mod-plus-right* [THEN iffD2], standard, simplified]

lemmas *push-mods'* = *zmod-zadd1-eq* [standard]

```
zmod-zmult-distrib [standard] zmod-zsub-distrib [standard]
```

```
zmod-uminus [symmetric, standard]
```

lemmas *push-mods* = *push-mods'* [THEN *eq-reflection*, standard]

lemmas *pull-mods* = *push-mods* [symmetric] *rdmods* [THEN *eq-reflection*, standard]

lemmas *mod-simps* =

```
zmod-zmult-self1 [THEN eq-reflection] zmod-zmult-self2 [THEN eq-reflection]
```

```
mod-mod-trivial [THEN eq-reflection]
```

lemma *nat-mod-eq*:

```
!!b. b < n ==> a mod n = b mod n ==> a mod n = (b :: nat)
```

```
by (induct a) auto
```

lemmas *nat-mod-eq'* = refl [THEN [2] *nat-mod-eq*]

lemma *nat-mod-lem*:

```
(0 :: nat) < n ==> b < n = (b mod n = b)
```

```
apply safe
```

```
apply (erule nat-mod-eq')
```

```
apply (erule subst)
```

```
apply (erule mod-less-divisor)
```

```
done
```

lemma *mod-nat-add*:

```
(x :: nat) < z ==> y < z ==>
```

```
(x + y) mod z = (if x + y < z then x + y else x + y - z)
```

```
apply (rule nat-mod-eq)
```

```
apply auto
```

```
apply (rule trans)
```

```
apply (rule le-mod-geq)
```

```
apply simp
```

```
apply (rule nat-mod-eq')
```

```
apply arith
```

```
done
```

lemma *mod-nat-sub*:

$(x :: \text{nat}) < z \implies (x - y) \bmod z = x - y$
by (*rule nat-mod-eq'*) *arith*

lemma *int-mod-lem*:

$(0 :: \text{int}) < n \implies (0 \leq b \ \& \ b < n) = (b \bmod n = b)$
apply *safe*
apply (*erule* (1) *mod-pos-pos-trivial*)
apply (*erule-tac* [!] *subst*)
apply *auto*
done

lemma *int-mod-eq*:

$(0 :: \text{int}) \leq b \implies b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$
by *clarsimp* (*rule mod-pos-pos-trivial*)

lemmas *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

lemma *int-mod-le*: $0 \leq a \implies 0 < (n :: \text{int}) \implies a \bmod n \leq a$

apply (*cases* $a < n$)
apply (*auto dest: mod-pos-pos-trivial pos-mod-bound* [**where** $a=a$])
done

lemma *int-mod-le'*: $0 \leq b - n \implies 0 < (n :: \text{int}) \implies b \bmod n \leq b - n$

by (*rule int-mod-le* [**where** $a = b - n$ **and** $n = n$, *simplified*])

lemma *int-mod-ge*: $a < n \implies 0 < (n :: \text{int}) \implies a \leq a \bmod n$

apply (*cases* $0 \leq a$)
apply (*erule* (1) *mod-pos-pos-trivial*)
apply *simp*
apply (*rule order-trans* [*OF* - *pos-mod-sign*])
apply *simp*
apply *assumption*
done

lemma *int-mod-ge'*: $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$

by (*rule int-mod-ge* [**where** $a = b + n$ **and** $n = n$, *simplified*])

lemma *mod-add-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
by (*auto intro: int-mod-eq*)

lemma *mod-sub-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
by (*auto intro: int-mod-eq*)

lemmas *zmde* = *zmod-zdiv-equality* [*THEN diff-eq-eq* [*THEN iffD2*], *symmetric*]
lemmas *mcl* = *mult-cancel-left* [*THEN iffD1*, *THEN make-pos-rule*]

lemma *zdiv-mult-self*: $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$
apply (*rule mcl*)
prefer 2
apply (*erule asm-rl*)
apply (*simp add: zmde ring-distrib*)
apply (*simp add: push-mods*)
done

lemma *eqne*: $\text{equiv } A \ r \implies X : A // r \implies X \sim = \{\}$
unfolding *equiv-def refl-def quotient-def Image-def* **by** *auto*

lemmas *Rep-Integ-ne* = *Integ.Rep-Integ*
[*THEN equiv-intrel* [*THEN eqne*, *simplified Integ-def* [*symmetric*]], *standard*]

lemmas *riq* = *Integ.Rep-Integ* [*simplified Integ-def*]
lemmas *intrel-refl* = *refl* [*THEN equiv-intrel-iff* [*THEN iffD1*], *standard*]
lemmas *Rep-Integ-equiv* = *quotient-eq-iff*
[*OF equiv-intrel riq riq*, *simplified Integ.Rep-Integ-inject*, *standard*]
lemmas *Rep-Integ-same* =
Rep-Integ-equiv [*THEN intrel-refl* [*THEN rev-iffD2*], *standard*]

lemma *RI-int*: $(a, 0) : \text{Rep-Integ } (\text{int } a)$
unfolding *int-def* **by** *auto*

lemmas *RI-intrel* [*simp*] = *UNIV-I* [*THEN quotientI*,
THEN Integ.Abs-Integ-inverse [*simplified Integ-def*], *standard*]

lemma *RI-minus*: $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (- x)$
apply (*rule-tac z=x in eq-Abs-Integ*)
apply (*clarsimp simp: minus*)
done

lemma *RI-add*:
 $(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$
 $(a + c, b + d) : \text{Rep-Integ } (x + y)$
apply (*rule-tac z=x in eq-Abs-Integ*)
apply (*rule-tac z=y in eq-Abs-Integ*)
apply (*clarsimp simp: add*)
done

lemma *mem-same*: $a : S \implies a = b \implies b : S$
by *fast*

```

lemma RI-eq-diff': (a, b) : Rep-Integ (int a - int b)
  apply (unfold diff-def)
  apply (rule mem-same)
  apply (rule RI-minus RI-add RI-int)+
  apply simp
  done

lemma RI-eq-diff: ((a, b) : Rep-Integ x) = (int a - int b = x)
  apply safe
  apply (rule Rep-Integ-same)
  prefer 2
  apply (erule asm-rl)
  apply (rule RI-eq-diff')+
  done

lemma mod-power-lem:
  a > 1 ==> a ^ n mod a ^ m = (if m <= n then 0 else (a :: int) ^ n)
  apply clarsimp
  apply safe
  apply (simp add: zdvd-iff-zmod-eq-0 [symmetric])
  apply (erule le-iff-add [THEN iffD1])
  apply (force simp: zpower-zadd-distrib)
  apply (rule mod-pos-pos-trivial)
  apply (simp add: zero-le-power)
  apply (rule power-strict-increasing)
  apply auto
  done

lemma min-pm [simp]: min a b + (a - b) = (a :: nat)
  by arith

lemmas min-pm1 [simp] = trans [OF add-commute min-pm]

lemma rev-min-pm [simp]: min b a + (a - b) = (a::nat)
  by simp

lemmas rev-min-pm1 [simp] = trans [OF add-commute rev-min-pm]

lemma pl-pl-rels:
  a + b = c + d ==>
  a >= c & b <= d | a <= c & b >= (d :: nat)
  apply (cut-tac n=a and m=c in nat-le-linear)
  apply (safe dest!: le-iff-add [THEN iffD1])
  apply arith+
  done

lemmas pl-pl-rels' = add-commute [THEN [2] trans, THEN pl-pl-rels]

lemma minus-eq: (m - k = m) = (k = 0 | m = (0 :: nat))

```

by *arith*

lemma *pl-pl-mm*: $(a :: \text{nat}) + b = c + d \implies a - c = d - b$
by *arith*

lemmas *pl-pl-mm'* = *add-commute* [THEN [2] *trans*, THEN *pl-pl-mm*]

lemma *min-minus* [*simp*] : $\text{min } m (m - k) = (m - k :: \text{nat})$
by *arith*

lemmas *min-minus'* [*simp*] = *trans* [OF *min-max.inf-commute min-minus*]

lemma *nat-no-eq-iff*:

$(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$
 $(\text{number-of } b = (\text{number-of } c :: \text{nat})) = (b = c)$

apply (*unfold nat-number-of-def*)

apply *safe*

apply (*drule* (2) *eq-nat-nat-iff* [THEN *iffD1*])

apply (*simp add: number-of-eq*)

done

lemmas *dme* = *box-equals* [OF *div-mod-equality add-0-right add-0-right*]

lemmas *dtle* = *xtr3* [OF *dme* [*symmetric*] *le-add1*]

lemmas *th2* = *order-trans* [OF *order-refl* [THEN [2] *mult-le-mono*] *dtle*]

lemma *td-gal*:

$0 < c \implies (a \geq b * c) = (a \text{ div } c \geq (b :: \text{nat}))$

apply *safe*

apply (*erule* (1) *xtr4* [OF *div-le-mono div-mult-self-is-m*])

apply (*erule th2*)

done

lemmas *td-gal-lt* = *td-gal* [*simplified le-def*, *simplified*]

lemma *div-mult-le*: $(a :: \text{nat}) \text{ div } b * b \leq a$

apply (*cases b*)

prefer 2

apply (*rule order-refl* [THEN *th2*])

apply *auto*

done

lemmas *sdl* = *split-div-lemma* [THEN *iffD1*, *symmetric*]

lemma *given-quot*: $f > (0 :: \text{nat}) \implies (f * l + (f - 1)) \text{ div } f = l$

by (*rule sdl*, *assumption*) (*simp (no-asm)*)

lemma *given-quot-alt*: $f > (0 :: \text{nat}) \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$

apply (*frule given-quot*)

apply (*rule trans*)

```

prefer 2
apply (erule asm-rl)
apply (rule-tac f=%n. n div f in arg-cong)
apply (simp add : mult-ac)
done

lemma diff-mod-le: (a::nat) < d ==> b dvd d ==> a - a mod b <= d - b
apply (unfold dvd-def)
apply clarify
apply (case-tac k)
apply clarsimp
apply clarify
apply (cases b > 0)
apply (drule mult-commute [THEN xtr1])
apply (frule (1) td-gal-lt [THEN iffD1])
apply (clarsimp simp: le-simps)
apply (rule mult-div-cancel [THEN [2] xtr4])
apply (rule mult-mono)
apply auto
done

lemma less-le-mult':
  w * c < b * c ==> 0 ≤ c ==> (w + 1) * c ≤ b * (c::int)
apply (rule mult-right-mono)
apply (rule zless-imp-add1-zle)
apply (erule (1) mult-right-less-imp-less)
apply assumption
done

lemmas less-le-mult = less-le-mult' [simplified left-distrib, simplified]

lemmas less-le-mult-minus = iffD2 [OF le-diff-eq less-le-mult,
  simplified left-diff-distrib, standard]

lemma lrlem':
assumes d: (i::nat) ≤ j ∨ m < j'
assumes R1: i * k ≤ j * k ==> R
assumes R2: Suc m * k' ≤ j' * k' ==> R
shows R using d
apply safe
apply (rule R1, erule mult-le-mono1)
apply (rule R2, erule Suc-le-eq [THEN iffD2 [THEN mult-le-mono1]])
done

lemma lrlem: (0::nat) < sc ==>
  (sc - n + (n + lb * n) <= m * n) = (sc + lb * n <= m * n)
apply safe
apply arith
apply (case-tac sc >= n)

```

```

  apply arith
  apply (insert linorder-le-less-linear [of m lb])
  apply (erule-tac k=n and k'=n in lrlem')
  apply arith
  apply simp
  done

lemma gen-minus: 0 < n ==> f n = f (Suc (n - 1))
  by auto

lemma mpl-lem: j <= (i :: nat) ==> k < j ==> i - j + k < i
  apply (induct i, clarsimp)
  apply (cases j, clarsimp)
  apply arith
  done

lemma nonneg-mod-div:
  0 <= a ==> 0 <= b ==> 0 <= (a mod b :: int) & 0 <= a div b
  apply (cases b = 0, clarsimp)
  apply (auto intro: pos-imp-zdiv-nonneg-iff [THEN iffD2])
  done

end

```

4 BinGeneral: Basic Definitions for Binary Integers

```
theory BinGeneral imports Num-Lemmas
```

```
begin
```

4.1 Recursion combinator for binary integers

```
lemma brlem: (bin = Numeral.Min) = (- bin + Numeral.pred 0 = 0)
  unfolding Min-def pred-def by arith
```

```
function
```

```
bin-rec' :: int * 'a * 'a * (int => bit => 'a => 'a) => 'a
```

```
where
```

```
bin-rec' (bin, f1, f2, f3) = (if bin = Numeral.Plus then f1
```

```
  else if bin = Numeral.Min then f2
```

```
  else case bin-rl bin of (w, b) => f3 w b (bin-rec' (w, f1, f2, f3)))
```

```
by pat-completeness auto
```

```
termination
```

```
apply (relation measure (nat o abs o fst))
```

```
apply simp
```

```

apply (simp add: Pls-def brlem)
apply (clarsimp simp: bin-rl-char pred-def)
apply (frule thin-rl [THEN refl [THEN bin-abs-lem [rule-format]]])
  apply (unfold Pls-def Min-def number-of-eq)
  prefer 2
  apply (erule asm-rl)
apply auto
done

```

constdefs

```

bin-rec :: 'a => 'a => (int => bit => 'a => 'a) => int => 'a
bin-rec f1 f2 f3 bin == bin-rec' (bin, f1, f2, f3)

```

lemma bin-rec-PM:

```

f = bin-rec f1 f2 f3 ==> f Numeral.Pls = f1 & f Numeral.Min = f2
apply safe
apply (unfold bin-rec-def)
apply (auto intro: bin-rec'.simps [THEN trans])
done

```

lemmas bin-rec-Pls = refl [THEN bin-rec-PM, THEN conjunct1, standard]

lemmas bin-rec-Min = refl [THEN bin-rec-PM, THEN conjunct2, standard]

lemma bin-rec-Bit:

```

f = bin-rec f1 f2 f3 ==> f3 Numeral.Pls bit.B0 f1 = f1 ==>
  f3 Numeral.Min bit.B1 f2 = f2 ==> f (w BIT b) = f3 w b (f w)
apply clarify
apply (unfold bin-rec-def)
apply (rule bin-rec'.simps [THEN trans])
apply auto
  apply (unfold Pls-def Min-def Bit-def)
  apply (cases b, auto)+
done

```

lemmas bin-rec-simps = refl [THEN bin-rec-Bit] bin-rec-Pls bin-rec-Min

4.2 Destructors for binary integers

consts

— corresponding operations analysing bins

```

bin-last :: int => bit
bin-rest :: int => int
bin-sign :: int => int
bin-nth :: int => nat => bool

```

primrec

```

Z : bin-nth w 0 = (bin-last w = bit.B1)
Suc : bin-nth w (Suc n) = bin-nth (bin-rest w) n

```

defs

bin-rest-def : $\text{bin-rest } w == \text{fst } (\text{bin-rl } w)$
bin-last-def : $\text{bin-last } w == \text{snd } (\text{bin-rl } w)$
bin-sign-def : $\text{bin-sign} == \text{bin-rec } \text{Numeral.Pls } \text{Numeral.Min } (\%_0 w \text{ b s. s})$

lemma *bin-rl*: $\text{bin-rl } w = (\text{bin-rest } w, \text{bin-last } w)$
unfolding *bin-rest-def bin-last-def* **by** *auto*

lemmas *bin-rl-simp* [*simp*] = *iffD1* [*OF bin-rl-char bin-rl*]

lemma *bin-rest-simps* [*simp*]:
bin-rest Numeral.Pls = *Numeral.Pls*
bin-rest Numeral.Min = *Numeral.Min*
bin-rest (w BIT b) = *w*
unfolding *bin-rest-def* **by** *auto*

lemma *bin-last-simps* [*simp*]:
bin-last Numeral.Pls = *bit.B0*
bin-last Numeral.Min = *bit.B1*
bin-last (w BIT b) = *b*
unfolding *bin-last-def* **by** *auto*

lemma *bin-sign-simps* [*simp*]:
bin-sign Numeral.Pls = *Numeral.Pls*
bin-sign Numeral.Min = *Numeral.Min*
bin-sign (w BIT b) = *bin-sign w*
unfolding *bin-sign-def* **by** (*auto simp: bin-rec-simps*)

lemma *bin-r-l-extras* [*simp*]:
bin-last 0 = *bit.B0*
bin-last (- 1) = *bit.B1*
bin-last -1 = *bit.B1*
bin-last 1 = *bit.B1*
bin-rest 1 = *0*
bin-rest 0 = *0*
bin-rest (- 1) = *- 1*
bin-rest -1 = *-1*
apply (*unfold number-of-Min*)
apply (*unfold Pls-def [symmetric] Min-def [symmetric]*)
apply (*unfold numeral-1-eq-1 [symmetric]*)
apply (*auto simp: number-of-eq*)
done

lemma *bin-last-mod*:
bin-last w = (*if w mod 2 = 0 then bit.B0 else bit.B1*)
apply (*case-tac w rule: bin-exhaust*)
apply (*case-tac b*)
apply *auto*
done

lemma *bin-rest-div*:
bin-rest $w = w \text{ div } 2$
apply (*case-tac* w *rule*: *bin-exhaust*)
apply (*rule trans*)
apply *clarsimp*
apply (*rule refl*)
apply (*drule trans*)
apply (*rule Bit-def*)
apply (*simp add*: *z1pdiv2 split*: *bit.split*)
done

lemma *Bit-div2* [*simp*]: $(w \text{ BIT } b) \text{ div } 2 = w$
unfolding *bin-rest-div* [*symmetric*] **by** *auto*

lemma *bin-nth-lem* [*rule-format*]:
ALL y . $\text{bin-nth } x = \text{bin-nth } y \dashrightarrow x = y$
apply (*induct* x *rule*: *bin-induct*)
apply *safe*
apply (*erule rev-mp*)
apply (*induct-tac* y *rule*: *bin-induct*)
apply *safe*
apply (*drule-tac* $x=0$ **in** *fun-cong*, *force*)
apply (*erule notE*, *rule ext*,
drule-tac $x=\text{Suc } x$ **in** *fun-cong*, *force*)
apply (*drule-tac* $x=0$ **in** *fun-cong*, *force*)
apply (*erule rev-mp*)
apply (*induct-tac* y *rule*: *bin-induct*)
apply *safe*
apply (*drule-tac* $x=0$ **in** *fun-cong*, *force*)
apply (*erule notE*, *rule ext*,
drule-tac $x=\text{Suc } x$ **in** *fun-cong*, *force*)
apply (*drule-tac* $x=0$ **in** *fun-cong*, *force*)
apply (*case-tac* y *rule*: *bin-exhaust*)
apply *clarify*
apply (*erule allE*)
apply (*erule impE*)
prefer 2
apply (*erule BIT-eqI*)
apply (*drule-tac* $x=0$ **in** *fun-cong*, *force*)
apply (*rule ext*)
apply (*drule-tac* $x=\text{Suc } ?x$ **in** *fun-cong*, *force*)
done

lemma *bin-nth-eq-iff*: $(\text{bin-nth } x = \text{bin-nth } y) = (x = y)$
by (*auto elim*: *bin-nth-lem*)

lemmas *bin-eqI = ext* [*THEN bin-nth-eq-iff* [*THEN iffD1*], *standard*]

lemma *bin-nth-Pls* [*simp*]: \sim *bin-nth Numeral.Pls* *n*
by (*induct n*) *auto*

lemma *bin-nth-Min* [*simp*]: *bin-nth Numeral.Min* *n*
by (*induct n*) *auto*

lemma *bin-nth-0-BIT*: *bin-nth* (*w BIT b*) 0 = (*b = bit.B1*)
by *auto*

lemma *bin-nth-Suc-BIT*: *bin-nth* (*w BIT b*) (*Suc n*) = *bin-nth w n*
by *auto*

lemma *bin-nth-minus* [*simp*]: $0 < n \implies$ *bin-nth* (*w BIT b*) *n* = *bin-nth w* (*n* - 1)
by (*cases n*) *auto*

lemmas *bin-nth-0 = bin-nth.simps(1)*
lemmas *bin-nth-Suc = bin-nth.simps(2)*

lemmas *bin-nth-simps =*
bin-nth-0 bin-nth-Suc bin-nth-Pls bin-nth-Min bin-nth-minus

lemma *bin-sign-rest* [*simp*]:
bin-sign (*bin-rest w*) = (*bin-sign w*)
by (*case-tac w rule: bin-exhaust*) *auto*

4.3 Truncating binary integers

consts

bintrunc :: *nat* => *int* => *int*

primrec

Z : *bintrunc* 0 *bin* = *Numeral.Pls*

Suc : *bintrunc* (*Suc n*) *bin* = *bintrunc n* (*bin-rest bin*) *BIT* (*bin-last bin*)

consts

sbintrunc :: *nat* => *int* => *int*

primrec

Z : *sbintrunc* 0 *bin* =

(*case bin-last bin of bit.B1 => Numeral.Min | bit.B0 => Numeral.Pls*)

Suc : *sbintrunc* (*Suc n*) *bin* = *sbintrunc n* (*bin-rest bin*) *BIT* (*bin-last bin*)

lemma *sign-bintr*:

!!*w*. *bin-sign* (*bintrunc n w*) = *Numeral.Pls*

by (*induct n*) *auto*

lemma *bintrunc-mod2p*:

!!*w*. *bintrunc n w* = (*w mod 2 ^ n* :: *int*)

apply (*induct n, clarsimp*)

apply (*simp add: bin-last-mod bin-rest-div Bit-def zmod-zmult2-eq*)

```

      cong: number-of-False-cong)
done

lemma sbintrunc-mod2p:
!!w. sbintrunc n w = ((w + 2 ^ n) mod 2 ^ (Suc n) - 2 ^ n :: int)
apply (induct n)
  apply clarsimp
  apply (subst zmod-zadd-left-eq)
  apply (simp add: bin-last-mod)
  apply (simp add: number-of-eq)
  apply clarsimp
  apply (simp add: bin-last-mod bin-rest-div Bit-def
    cong: number-of-False-cong)
  apply (clarsimp simp: zmod-zmult-zmult1 [symmetric]
    zmod-zdiv-equality [THEN diff-eq-eq [THEN iffD2 [THEN sym]]])
  apply (rule trans [symmetric, OF - emep1])
  apply auto
  apply (auto simp: even-def)
done

```

4.4 Simplifications for (s)bintrunc

```

lemma bit-bool:
(b = (b' = bit.B1)) = (b' = (if b then bit.B1 else bit.B0))
by (cases b') auto

```

```

lemmas bit-bool1 [simp] = refl [THEN bit-bool [THEN iffD1], symmetric]

```

```

lemma bin-sign-lem:
!!bin. (bin-sign (sbintrunc n bin) = Numeral.Min) = bin-nth bin n
apply (induct n)
  apply (case-tac bin rule: bin-exhaust, case-tac b, auto)+
done

```

```

lemma nth-bintr:
!!w m. bin-nth (bintrunc m w) n = (n < m & bin-nth w n)
apply (induct n)
  apply (case-tac m, auto)[1]
  apply (case-tac m, auto)[1]
done

```

```

lemma nth-sbintr:
!!w m. bin-nth (sbintrunc m w) n =
  (if n < m then bin-nth w n else bin-nth w m)
apply (induct n)
  apply (case-tac m, simp-all split: bit.splits)[1]
  apply (case-tac m, simp-all split: bit.splits)[1]
done

```

lemma *bin-nth-Bit*:

bin-nth (w BIT b) n = (n = 0 & b = bit.B1 | (EX m. n = Suc m & bin-nth w m))
by (*cases n*) *auto*

lemma *bintrunc-bintrunc-l*:

n <= m ==> (bintrunc m (bintrunc n w) = bintrunc n w)
by (*rule bin-eqI*) (*auto simp add : nth-bintr*)

lemma *sbintrunc-sbintrunc-l*:

n <= m ==> (sbintrunc m (sbintrunc n w) = sbintrunc n w)
by (*rule bin-eqI*) (*auto simp: nth-sbintr min-def*)

lemma *bintrunc-bintrunc-ge*:

n <= m ==> (bintrunc n (bintrunc m w) = bintrunc n w)
by (*rule bin-eqI*) (*auto simp: nth-bintr*)

lemma *bintrunc-bintrunc-min [simp]*:

bintrunc m (bintrunc n w) = bintrunc (min m n) w
apply (*unfold min-def*)
apply (*rule bin-eqI*)
apply (*auto simp: nth-bintr*)
done

lemma *sbintrunc-sbintrunc-min [simp]*:

sbintrunc m (sbintrunc n w) = sbintrunc (min m n) w
apply (*unfold min-def*)
apply (*rule bin-eqI*)
apply (*auto simp: nth-sbintr*)
done

lemmas *bintrunc-Pls =*

bintrunc.Suc [where bin=Numeral.Pls, simplified bin-last-simps bin-rest-simps, standard]

lemmas *bintrunc-Min [simp] =*

bintrunc.Suc [where bin=Numeral.Min, simplified bin-last-simps bin-rest-simps, standard]

lemmas *bintrunc-BIT [simp] =*

bintrunc.Suc [where bin=w BIT b, simplified bin-last-simps bin-rest-simps, standard]

lemmas *bintrunc-Sucs = bintrunc-Pls bintrunc-Min bintrunc-BIT*

lemmas *sbintrunc-Suc-Pls =*

sbintrunc.Suc [where bin=Numeral.Pls, simplified bin-last-simps bin-rest-simps, standard]

lemmas *sbintrunc-Suc-Min* =
sbintrunc.Suc [**where** *bin=Numeral.Min*, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *sbintrunc-Suc-BIT* [*simp*] =
sbintrunc.Suc [**where** *bin=w BIT b*, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls* *sbintrunc-Suc-Min* *sbintrunc-Suc-BIT*

lemmas *sbintrunc-Pls* =
sbintrunc.Z [**where** *bin=Numeral.Pl*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-Min* =
sbintrunc.Z [**where** *bin=Numeral.Min*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-0-BIT-B0* [*simp*] =
sbintrunc.Z [**where** *bin=w BIT bit.B0*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-0-BIT-B1* [*simp*] =
sbintrunc.Z [**where** *bin=w BIT bit.B1*,
simplified bin-last-simps bin-rest-simps bit.simps, *standard*]

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls *sbintrunc-Min* *sbintrunc-0-BIT-B0* *sbintrunc-0-BIT-B1*

lemmas *bintrunc-simps* = *bintrunc.Z* *bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps* *sbintrunc-Sucs*

lemma *bintrunc-minus*:
 $0 < n ==> \text{bintrunc } (\text{Suc } (n - 1)) w = \text{bintrunc } n w$
by *auto*

lemma *sbintrunc-minus*:
 $0 < n ==> \text{sbintrunc } (\text{Suc } (n - 1)) w = \text{sbintrunc } n w$
by *auto*

lemmas *bintrunc-minus-simps* =
bintrunc-Sucs [*THEN* [2] *bintrunc-minus* [*symmetric*, *THEN trans*], *standard*]

lemmas *sbintrunc-minus-simps* =
sbintrunc-Sucs [*THEN* [2] *sbintrunc-minus* [*symmetric*, *THEN trans*], *standard*]

lemma *bintrunc-n-Pls* [*simp*]:
 $\text{bintrunc } n \text{ Numeral.Pl} = \text{Numeral.Pl}$
by (*induct n*) *auto*

lemma *sbintrunc-n-PM* [simp]:

sbintrunc n Numeral.Pls = Numeral.Pls
sbintrunc n Numeral.Min = Numeral.Min
by (induct n) auto

lemmas *thobini1* = arg-cong [where $f = \%w. w \text{ BIT } b$, standard]

lemmas *bintrunc-BIT-I* = trans [OF *bintrunc-BIT thobini1*]

lemmas *bintrunc-Min-I* = trans [OF *bintrunc-Min thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps* [THEN *thobini1* [THEN [2] trans], standard]

lemmas *bintrunc-Pls-minus-I* = *bmsts*(1)

lemmas *bintrunc-Min-minus-I* = *bmsts*(2)

lemmas *bintrunc-BIT-minus-I* = *bmsts*(3)

lemma *bintrunc-0-Min*: *bintrunc 0 Numeral.Min = Numeral.Pls*

by auto

lemma *bintrunc-0-BIT*: *bintrunc 0 (w BIT b) = Numeral.Pls*

by auto

lemma *bintrunc-Suc-lem*:

bintrunc (Suc n) x = y ==> m = Suc n ==> bintrunc m x = y

by auto

lemmas *bintrunc-Suc-Ialts* =

bintrunc-Min-I bintrunc-BIT-I [THEN *bintrunc-Suc-lem*, standard]

lemmas *sbintrunc-BIT-I* = trans [OF *sbintrunc-Suc-BIT thobini1*]

lemmas *sbintrunc-Suc-Is* =

sbintrunc-Sucs [THEN *thobini1* [THEN [2] trans], standard]

lemmas *sbintrunc-Suc-minus-Is* =

sbintrunc-minus-simps [THEN *thobini1* [THEN [2] trans], standard]

lemma *sbintrunc-Suc-lem*:

sbintrunc (Suc n) x = y ==> m = Suc n ==> sbintrunc m x = y

by auto

lemmas *sbintrunc-Suc-Ialts* =

sbintrunc-Suc-Is [THEN *sbintrunc-Suc-lem*, standard]

lemma *sbintrunc-bintrunc-lt*:

$m > n ==> sbintrunc n (bintrunc m w) = sbintrunc n w$

by (rule *bin-eqI*) (auto simp: *nth-sbintr nth-bintr*)

lemma *bintrunc-sbintrunc-le*:

$m \leq Suc n ==> bintrunc m (sbintrunc n w) = bintrunc m w$

```

apply (rule bin-eqI)
apply (auto simp: nth-sbintr nth-bintr)
apply (subgoal-tac x=n, safe, arith+)[1]
apply (subgoal-tac x=n, safe, arith+)[1]
done

```

```

lemmas bintrunc-sbintrunc [simp] = order-refl [THEN bintrunc-sbintrunc-le]
lemmas sbintrunc-bintrunc [simp] = lessI [THEN sbintrunc-bintrunc-lt]
lemmas bintrunc-bintrunc [simp] = order-refl [THEN bintrunc-bintrunc-l]
lemmas sbintrunc-sbintrunc [simp] = order-refl [THEN sbintrunc-sbintrunc-l]

```

```

lemma bintrunc-sbintrunc' [simp]:
   $0 < n \implies \text{bintrunc } n (\text{sbintrunc } (n - 1) w) = \text{bintrunc } n w$ 
by (cases n) (auto simp del: bintrunc.Suc)

```

```

lemma sbintrunc-bintrunc' [simp]:
   $0 < n \implies \text{sbintrunc } (n - 1) (\text{bintrunc } n w) = \text{sbintrunc } (n - 1) w$ 
by (cases n) (auto simp del: bintrunc.Suc)

```

```

lemma bin-sbin-eq-iff:
   $\text{bintrunc } (\text{Suc } n) x = \text{bintrunc } (\text{Suc } n) y \iff$ 
   $\text{sbintrunc } n x = \text{sbintrunc } n y$ 
apply (rule iffI)
apply (rule box-equals [OF - sbintrunc-bintrunc sbintrunc-bintrunc])
apply simp
apply (rule box-equals [OF - bintrunc-sbintrunc bintrunc-sbintrunc])
apply simp
done

```

```

lemma bin-sbin-eq-iff':
   $0 < n \implies \text{bintrunc } n x = \text{bintrunc } n y \iff$ 
   $\text{sbintrunc } (n - 1) x = \text{sbintrunc } (n - 1) y$ 
by (cases n) (simp-all add: bin-sbin-eq-iff del: bintrunc.Suc)

```

```

lemmas bintrunc-sbintruncS0 [simp] = bintrunc-sbintrunc' [unfolded One-nat-def]
lemmas sbintrunc-bintruncS0 [simp] = sbintrunc-bintrunc' [unfolded One-nat-def]

```

```

lemmas bintrunc-bintrunc-l' = le-add1 [THEN bintrunc-bintrunc-l]
lemmas sbintrunc-sbintrunc-l' = le-add1 [THEN sbintrunc-sbintrunc-l]

```

```

lemmas nat-non0-gr =
  trans [OF iszero-def [THEN Not-eq-iff [THEN iffD2]] refl, standard]

```

```

lemmas bintrunc-pred-simps [simp] =
  bintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

```

```

lemmas sbintrunc-pred-simps [simp] =

```

sbintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

lemma *no-bintr-alt*:

number-of (bintrunc n w) = $w \bmod 2^n$
by (*simp add*: *number-of-eq bintrunc-mod2p*)

lemma *no-bintr-alt1*: bintrunc $n = (\%w. w \bmod 2^n :: int)$

by (*rule ext*) (*rule bintrunc-mod2p*)

lemma *range-bintrunc*: range (bintrunc n) = $\{i. 0 \leq i \ \& \ i < 2^n\}$

apply (*unfold no-bintr-alt1*)

apply (*auto simp add*: *image-iff*)

apply (*rule exI*)

apply (*auto intro*: *int-mod-lem* [THEN *iffD1*, *symmetric*])

done

lemma *no-bintr*:

number-of (bintrunc n w) = (*number-of* $w \bmod 2^n :: int$)

by (*simp add* : *bintrunc-mod2p number-of-eq*)

lemma *no-sbintr-alt2*:

sbintrunc $n = (\%w. (w + 2^n) \bmod 2^{Suc\ n} - 2^n :: int)$

by (*rule ext*) (*simp add* : *sbintrunc-mod2p*)

lemma *no-sbintr*:

number-of (*sbintrunc* n w) =

(*number-of* $w + 2^n \bmod 2^{Suc\ n} - 2^n :: int$)

by (*simp add* : *no-sbintr-alt2 number-of-eq*)

lemma *range-sbintrunc*:

range (*sbintrunc* n) = $\{i. -(2^n) \leq i \ \& \ i < 2^n\}$

apply (*unfold no-sbintr-alt2*)

apply (*auto simp add*: *image-iff eq-diff-eq*)

apply (*rule exI*)

apply (*auto intro*: *int-mod-lem* [THEN *iffD1*, *symmetric*])

done

lemma *sb-inc-lem*:

($a :: int$) + $2^k < 0 \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$

apply (*erule int-mod-ge'* [**where** $n = 2^{(Suc\ k)}$ **and** $b = a + 2^k$, *simplified zless2p*])

apply (*rule TrueI*)

done

lemma *sb-inc-lem'*:

($a :: int$) < $-(2^k) \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$

by (*rule iffD1* [*OF less-diff-eq*, THEN *sb-inc-lem*, *simplified OrderedGroup.diff-0*])

lemma *sbintrunc-inc*:

$x < - (2^n) \implies x + 2^{(Suc\ n)} \leq sbintrunc\ n\ x$
unfolding *no-sbintr-alt2* **by** (*drule sb-inc-lem'*) *simp*

lemma *sb-dec-lem*:

$(0::int) \leq - (2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$
by (*rule int-mod-le'* [**where** $n = 2^k$ **and** $b = a + 2^k$,
simplified zless2p, OF - TrueI, simplified])

lemma *sb-dec-lem'*:

$(2::int)^k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$
by (*rule iffD1* [*OF diff-le-eq'*, *THEN sb-dec-lem, simplified*])

lemma *sbintrunc-dec*:

$x \geq (2^n) \implies x - 2^{(Suc\ n)} \geq sbintrunc\ n\ x$
unfolding *no-sbintr-alt2* **by** (*drule sb-dec-lem'*) *simp*

lemmas *zmod-uminus'* = *zmod-uminus* [**where** $b=c$, *standard*]

lemmas *zpower-zmod'* = *zpower-zmod* [**where** $m=c$ **and** $y=k$, *standard*]

lemmas *brdmod1s'* [*symmetric*] =

zmod-zadd-left-eq zmod-zadd-right-eq
zmod-zsub-left-eq zmod-zsub-right-eq
zmod-zmult1-eq zmod-zmult1-eq-rev

lemmas *brdmods'* [*symmetric*] =

zpower-zmod' [*symmetric*]
trans [*OF zmod-zadd-left-eq zmod-zadd-right-eq*]
trans [*OF zmod-zsub-left-eq zmod-zsub-right-eq*]
trans [*OF zmod-zmult1-eq zmod-zmult1-eq-rev*]
zmod-uminus' [*symmetric*]
zmod-zadd-left-eq [**where** $b = 1$]
zmod-zsub-left-eq [**where** $b = 1$]

lemmas *bintr-arith1s* =

brdmod1s' [**where** $c=2^n$, *folded pred-def succ-def bintrunc-mod2p, standard*]

lemmas *bintr-ariths* =

brdmods' [**where** $c=2^n$, *folded pred-def succ-def bintrunc-mod2p, standard*]

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [*OF zless2p, standard*]

lemma *bintr-ge0*: $(0 :: int) \leq \text{number-of} (bintrunc\ n\ w)$

by (*simp add : no-bintr m2pths*)

lemma *bintr-lt2p*: $\text{number-of} (bintrunc\ n\ w) < (2^n :: int)$

by (*simp add : no-bintr m2pths*)

lemma *bintr-Min*:

$\text{number-of} (bintrunc\ n\ \text{Numeral.Min}) = (2^n :: int) - 1$

by (*simp add : no-bintr m1mod2k*)

lemma *sbintr-ge*: $(- (2 \wedge n) :: \text{int}) \leq \text{number-of } (\text{sbintrunc } n \ w)$
by (*simp add : no-sbintr m2pths*)

lemma *sbintr-lt*: $\text{number-of } (\text{sbintrunc } n \ w) < (2 \wedge n :: \text{int})$
by (*simp add : no-sbintr m2pths*)

lemma *bintrunc-Suc*:
 $\text{bintrunc } (\text{Suc } n) \ \text{bin} = \text{bintrunc } n \ (\text{bin-rest } \text{bin}) \ \text{BIT } \text{bin-last } \text{bin}$
by (*case-tac bin rule: bin-exhaust*) *auto*

lemma *sign-Pls-ge-0*:
 $(\text{bin-sign } \text{bin} = \text{Numeral.Pls}) = (\text{number-of } \text{bin} \geq (0 :: \text{int}))$
by (*induct bin rule: bin-induct*) *auto*

lemma *sign-Min-lt-0*:
 $(\text{bin-sign } \text{bin} = \text{Numeral.Min}) = (\text{number-of } \text{bin} < (0 :: \text{int}))$
by (*induct bin rule: bin-induct*) *auto*

lemmas *sign-Min-neg* = *trans [OF sign-Min-lt-0 neg-def [symmetric]]*

lemma *bin-rest-trunc*:
 $!!\text{bin. } (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bintrunc } (n - 1) \ (\text{bin-rest } \text{bin})$
by (*induct n*) *auto*

lemma *bin-rest-power-trunc* [*rule-format*] :
 $(\text{bin-rest } \wedge k) \ (\text{bintrunc } n \ \text{bin}) =$
 $\text{bintrunc } (n - k) \ ((\text{bin-rest } \wedge k) \ \text{bin})$
by (*induct k*) (*auto simp: bin-rest-trunc*)

lemma *bin-rest-trunc-i*:
 $\text{bintrunc } n \ (\text{bin-rest } \text{bin}) = \text{bin-rest } (\text{bintrunc } (\text{Suc } n) \ \text{bin})$
by *auto*

lemma *bin-rest-strunc*:
 $!!\text{bin. } \text{bin-rest } (\text{sbintrunc } (\text{Suc } n) \ \text{bin}) = \text{sbintrunc } n \ (\text{bin-rest } \text{bin})$
by (*induct n*) *auto*

lemma *bintrunc-rest* [*simp*]:
 $!!\text{bin. } \text{bintrunc } n \ (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bin-rest } (\text{bintrunc } n \ \text{bin})$
apply (*induct n, simp*)
apply (*case-tac bin rule: bin-exhaust*)
apply (*auto simp: bintrunc-bintrunc-l*)
done

lemma *sbintrunc-rest* [*simp*]:
 $!!\text{bin. } \text{sbintrunc } n \ (\text{bin-rest } (\text{sbintrunc } n \ \text{bin})) = \text{bin-rest } (\text{sbintrunc } n \ \text{bin})$
apply (*induct n, simp*)
apply (*case-tac bin rule: bin-exhaust*)

apply (*auto simp: bintrunc-bintrunc-l split: bit.splits*)
done

lemma *bintrunc-rest'*:
 $\text{bintrunc } n \text{ o bin-rest o bintrunc } n = \text{bin-rest o bintrunc } n$
by (*rule ext*) *auto*

lemma *sbintrunc-rest'*:
 $\text{sbintrunc } n \text{ o bin-rest o sbintrunc } n = \text{bin-rest o sbintrunc } n$
by (*rule ext*) *auto*

lemma *rco-lem*:
 $f \text{ o } g \text{ o } f = g \text{ o } f \implies f \text{ o } (g \text{ o } f) ^ n = g ^ n \text{ o } f$
apply (*rule ext*)
apply (*induct-tac n*)
apply (*simp-all (no-asm)*)
apply (*drule fun-cong*)
apply (*unfold o-def*)
apply (*erule trans*)
apply *simp*
done

lemma *rco-alt*: $(f \text{ o } g) ^ n \text{ o } f = f \text{ o } (g \text{ o } f) ^ n$
apply (*rule ext*)
apply (*induct n*)
apply (*simp-all add: o-def*)
done

lemmas *rco-bintr = bintrunc-rest'*
[*THEN rco-lem [THEN fun-cong], unfolded o-def*]
lemmas *rco-sbintr = sbintrunc-rest'*
[*THEN rco-lem [THEN fun-cong], unfolded o-def*]

4.5 Splitting and concatenation

consts

bin-split :: $\text{nat} \Rightarrow \text{int} \Rightarrow \text{int} * \text{int}$

primrec

Z : $\text{bin-split } 0 \ w = (w, \text{Numeral.Pls})$

Suc : $\text{bin-split } (\text{Suc } n) \ w = (\text{let } (w1, w2) = \text{bin-split } n \ (\text{bin-rest } w) \text{ in } (w1, w2 \ \text{BIT } \text{bin-last } w))$

consts

bin-cat :: $\text{int} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int}$

primrec

Z : $\text{bin-cat } w \ 0 \ v = w$

Suc : $\text{bin-cat } w \ (\text{Suc } n) \ v = \text{bin-cat } w \ n \ (\text{bin-rest } v) \ \text{BIT } \text{bin-last } v$

4.6 Miscellaneous lemmas

lemmas *funpow-minus-simp* =

trans [OF gen-minus [where f = power f] funpow-Suc, standard]

lemmas *funpow-pred-simp [simp]* =

funpow-minus-simp [of number-of bin, simplified nobm1, standard]

lemmas *replicate-minus-simp* =

trans [OF gen-minus [where f = %n. replicate n x] replicate.replicate-Suc, standard]

lemmas *replicate-pred-simp [simp]* =

replicate-minus-simp [of number-of bin, simplified nobm1, standard]

lemmas *power-Suc-no [simp]* = *power-Suc [of number-of a, standard]*

lemmas *power-minus-simp* =

trans [OF gen-minus [where f = power f] power-Suc, standard]

lemmas *power-pred-simp* =

power-minus-simp [of number-of bin, simplified nobm1, standard]

lemmas *power-pred-simp-no [simp]* = *power-pred-simp [where f = number-of f, standard]*

lemma *list-exhaust-size-gt0:*

assumes *y: $\bigwedge a \text{ list. } y = a \# \text{ list} \implies P$*

shows *$0 < \text{length } y \implies P$*

apply *(cases y, simp)*

apply *(rule y)*

apply *fastsimp*

done

lemma *list-exhaust-size-eq0:*

assumes *y: $y = [] \implies P$*

shows *$\text{length } y = 0 \implies P$*

apply *(cases y)*

apply *(rule y, simp)*

apply *simp*

done

lemma *size-Cons-lem-eq:*

y = xa # list ==> size y = Suc k ==> size list = k

by *auto*

lemma *size-Cons-lem-eq-bin:*

y = xa # list ==> size y = number-of (Numeral.succ k) ==>

size list = number-of k

by *(auto simp: pred-def succ-def split add : split-if-asm)*

```

lemmas ls-splits =
  prod.split split-split prod.split-asm split-split-asm split-if-asm

```

```

lemma not-B1-is-B0:  $y \neq \text{bit}.B1 \implies y = \text{bit}.B0$ 
  by (cases y) auto

```

```

lemma B1-ass-B0:
  assumes  $y: y = \text{bit}.B0 \implies y = \text{bit}.B1$ 
  shows  $y = \text{bit}.B1$ 
  apply (rule classical)
  apply (drule not-B1-is-B0)
  apply (erule y)
  done

```

— simplifications for specific word lengths

```

lemmas n2s-ths [THEN eq-reflection] = add-2-eq-Suc add-2-eq-Suc'

```

```

lemmas s2n-ths = n2s-ths [symmetric]

```

end

5 BitSyntax: Syntactic class for bitwise operations

```

theory BitSyntax
imports Main
begin

```

```

class bit = type +
  fixes bitNOT :: 'a  $\Rightarrow$  'a
    and bitAND :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a
    and bitOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a
    and bitXOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a

```

We want the bitwise operations to bind slightly weaker than + and −, but ~~ to bind slightly stronger than *.

notation

```

bitNOT (NOT - [70] 71) and
bitAND (infixr AND 64) and
bitOR (infixr OR 59) and
bitXOR (infixr XOR 59)

```

Testing and shifting operations.

consts

```

test-bit :: 'a::bit  $\Rightarrow$  nat  $\Rightarrow$  bool (infixl !! 100)
lsb      :: 'a::bit  $\Rightarrow$  bool

```

```

msb      :: 'a::bit ⇒ bool
set-bit  :: 'a::bit ⇒ nat ⇒ bool ⇒ 'a
set-bits :: (nat ⇒ bool) ⇒ 'a::bit (binder BITS 10)
shiftr   :: 'a::bit ⇒ nat ⇒ 'a (infixl << 55)
shiftr   :: 'a::bit ⇒ nat ⇒ 'a (infixl >> 55)

```

5.1 Bitwise operations on type *bit*

instance *bit* :: *bit*

```

NOT-bit-def: NOT x ≡ case x of bit.B0 ⇒ bit.B1 | bit.B1 ⇒ bit.B0
AND-bit-def: x AND y ≡ case x of bit.B0 ⇒ bit.B0 | bit.B1 ⇒ y
OR-bit-def:  x OR y  :: bit ≡ NOT (NOT x AND NOT y)
XOR-bit-def: x XOR y :: bit ≡ (x AND NOT y) OR (NOT x AND y)
..

```

lemma *bit-simps* [*simp*]:

```

NOT bit.B0 = bit.B1
NOT bit.B1 = bit.B0
bit.B0 AND y = bit.B0
bit.B1 AND y = y
bit.B0 OR y = y
bit.B1 OR y = bit.B1
bit.B0 XOR y = y
bit.B1 XOR y = NOT y
by (simp-all add: NOT-bit-def AND-bit-def OR-bit-def
      XOR-bit-def split: bit.split)

```

lemma *bit-extra-simps* [*simp*]:

```

x AND bit.B0 = bit.B0
x AND bit.B1 = x
x OR bit.B1 = bit.B1
x OR bit.B0 = x
x XOR bit.B1 = NOT x
x XOR bit.B0 = x
by (cases x, auto)+

```

lemma *bit-ops-comm*:

```

(x::bit) AND y = y AND x
(x::bit) OR y = y OR x
(x::bit) XOR y = y XOR x
by (cases y, auto)+

```

lemma *bit-ops-same* [*simp*]:

```

(x::bit) AND x = x
(x::bit) OR x = x
(x::bit) XOR x = bit.B0
by (cases x, auto)+

```

lemma *bit-not-not* [*simp*]: $NOT (NOT (x::bit)) = x$

```

  by (cases x) auto
end

```

6 BinOperations: Bitwise Operations on Binary Integers

```

theory BinOperations imports BinGeneral BitSyntax

```

```

begin

```

6.1 Logical operations

bit-wise logical operations on the int type

```

instance int :: bit
  int-not-def: bitNOT  $\equiv$  bin-rec Numeral.Min Numeral.Pls
    ( $\lambda w b s. s \text{ BIT } (\text{NOT } b)$ )
  int-and-def: bitAND  $\equiv$  bin-rec ( $\lambda x. \text{Numeral.Pls}$ ) ( $\lambda y. y$ )
    ( $\lambda w b s y. s \text{ (bin-rest } y) \text{ BIT } (b \text{ AND } \text{bin-last } y)$ )
  int-or-def: bitOR  $\equiv$  bin-rec ( $\lambda x. x$ ) ( $\lambda y. \text{Numeral.Min}$ )
    ( $\lambda w b s y. s \text{ (bin-rest } y) \text{ BIT } (b \text{ OR } \text{bin-last } y)$ )
  int-xor-def: bitXOR  $\equiv$  bin-rec ( $\lambda x. x$ ) bitNOT
    ( $\lambda w b s y. s \text{ (bin-rest } y) \text{ BIT } (b \text{ XOR } \text{bin-last } y)$ )
..

```

```

lemma int-not-simps [simp]:
  NOT Numeral.Pls = Numeral.Min
  NOT Numeral.Min = Numeral.Pls
  NOT (w BIT b) = (NOT w) BIT (NOT b)
  by (unfold int-not-def) (auto intro: bin-rec-simps)

```

```

lemma int-xor-Pls [simp]:
  Numeral.Pls XOR x = x
  unfolding int-xor-def by (simp add: bin-rec-PM)

```

```

lemma int-xor-Min [simp]:
  Numeral.Min XOR x = NOT x
  unfolding int-xor-def by (simp add: bin-rec-PM)

```

```

lemma int-xor-Bits [simp]:
  (x BIT b) XOR (y BIT c) = (x XOR y) BIT (b XOR c)
  apply (unfold int-xor-def)
  apply (rule bin-rec-simps (1) [THEN fun-cong, THEN trans])
  apply (rule ext, simp)
  prefer 2
  apply simp
  apply (rule ext)

```

apply (*simp add: int-not-simps [symmetric]*)
done

lemma *int-xor-x-simps'*:
 $w \text{ XOR } (\text{Numeral.Pls BIT bit.B0}) = w$
 $w \text{ XOR } (\text{Numeral.Min BIT bit.B1}) = \text{NOT } w$
apply (*induct w rule: bin-induct*)
apply *simp-all[4]*
apply (*unfold int-xor-Bits*)
apply *clarsimp+*
done

lemmas *int-xor-extra-simps [simp] = int-xor-x-simps' [simplified arith-simps]*

lemma *int-or-Pls [simp]*:
 $\text{Numeral.Pls OR } x = x$
by (*unfold int-or-def*) (*simp add: bin-rec-PM*)

lemma *int-or-Min [simp]*:
 $\text{Numeral.Min OR } x = \text{Numeral.Min}$
by (*unfold int-or-def*) (*simp add: bin-rec-PM*)

lemma *int-or-Bits [simp]*:
 $(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \text{ OR } c)$
unfolding *int-or-def* **by** (*simp add: bin-rec-simps*)

lemma *int-or-x-simps'*:
 $w \text{ OR } (\text{Numeral.Pls BIT bit.B0}) = w$
 $w \text{ OR } (\text{Numeral.Min BIT bit.B1}) = \text{Numeral.Min}$
apply (*induct w rule: bin-induct*)
apply *simp-all[4]*
apply (*unfold int-or-Bits*)
apply *clarsimp+*
done

lemmas *int-or-extra-simps [simp] = int-or-x-simps' [simplified arith-simps]*

lemma *int-and-Pls [simp]*:
 $\text{Numeral.Pls AND } x = \text{Numeral.Pls}$
unfolding *int-and-def* **by** (*simp add: bin-rec-PM*)

lemma *int-and-Min [simp]*:
 $\text{Numeral.Min AND } x = x$
unfolding *int-and-def* **by** (*simp add: bin-rec-PM*)

lemma *int-and-Bits [simp]*:
 $(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \text{ AND } c)$
unfolding *int-and-def* **by** (*simp add: bin-rec-simps*)

lemma *int-and-x-simps'*:

```

w AND (Numeral.Pls BIT bit.B0) = Numeral.Pls
w AND (Numeral.Min BIT bit.B1) = w
apply (induct w rule: bin-induct)
  apply simp-all[4]
  apply (unfold int-and-Bits)
  apply clarsimp+
done

```

lemmas *int-and-extra-simps* [simp] = *int-and-x-simps'* [simplified arith-simps]

lemma *bin-ops-comm*:

```

shows
  int-and-comm: !!y::int. x AND y = y AND x and
  int-or-comm: !!y::int. x OR y = y OR x and
  int-xor-comm: !!y::int. x XOR y = y XOR x
apply (induct x rule: bin-induct)
  apply simp-all[6]
  apply (case-tac y rule: bin-exhaust, simp add: bit-ops-comm)+
done

```

lemma *bin-ops-same* [simp]:

```

(x::int) AND x = x
(x::int) OR x = x
(x::int) XOR x = Numeral.Pls
by (induct x rule: bin-induct) auto

```

lemma *int-not-not* [simp]: NOT (NOT (x::int)) = x

```

by (induct x rule: bin-induct) auto

```

lemmas *bin-log-esimps* =

```

int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min

```

lemma *bbw-ao-absorb*:

```

!!y::int. x AND (y OR x) = x & x OR (y AND x) = x
apply (induct x rule: bin-induct)
  apply auto
  apply (case-tac [!] y rule: bin-exhaust)
  apply auto
  apply (case-tac [!] bit)
  apply auto
done

```

lemma *bbw-ao-absorbs-other*:

```

x AND (x OR y) = x ∧ (y AND x) OR x = (x::int)
(y OR x) AND x = x ∧ x OR (x AND y) = (x::int)
(x OR y) AND x = x ∧ (x AND y) OR x = (x::int)
apply (auto simp: bbw-ao-absorb int-or-comm)
  apply (subst int-or-comm)
  apply (simp add: bbw-ao-absorb)
  apply (subst int-and-comm)
  apply (subst int-or-comm)
  apply (simp add: bbw-ao-absorb)
  apply (subst int-and-comm)
  apply (simp add: bbw-ao-absorb)
done

```

lemmas bbw-ao-absorbs [simp] = bbw-ao-absorb bbw-ao-absorbs-other

lemma int-xor-not:

```

!!y::int. (NOT x) XOR y = NOT (x XOR y) &
  x XOR (NOT y) = NOT (x XOR y)
apply (induct x rule: bin-induct)
  apply auto
  apply (case-tac y rule: bin-exhaust, auto,
    case-tac b, auto)+
done

```

lemma bbw-assocs':

```

!!y z::int. (x AND y) AND z = x AND (y AND z) &
  (x OR y) OR z = x OR (y OR z) &
  (x XOR y) XOR z = x XOR (y XOR z)
apply (induct x rule: bin-induct)
  apply (auto simp: int-xor-not)
  apply (case-tac [!] y rule: bin-exhaust)
  apply (case-tac [!] z rule: bin-exhaust)
  apply (case-tac [!] bit)
  apply (case-tac [!] b)
  apply auto
done

```

lemma int-and-assoc:

```

(x AND y) AND (z::int) = x AND (y AND z)
by (simp add: bbw-assocs')
```

lemma int-or-assoc:

```

(x OR y) OR (z::int) = x OR (y OR z)
by (simp add: bbw-assocs')
```

lemma int-xor-assoc:

```

(x XOR y) XOR (z::int) = x XOR (y XOR z)
by (simp add: bbw-assocs')
```

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [*simp*]:

$(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$

$(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$

$(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$

apply (*auto simp: bbw-assocs [symmetric]*)

apply (*auto simp: bin-ops-comm*)

done

lemma *bbw-not-dist*:

$!!y::int. \text{ NOT } (x \text{ OR } y) = (\text{ NOT } x) \text{ AND } (\text{ NOT } y)$

$!!y::int. \text{ NOT } (x \text{ AND } y) = (\text{ NOT } x) \text{ OR } (\text{ NOT } y)$

apply (*induct x rule: bin-induct*)

apply *auto*

apply (*case-tac [!] y rule: bin-exhaust*)

apply (*case-tac [!] bit, auto*)

done

lemma *bbw-oa-dist*:

$!!y z::int. (x \text{ AND } y) \text{ OR } z =$

$(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$

apply (*induct x rule: bin-induct*)

apply *auto*

apply (*case-tac y rule: bin-exhaust*)

apply (*case-tac z rule: bin-exhaust*)

apply (*case-tac ba, auto*)

done

lemma *bbw-ao-dist*:

$!!y z::int. (x \text{ OR } y) \text{ AND } z =$

$(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$

apply (*induct x rule: bin-induct*)

apply *auto*

apply (*case-tac y rule: bin-exhaust*)

apply (*case-tac z rule: bin-exhaust*)

apply (*case-tac ba, auto*)

done

lemma *plus-and-or* [*rule-format*]:

$ALL y::int. (x \text{ AND } y) + (x \text{ OR } y) = x + y$

apply (*induct x rule: bin-induct*)

apply *clarsimp*

apply *clarsimp*

apply *clarsimp*

apply (*case-tac y rule: bin-exhaust*)

apply *clarsimp*

```

apply (unfold Bit-def)
apply clarsimp
apply (erule-tac  $x = x$  in allE)
apply (simp split: bit.split)
done

```

```

lemma le-int-or:
!!x. bin-sign y = Numeral.Pls ==> x <= x OR y
apply (induct y rule: bin-induct)
  apply clarsimp
  apply clarsimp
apply (case-tac x rule: bin-exhaust)
apply (case-tac b)
  apply (case-tac [!] bit)
  apply (auto simp: less-eq-numeral-code)
done

```

```

lemmas int-and-le =
  xtr3 [OF bbw-ao-absorbs (2) [THEN conjunct2, symmetric] le-int-or]

```

```

lemma bin-nth-ops:
!!x y. bin-nth (x AND y) n = (bin-nth x n & bin-nth y n)
!!x y. bin-nth (x OR y) n = (bin-nth x n | bin-nth y n)
!!x y. bin-nth (x XOR y) n = (bin-nth x n ~ bin-nth y n)
!!x. bin-nth (NOT x) n = (~ bin-nth x n)
apply (induct n)
  apply safe
    apply (case-tac [!] x rule: bin-exhaust)
    apply simp-all
    apply (case-tac [!] y rule: bin-exhaust)
    apply simp-all
  apply (auto dest: not-B1-is-B0 intro: B1-ass-B0)
done

```

```

lemma bin-add-not:  $x + NOT x = Numeral.Min$ 
apply (induct x rule: bin-induct)
  apply clarsimp
  apply clarsimp
apply (case-tac bit, auto)
done

```

```

lemma bin-trunc-ao:
!!x y. (bintrunc n x) AND (bintrunc n y) = bintrunc n (x AND y)
!!x y. (bintrunc n x) OR (bintrunc n y) = bintrunc n (x OR y)
apply (induct n)
  apply auto

```

```

apply (case-tac [!] x rule: bin-exhaust)
apply (case-tac [!] y rule: bin-exhaust)
apply auto
done

```

lemma *bin-trunc-xor*:

```

!!x y. bintrunc n (bintrunc n x XOR bintrunc n y) =
      bintrunc n (x XOR y)
apply (induct n)
apply auto
apply (case-tac [!] x rule: bin-exhaust)
apply (case-tac [!] y rule: bin-exhaust)
apply auto
done

```

lemma *bin-trunc-not*:

```

!!x. bintrunc n (NOT (bintrunc n x)) = bintrunc n (NOT x)
apply (induct n)
apply auto
apply (case-tac [!] x rule: bin-exhaust)
apply auto
done

```

lemma *bintr-bintr-i*:

```

x = bintrunc n y ==> bintrunc n x = bintrunc n y
by auto

```

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

6.2 Setting and clearing bits

consts

```

bin-sc :: nat => bit => int => int

```

primrec

```

Z : bin-sc 0 b w = bin-rest w BIT b

```

```

Suc :

```

```

  bin-sc (Suc n) b w = bin-sc n b (bin-rest w) BIT bin-last w

```

lemma *bin-nth-sc* [simp]:

```

!!w. bin-nth (bin-sc n b w) n = (b = bit.B1)
by (induct n) auto

```

lemma *bin-sc-sc-same* [simp]:

```

!!w. bin-sc n c (bin-sc n b w) = bin-sc n c w

```

by (*induct n*) *auto*

lemma *bin-sc-sc-diff*:

!!*w m. m* \sim *n* \implies

bin-sc m c (bin-sc n b w) = bin-sc n b (bin-sc m c w)

apply (*induct n*)

apply (*case-tac* [!] *m*)

apply *auto*

done

lemma *bin-nth-sc-gen*:

!!*w m. bin-nth (bin-sc n b w) m = (if m = n then b = bit.B1 else bin-nth w m)*

by (*induct n*) (*case-tac* [!] *m, auto*)

lemma *bin-sc-nth [simp]*:

!!*w. (bin-sc n (If (bin-nth w n) bit.B1 bit.B0) w) = w*

by (*induct n*) *auto*

lemma *bin-sign-sc [simp]*:

!!*w. bin-sign (bin-sc n b w) = bin-sign w*

by (*induct n*) *auto*

lemma *bin-sc-bintr [simp]*:

!!*w m. bintrunc m (bin-sc n x (bintrunc m (w))) = bintrunc m (bin-sc n x w)*

apply (*induct n*)

apply (*case-tac* [!] *w rule: bin-exhaust*)

apply (*case-tac* [!] *m, auto*)

done

lemma *bin-clr-le*:

!!*w. bin-sc n bit.B0 w* \leq *w*

apply (*induct n*)

apply (*case-tac* [!] *w rule: bin-exhaust*)

apply *auto*

apply (*unfold Bit-def*)

apply (*simp-all split: bit.split*)

done

lemma *bin-set-ge*:

!!*w. bin-sc n bit.B1 w* \geq *w*

apply (*induct n*)

apply (*case-tac* [!] *w rule: bin-exhaust*)

apply *auto*

apply (*unfold Bit-def*)

apply (*simp-all split: bit.split*)

done

lemma *bintr-bin-clr-le*:

!!*w m. bintrunc n (bin-sc m bit.B0 w) \leq bintrunc n w*

```

apply (induct n)
apply simp
apply (case-tac w rule: bin-exhaust)
apply (case-tac m)
apply auto
apply (unfold Bit-def)
apply (simp-all split: bit.split)
done

```

```

lemma bintr-bin-set-ge:
  !!w m. bintrunc n (bin-sc m bit.B1 w) >= bintrunc n w
apply (induct n)
apply simp
apply (case-tac w rule: bin-exhaust)
apply (case-tac m)
apply auto
apply (unfold Bit-def)
apply (simp-all split: bit.split)
done

```

```

lemma bin-sc-FP [simp]: bin-sc n bit.B0 Numeral.Pls = Numeral.Pls
  by (induct n) auto

```

```

lemma bin-sc-TM [simp]: bin-sc n bit.B1 Numeral.Min = Numeral.Min
  by (induct n) auto

```

```

lemmas bin-sc-simps = bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP

```

```

lemma bin-sc-minus:
  0 < n ==> bin-sc (Suc (n - 1)) b w = bin-sc n b w
  by auto

```

```

lemmas bin-sc-Suc-minus =
  trans [OF bin-sc-minus [symmetric] bin-sc.Suc, standard]

```

```

lemmas bin-sc-Suc-pred [simp] =
  bin-sc-Suc-minus [of number-of bin, simplified nobm1, standard]

```

6.3 Operations on lists of booleans

consts

```

bin-to-bl :: nat => int => bool list
bin-to-bl-aux :: nat => int => bool list => bool list
bl-to-bin :: bool list => int
bl-to-bin-aux :: int => bool list => int

```

```

bl-of-nth :: nat => (nat => bool) => bool list

```

primrec

$Nil : bl-to-bin-aux\ w\ [] = w$
 $Cons : bl-to-bin-aux\ w\ (b\ \# bs) =$
 $bl-to-bin-aux\ (w\ BIT\ (if\ b\ then\ bit.B1\ else\ bit.B0))\ bs$

primrec

$Z : bin-to-bl-aux\ 0\ w\ bl = bl$
 $Suc : bin-to-bl-aux\ (Suc\ n)\ w\ bl =$
 $bin-to-bl-aux\ n\ (bin-rest\ w)\ ((bin-last\ w = bit.B1) \# bl)$

defs

$bin-to-bl-def : bin-to-bl\ n\ w == bin-to-bl-aux\ n\ w\ []$
 $bl-to-bin-def : bl-to-bin\ bs == bl-to-bin-aux\ Numeral.Pls\ bs$

primrec

$Suc : bl-of-nth\ (Suc\ n)\ f = f\ n\ \# bl-of-nth\ n\ f$
 $Z : bl-of-nth\ 0\ f = []$

consts

$takefill :: 'a\ =>\ nat\ =>\ 'a\ list\ =>\ 'a\ list$
 $app2 :: ('a\ =>\ 'b\ =>\ 'c)\ =>\ 'a\ list\ =>\ 'b\ list\ =>\ 'c\ list$

— takefill - like take but if argument list too short,
— extends result to get requested length

primrec

$Z : takefill\ fill\ 0\ xs = []$
 $Suc : takefill\ fill\ (Suc\ n)\ xs = ($
 $case\ xs\ of\ [] => fill\ \# takefill\ fill\ n\ xs$
 $| y\ \# ys => y\ \# takefill\ fill\ n\ ys)$

defs

$app2-def : app2\ f\ as\ bs == map\ (split\ f)\ (zip\ as\ bs)$

6.4 Splitting and concatenation

— rcat and rsplit

consts

$bin-rcat :: nat\ =>\ int\ list\ =>\ int$
 $bin-rsplit-aux :: nat\ * int\ list\ * nat\ * int\ =>\ int\ list$
 $bin-rsplit :: nat\ =>\ (nat\ * int)\ =>\ int\ list$
 $bin-rsplitl-aux :: nat\ * int\ list\ * nat\ * int\ =>\ int\ list$
 $bin-rsplitl :: nat\ =>\ (nat\ * int)\ =>\ int\ list$

recdef bin-rsplit-aux measure (fst o snd o snd)

$bin-rsplit-aux\ (n,\ bs,\ (m,\ c)) =$
 $(if\ m = 0\ | n = 0\ then\ bs\ else$
 $let\ (a,\ b) = bin-split\ n\ c$
 $in\ bin-rsplit-aux\ (n,\ b\ \# bs,\ (m - n,\ a)))$

recdef bin-rsplitl-aux measure (fst o snd o snd)

```

bin-rsplitl-aux (n, bs, (m, c)) =
  (if m = 0 | n = 0 then bs else
   let (a, b) = bin-split (min m n) c
   in bin-rsplitl-aux (n, b # bs, (m - n, a)))

```

defs

```

bin-rcat-def : bin-rcat n bs == foldl (%u v. bin-cat u n v) Numeral.Pls bs
bin-rsplit-def : bin-rsplit n w == bin-rsplit-aux (n, [], w)
bin-rsplitl-def : bin-rsplitl n w == bin-rsplitl-aux (n, [], w)

```

```

declare bin-rsplit-aux.simps [simp del]
declare bin-rsplitl-aux.simps [simp del]

```

lemma *bin-sign-cat*:

```

!!y. bin-sign (bin-cat x n y) = bin-sign x
by (induct n) auto

```

lemma *bin-cat-Suc-Bit*:

```

bin-cat w (Suc n) (v BIT b) = bin-cat w n v BIT b
by auto

```

lemma *bin-nth-cat*:

```

!!n y. bin-nth (bin-cat x k y) n =
  (if n < k then bin-nth y n else bin-nth x (n - k))
apply (induct k)
apply clarsimp
apply (case-tac n, auto)
done

```

lemma *bin-nth-split*:

```

!!b c. bin-split n c = (a, b) ==>
  (ALL k. bin-nth a k = bin-nth c (n + k)) &
  (ALL k. bin-nth b k = (k < n & bin-nth c k))
apply (induct n)
apply clarsimp
apply (clarsimp simp: Let-def split: ls-splits)
apply (case-tac k)
apply auto
done

```

lemma *bin-cat-assoc*:

```

!!z. bin-cat (bin-cat x m y) n z = bin-cat x (m + n) (bin-cat y n z)
by (induct n) auto

```

lemma *bin-cat-assoc-sym*: !!z m.

```

bin-cat x m (bin-cat y n z) = bin-cat (bin-cat x (m - n) y) (min m n) z
apply (induct n, clarsimp)

```

apply (*case-tac m, auto*)
done

lemma *bin-cat-Pls* [*simp*]:
 !!*w. bin-cat Numeral.Pls n w = bintrunc n w*
by (*induct n auto*)

lemma *bintr-cat1*:
 !!*b. bintrunc (k + n) (bin-cat a n b) = bin-cat (bintrunc k a) n b*
by (*induct n auto*)

lemma *bintr-cat*: *bintrunc m (bin-cat a n b) =*
bin-cat (bintrunc (m - n) a) n (bintrunc (min m n) b)
by (*rule bin-eqI*) (*auto simp: bin-nth-cat nth-bintr*)

lemma *bintr-cat-same* [*simp*]:
bintrunc n (bin-cat a n b) = bintrunc n b
by (*auto simp add : bintr-cat*)

lemma *cat-bintr* [*simp*]:
 !!*b. bin-cat a n (bintrunc n b) = bin-cat a n b*
by (*induct n auto*)

lemma *split-bintrunc*:
 !!*b c. bin-split n c = (a, b) ==> b = bintrunc n c*
by (*induct n*) (*auto simp: Let-def split: ls-splits*)

lemma *bin-cat-split*:
 !!*v w. bin-split n w = (u, v) ==> w = bin-cat u n v*
by (*induct n*) (*auto simp: Let-def split: ls-splits*)

lemma *bin-split-cat*:
 !!*w. bin-split n (bin-cat v n w) = (v, bintrunc n w)*
by (*induct n auto*)

lemma *bin-split-Pls* [*simp*]:
bin-split n Numeral.Pls = (Numeral.Pls, Numeral.Pls)
by (*induct n*) (*auto simp: Let-def split: ls-splits*)

lemma *bin-split-Min* [*simp*]:
bin-split n Numeral.Min = (Numeral.Min, bintrunc n Numeral.Min)
by (*induct n*) (*auto simp: Let-def split: ls-splits*)

lemma *bin-split-trunc*:
 !!*m b c. bin-split (min m n) c = (a, b) ==>*
bin-split n (bintrunc m c) = (bintrunc (m - n) a, b)
apply (*induct n, clarsimp*)
apply (*simp add: bin-rest-trunc Let-def split: ls-splits*)
apply (*case-tac m*)

```

apply (auto simp: Let-def split: ls-splits)
done

```

lemma *bin-split-trunc1*:

```

!!m b c. bin-split n c = (a, b) ==>
  bin-split n (bintrunc m c) = (bintrunc (m - n) a, bintrunc m b)
apply (induct n, clarsimp)
apply (simp add: bin-rest-trunc Let-def split: ls-splits)
apply (case-tac m)
apply (auto simp: Let-def split: ls-splits)
done

```

lemma *bin-cat-num*:

```

!!b. bin-cat a n b = a * 2 ^ n + bintrunc n b
apply (induct n, clarsimp)
apply (simp add: Bit-def cong: number-of-False-cong)
done

```

lemma *bin-split-num*:

```

!!b. bin-split n b = (b div 2 ^ n, b mod 2 ^ n)
apply (induct n, clarsimp)
apply (simp add: bin-rest-div zdiv-zmult2-eq)
apply (case-tac b rule: bin-exhaust)
apply simp
apply (simp add: Bit-def zmod-zmult-zmult1 p1mod22k
  split: bit.split
  cong: number-of-False-cong)
done

```

6.5 Miscellaneous lemmas

lemma *nth-2p-bin*:

```

!!m. bin-nth (2 ^ n) m = (m = n)
apply (induct n)
apply clarsimp
apply safe
apply (case-tac m)
apply (auto simp: trans [OF numeral-1-eq-1 [symmetric] number-of-eq])
apply (case-tac m)
apply (auto simp: Bit-B0-2t [symmetric])
done

```

lemma *ex-eq-or*:

```

(EX m. n = Suc m & (m = k | P m)) = (n = Suc k | (EX m. n = Suc m & P
m))
by auto

```

end

7 BinBoolList: Bool lists and integers

theory *BinBoolList* imports *BinOperations* begin

7.1 Arithmetic in terms of bool lists

consts

rbl-succ :: *bool list* => *bool list*
rbl-pred :: *bool list* => *bool list*
rbl-add :: *bool list* => *bool list* => *bool list*
rbl-mult :: *bool list* => *bool list* => *bool list*

primrec

Nil: *rbl-succ Nil* = *Nil*
Cons: *rbl-succ* (*x # xs*) = (if *x* then *False # rbl-succ xs* else *True # xs*)

primrec

Nil : *rbl-pred Nil* = *Nil*
Cons : *rbl-pred* (*x # xs*) = (if *x* then *False # xs* else *True # rbl-pred xs*)

primrec

Nil : *rbl-add Nil x* = *Nil*
Cons : *rbl-add* (*y # ys*) *x* = (let *ws* = *rbl-add ys (tl x)* in
(y ~ = hd x) # (if hd x & y then rbl-succ ws else ws))

primrec

Nil : *rbl-mult Nil x* = *Nil*
Cons : *rbl-mult* (*y # ys*) *x* = (let *ws* = *False # rbl-mult ys x* in
if *y* then *rbl-add ws x* else *ws*)

lemma *tl-take*: *tl (take n l) = take (n - 1) (tl l)*

apply (*cases n, clarsimp*)
apply (*cases l, auto*)
done

lemma *take-butlast* [*rule-format*] :

ALL n. n < length l --> take n (butlast l) = take n l
apply (*induct l, clarsimp*)
apply *clarsimp*
apply (*case-tac n*)
apply *auto*
done

lemma *butlast-take* [*rule-format*] :

ALL n. n <= length l --> butlast (take n l) = take (n - 1) l

```

apply (induct l, clarsimp)
apply clarsimp
apply (case-tac n)
apply safe
apply simp-all
apply (case-tac nat)
apply auto
done

```

lemma *butlast-drop* [rule-format] :
 ALL n. *butlast* (drop n l) = drop n (*butlast* l)
apply (induct l)
apply clarsimp
apply clarsimp
apply safe
apply ((case-tac n, auto)[1])+
done

lemma *butlast-power*:
 (*butlast* ^ n) bl = take (length bl - n) bl
by (induct n) (auto simp: *butlast-take*)

lemma *bin-to-bl-aux-Pls-minus-simp*:
 0 < n ==> *bin-to-bl-aux* n Numeral.Pls bl =
bin-to-bl-aux (n - 1) Numeral.Pls (False # bl)
by (cases n) auto

lemma *bin-to-bl-aux-Min-minus-simp*:
 0 < n ==> *bin-to-bl-aux* n Numeral.Min bl =
bin-to-bl-aux (n - 1) Numeral.Min (True # bl)
by (cases n) auto

lemma *bin-to-bl-aux-Bit-minus-simp*:
 0 < n ==> *bin-to-bl-aux* n (w BIT b) bl =
bin-to-bl-aux (n - 1) w ((b = bit.B1) # bl)
by (cases n) auto

```

declare bin-to-bl-aux-Pls-minus-simp [simp]
bin-to-bl-aux-Min-minus-simp [simp]
bin-to-bl-aux-Bit-minus-simp [simp]

```

lemma *bl-to-bin-aux-append* [rule-format] :
 ALL w. *bl-to-bin-aux* w (bs @ cs) = *bl-to-bin-aux* (*bl-to-bin-aux* w bs) cs
by (induct bs) auto

lemma *bin-to-bl-aux-append* [rule-format] :
 ALL w bs. *bin-to-bl-aux* n w bs @ cs = *bin-to-bl-aux* n w (bs @ cs)

by (*induct n*) *auto*

lemma *bl-to-bin-append*:

bl-to-bin (bs @ cs) = bl-to-bin-aux (bl-to-bin bs) cs

unfolding *bl-to-bin-def* **by** (*rule bl-to-bin-aux-append*)

lemma *bin-to-bl-aux-alt*:

bin-to-bl-aux n w bs = bin-to-bl n w @ bs

unfolding *bin-to-bl-def* **by** (*simp add : bin-to-bl-aux-append*)

lemma *bin-to-bl-0*: *bin-to-bl 0 bs = []*

unfolding *bin-to-bl-def* **by** *auto*

lemma *size-bin-to-bl-aux* [*rule-format*] :

ALL w bs. size (bin-to-bl-aux n w bs) = n + length bs

by (*induct n*) *auto*

lemma *size-bin-to-bl*: *size (bin-to-bl n w) = n*

unfolding *bin-to-bl-def* **by** (*simp add : size-bin-to-bl-aux*)

lemma *bin-bl-bin'* [*rule-format*] :

ALL w bs. bl-to-bin (bin-to-bl-aux n w bs) =

bl-to-bin-aux (bintrunc n w) bs

by (*induct n*) (*auto simp add : bl-to-bin-def*)

lemma *bin-bl-bin*: *bl-to-bin (bin-to-bl n w) = bintrunc n w*

unfolding *bin-to-bl-def bin-bl-bin'* **by** *auto*

lemma *bl-bin-bl'* [*rule-format*] :

ALL w n. bin-to-bl (n + length bs) (bl-to-bin-aux w bs) =

bin-to-bl-aux n w bs

apply (*induct bs*)

apply *auto*

apply (*simp-all only : add-Suc [symmetric]*)

apply (*auto simp add : bin-to-bl-def*)

done

lemma *bl-bin-bl*: *bin-to-bl (length bs) (bl-to-bin bs) = bs*

unfolding *bl-to-bin-def*

apply (*rule box-equals*)

apply (*rule bl-bin-bl'*)

prefer 2

apply (*rule bin-to-bl-aux.Z*)

apply *simp*

done

declare

bin-to-bl-0 [*simp*]

size-bin-to-bl [*simp*]

bin-bl-bin [simp]
bl-bin-bl [simp]

lemma *bl-to-bin-inj*:

bl-to-bin bs = bl-to-bin cs ==> length bs = length cs ==> bs = cs
apply (*rule-tac box-equals*)
defer
apply (*rule bl-bin-bl*)
apply (*rule bl-bin-bl*)
apply *simp*
done

lemma *bl-to-bin-False*: *bl-to-bin (False # bl) = bl-to-bin bl*
unfolding *bl-to-bin-def* **by** *auto*

lemma *bl-to-bin-Nil*: *bl-to-bin [] = Numeral.Pls*
unfolding *bl-to-bin-def* **by** *auto*

lemma *bin-to-bl-Pls-aux* [rule-format] :
ALL bl. bin-to-bl-aux n Numeral.Pls bl = replicate n False @ bl
by (*induct n*) (*auto simp: replicate-app-Cons-same*)

lemma *bin-to-bl-Pls*: *bin-to-bl n Numeral.Pls = replicate n False*
unfolding *bin-to-bl-def* **by** (*simp add : bin-to-bl-Pls-aux*)

lemma *bin-to-bl-Min-aux* [rule-format] :
ALL bl. bin-to-bl-aux n Numeral.Min bl = replicate n True @ bl
by (*induct n*) (*auto simp: replicate-app-Cons-same*)

lemma *bin-to-bl-Min*: *bin-to-bl n Numeral.Min = replicate n True*
unfolding *bin-to-bl-def* **by** (*simp add : bin-to-bl-Min-aux*)

lemma *bl-to-bin-rep-F*:
bl-to-bin (replicate n False @ bl) = bl-to-bin bl
apply (*simp add: bin-to-bl-Pls-aux [symmetric] bin-bl-bin'*)
apply (*simp add: bl-to-bin-def*)
done

lemma *bin-to-bl-trunc*:
n <= m ==> bin-to-bl n (bintrunc m w) = bin-to-bl n w
by (*auto intro: bl-to-bin-inj*)

declare

bin-to-bl-trunc [simp]
bl-to-bin-False [simp]
bl-to-bin-Nil [simp]

lemma *bin-to-bl-aux-bintr* [rule-format] :

ALL m bin bl. bin-to-bl-aux n (bintrunc m bin) bl =

```

    replicate (n - m) False @ bin-to-bl-aux (min n m) bin bl
  apply (induct-tac n)
  apply clarsimp
  apply clarsimp
  apply (case-tac m)
  apply (clarsimp simp: bin-to-bl-Pls-aux)
  apply (erule thin-rl)
  apply (induct-tac n)
  apply auto
done

```

lemmas *bin-to-bl-bintr* =
bin-to-bl-aux-bintr [where *bl* = [], folded *bin-to-bl-def*]

lemma *bl-to-bin-rep-False*: *bl-to-bin* (replicate *n* False) = *Numeral.Pls*
 by (induct *n*) auto

lemma *len-bin-to-bl-aux* [rule-format] :
 ALL *w bs*. length (bin-to-bl-aux *n w bs*) = *n* + length *bs*
 by (induct *n*) auto

lemma *len-bin-to-bl* [simp]: length (bin-to-bl *n w*) = *n*
 unfolding *bin-to-bl-def len-bin-to-bl-aux* by auto

lemma *sign-bl-bin'* [rule-format] :
 ALL *w*. bin-sign (bl-to-bin-aux *w bs*) = bin-sign *w*
 by (induct *bs*) auto

lemma *sign-bl-bin*: bin-sign (bl-to-bin *bs*) = *Numeral.Pls*
 unfolding *bl-to-bin-def* by (simp add : *sign-bl-bin'*)

lemma *bl-sbin-sign-aux* [rule-format] :
 ALL *w bs*. hd (bin-to-bl-aux (Suc *n*) *w bs*) =
 (bin-sign (sbintrunc *n w*) = *Numeral.Min*)
 apply (induct *n*)
 apply clarsimp
 apply (case-tac *w* rule: bin-exhaust)
 apply (simp split add : bit.split)
 apply clarsimp
 done

lemma *bl-sbin-sign*:
 hd (bin-to-bl (Suc *n*) *w*) = (bin-sign (sbintrunc *n w*) = *Numeral.Min*)
 unfolding *bin-to-bl-def* by (rule *bl-sbin-sign-aux*)

lemma *bin-nth-of-bl-aux* [rule-format] :
 ALL *w*. bin-nth (bl-to-bin-aux *w bl*) *n* =
 (*n* < size *bl* & rev *bl* ! *n* | *n* >= length *bl* & bin-nth *w* (*n* - size *bl*))
 apply (induct-tac *bl*)

```

apply clarsimp
apply clarsimp
apply (cut-tac  $x=n$  and  $y=size\ list$  in linorder-less-linear)
apply (erule disjE, simp add: nth-append)+
apply (simp add: nth-append)
done

```

lemma *bin-nth-of-bl*: $bin-nth\ (bl-to-bin\ bl)\ n = (n < length\ bl \ \&\ rev\ bl\ !\ n)$
unfolding *bl-to-bin-def* **by** (*simp* *add : bin-nth-of-bl-aux*)

```

lemma bin-nth-bl [rule-format] :  $ALL\ m\ w.\ n < m \ \longrightarrow$   

 $bin-nth\ w\ n = nth\ (rev\ (bin-to-bl\ m\ w))\ n$   

apply (induct  $n$ )  

apply clarsimp  

apply (case-tac  $m$ , clarsimp)  

apply (clarsimp simp: bin-to-bl-def)  

apply (simp add: bin-to-bl-aux-alt)  

apply clarsimp  

apply (case-tac  $m$ , clarsimp)  

apply (clarsimp simp: bin-to-bl-def)  

apply (simp add: bin-to-bl-aux-alt)  

done

```

```

lemma nth-rev [rule-format] :  

 $n < length\ xs \ \longrightarrow rev\ xs\ !\ n = xs\ !\ (length\ xs - 1 - n)$   

apply (induct-tac  $xs$ )  

apply simp  

apply (clarsimp simp add : nth-append nth.simps split add : nat.split)  

apply (rule-tac  $f = \%n.\ list\ !\ n$  in arg-cong)  

apply arith  

done

```

lemmas *nth-rev-alt = nth-rev* [**where** $xs = rev\ ys$, *simplified*, *standard*]

```

lemma nth-bin-to-bl-aux [rule-format] :  

 $ALL\ w\ n\ bl.\ n < m + length\ bl \ \longrightarrow (bin-to-bl-aux\ m\ w\ bl)\ !\ n =$   

 $(if\ n < m\ then\ bin-nth\ w\ (m - 1 - n)\ else\ bl\ !\ (n - m))$   

apply (induct-tac  $m$ )  

apply clarsimp  

apply clarsimp  

apply (case-tac  $w$  rule: bin-exhaust)  

apply clarsimp  

apply (case-tac  $na - n$ )  

apply arith  

apply simp  

apply (rule-tac  $f = \%n.\ bl\ !\ n$  in arg-cong)  

apply arith  

done

```

lemma *nth-bin-to-bl*: $n < m \implies (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$
unfolding *bin-to-bl-def* **by** (*simp add : nth-bin-to-bl-aux*)

lemma *bl-to-bin-lt2p-aux* [*rule-format*] :
*ALL w. bl-to-bin-aux w bs < (w + 1) * (2 ^ length bs)*
apply (*induct bs*)
apply *clarsimp*
apply *clarsimp*
apply *safe*
apply (*erule alle, erule xtr8 [rotated]*),
simp add: Bit-def ring-simps cong add : number-of-False-cong)
done

lemma *bl-to-bin-lt2p*: *bl-to-bin bs < (2 ^ length bs)*
apply (*unfold bl-to-bin-def*)
apply (*rule xtr1*)
prefer 2
apply (*rule bl-to-bin-lt2p-aux*)
apply *simp*
done

lemma *bl-to-bin-ge2p-aux* [*rule-format*] :
*ALL w. bl-to-bin-aux w bs >= w * (2 ^ length bs)*
apply (*induct bs*)
apply *clarsimp*
apply *clarsimp*
apply *safe*
apply (*erule alle, erule less-eq-less.order-trans [rotated]*),
simp add: Bit-def ring-simps cong add : number-of-False-cong)
done

lemma *bl-to-bin-ge0*: *bl-to-bin bs >= 0*
apply (*unfold bl-to-bin-def*)
apply (*rule xtr4*)
apply (*rule bl-to-bin-ge2p-aux*)
apply *simp*
done

lemma *butlast-rest-bin*:
butlast (bin-to-bl n w) = bin-to-bl (n - 1) (bin-rest w)
apply (*unfold bin-to-bl-def*)
apply (*cases w rule: bin-exhaust*)
apply (*cases n, clarsimp*)
apply *clarsimp*
apply (*auto simp add: bin-to-bl-aux-alt*)
done

lemmas *butlast-bin-rest = butlast-rest-bin*
[*where w=bl-to-bin bl and n=length bl, simplified, standard*]

lemma *butlast-rest-bl2bin-aux* [rule-format] :
 ALL w. bl $\sim = []$ -->
 bl-to-bin-aux w (butlast bl) = bin-rest (bl-to-bin-aux w bl)
 by (induct bl) auto

lemma *butlast-rest-bl2bin*:
 bl-to-bin (butlast bl) = bin-rest (bl-to-bin bl)
 apply (unfold bl-to-bin-def)
 apply (cases bl)
 apply (auto simp add: butlast-rest-bl2bin-aux)
 done

lemma *trunc-bl2bin-aux* [rule-format] :
 ALL w. bintrunc m (bl-to-bin-aux w bl) =
 bl-to-bin-aux (bintrunc (m - length bl) w) (drop (length bl - m) bl)
 apply (induct-tac bl)
 apply clarsimp
 apply clarsimp
 apply safe
 apply (case-tac m - size list)
 apply (simp add : diff-is-0-eq [THEN iffD1, THEN Suc-diff-le])
 apply simp
 apply (rule-tac f = %nat. bl-to-bin-aux (bintrunc nat w BIT bit.B1) list
 in arg-cong)
 apply simp
 apply (case-tac m - size list)
 apply (simp add: diff-is-0-eq [THEN iffD1, THEN Suc-diff-le])
 apply simp
 apply (rule-tac f = %nat. bl-to-bin-aux (bintrunc nat w BIT bit.B0) list
 in arg-cong)
 apply simp
 done

lemma *trunc-bl2bin*:
 bintrunc m (bl-to-bin bl) = bl-to-bin (drop (length bl - m) bl)
 unfolding bl-to-bin-def by (simp add : trunc-bl2bin-aux)

lemmas *trunc-bl2bin-len* [simp] =
 trunc-bl2bin [of length bl bl, simplified, standard]

lemma *bl2bin-drop*:
 bl-to-bin (drop k bl) = bintrunc (length bl - k) (bl-to-bin bl)
 apply (rule trans)
 prefer 2
 apply (rule trunc-bl2bin [symmetric])
 apply (cases k <= length bl)
 apply auto
 done

lemma *nth-rest-power-bin* [rule-format] :
 ALL n . $\text{bin-nth } ((\text{bin-rest } ^k) w) n = \text{bin-nth } w (n + k)$
apply (induct k , clarsimp)
apply clarsimp
apply (simp only: *bin-nth.Suc* [symmetric] *add-Suc*)
done

lemma *take-rest-power-bin*:
 $m \leq n \implies \text{take } m (\text{bin-to-bl } n w) = \text{bin-to-bl } m ((\text{bin-rest } ^{(n - m)}) w)$
apply (rule *nth-equalityI*)
apply simp
apply (clarsimp simp add: *nth-bin-to-bl nth-rest-power-bin*)
done

lemma *hd-butlast*: $\text{size } xs > 1 \implies \text{hd } (\text{butlast } xs) = \text{hd } xs$
by (cases xs) auto

lemma *last-bin-last'* [rule-format] :
 ALL w . $\text{size } xs > 0 \dashrightarrow \text{last } xs = (\text{bin-last } (\text{bl-to-bin-aux } w xs) = \text{bit.B1})$
by (induct xs) auto

lemma *last-bin-last*:
 $\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last } (\text{bl-to-bin } xs) = \text{bit.B1})$
unfolding *bl-to-bin-def* **by** (erule *last-bin-last'*)

lemma *bin-last-last*:
 $\text{bin-last } w = (\text{if } \text{last } (\text{bin-to-bl } (\text{Suc } n) w) \text{ then } \text{bit.B1} \text{ else } \text{bit.B0})$
apply (unfold *bin-to-bl-def*)
apply simp
apply (auto simp add: *bin-to-bl-aux-alt*)
done

lemma *app2-Nil* [simp]: $\text{app2 } f [] ys = []$
unfolding *app2-def* **by** auto

lemma *app2-Cons* [simp]:
 $\text{app2 } f (x \# xs) (y \# ys) = f x y \# \text{app2 } f xs ys$
unfolding *app2-def* **by** auto

lemma *bl-xor-aux-bin* [rule-format] : ALL $v w bs cs$.
 $\text{app2 } (\%x y. x \sim y) (\text{bin-to-bl-aux } n v bs) (\text{bin-to-bl-aux } n w cs) =$
 $\text{bin-to-bl-aux } n (v \text{ XOR } w) (\text{app2 } (\%x y. x \sim y) bs cs)$
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: *bin-exhaust*)

```

apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

```

```

lemma bl-or-aux-bin [rule-format] : ALL v w bs cs.
  app2 (op | ) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v OR w) (app2 (op | ) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

```

```

lemma bl-and-aux-bin [rule-format] : ALL v w bs cs.
  app2 (op & ) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v AND w) (app2 (op & ) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

```

```

lemma bl-not-aux-bin [rule-format] :
  ALL w cs. map Not (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (NOT w) (map Not cs)
apply (induct n)
apply clarsimp
apply clarsimp
apply (case-tac w rule: bin-exhaust)
apply (case-tac b)
apply auto
done

```

```

lemmas bl-not-bin = bl-not-aux-bin
  [where cs = [], unfolded bin-to-bl-def [symmetric] map.simps]

```

```

lemmas bl-and-bin = bl-and-aux-bin [where bs=[] and cs=[],
  unfolded app2-Nil, folded bin-to-bl-def]

```

lemmas $bl\text{-}or\text{-}bin = bl\text{-}or\text{-}aux\text{-}bin$ [where $bs=[]$ and $cs=[]$,
unfolded app2-Nil, folded bin-to-bl-def]

lemmas $bl\text{-}xor\text{-}bin = bl\text{-}xor\text{-}aux\text{-}bin$ [where $bs=[]$ and $cs=[]$,
unfolded app2-Nil, folded bin-to-bl-def]

lemma $drop\text{-}bin2bl\text{-}aux$ [rule-format] :
 ALL m bin bs . $drop\ m\ (bin\text{-}to\text{-}bl\text{-}aux\ n\ bin\ bs) =$
 $bin\text{-}to\text{-}bl\text{-}aux\ (n - m)\ bin\ (drop\ (m - n)\ bs)$
apply (*induct* n , *clarsimp*)
apply *clarsimp*
apply (*case-tac* bin *rule: bin-exhaust*)
apply (*case-tac* $m \leq n$, *simp*)
apply (*case-tac* $m - n$, *simp*)
apply *simp*
apply (*rule-tac* $f = \%nat.$ *drop nat bs in arg-cong*)
apply *simp*
done

lemma $drop\text{-}bin2bl$: $drop\ m\ (bin\text{-}to\text{-}bl\ n\ bin) = bin\text{-}to\text{-}bl\ (n - m)\ bin$
unfolding $bin\text{-}to\text{-}bl\text{-}def$ **by** (*simp add: drop-bin2bl-aux*)

lemma $take\text{-}bin2bl\text{-}lem1$ [rule-format] :
 ALL w bs . $take\ m\ (bin\text{-}to\text{-}bl\text{-}aux\ m\ w\ bs) = bin\text{-}to\text{-}bl\ m\ w$
apply (*induct* m , *clarsimp*)
apply *clarsimp*
apply (*simp add: bin-to-bl-aux-alt*)
apply (*simp add: bin-to-bl-def*)
apply (*simp add: bin-to-bl-aux-alt*)
done

lemma $take\text{-}bin2bl\text{-}lem$ [rule-format] :
 ALL w bs . $take\ m\ (bin\text{-}to\text{-}bl\text{-}aux\ (m + n)\ w\ bs) =$
 $take\ m\ (bin\text{-}to\text{-}bl\ (m + n)\ w)$
apply (*induct* n)
apply *clarify*
apply (*simp-all (no-asm) add: bin-to-bl-def take-bin2bl-lem1*)
apply *simp*
done

lemma $bin\text{-}split\text{-}take$ [rule-format] :
 ALL b c . $bin\text{-}split\ n\ c = (a, b) \text{ --->}$
 $bin\text{-}to\text{-}bl\ m\ a = take\ m\ (bin\text{-}to\text{-}bl\ (m + n)\ c)$
apply (*induct* n)
apply *clarsimp*
apply (*clarsimp simp: Let-def split: ls-splits*)
apply (*simp add: bin-to-bl-def*)
apply (*simp add: take-bin2bl-lem*)
done

lemma *bin-split-take1*:

$k = m + n \implies \text{bin-split } n \ c = (a, b) \implies$
 $\text{bin-to-bl } m \ a = \text{take } m \ (\text{bin-to-bl } k \ c)$
by (*auto elim: bin-split-take*)

lemma *nth-takefill* [*rule-format*] : *ALL m l. m < n -->*

takefill fill n l ! m = (if m < length l then l ! m else fill)
apply (*induct n, clarsimp*)
apply *clarsimp*
apply (*case-tac m*)
apply (*simp split: list.split*)
apply *clarsimp*
apply (*erule allE*)
apply (*erule (1) impE*)
apply (*simp split: list.split*)
done

lemma *takefill-alt* [*rule-format*] :

ALL l. takefill fill n l = take n l @ replicate (n - length l) fill
by (*induct n*) (*auto split: list.split*)

lemma *takefill-replicate* [*simp*]:

takefill fill n (replicate m fill) = replicate n fill
by (*simp add : takefill-alt replicate-add [symmetric]*)

lemma *takefill-le'* [*rule-format*] :

ALL l n. n = m + k --> takefill x m (takefill x n l) = takefill x m l
by (*induct m*) (*auto split: list.split*)

lemma *length-takefill* [*simp*]: *length (takefill fill n l) = n*

by (*simp add : takefill-alt*)

lemma *take-takefill'*:

!!w n. n = k + m ==> take k (takefill fill n w) = takefill fill k w
by (*induct k*) (*auto split add : list.split*)

lemma *drop-takefill*:

!!w. drop k (takefill fill (m + k) w) = takefill fill m (drop k w)
by (*induct k*) (*auto split add : list.split*)

lemma *takefill-le* [*simp*]:

$m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
by (*auto simp: le-iff-add takefill-le'*)

lemma *take-takefill* [*simp*]:

$m \leq n \implies \text{take } m \ (\text{takefill } fill \ n \ w) = \text{takefill } fill \ m \ w$
by (*auto simp: le-iff-add take-takefill'*)

lemma *takefill-append*:

takefill fill (m + length xs) (xs @ w) = xs @ (takefill fill m w)
by (*induct xs*) *auto*

lemma *takefill-same'*:

l = length xs ==> takefill fill l xs = xs
by *clarify (induct xs, auto)*

lemmas *takefill-same [simp] = takefill-same' [OF refl]*

lemma *takefill-bintrunc*:

takefill False n bl = rev (bin-to-bl n (bl-to-bin (rev bl)))
apply (*rule nth-equalityI*)
apply *simp*
apply (*clarsimp simp: nth-takefill nth-rev nth-bin-to-bl bin-nth-of-bl*)
done

lemma *bl-bin-bl-rtf*:

bin-to-bl n (bl-to-bin bl) = rev (takefill False n (rev bl))
by (*simp add : takefill-bintrunc*)

lemmas *bl-bin-bl-rep-drop =*

bl-bin-bl-rtf [simplified takefill-alt,
simplified, simplified rev-take, simplified]

lemma *tf-rev*:

n + k = m + length bl ==> takefill x m (rev (takefill y n bl)) =
rev (takefill y m (rev (takefill x k (rev bl))))
apply (*rule nth-equalityI*)
apply (*auto simp add: nth-takefill nth-rev*)
apply (*rule-tac f = %n. bl ! n in arg-cong*)
apply *arith*
done

lemma *takefill-minus*:

0 < n ==> takefill fill (Suc (n - 1)) w = takefill fill n w
by *auto*

lemmas *takefill-Suc-cases =*

list.cases [THEN takefill.Suc [THEN trans], standard]

lemmas *takefill-Suc-Nil = takefill-Suc-cases (1)*

lemmas *takefill-Suc-Cons = takefill-Suc-cases (2)*

lemmas *takefill-minus-simps = takefill-Suc-cases [THEN [2]*

takefill-minus [symmetric, THEN trans], standard]

lemmas *takefill-pred-simps [simp] =*

takefill-minus-simps [where n=number-of bin, simplified nobm1, standard]

lemma *bl-to-bin-aux-cat*:

```
!!nv v. bl-to-bin-aux (bin-cat w nv v) bs =
  bin-cat w (nv + length bs) (bl-to-bin-aux v bs)
apply (induct bs)
apply simp
apply (simp add: bin-cat-Suc-Bit [symmetric] del: bin-cat.simps)
done
```

lemma *bin-to-bl-aux-cat*:

```
!!w bs. bin-to-bl-aux (nv + nw) (bin-cat v nw w) bs =
  bin-to-bl-aux nv v (bin-to-bl-aux nw w bs)
by (induct nw) auto
```

lemmas *bl-to-bin-aux-alt* =

```
bl-to-bin-aux-cat [where nv = 0 and v = Numeral.Pls,
  simplified bl-to-bin-def [symmetric], simplified]
```

lemmas *bin-to-bl-cat* =

```
bin-to-bl-aux-cat [where bs = [], folded bin-to-bl-def]
```

lemmas *bl-to-bin-aux-app-cat* =

```
trans [OF bl-to-bin-aux-append bl-to-bin-aux-alt]
```

lemmas *bin-to-bl-aux-cat-app* =

```
trans [OF bin-to-bl-aux-cat bin-to-bl-aux-alt]
```

lemmas *bl-to-bin-app-cat* = *bl-to-bin-aux-app-cat*

```
[where w = Numeral.Pls, folded bl-to-bin-def]
```

lemmas *bin-to-bl-cat-app* = *bin-to-bl-aux-cat-app*

```
[where bs = [], folded bin-to-bl-def]
```

lemma *bl-to-bin-app-cat-alt*:

```
bin-cat (bl-to-bin cs) n w = bl-to-bin (cs @ bin-to-bl n w)
by (simp add : bl-to-bin-app-cat)
```

lemma *mask-lem*: (*bl-to-bin* (True # replicate n False)) =

```
Numeral.succ (bl-to-bin (replicate n True))
apply (unfold bl-to-bin-def)
apply (induct n)
apply simp
apply (simp only: Suc-eq-add-numeral-1 replicate-add
  append-Cons [symmetric] bl-to-bin-aux-append)
apply simp
done
```

lemma *length-bl-of-nth* [*simp*]: $\text{length } (\text{bl-of-nth } n \ f) = n$
by (*induct n*) *auto*

lemma *nth-bl-of-nth* [*simp*]:
 $m < n \implies \text{rev } (\text{bl-of-nth } n \ f) ! m = f \ m$
apply (*induct n*)
apply *simp*
apply (*clarsimp simp add : nth-append*)
apply (*rule-tac f = f in arg-cong*)
apply *simp*
done

lemma *bl-of-nth-inj*:
 $(!!k. k < n \implies f \ k = g \ k) \implies \text{bl-of-nth } n \ f = \text{bl-of-nth } n \ g$
by (*induct n*) *auto*

lemma *bl-of-nth-nth-le* [*rule-format*] : *ALL xs.*
 $\text{length } xs \geq n \dashrightarrow \text{bl-of-nth } n \ (\text{nth } (\text{rev } xs)) = \text{drop } (\text{length } xs - n) \ xs$
apply (*induct n, clarsimp*)
apply *clarsimp*
apply (*rule trans [OF - hd-Cons-tl]*)
apply (*frule Suc-le-lessD*)
apply (*simp add: nth-rev trans [OF drop-Suc drop-tl, symmetric]*)
apply (*subst hd-drop-conv-nth*)
apply *force*
apply *simp-all*
apply (*rule-tac f = %n. drop n xs in arg-cong*)
apply *simp*
done

lemmas *bl-of-nth-nth* [*simp*] = *order-refl [THEN bl-of-nth-nth-le, simplified]*

lemma *size-rbl-pred*: $\text{length } (\text{rbl-pred } bl) = \text{length } bl$
by (*induct bl*) *auto*

lemma *size-rbl-succ*: $\text{length } (\text{rbl-succ } bl) = \text{length } bl$
by (*induct bl*) *auto*

lemma *size-rbl-add*:
 $!!cl. \text{length } (\text{rbl-add } bl \ cl) = \text{length } bl$
by (*induct bl*) (*auto simp: Let-def size-rbl-succ*)

lemma *size-rbl-mult*:
 $!!cl. \text{length } (\text{rbl-mult } bl \ cl) = \text{length } bl$
by (*induct bl*) (*auto simp add : Let-def size-rbl-add*)

lemmas *rbl-sizes* [*simp*] =

size-rbl-pred size-rbl-succ size-rbl-add size-rbl-mult

lemmas *rbl-Nils* =
rbl-pred.Nil rbl-succ.Nil rbl-add.Nil rbl-mult.Nil

lemma *rbl-pred*:
 !!*bin*. *rbl-pred* (*rev* (*bin-to-bl* *n bin*)) = *rev* (*bin-to-bl* *n* (*Numeral.pred bin*))
apply (*induct n, simp*)
apply (*unfold bin-to-bl-def*)
apply *clarsimp*
apply (*case-tac bin rule: bin-exhaust*)
apply (*case-tac b*)
apply (*clarsimp simp: bin-to-bl-aux-alt*) +
done

lemma *rbl-succ*:
 !!*bin*. *rbl-succ* (*rev* (*bin-to-bl* *n bin*)) = *rev* (*bin-to-bl* *n* (*Numeral.succ bin*))
apply (*induct n, simp*)
apply (*unfold bin-to-bl-def*)
apply *clarsimp*
apply (*case-tac bin rule: bin-exhaust*)
apply (*case-tac b*)
apply (*clarsimp simp: bin-to-bl-aux-alt*) +
done

lemma *rbl-add*:
 !!*bina binb*. *rbl-add* (*rev* (*bin-to-bl* *n bina*)) (*rev* (*bin-to-bl* *n binb*)) =
rev (*bin-to-bl* *n* (*bina + binb*))
apply (*induct n, simp*)
apply (*unfold bin-to-bl-def*)
apply *clarsimp*
apply (*case-tac bina rule: bin-exhaust*)
apply (*case-tac binb rule: bin-exhaust*)
apply (*case-tac b*)
apply (*case-tac [!] ba*)
apply (*auto simp: rbl-succ succ-def bin-to-bl-aux-alt Let-def add-ac*)
done

lemma *rbl-add-app2*:
 !!*blb*. *length blb* >= *length bla* ==>
rbl-add bla (*blb @ blc*) = *rbl-add bla blb*
apply (*induct bla, simp*)
apply *clarsimp*
apply (*case-tac blb, clarsimp*)
apply (*clarsimp simp: Let-def*)
done

lemma *rbl-add-take2*:
 !!*blb*. *length blb* >= *length bla* ==>

```

  rbl-add bla (take (length bla) bbb) = rbl-add bla bbb
apply (induct bla, simp)
apply clarsimp
apply (case-tac bbb, clarsimp)
apply (clarsimp simp: Let-def)
done

```

lemma *rbl-add-long*:

```

  m >= n ==> rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
    rev (bin-to-bl n (bina + binb))
apply (rule box-equals [OF - rbl-add-take2 rbl-add])
apply (rule-tac f = rbl-add (rev (bin-to-bl n bina)) in arg-cong)
apply (rule rev-swap [THEN iffD1])
apply (simp add: rev-take drop-bin2bl)
apply simp
done

```

lemma *rbl-mult-app2*:

```

  !!bbb. length bbb >= length bla ==>
    rbl-mult bla (bbb @ bbb) = rbl-mult bla bbb
apply (induct bla, simp)
apply clarsimp
apply (case-tac bbb, clarsimp)
apply (clarsimp simp: Let-def rbl-add-app2)
done

```

lemma *rbl-mult-take2*:

```

  length bbb >= length bla ==>
    rbl-mult bla (take (length bla) bbb) = rbl-mult bla bbb
apply (rule trans)
apply (rule rbl-mult-app2 [symmetric])
apply simp
apply (rule-tac f = rbl-mult bla in arg-cong)
apply (rule append-take-drop-id)
done

```

lemma *rbl-mult-gt1*:

```

  m >= length bl ==> rbl-mult bl (rev (bin-to-bl m binb)) =
    rbl-mult bl (rev (bin-to-bl (length bl) binb))
apply (rule trans)
apply (rule rbl-mult-take2 [symmetric])
apply simp-all
apply (rule-tac f = rbl-mult bl in arg-cong)
apply (rule rev-swap [THEN iffD1])
apply (simp add: rev-take drop-bin2bl)
done

```

lemma *rbl-mult-gt*:

```

  m > n ==> rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =

```

```

  rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb))
  by (auto intro: trans [OF rbl-mult-gt1])

```

lemmas *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

lemma *rbbl-Cons*:

```

  b # rev (bin-to-bl n x) = rev (bin-to-bl (Suc n) (x BIT If b bit.B1 bit.B0))
  apply (unfold bin-to-bl-def)
  apply simp
  apply (simp add: bin-to-bl-aux-alt)
  done

```

lemma *rbl-mult*: $!!bina\ binb.$

```

  rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =
  rev (bin-to-bl n (bina * binb))
  apply (induct n)
  apply simp
  apply (unfold bin-to-bl-def)
  apply clarsimp
  apply (case-tac bina rule: bin-exhaust)
  apply (case-tac binb rule: bin-exhaust)
  apply (case-tac b)
  apply (case-tac [!] ba)
  apply (auto simp: bin-to-bl-aux-alt Let-def)
  apply (auto simp: rbbl-Cons rbl-mult-Suc rbl-add)
  done

```

lemma *rbl-add-split*:

```

  P (rbl-add (y # ys) (x # xs)) =
  (ALL ws. length ws = length ys --> ws = rbl-add ys xs -->
  (y --> ((x --> P (False # rbl-succ ws)) & (~ x --> P (True # ws))))
  &
  (~ y --> P (x # ws)))
  apply (auto simp add: Let-def)
  apply (case-tac [!] y)
  apply auto
  done

```

lemma *rbl-mult-split*:

```

  P (rbl-mult (y # ys) xs) =
  (ALL ws. length ws = Suc (length ys) --> ws = False # rbl-mult ys xs -->
  (y --> P (rbl-add ws xs)) & (~ y --> P ws))
  by (clarsimp simp add : Let-def)

```

lemma *and-len*: $xs = ys ==> xs = ys \ \& \ length\ xs = length\ ys$

by *auto*

lemma *size-if*: $size\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ size\ xs\ else\ size\ ys)$

by *auto*

lemma *tl-if*: $tl (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ tl\ xs\ else\ tl\ ys)$
by *auto*

lemma *hd-if*: $hd (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ hd\ xs\ else\ hd\ ys)$
by *auto*

lemma *if-Not-x*: $(if\ p\ then\ \sim\ x\ else\ x) = (p = (\sim\ x))$
by *auto*

lemma *if-x-Not*: $(if\ p\ then\ x\ else\ \sim\ x) = (p = x)$
by *auto*

lemma *if-same-and*: $(If\ p\ x\ y\ \&\ If\ p\ u\ v) = (if\ p\ then\ x\ \&\ u\ else\ y\ \&\ v)$
by *auto*

lemma *if-same-eq*: $(If\ p\ x\ y = If\ p\ u\ v) = (if\ p\ then\ x = (u)\ else\ y = (v))$
by *auto*

lemma *if-same-eq-not*:
 $(If\ p\ x\ y = (\sim\ If\ p\ u\ v)) = (if\ p\ then\ x = (\sim\ u)\ else\ y = (\sim\ v))$
by *auto*

lemma *if-Cons*: $(if\ p\ then\ x\ \#\ xs\ else\ y\ \#\ ys) = If\ p\ x\ y\ \#\ If\ p\ xs\ ys$
by *auto*

lemma *if-single*:
 $(if\ xc\ then\ [xab]\ else\ [an]) = [if\ xc\ then\ xab\ else\ an]$
by *auto*

lemma *if-bool-simps*:
 $If\ p\ True\ y = (p \mid y) \ \&\ If\ p\ False\ y = (\sim\ p \ \&\ y) \ \&$
 $If\ p\ y\ True = (p \dashrightarrow y) \ \&\ If\ p\ y\ False = (p \ \&\ y)$
by *auto*

lemmas *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

lemmas *seqr* = *eq-reflection* [where $x = size\ w$, *standard*]

lemmas *tl-Nil* = *tl.simps* (1)
lemmas *tl-Cons* = *tl.simps* (2)

7.2 Repeated splitting or concatenation

lemma *sclem*:
 $size\ (concat\ (map\ (bin\ to\ bl\ n)\ xs)) = length\ xs * n$
by (*induct xs*) *auto*

```

lemma bin-cat-foldl-lem [rule-format] :
  ALL x. foldl (%u. bin-cat u n) x xs =
    bin-cat x (size xs * n) (foldl (%u. bin-cat u n) y xs)
apply (induct xs)
apply simp
apply clarify
apply (simp (no-asm))
apply (frule asm-rl)
apply (drule spec)
apply (erule trans)
apply (drule-tac x = bin-cat y n a in spec)
apply (simp add : bin-cat-assoc-sym min-def)
done

lemma bin-rcat-bl:
  (bin-rcat n wl) = bl-to-bin (concat (map (bin-to-bl n) wl))
apply (unfold bin-rcat-def)
apply (rule sym)
apply (induct wl)
apply (auto simp add : bl-to-bin-append)
apply (simp add : bl-to-bin-aux-alt sclem)
apply (simp add : bin-cat-foldl-lem [symmetric])
done

lemmas bin-rsplit-aux-simps = bin-rsplit-aux.simps bin-rsplitl-aux.simps
lemmas rsplit-aux-simps = bin-rsplit-aux-simps

lemmas th-if-simp1 = split-if [where P = op = l,
  THEN iffD1, THEN conjunct1, THEN mp, standard]
lemmas th-if-simp2 = split-if [where P = op = l,
  THEN iffD1, THEN conjunct2, THEN mp, standard]

lemmas rsplit-aux-simp1s = rsplit-aux-simps [THEN th-if-simp1]

lemmas rsplit-aux-simp2ls = rsplit-aux-simps [THEN th-if-simp2]

lemmas bin-rsplit-aux-simp2s [simp] = rsplit-aux-simp2ls [unfolded Let-def]
lemmas rbscl = bin-rsplit-aux-simp2s (2)

lemmas rsplit-aux-0-simps [simp] =
  rsplit-aux-simp1s [OF disjI1] rsplit-aux-simp1s [OF disjI2]

lemma bin-rsplit-aux-append:
  bin-rsplit-aux (n, bs @ cs, m, c) = bin-rsplit-aux (n, bs, m, c) @ cs
apply (rule-tac u=n and v=bs and w=m and x=c in bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (subst bin-rsplit-aux.simps)
apply (clarsimp split: ls-splits)

```

done

lemma *bin-rsplitl-aux-append*:

bin-rsplitl-aux (n , bs @ cs , m , c) = *bin-rsplitl-aux* (n , bs , m , c) @ cs
apply (*rule-tac* $u=n$ **and** $v=bs$ **and** $w=m$ **and** $x=c$ **in** *bin-rsplitl-aux.induct*)
apply (*subst* *bin-rsplitl-aux.simps*)
apply (*subst* *bin-rsplitl-aux.simps*)
apply (*clarsimp* *split: ls-splits*)
done

lemmas *rsplit-aux-apps* [where $bs = []$] =
bin-rsplit-aux-append *bin-rsplitl-aux-append*

lemmas *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplitl-def*

lemmas *rsplit-aux-alts* = *rsplit-aux-apps*
[*unfolded* *append-Nil* *rsplit-def-auxs* [*symmetric*]]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) w = \text{bin-split } n w$
by *auto*

lemmas *bin-split-minus-simp* =
bin-split.Suc [THEN [2] *bin-split-minus* [*symmetric*, THEN *trans*], *standard*]

lemma *bin-split-pred-simp* [*simp*]:

$(0::\text{nat}) < \text{number-of bin} \implies$
bin-split (*number-of bin*) $w =$
(*let* ($w1$, $w2$) = *bin-split* (*number-of* (*Numeral.pred bin*)) (*bin-rest w*)
in ($w1$, $w2$ BIT *bin-last w*))
by (*simp* *only: nobm1 bin-split-minus-simp*)

declare *bin-split-pred-simp* [*simp*]

lemma *bin-rsplit-aux-simp-alt*:

bin-rsplit-aux (n , bs , m , c) =
(*if* $m = 0 \vee n = 0$
then bs
else *let* (a , b) = *bin-split* n c *in* *bin-rsplit* n ($m - n$, a) @ $b \# bs$)
apply (*rule* *trans*)
apply (*subst* *bin-rsplit-aux.simps*, *rule* *refl*)
apply (*simp* *add: rsplit-aux-alts*)
done

lemmas *bin-rsplit-simp-alt* =
trans [OF *bin-rsplit-def* [THEN *meta-eq-to-obj-eq*]
bin-rsplit-aux-simp-alt, *standard*]

lemmas *bthrs* = *bin-rsplit-simp-alt* [THEN [2] *trans*]

lemma *bin-rsplit-size-sign'* [rule-format] :
 $n > 0 \implies (ALL\ nw\ w.\ rev\ sw = bin-rsplit\ n\ (nw, w) \dashrightarrow$
 $(ALL\ v:\ set\ sw.\ bintrunc\ n\ v = v))$
apply (induct sw)
apply clarsimp
apply clarsimp
apply (drule bthrs)
apply (simp (no-asm-use) add: Let-def split: ls-splits)
apply clarify
apply (erule impE, rule exI, erule exI)
apply (drule split-bintrunc)
apply simp
done

lemmas *bin-rsplit-size-sign = bin-rsplit-size-sign'* [OF asm-rl
rev-rev-ident [THEN trans] set-rev [THEN equalityD2 [THEN subsetD]],
standard]

lemma *bin-nth-rsplit* [rule-format] :
 $n > 0 \implies m < n \implies (ALL\ w\ k\ nw.\ rev\ sw = bin-rsplit\ n\ (nw, w) \dashrightarrow$
 $k < size\ sw \dashrightarrow bin-nth\ (sw\ !\ k)\ m = bin-nth\ w\ (k * n + m))$
apply (induct sw)
apply clarsimp
apply clarsimp
apply (drule bthrs)
apply (simp (no-asm-use) add: Let-def split: ls-splits)
apply clarify
apply (erule allE, erule impE, erule exI)
apply (case-tac k)
apply clarsimp
prefer 2
apply clarsimp
apply (erule allE)
apply (erule (1) impE)
apply (drule bin-nth-split, erule conjE, erule allE,
erule trans, simp add : add-ac)+
done

lemma *bin-rsplit-all*:
 $0 < nw \implies nw \leq n \implies bin-rsplit\ n\ (nw, w) = [bintrunc\ n\ w]$
unfolding bin-rsplit-def
by (clarsimp dest!: split-bintrunc simp: rsplit-aux-simp2ls split: ls-splits)

lemma *bin-rsplit-l* [rule-format] :
 $ALL\ bin.\ bin-rsplitl\ n\ (m, bin) = bin-rsplit\ n\ (m, bintrunc\ m\ bin)$
apply (rule-tac a = m in wf-less-than [THEN wf-induct])
apply (simp (no-asm) add : bin-rsplitl-def bin-rsplit-def)
apply (rule allI)
apply (subst bin-rsplitl-aux.simps)

```

apply (subst bin-rsplit-aux.simps)
apply (clarsimp simp: rsplit-aux-alts Let-def split: ls-splits)
apply (drule bin-split-trunc)
apply (drule sym [THEN trans], assumption)
apply fast
done

```

```

lemma bin-rsplit-rcat [rule-format] :
   $n > 0 \longrightarrow \text{bin-rsplit } n (n * \text{size } ws, \text{bin-rcat } n \text{ } ws) = \text{map } (\text{bintrunc } n) \text{ } ws$ 
apply (unfold bin-rsplit-def bin-rcat-def)
apply (rule-tac  $xs = ws$  in rev-induct)
apply clarsimp
apply clarsimp
apply (clarsimp simp add: bin-split-cat rsplit-aux-alts)
done

```

```

lemma bin-rsplit-aux-len-le [rule-format] :
  ALL  $ws \ m. n \neq 0 \longrightarrow ws = \text{bin-rsplit-aux } (n, bs, nw, w) \longrightarrow$ 
  ( $\text{length } ws \leq m$ ) = ( $nw + \text{length } bs * n \leq m * n$ )
apply (rule-tac  $u=n$  and  $v=bs$  and  $w=nw$  and  $x=w$  in bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (simp add: lrlem Let-def split: ls-splits )
done

```

```

lemma bin-rsplit-len-le:
   $n \neq 0 \longrightarrow ws = \text{bin-rsplit } n (nw, w) \longrightarrow (\text{length } ws \leq m) = (nw \leq m * n)$ 
unfolding bin-rsplit-def by (clarsimp simp add : bin-rsplit-aux-len-le)

```

```

lemma bin-rsplit-aux-len [rule-format] :
   $n \neq 0 \longrightarrow \text{length } (\text{bin-rsplit-aux } (n, cs, nw, w)) =$ 
  ( $nw + n - 1$ ) div  $n + \text{length } cs$ 
apply (rule-tac  $u=n$  and  $v=cs$  and  $w=nw$  and  $x=w$  in bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (clarsimp simp: Let-def split: ls-splits)
apply (erule thin-rl)
apply (case-tac  $m \leq n$ )
prefer 2
apply (simp add: div-add-self2 [symmetric])
apply (case-tac  $m$ , clarsimp)
apply (simp add: div-add-self2)
done

```

```

lemma bin-rsplit-len:
   $n \neq 0 \implies \text{length } (\text{bin-rsplit } n (nw, w)) = (nw + n - 1) \text{ div } n$ 
unfolding bin-rsplit-def by (clarsimp simp add : bin-rsplit-aux-len)

```

```

lemma bin-rsplit-aux-len-indep [rule-format] :
   $n \neq 0 \implies (\text{ALL } v \text{ } bs. \text{length } bs = \text{length } cs \longrightarrow$ 

```

```

    length (bin-rsplit-aux (n, bs, nw, v)) =
      length (bin-rsplit-aux (n, cs, nw, w))
  apply (rule-tac u=n and v=cs and w=nw and x=w in bin-rsplit-aux.induct)
  apply clarsimp
  apply (simp (no-asm-simp) add: bin-rsplit-aux-simp-alt Let-def
            split: ls-splits)

  apply clarify
  apply (erule allE)+
  apply (erule impE)
  apply (fast elim!: sym)
  apply (simp (no-asm-use) add: rsplit-aux-alts)
  apply (erule impE)
  apply (rule-tac x=ba # bs in exI)
  apply auto
  done

lemma bin-rsplit-len-indep:
  n≠0 ==> length (bin-rsplit n (nw, v)) = length (bin-rsplit n (nw, w))
  apply (unfold bin-rsplit-def)
  apply (erule bin-rsplit-aux-len-indep)
  apply (rule refl)
  done

end

```

8 TdThs: Type Definition Theorems

```
theory TdThs imports Main begin
```

9 More lemmas about normal type definitions

```

lemma
  tdD1: type-definition Rep Abs A ==> ∀x. Rep x ∈ A and
  tdD2: type-definition Rep Abs A ==> ∀x. Abs (Rep x) = x and
  tdD3: type-definition Rep Abs A ==> ∀y. y ∈ A → Rep (Abs y) = y
  by (auto simp: type-definition-def)

```

```

lemma td-nat-int:
  type-definition int nat (Collect (op <= 0))
  unfolding type-definition-def by auto

```

```

context type-definition
begin

```

```

lemmas Rep' [iff] = Rep [simplified]

```

declare *Rep-inverse* [simp] *Rep-inject* [simp]

lemma *Abs-eqD*: $Abs\ x = Abs\ y \implies x \in A \implies y \in A \implies x = y$
by (*simp add: Abs-inject*)

lemma *Abs-inverse'*:
 $r : A \implies Abs\ r = a \implies Rep\ a = r$
by (*safe elim!: Abs-inverse*)

lemma *Rep-comp-inverse*:
 $Rep\ o\ f = g \implies Abs\ o\ g = f$
using *Rep-inverse* **by** (*auto intro: ext*)

lemma *Rep-eqD* [*elim!*]: $Rep\ x = Rep\ y \implies x = y$
by *simp*

lemma *Rep-inverse'*: $Rep\ a = r \implies Abs\ r = a$
by (*safe intro!: Rep-inverse*)

lemma *comp-Abs-inverse*:
 $f\ o\ Abs = g \implies g\ o\ Rep = f$
using *Rep-inverse* **by** (*auto intro: ext*)

lemma *set-Rep*:
 $A = range\ Rep$
proof (*rule set-ext*)
fix x
show $(x \in A) = (x \in range\ Rep)$
by (*auto dest: Abs-inverse [of x, symmetric]*)
qed

lemma *set-Rep-Abs*: $A = range\ (Rep\ o\ Abs)$
proof (*rule set-ext*)
fix x
show $(x \in A) = (x \in range\ (Rep\ o\ Abs))$
by (*auto dest: Abs-inverse [of x, symmetric]*)
qed

lemma *Abs-inj-on*: *inj-on* $Abs\ A$
unfolding *inj-on-def*
by (*auto dest: Abs-inject [THEN iffD1]*)

lemma *image*: $Abs\ ` A = UNIV$
by (*auto intro!: image-eqI*)

lemmas *td-thm* = *type-definition-axioms*

lemma *fns1*:
 $Rep\ o\ fa = fr\ o\ Rep \mid fa\ o\ Abs = Abs\ o\ fr \implies Abs\ o\ fr\ o\ Rep = fa$

by (*auto dest: Rep-comp-inverse elim: comp-Abs-inverse simp: o-assoc*)

lemmas *fns1a* = *disjI1* [*THEN fns1*]

lemmas *fns1b* = *disjI2* [*THEN fns1*]

lemma *fns4*:

Rep o fa o Abs = fr ==>

Rep o fa = fr o Rep & fa o Abs = Abs o fr

by (*auto intro!: ext*)

end

interpretation *nat-int*: *type-definition* [*int nat Collect (op <= 0)*]

by (*rule td-nat-int*)

— resetting to the default nat induct and cases rules

declare *Nat.induct* [*case-names 0 Suc, induct type*]

declare *Nat.exhaust* [*case-names 0 Suc, cases type*]

9.1 Extended form of type definition predicate

lemma *td-conds*:

norm o norm = norm ==> (fr o norm = norm o fr) =

(norm o fr o norm = fr o norm & norm o fr o norm = norm o fr)

apply *safe*

apply (*simp-all add: o-assoc [symmetric]*)

apply (*simp-all add: o-assoc*)

done

lemma *fn-comm-power*:

fa o tr = tr o fr ==> fa ^ n o tr = tr o fr ^ n

apply (*rule ext*)

apply (*induct n*)

apply (*auto dest: fun-cong*)

done

lemmas *fn-comm-power'* =

ext [THEN fn-comm-power, THEN fun-cong, unfolded o-def, standard]

locale *td-ext* = *type-definition* +

fixes *norm*

assumes *eq-norm*: $\bigwedge x. \text{Rep} (\text{Abs } x) = \text{norm } x$

begin

lemma *Abs-norm* [*simp*]:

Abs (norm x) = Abs x

using *eq-norm* [*of x*] **by** (*auto elim: Rep-inverse'*)

```

lemma td-th:
   $g \circ \text{Abs} = f \implies f (\text{Rep } x) = g x$ 
  by (drule comp-Abs-inverse [symmetric]) simp

lemma eq-norm':  $\text{Rep} \circ \text{Abs} = \text{norm}$ 
  by (auto simp: eq-norm intro!: ext)

lemma norm-Rep [simp]:  $\text{norm} (\text{Rep } x) = \text{Rep } x$ 
  by (auto simp: eq-norm' intro: td-th)

lemmas td = td-thm

lemma set-iff-norm:  $w : A \iff w = \text{norm } w$ 
  by (auto simp: set-Rep-Abs eq-norm' eq-norm [symmetric])

lemma inverse-norm:
   $(\text{Abs } n = w) = (\text{Rep } w = \text{norm } n)$ 
  apply (rule iffI)
  apply (clarsimp simp add: eq-norm)
  apply (simp add: eq-norm' [symmetric])
  done

lemma norm-eq-iff:
   $(\text{norm } x = \text{norm } y) = (\text{Abs } x = \text{Abs } y)$ 
  by (simp add: eq-norm' [symmetric])

lemma norm-comps:
   $\text{Abs} \circ \text{norm} = \text{Abs}$ 
   $\text{norm} \circ \text{Rep} = \text{Rep}$ 
   $\text{norm} \circ \text{norm} = \text{norm}$ 
  by (auto simp: eq-norm' [symmetric] o-def)

lemmas norm-norm [simp] = norm-comps

lemma fns5:
   $\text{Rep} \circ \text{fa} \circ \text{Abs} = \text{fr} \implies$ 
   $\text{fr} \circ \text{norm} = \text{fr} \ \& \ \text{norm} \circ \text{fr} = \text{fr}$ 
  by (fold eq-norm') (auto intro!: ext)

lemma fns2:
   $\text{Abs} \circ \text{fr} \circ \text{Rep} = \text{fa} \implies$ 
   $(\text{norm} \circ \text{fr} \circ \text{norm} = \text{fr} \circ \text{norm}) = (\text{Rep} \circ \text{fa} = \text{fr} \circ \text{Rep})$ 
  apply (fold eq-norm')
  apply safe
  prefer 2
  apply (simp add: o-assoc)
  apply (rule ext)
  apply (drule-tac x=Rep x in fun-cong)

```

```

apply auto
done

```

```

lemma fns3:
  Abs o fr o Rep = fa ==>
  (norm o fr o norm = norm o fr) = (fa o Abs = Abs o fr)
apply (fold eq-norm ^)
apply safe
prefer 2
apply (simp add: o-assoc [symmetric])
apply (rule ext)
apply (drule fun-cong)
apply simp
done

```

```

lemma fns:
  fr o norm = norm o fr ==>
  (fa o Abs = Abs o fr) = (Rep o fa = fr o Rep)
apply safe
apply (frule fns1b)
prefer 2
apply (frule fns1a)
apply (rule fns3 [THEN iffD1])
prefer 3
apply (rule fns2 [THEN iffD1])
apply (simp-all add: o-assoc [symmetric])
apply (simp-all add: o-assoc)
done

```

```

lemma range-norm:
  range (Rep o Abs) = A
by (simp add: set-Rep-Abs)

```

```

end

```

```

lemmas td-ext-def' =
  td-ext-def [unfolded type-definition-def td-ext-axioms-def]

```

```

end

```

10 WordDefinition: Definition of Word Type

```

theory WordDefinition imports Size BinBoolList TdThs begin

```

```

typedef (open word) 'a word'
  = {(0::int) ..< 2^len-of TYPE('a::len0)} by auto

```

```

instance word :: (len0) number ..
instance word :: (type) minus ..
instance word :: (type) plus ..
instance word :: (type) one ..
instance word :: (type) zero ..
instance word :: (type) times ..
instance word :: (type) Divides.div ..
instance word :: (type) power ..
instance word :: (type) ord ..
instance word :: (type) size ..
instance word :: (type) inverse ..
instance word :: (type) bit ..

```

10.1 Type conversions and casting

constdefs

— representation of words using unsigned or signed bins, only difference in these is the type class

```

word-of-int :: int => 'a :: len0 word
word-of-int w == Abs-word (bintrunc (len-of TYPE ('a)) w)

```

— uint and sint cast a word to an integer, uint treats the word as unsigned, sint treats the most-significant-bit as a sign bit

```

uint :: 'a :: len0 word => int
uint w == Rep-word w
sint :: 'a :: len word => int
sint-uint: sint w == sbintrunc (len-of TYPE ('a) - 1) (uint w)
unat :: 'a :: len0 word => nat
unat w == nat (uint w)

```

— the sets of integers representing the words

```

uints :: nat => int set
uints n == range (bintrunc n)
sints :: nat => int set
sints n == range (sbintrunc (n - 1))
unats :: nat => nat set
unats n == {i. i < 2 ^ n}
norm-sint :: nat => int => int
norm-sint n w == (w + 2 ^ (n - 1)) mod 2 ^ n - 2 ^ (n - 1)

```

— cast a word to a different length

```

scast :: 'a :: len word => 'b :: len word
scast w == word-of-int (sint w)
ucast :: 'a :: len0 word => 'b :: len0 word
ucast w == word-of-int (uint w)

```

— whether a cast (or other) function is to a longer or shorter length

```

source-size :: ('a :: len0 word => 'b) => nat
source-size c == let arb = arbitrary ; x = c arb in size arb

```

```

target-size :: ('a => 'b :: len0 word) => nat
target-size c == size (c arbitrary)
is-up :: ('a :: len0 word => 'b :: len0 word) => bool
is-up c == source-size c <= target-size c
is-down :: ('a :: len0 word => 'b :: len0 word) => bool
is-down c == target-size c <= source-size c

```

constdefs

```

of-bl :: bool list => 'a :: len0 word
of-bl bl == word-of-int (bl-to-bin bl)
to-bl :: 'a :: len0 word => bool list
to-bl w ==
bin-to-bl (len-of TYPE ('a)) (uint w)

```

```

word-reverse :: 'a :: len0 word => 'a word
word-reverse w == of-bl (rev (to-bl w))

```

defs (overloaded)

```

word-size: size (w :: 'a :: len0 word) == len-of TYPE('a)
word-number-of-def: number-of w == word-of-int w

```

constdefs

```

word-int-case :: (int => 'b) => ('a :: len0 word) => 'b
word-int-case f w == f (uint w)

```

syntax

```

of-int :: int => 'a

```

translations

```

case x of of-int y => b == word-int-case (%y. b) x

```

10.2 Arithmetic operations**defs (overloaded)**

```

word-1-wi: (1 :: ('a :: len0) word) == word-of-int 1
word-0-wi: (0 :: ('a :: len0) word) == word-of-int 0

```

```

word-le-def: a <= b == uint a <= uint b
word-less-def: x < y == x <= y & x ~ = (y :: 'a :: len0 word)

```

constdefs

```

word-succ :: 'a :: len0 word => 'a word
word-succ a == word-of-int (Numeral.succ (uint a))

```

```

word-pred :: 'a :: len0 word => 'a word
word-pred a == word-of-int (Numeral.pred (uint a))

```

```

udvd :: 'a::len word => 'a::len word => bool (infixl udvd 50)
a udvd b == EX n>=0. uint b = n * uint a

```

word-sle :: 'a :: len word => 'a word => bool ((-/ <=s -) [50, 51] 50)
a <=s b == *sint a <= sint b*

word-sless :: 'a :: len word => 'a word => bool ((-/ <s -) [50, 51] 50)
(x <s y) == *(x <=s y & x ~ = y)*

consts

word-power :: 'a :: len0 word => nat => 'a word

primrec

word-power a 0 = 1

word-power a (Suc n) = *a * word-power a n*

defs (overloaded)

word-pow: *power* == *word-power*

word-add-def: *a + b* == *word-of-int (uint a + uint b)*

word-sub-wi: *a - b* == *word-of-int (uint a - uint b)*

word-minus-def: *- a* == *word-of-int (- uint a)*

word-mult-def: *a * b* == *word-of-int (uint a * uint b)*

word-div-def: *a div b* == *word-of-int (uint a div uint b)*

word-mod-def: *a mod b* == *word-of-int (uint a mod uint b)*

10.3 Bit-wise operations**defs (overloaded)**

word-and-def:

(a::'a::len0 word) AND b == *word-of-int (uint a AND uint b)*

word-or-def:

(a::'a::len0 word) OR b == *word-of-int (uint a OR uint b)*

word-xor-def:

(a::'a::len0 word) XOR b == *word-of-int (uint a XOR uint b)*

word-not-def:

NOT (a::'a::len0 word) == *word-of-int (NOT (uint a))*

word-test-bit-def:

test-bit (a::'a::len0 word) == *bin-nth (uint a)*

word-set-bit-def:

set-bit (a::'a::len0 word) n x ==

word-of-int (bin-sc n (If x bit.B1 bit.B0) (uint a))

word-set-bits-def:

(BITS n. f n)::'a::len0 word == *of-bl (bl-of-nth (len-of TYPE ('a)) f)*

word-lsb-def:

lsb (a::'a::len0 word) == *bin-last (uint a) = bit.B1*

word-msb-def:
msb (*a*::*'a*::*len* *word*) == *bin-sign* (*sint* *a*) = *Numeral.Min*

constdefs

setBit :: *'a* :: *len0* *word* => *nat* => *'a* *word*
setBit *w* *n* == *set-bit* *w* *n* *True*

clearBit :: *'a* :: *len0* *word* => *nat* => *'a* *word*
clearBit *w* *n* == *set-bit* *w* *n* *False*

10.4 Shift operations**constdefs**

shiffl1 :: *'a* :: *len0* *word* => *'a* *word*
shiffl1 *w* == *word-of-int* (*uint* *w* *BIT* *bit.B0*)

— shift right as unsigned or as signed, ie logical or arithmetic

shiftr1 :: *'a* :: *len0* *word* => *'a* *word*
shiftr1 *w* == *word-of-int* (*bin-rest* (*uint* *w*))

sshiftr1 :: *'a* :: *len* *word* => *'a* *word*
sshiftr1 *w* == *word-of-int* (*bin-rest* (*sint* *w*))

bshiftr1 :: *bool* => *'a* :: *len* *word* => *'a* *word*
bshiftr1 *b* *w* == *of-bl* (*b* # *butlast* (*to-bl* *w*))

sshiftr :: *'a* :: *len* *word* => *nat* => *'a* *word* (**infixl** >>> 55)
w >>> *n* == (*sshiftr1* ^ *n*) *w*

mask :: *nat* => *'a*::*len* *word*
mask *n* == (*1* << *n*) - *1*

revcast :: *'a* :: *len0* *word* => *'b* :: *len0* *word*
revcast *w* == *of-bl* (*takefill* *False* (*len-of* *TYPE*('b)) (*to-bl* *w*))

slice1 :: *nat* => *'a* :: *len0* *word* => *'b* :: *len0* *word*
slice1 *n* *w* == *of-bl* (*takefill* *False* *n* (*to-bl* *w*))

slice :: *nat* => *'a* :: *len0* *word* => *'b* :: *len0* *word*
slice *n* *w* == *slice1* (*size* *w* - *n*) *w*

defs (overloaded)

shiffl-def: (*w*::*'a*::*len0* *word*) << *n* == (*shiffl1* ^ *n*) *w*
shiftr-def: (*w*::*'a*::*len0* *word*) >> *n* == (*shiftr1* ^ *n*) *w*

10.5 Rotation**constdefs**

```

rotater1 :: 'a list => 'a list
rotater1 ys ==
  case ys of [] => [] | x # xs => last ys # butlast ys

rotater :: nat => 'a list => 'a list
rotater n == rotater1 ^ n

word-rotr :: nat => 'a :: len0 word => 'a :: len0 word
word-rotr n w == of-bl (rotater n (to-bl w))

word-rotl :: nat => 'a :: len0 word => 'a :: len0 word
word-rotl n w == of-bl (rotate n (to-bl w))

word-roti :: int => 'a :: len0 word => 'a :: len0 word
word-roti i w == if i >= 0 then word-rotr (nat i) w
                else word-rotl (nat (- i)) w

```

10.6 Split and cat operations

constdefs

```

word-cat :: 'a :: len0 word => 'b :: len0 word => 'c :: len0 word
word-cat a b == word-of-int (bin-cat (uint a) (len-of TYPE ('b)) (uint b))

word-split :: 'a :: len0 word => ('b :: len0 word) * ('c :: len0 word)
word-split a ==
  case bin-split (len-of TYPE ('c)) (uint a) of
    (u, v) => (word-of-int u, word-of-int v)

word-rcat :: 'a :: len0 word list => 'b :: len0 word
word-rcat ws ==
  word-of-int (bin-rcat (len-of TYPE ('a)) (map uint ws))

word-rsplit :: 'a :: len0 word => 'b :: len word list
word-rsplit w ==
  map word-of-int (bin-rsplit (len-of TYPE ('b)) (len-of TYPE ('a), uint w))

```

constdefs

```

— Largest representable machine integer.
max-word :: 'a::len word
max-word ≡ word-of-int (2^len-of TYPE('a) - 1)

```

consts

```

of-bool :: bool => 'a::len word

```

primrec

```

of-bool False = 0
of-bool True = 1

```

lemmas *of-nth-def* = *word-set-bits-def*

lemmas *word-size-gt-0* [iff] =

xtr1 [OF *word-size* [THEN *meta-eq-to-obj-eq*] *len-gt-0*, *standard*]

lemmas *lens-gt-0* = *word-size-gt-0 len-gt-0*

lemmas *lens-not-0* [iff] = *lens-gt-0* [THEN *gr-implies-not0*, *standard*]

lemma *uints-num*: $uints\ n = \{i. 0 \leq i \wedge i < 2 \wedge n\}$

by (*simp add: uints-def range-bintrunc*)

lemma *sints-num*: $sints\ n = \{i. -(2 \wedge (n - 1)) \leq i \wedge i < 2 \wedge (n - 1)\}$

by (*simp add: sints-def range-sbintrunc*)

lemmas *atLeastLessThan-alt* = *atLeastLessThan-def* [*unfolded*

atLeast-def lessThan-def Collect-conj-eq [*symmetric*]

lemma *mod-in-reps*: $m > 0 \implies y \bmod m : \{0::int ..< m\}$

unfolding *atLeastLessThan-alt* **by** *auto*

lemma

Rep-word-0:0 <= *Rep-word* *x* **and**

Rep-word-lt: *Rep-word* ($x::'a::len0$ *word*) < $2 \wedge len\text{-of}\ TYPE('a)$

by (*auto simp: Rep-word* [*simplified*])

lemma *Rep-word-mod-same*:

Rep-word $x \bmod 2 \wedge len\text{-of}\ TYPE('a) = Rep\text{-word}$ ($x::'a::len0$ *word*)

by (*simp add: int-mod-eq Rep-word-lt Rep-word-0*)

lemma *td-ext-uint*:

td-ext ($uint :: 'a\ word \implies int$) *word-of-int* ($uints$ ($len\text{-of}\ TYPE('a::len0)$))
 ($\%w::int. w \bmod 2 \wedge len\text{-of}\ TYPE('a)$)

apply (*unfold td-ext-def'*)

apply (*simp add: uints-num uint-def word-of-int-def bintrunc-mod2p*)

apply (*simp add: Rep-word-mod-same Rep-word-0 Rep-word-lt*

word.Rep-word-inverse word.Abs-word-inverse int-mod-lem)

done

lemmas *int-word-uint* = *td-ext-uint* [THEN *td-ext.eq-norm*, *standard*]

interpretation *word-uint*:

td-ext [$uint::'a::len0$ *word* $\implies int$

word-of-int

$uints$ ($len\text{-of}\ TYPE('a::len0)$)

$\lambda w. w \bmod 2 \wedge len\text{-of}\ TYPE('a::len0)$]

by (*rule td-ext-uint*)

lemmas *td-uint* = *word-uint.td-thm*

lemmas *td-ext-ubin* = *td-ext-uint*

[*simplified len-gt-0 no-bintr-alt1 [symmetric]*]

interpretation *word-ubin*:

td-ext [*uint*::'a::len0 *word* \Rightarrow *int*
word-of-int
uints (*len-of TYPE*('a::len0))
bintrunc (*len-of TYPE*('a::len0))]
by (*rule td-ext-ubin*)

lemma *sint-sbintrunc'*:

sint (*word-of-int bin* :: 'a *word*) =
(*sbintrunc* (*len-of TYPE* ('a :: len) - 1) *bin*)
unfolding *sint-uint*
by (*auto simp: word-ubin.eq-norm sbintrunc-bintrunc-lt*)

lemma *uint-sint*:

uint w = *bintrunc* (*len-of TYPE*('a)) (*sint* (*w* :: 'a :: len *word*))
unfolding *sint-uint* **by** (*auto simp: bintrunc-sbintrunc-le*)

lemma *bintr-uint'*:

n \geq *size w* \implies *bintrunc n* (*uint w*) = *uint w*
apply (*unfold word-size*)
apply (*subst word-ubin.norm-Rep [symmetric]*)
apply (*simp only: bintrunc-bintrunc-min word-size min-def*)
apply *simp*
done

lemma *wi-bintr'*:

wb = *word-of-int bin* \implies *n* \geq *size wb* \implies
word-of-int (*bintrunc n bin*) = *wb*
unfolding *word-size*
by (*clarsimp simp add : word-ubin.norm-eq-iff [symmetric] min-def*)

lemmas *bintr-uint* = *bintr-uint'* [*unfolded word-size*]

lemmas *wi-bintr* = *wi-bintr'* [*unfolded word-size*]

lemma *td-ext-sbin*:

td-ext (*sint* :: 'a *word* \implies *int*) *word-of-int* (*sints* (*len-of TYPE*('a::len)))
(*sbintrunc* (*len-of TYPE*('a) - 1))
apply (*unfold td-ext-def' sint-uint*)
apply (*simp add : word-ubin.eq-norm*)
apply (*cases len-of TYPE*('a))
apply (*auto simp add : sints-def*)
apply (*rule sym [THEN trans]*)
apply (*rule word-ubin.Abs-norm*)
apply (*simp only: bintrunc-sbintrunc*)
apply (*drule sym*)
apply *simp*
done

lemmas *td-ext-sint = td-ext-sbin*
[simplified len-gt-0 no-sbintr-alt2 Suc-pred' [symmetric]]

interpretation *word-sint:*

td-ext [sint :: 'a::len word => int
word-of-int
sints (len-of TYPE('a::len))
%w. (w + 2^(len-of TYPE('a::len) - 1)) mod 2^len-of TYPE('a::len) -
2 ^ (len-of TYPE('a::len) - 1)]
by (*rule td-ext-sint*)

interpretation *word-sbin:*

td-ext [sint :: 'a::len word => int
word-of-int
sints (len-of TYPE('a::len))
sbintrunc (len-of TYPE('a::len) - 1)]
by (*rule td-ext-sbin*)

lemmas *int-word-sint = td-ext-sint [THEN td-ext.eq-norm, standard]*

lemmas *td-sint = word-sint.td*

lemma *word-number-of-alt: number-of b == word-of-int (number-of b)*
unfolding *word-number-of-def* **by** (*simp add: number-of-eq*)

lemma *word-no-wi: number-of = word-of-int*
by (*auto simp: word-number-of-def intro: ext*)

lemma *to-bl-def':*
(to-bl :: 'a :: len0 word => bool list) =
bin-to-bl (len-of TYPE('a)) o uint
by (*auto simp: to-bl-def intro: ext*)

lemmas *word-reverse-no-def [simp] = word-reverse-def [of number-of w, standard]*

lemmas *uints-mod = uints-def [unfolded no-bintr-alt1]*

lemma *uint-bintrunc: uint (number-of bin :: 'a word) =*
number-of (bintrunc (len-of TYPE ('a :: len0)) bin)
unfolding *word-number-of-def number-of-eq*
by (*auto intro: word-ubin.eq-norm*)

lemma *sint-sbintrunc: sint (number-of bin :: 'a word) =*
number-of (sbintrunc (len-of TYPE ('a :: len) - 1) bin)
unfolding *word-number-of-def number-of-eq*
by (*subst word-sbin.eq-norm*) *simp*

lemma *unat-bintrunc*:

unat (*number-of* *bin* :: 'a :: len0 word) =
number-of (*bintrunc* (*len-of* *TYPE*('a)) *bin*)
unfolding *unat-def* *nat-number-of-def*
by (*simp only*: *uint-bintrunc*)

declare

uint-bintrunc [*simp*]
sint-sbintrunc [*simp*]
unat-bintrunc [*simp*]

lemma *size-0-eq*: *size* (*w* :: 'a :: len0 word) = 0 ==> *v* = *w*

apply (*unfold* *word-size*)
apply (*rule* *word-uint.Rep-eqD*)
apply (*rule* *box-equals*)
defer
apply (*rule* *word-ubin.norm-Rep*)+
apply *simp*
done

lemmas *uint-lem* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]

lemmas *sint-lem* = *word-sint.Rep* [*unfolded sints-num mem-Collect-eq*]

lemmas *uint-ge-0* [*iff*] = *uint-lem* [*THEN conjunct1, standard*]

lemmas *uint-lt2p* [*iff*] = *uint-lem* [*THEN conjunct2, standard*]

lemmas *sint-ge* = *sint-lem* [*THEN conjunct1, standard*]

lemmas *sint-lt* = *sint-lem* [*THEN conjunct2, standard*]

lemma *sign-uint-Pls* [*simp*]:

bin-sign (*uint* *x*) = *Numeral.Pl*s
by (*simp add*: *sign-Pls-ge-0 number-of-eq*)

lemmas *uint-m2p-neg* = *iffD2* [*OF diff-less-0-iff-less uint-lt2p, standard*]

lemmas *uint-m2p-not-non-neg* =
iffD2 [*OF linorder-not-le uint-m2p-neg, standard*]

lemma *lt2p-lem*:

len-of *TYPE*('a) <= *n* ==> *uint* (*w* :: 'a :: len0 word) < 2 ^ *n*
by (*rule* *xtr8* [*OF - uint-lt2p*]) *simp*

lemmas *uint-le-0-iff* [*simp*] =

uint-ge-0 [*THEN leD, THEN linorder-antisym-conv1, standard*]

lemma *uint-nat*: *uint* *w* == *int* (*unat* *w*)

unfolding *unat-def* **by** *auto*

lemma *uint-number-of*:

uint (*number-of* *b* :: 'a :: len0 word) = *number-of* *b* mod 2 ^ *len-of* *TYPE*('a)
unfolding *word-number-of-alt*

by (*simp only: int-word-uint*)

lemma *unat-number-of*:

bin-sign b = Numeral.Pls ==>

unat (number-of b::'a::len0 word) = number-of b mod 2 ^ len-of TYPE ('a)

apply (*unfold unat-def*)

apply (*clarsimp simp only: uint-number-of*)

apply (*rule nat-mod-distrib [THEN trans]*)

apply (*erule sign-Pls-ge-0 [THEN iffD1]*)

apply (*simp-all add: nat-power-eq*)

done

lemma *sint-number-of*: *sint (number-of b :: 'a :: len word) = (number-of b +*

2 ^ (len-of TYPE('a) - 1)) mod 2 ^ len-of TYPE('a) -

2 ^ (len-of TYPE('a) - 1)

unfolding *word-number-of-alt* **by** (*rule int-word-sint*)

lemma *word-of-int-bin* [*simp*] :

(word-of-int (number-of bin) :: 'a :: len0 word) = (number-of bin)

unfolding *word-number-of-alt* **by** *auto*

lemma *word-int-case-wi*:

word-int-case f (word-of-int i :: 'b word) =

f (i mod 2 ^ len-of TYPE('b::len0))

unfolding *word-int-case-def* **by** (*simp add: word-uint.eq-norm*)

lemma *word-int-split*:

P (word-int-case f x) =

(ALL i. x = (word-of-int i :: 'b :: len0 word) &

0 <= i & i < 2 ^ len-of TYPE('b) --> P (f i))

unfolding *word-int-case-def*

by (*auto simp: word-uint.eq-norm int-mod-eq'*)

lemma *word-int-split-asm*:

P (word-int-case f x) =

(~ (EX n. x = (word-of-int n :: 'b::len0 word) &

0 <= n & n < 2 ^ len-of TYPE('b::len0) & ~ P (f n)))

unfolding *word-int-case-def*

by (*auto simp: word-uint.eq-norm int-mod-eq'*)

lemmas *uint-range'* =

word-uint.Rep [unfolded uints-num mem-Collect-eq, standard]

lemmas *sint-range'* = *word-sint.Rep [unfolded One-nat-def*

sints-num mem-Collect-eq, standard]

lemma *uint-range-size*: *0 <= uint w & uint w < 2 ^ size w*

unfolding *word-size* **by** (*rule uint-range'*)

lemma *sint-range-size*:

– $(2 \wedge (\text{size } w - \text{Suc } 0)) \leq \text{sint } w \ \& \ \text{sint } w < 2 \wedge (\text{size } w - \text{Suc } 0)$
unfolding *word-size* **by** (*rule sint-range'*)

lemmas *sint-above-size = sint-range-size*
 [THEN *conjunct2*, THEN [2] *xtr8*, folded *One-nat-def*, *standard*]

lemmas *sint-below-size = sint-range-size*
 [THEN *conjunct1*, THEN [2] *order-trans*, folded *One-nat-def*, *standard*]

lemma *test-bit-eq-iff*: $(\text{test-bit } (u::'a::\text{len0 } \text{word}) = \text{test-bit } v) = (u = v)$
unfolding *word-test-bit-def* **by** (*simp add: bin-nth-eq-iff*)

lemma *test-bit-size* [rule-format] : $(w::'a::\text{len0 } \text{word}) !! n \longrightarrow n < \text{size } w$
apply (*unfold word-test-bit-def*)
apply (*subst word-ubin.norm-Rep [symmetric]*)
apply (*simp only: nth-bintr word-size*)
apply *fast*
done

lemma *word-eqI* [rule-format] :
fixes $u :: 'a::\text{len0 } \text{word}$
shows $(\text{ALL } n. n < \text{size } u \longrightarrow u !! n = v !! n) \implies u = v$
apply (*rule test-bit-eq-iff [THEN iffD1]*)
apply (*rule ext*)
apply (*erule allE*)
apply (*erule impCE*)
prefer 2
apply *assumption*
apply (*auto dest!: test-bit-size simp add: word-size*)
done

lemmas *word-eqD = test-bit-eq-iff* [THEN *iffD2*, THEN *fun-cong*, *standard*]

lemma *test-bit-bin'*: $w !! n = (n < \text{size } w \ \& \ \text{bin-nth } (\text{uint } w) \ n)$
unfolding *word-test-bit-def word-size*
by (*simp add: nth-bintr [symmetric]*)

lemmas *test-bit-bin = test-bit-bin'* [unfolding *word-size*]

lemma *bin-nth-uint-imp'*: $\text{bin-nth } (\text{uint } w) \ n \longrightarrow n < \text{size } w$
apply (*unfold word-size*)
apply (*rule impI*)
apply (*rule nth-bintr [THEN iffD1, THEN conjunct1]*)
apply (*subst word-ubin.norm-Rep*)
apply *assumption*
done

lemma *bin-nth-sint'*:
 $n \geq \text{size } w \longrightarrow \text{bin-nth } (\text{sint } w) \ n = \text{bin-nth } (\text{sint } w) \ (\text{size } w - 1)$

```

apply (rule impI)
apply (subst word-sbin.norm-Rep [symmetric])
apply (simp add : nth-sbintr word-size)
apply auto
done

```

lemmas $bin\text{-}nth\text{-}uint\text{-}imp = bin\text{-}nth\text{-}uint\text{-}imp'$ [rule-format, unfolded word-size]

lemmas $bin\text{-}nth\text{-}sint = bin\text{-}nth\text{-}sint'$ [rule-format, unfolded word-size]

```

lemma td-bl:
  type-definition (to-bl :: 'a::len0 word => bool list)
    of-bl
    {bl. length bl = len-of TYPE('a)}
apply (unfold type-definition-def of-bl-def to-bl-def)
apply (simp add: word-ubin.eq-norm)
apply safe
apply (drule sym)
apply simp
done

```

```

interpretation word-bl:
  type-definition [to-bl :: 'a::len0 word => bool list]
    of-bl
    {bl. length bl = len-of TYPE('a::len0)}
by (rule td-bl)

```

lemma word-size-bl: $size\ w == size\ (to\text{-}bl\ w)$
unfolding word-size **by** auto

```

lemma to-bl-use-of-bl:
  (to-bl w = bl) = (w = of-bl bl  $\wedge$  length bl = length (to-bl w))
by (fastsimp elim!: word-bl.Abs-inverse [simplified])

```

lemma to-bl-word-rev: $to\text{-}bl\ (word\text{-}reverse\ w) = rev\ (to\text{-}bl\ w)$
unfolding word-reverse-def **by** (simp add: word-bl.Abs-inverse)

lemma word-rev-rev [simp] : $word\text{-}reverse\ (word\text{-}reverse\ w) = w$
unfolding word-reverse-def **by** (simp add : word-bl.Abs-inverse)

lemma word-rev-gal: $word\text{-}reverse\ w = u ==> word\text{-}reverse\ u = w$
by auto

lemmas $word\text{-}rev\text{-}gal' = sym$ [THEN word-rev-gal, symmetric, standard]

lemmas $length\text{-}bl\text{-}gt\text{-}0$ [iff] = xtr1 [OF word-bl.Rep' len-gt-0, standard]

lemmas $bl\text{-}not\text{-}Nil$ [iff] =

$length\text{-}bl\text{-}gt\text{-}0$ [THEN length-greater-0-conv [THEN iffD1], standard]

lemmas $length\text{-}bl\text{-}neq\text{-}0$ [iff] = $length\text{-}bl\text{-}gt\text{-}0$ [THEN gr-implies-not0]

lemma *hd-bl-sign-sint*: $hd (to-bl w) = (bin-sign (sint w) = Numeral.Min)$
apply (*unfold to-bl-def sint-uint*)
apply (*rule trans [OF - bl-sbin-sign]*)
apply *simp*
done

lemma *of-bl-drop'*:
 $lend = length bl - len-of TYPE ('a :: len0) ==>$
 $of-bl (drop lend bl) = (of-bl bl :: 'a word)$
apply (*unfold of-bl-def*)
apply (*clarsimp simp add : trunc-bl2bin [symmetric]*)
done

lemmas *of-bl-no = of-bl-def [folded word-number-of-def]*

lemma *test-bit-of-bl*:
 $(of-bl bl :: 'a :: len0 word) !! n = (rev bl ! n \wedge n < len-of TYPE('a) \wedge n < length bl)$
apply (*unfold of-bl-def word-test-bit-def*)
apply (*auto simp add : word-size word-ubin.eq-norm nth-bintr bin-nth-of-bl*)
done

lemma *no-of-bl*:
 $(number-of bin :: 'a :: len0 word) = of-bl (bin-to-bl (len-of TYPE('a)) bin)$
unfolding *word-size of-bl-no by (simp add : word-number-of-def)*

lemma *uint-bl*: $to-bl w == bin-to-bl (size w) (uint w)$
unfolding *word-size to-bl-def by auto*

lemma *to-bl-bin*: $bl-to-bin (to-bl w) = uint w$
unfolding *uint-bl by (simp add : word-size)*

lemma *to-bl-of-bin*:
 $to-bl (word-of-int bin :: 'a :: len0 word) = bin-to-bl (len-of TYPE('a)) bin$
unfolding *uint-bl by (clarsimp simp add : word-ubin.eq-norm word-size)*

lemmas *to-bl-no-bin [simp] = to-bl-of-bin [folded word-number-of-def]*

lemma *to-bl-to-bin [simp]* : $bl-to-bin (to-bl w) = uint w$
unfolding *uint-bl by (simp add : word-size)*

lemmas *uint-bl-bin [simp] = trans [OF bin-bl-bin word-ubin.norm-Rep, standard]*

lemmas *num-AB-u [simp] = word-uint.Rep-inverse*
[unfolded o-def word-number-of-def [symmetric], standard]

lemmas *num-AB-s [simp] = word-sint.Rep-inverse*
[unfolded o-def word-number-of-def [symmetric], standard]

```

lemma wints-unats: wints n = int ‘ unats n
  apply (unfold unats-def wints-num)
  apply safe
  apply (rule-tac image-eqI)
  apply (erule-tac nat-0-le [symmetric])
  apply auto
  apply (erule-tac nat-less-iff [THEN iffD2])
  apply (rule-tac [2] zless-nat-eq-int-zless [THEN iffD1])
  apply (auto simp add : nat-power-eq int-power)
  done

lemma unats-wints: unats n = nat ‘ wints n
  by (auto simp add : wints-unats image-iff)

lemmas bintr-num = word-ubin.norm-eq-iff
  [symmetric, folded word-number-of-def, standard]
lemmas sbintr-num = word-sbin.norm-eq-iff
  [symmetric, folded word-number-of-def, standard]

lemmas num-of-bintr = word-ubin.Abs-norm [folded word-number-of-def, stan-
dard]
lemmas num-of-sbintr = word-sbin.Abs-norm [folded word-number-of-def, stan-
dard]

lemma num-of-bintr':
  bintrunc (len-of TYPE('a :: len0)) a = b ==>
    number-of a = (number-of b :: 'a word)
  apply safe
  apply (rule-tac num-of-bintr [symmetric])
  done

lemma num-of-sbintr':
  sbintrunc (len-of TYPE('a :: len) - 1) a = b ==>
    number-of a = (number-of b :: 'a word)
  apply safe
  apply (rule-tac num-of-sbintr [symmetric])
  done

lemmas num-abs-bintr = sym [THEN trans,
  OF num-of-bintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]
lemmas num-abs-sbintr = sym [THEN trans,
  OF num-of-sbintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]

lemma ucast-id: ucast w = w

```

unfolding *ucast-def* **by** *auto*

lemma *scast-id*: *scast w = w*
unfolding *scast-def* **by** *auto*

lemma *ucast-bl*: *ucast w == of-bl (to-bl w)*
unfolding *ucast-def of-bl-def uint-bl*
by (*auto simp add : word-size*)

lemma *nth-ucast*:
(ucast w :: 'a :: len0 word) !! n = (w !! n & n < len-of TYPE('a))
apply (*unfold ucast-def test-bit-bin*)
apply (*simp add: word-ubin.eq-norm nth-bintr word-size*)
apply (*fast elim!: bin-nth-uint-imp*)
done

lemma *ucast-bintr [simp]*:
ucast (number-of w :: 'a :: len0 word) =
number-of (bintrunc (len-of TYPE('a)) w)
unfolding *ucast-def* **by** *simp*

lemma *scast-sbintr [simp]*:
scast (number-of w :: 'a :: len word) =
number-of (sbintrunc (len-of TYPE('a) - Suc 0) w)
unfolding *scast-def* **by** *simp*

lemmas *source-size = source-size-def [unfolded Let-def word-size]*
lemmas *target-size = target-size-def [unfolded Let-def word-size]*
lemmas *is-down = is-down-def [unfolded source-size target-size]*
lemmas *is-up = is-up-def [unfolded source-size target-size]*

lemmas *is-up-down =*
trans [OF is-up [THEN meta-eq-to-obj-eq]
is-down [THEN meta-eq-to-obj-eq, symmetric],
standard]

lemma *down-cast-same'*: *uc = ucast ==> is-down uc ==> uc = scast*
apply (*unfold is-down*)
apply *safe*
apply (*rule ext*)
apply (*unfold ucast-def scast-def uint-sint*)
apply (*rule word-ubin.norm-eq-iff [THEN iffD1]*)
apply *simp*
done

lemma *word-rev-tf'*:
r = to-bl (of-bl bl) ==> r = rev (takefill False (length r) (rev bl))

unfolding *of-bl-def uint-bl*
by (*clarsimp simp add: bl-bin-bl-rtf word-ubin.eq-norm word-size*)

lemmas *word-rev-tf = refl [THEN word-rev-tf', unfolded word-bl.Rep', standard]*

lemmas *word-rep-drop = word-rev-tf [simplified takefill-alt, simplified, simplified rev-take, simplified]*

lemma *to-bl-ucast:*

to-bl (ucast (w::'b::len0 word) ::'a::len0 word) =
replicate (len-of TYPE('a) - len-of TYPE('b)) False @
drop (len-of TYPE('b) - len-of TYPE('a)) (to-bl w)
apply (*unfold ucast-bl*)
apply (*rule trans*)
apply (*rule word-rep-drop*)
apply *simp*
done

lemma *ucast-up-app':*

uc = ucast ==> source-size uc + n = target-size uc ==>
to-bl (uc w) = replicate n False @ (to-bl w)
apply (*auto simp add : source-size target-size to-bl-ucast*)
apply (*rule-tac f = %n. replicate n False in arg-cong*)
apply *simp*
done

lemma *ucast-down-drop':*

uc = ucast ==> source-size uc = target-size uc + n ==>
to-bl (uc w) = drop n (to-bl w)
by (*auto simp add : source-size target-size to-bl-ucast*)

lemma *scast-down-drop':*

sc = scast ==> source-size sc = target-size sc + n ==>
to-bl (sc w) = drop n (to-bl w)
apply (*subgoal-tac sc = ucast*)
apply *safe*
apply *simp*
apply (*erule refl [THEN ucast-down-drop']*)
apply (*rule refl [THEN down-cast-same', symmetric]*)
apply (*simp add : source-size target-size is-down*)
done

lemma *sint-up-scast':*

sc = scast ==> is-up sc ==> sint (sc w) = sint w
apply (*unfold is-up*)
apply *safe*
apply (*simp add: scast-def word-sbin.eq-norm*)
apply (*rule box-equals*)
prefer 3

```

  apply (rule word-sbin.norm-Rep)
  apply (rule sbintrunc-sbintrunc-l)
  defer
  apply (subst word-sbin.norm-Rep)
  apply (rule refl)
  apply simp
done

```

```

lemma uint-up-ucast':
  uc = ucast ==> is-up uc ==> uint (uc w) = uint w
  apply (unfold is-up)
  apply safe
  apply (rule bin-eqI)
  apply (fold word-test-bit-def)
  apply (auto simp add: nth-ucast)
  apply (auto simp add: test-bit-bin)
done

```

```

lemmas down-cast-same = refl [THEN down-cast-same]
lemmas ucast-up-app = refl [THEN ucast-up-app]
lemmas ucast-down-drop = refl [THEN ucast-down-drop]
lemmas scast-down-drop = refl [THEN scast-down-drop]
lemmas uint-up-ucast = refl [THEN uint-up-ucast]
lemmas sint-up-scast = refl [THEN sint-up-scast]

```

```

lemma ucast-up-ucast': uc = ucast ==> is-up uc ==> ucast (uc w) = ucast w
  apply (simp (no-asm) add: ucast-def)
  apply (clarsimp simp add: uint-up-ucast)
done

```

```

lemma scast-up-scast': sc = scast ==> is-up sc ==> scast (sc w) = scast w
  apply (simp (no-asm) add: scast-def)
  apply (clarsimp simp add: sint-up-scast)
done

```

```

lemma ucast-of-bl-up':
  w = of-bl bl ==> size bl <= size w ==> ucast w = of-bl bl
  by (auto simp add : nth-ucast word-size test-bit-of-bl intro!: word-eqI)

```

```

lemmas ucast-up-ucast = refl [THEN ucast-up-ucast]
lemmas scast-up-scast = refl [THEN scast-up-scast]
lemmas ucast-of-bl-up = refl [THEN ucast-of-bl-up]

```

```

lemmas ucast-up-ucast-id = trans [OF ucast-up-ucast ucast-id]
lemmas scast-up-scast-id = trans [OF scast-up-scast scast-id]

```

```

lemmas isduu = is-up-down [where c = ucast, THEN iffD2]
lemmas isdus = is-up-down [where c = scast, THEN iffD2]
lemmas ucast-down-ucast-id = isduu [THEN ucast-up-ucast-id]

```

lemmas *scast-down-scast-id* = *isdus* [*THEN* *ucast-up-ucast-id*]

lemma *up-ucast-surj*:

is-up (*ucast* :: 'b::len0 word => 'a::len0 word) ==>
surj (*ucast* :: 'a word => 'b word)
by (*rule* *surjI*, *erule* *ucast-up-ucast-id*)

lemma *up-scast-surj*:

is-up (*scast* :: 'b::len word => 'a::len word) ==>
surj (*scast* :: 'a word => 'b word)
by (*rule* *surjI*, *erule* *scast-up-scast-id*)

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word => 'a::len word) ==>
inj-on (*ucast* :: 'a word => 'b word) *A*
by (*rule* *inj-on-inverseI*, *erule* *scast-down-scast-id*)

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word => 'a::len0 word) ==>
inj-on (*ucast* :: 'a word => 'b word) *A*
by (*rule* *inj-on-inverseI*, *erule* *ucast-down-ucast-id*)

lemma *of-bl-append-same*: *of-bl* (*X* @ *to-bl* *w*) = *w*

by (*rule* *word-bl.Rep-eqD*) (*simp* *add*: *word-rep-drop*)

lemma *ucast-down-no'*:

uc = *ucast* ==> *is-down* *uc* ==> *uc* (*number-of* *bin*) = *number-of* *bin*
apply (*unfold* *word-number-of-def* *is-down*)
apply (*clarsimp* *simp* *add*: *ucast-def* *word-ubin.eq-norm*)
apply (*rule* *word-ubin.norm-eq-iff* [*THEN* *iffD1*])
apply (*erule* *bintrunc-bintrunc-ge*)
done

lemmas *ucast-down-no* = *ucast-down-no'* [*OF* *refl*]

lemma *ucast-down-bl'*: *uc* = *ucast* ==> *is-down* *uc* ==> *uc* (*of-bl* *bl*) = *of-bl* *bl*

unfolding *of-bl-no* **by** *clarify* (*erule* *ucast-down-no*)

lemmas *ucast-down-bl* = *ucast-down-bl'* [*OF* *refl*]

lemmas *slice-def'* = *slice-def* [*unfolded* *word-size*]

lemmas *test-bit-def'* = *word-test-bit-def* [*THEN* *meta-eq-to-obj-eq*, *THEN* *fun-cong*]

lemmas *word-log-defs* = *word-and-def* *word-or-def* *word-xor-def* *word-not-def*

lemmas *word-log-bin-defs* = *word-log-defs*

end

11 WordArith: Word Arithmetic

theory *WordArith* **imports** *WordDefinition* **begin**

lemma *word-less-alt*: $(a < b) = (uint\ a < uint\ b)$
unfolding *word-less-def word-le-def*
by (*auto simp del: word-uint.Rep-inject*
simp: word-uint.Rep-inject [symmetric])

lemma *signed-linorder*: *linorder word-sle word-sless*
apply *unfold-locales*
apply (*unfold word-sle-def word-sless-def*)
by *auto*

interpretation *signed*: *linorder [word-sle word-sless]*
by (*rule signed-linorder*)

lemmas *word-arith-wis* [*THEN meta-eq-to-obj-eq*] =
word-add-def word-mult-def word-minus-def
word-succ-def word-pred-def word-0-wi word-1-wi

lemma *udvdI*:
 $0 \leq n \implies uint\ b = n * uint\ a \implies a\ udvd\ b$
by (*auto simp: udvd-def*)

lemmas *word-div-no* [*simp*] =
word-div-def [of number-of a number-of b, standard]

lemmas *word-mod-no* [*simp*] =
word-mod-def [of number-of a number-of b, standard]

lemmas *word-less-no* [*simp*] =
word-less-def [of number-of a number-of b, standard]

lemmas *word-le-no* [*simp*] =
word-le-def [of number-of a number-of b, standard]

lemmas *word-sless-no* [*simp*] =
word-sless-def [of number-of a number-of b, standard]

lemmas *word-sle-no* [*simp*] =
word-sle-def [of number-of a number-of b, standard]

lemmas *word-0-wi-Pls* = *word-0-wi* [*folded Pls-def*]

lemmas *word-0-no* = *word-0-wi-Pls* [*folded word-no-wi*]

lemma *int-one-bin*: $(1 :: int) == (Numeral.Pls\ BIT\ bit.B1)$

unfolding *Pls-def Bit-def* **by** *auto*

lemma *word-1-no*:

$(1 :: 'a :: \text{len0 } \text{word}) == \text{number-of } (\text{Numeral.Pls BIT bit.B1})$

unfolding *word-1-wi word-number-of-def int-one-bin* **by** *auto*

lemma *word-m1-wi*: $-1 == \text{word-of-int } -1$

by (*rule word-number-of-alt*)

lemma *word-m1-wi-Min*: $-1 = \text{word-of-int } \text{Numeral.Min}$

by (*simp add: word-m1-wi number-of-eq*)

lemma *word-0-bl*: $\text{of-bl } [] = 0$

unfolding *word-0-wi of-bl-def* **by** (*simp add : Pls-def*)

lemma *word-1-bl*: $\text{of-bl } [\text{True}] = 1$

unfolding *word-1-wi of-bl-def*

by (*simp add : bl-to-bin-def Bit-def Pls-def*)

lemma *uint-0 [simp]*: $(\text{uint } 0 = 0)$

unfolding *word-0-wi*

by (*simp add: word-ubin.eq-norm Pls-def [symmetric]*)

lemma *of-bl-0 [simp]*: $\text{of-bl } (\text{replicate } n \text{ False}) = 0$

by (*simp add : word-0-wi of-bl-def bl-to-bin-rep-False Pls-def*)

lemma *to-bl-0*:

$\text{to-bl } (0 :: 'a :: \text{len0 } \text{word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ False}$

unfolding *uint-bl*

by (*simp add : word-size bin-to-bl-Pls Pls-def [symmetric]*)

lemma *uint-0-iff*: $(\text{uint } x = 0) = (x = 0)$

by (*auto intro!: word-uint.Rep-eqD*)

lemma *unat-0-iff*: $(\text{unat } x = 0) = (x = 0)$

unfolding *unat-def* **by** (*auto simp add : nat-eq-iff uint-0-iff*)

lemma *unat-0 [simp]*: $\text{unat } 0 = 0$

unfolding *unat-def* **by** *auto*

lemma *size-0-same'*: $\text{size } w = 0 ==> w = (v :: 'a :: \text{len0 } \text{word})$

apply (*unfold word-size*)

apply (*rule box-equals*)

defer

apply (*rule word-uint.Rep-inverse*)+

apply (*rule word-ubin.norm-eq-iff [THEN iffD1]*)

apply *simp*

done

```

lemmas size-0-same = size-0-same' [folded word-size]

lemmas unat-eq-0 = unat-0-iff
lemmas unat-eq-zero = unat-0-iff

lemma unat-gt-0: (0 < unat x) = (x ~ = 0)
by (auto simp: unat-0-iff [symmetric])

lemma ucast-0 [simp] : ucast 0 = 0
unfolding ucast-def
by simp (simp add: word-0-wi)

lemma sint-0 [simp] : sint 0 = 0
unfolding sint-wint
by (simp add: Pls-def [symmetric])

lemma scast-0 [simp] : scast 0 = 0
apply (unfold scast-def)
apply simp
apply (simp add: word-0-wi)
done

lemma sint-n1 [simp] : sint -1 = -1
apply (unfold word-m1-wi-Min)
apply (simp add: word-sbin.eq-norm)
apply (unfold Min-def number-of-eq)
apply simp
done

lemma scast-n1 [simp] : scast -1 = -1
apply (unfold scast-def sint-n1)
apply (unfold word-number-of-alt)
apply (rule refl)
done

lemma uint-1 [simp] : uint (1 :: 'a :: len word) = 1
unfolding word-1-wi
by (simp add: word-ubin.eq-norm int-one-bin bintrunc-minus-simps)

lemma unat-1 [simp] : unat (1 :: 'a :: len word) = 1
by (unfold unat-def uint-1) auto

lemma ucast-1 [simp] : ucast (1 :: 'a :: len word) = 1
unfolding ucast-def word-1-wi
by (simp add: word-ubin.eq-norm int-one-bin bintrunc-minus-simps)

lemmas arths =

```

bintr-ariths [THEN *word-ubin.norm-eq-iff* [THEN *iffD1*],
folded *word-ubin.eq-norm*, *standard*]

lemma *wi-homs*:

shows

wi-hom-add: $\text{word-of-int } a + \text{word-of-int } b = \text{word-of-int } (a + b)$ **and**
wi-hom-mult: $\text{word-of-int } a * \text{word-of-int } b = \text{word-of-int } (a * b)$ **and**
wi-hom-neg: $-\text{word-of-int } a = \text{word-of-int } (-a)$ **and**
wi-hom-succ: $\text{word-succ } (\text{word-of-int } a) = \text{word-of-int } (\text{Numeral.succ } a)$ **and**
wi-hom-pred: $\text{word-pred } (\text{word-of-int } a) = \text{word-of-int } (\text{Numeral.pred } a)$
by (*auto simp*: *word-arith-wis arths*)

lemmas *wi-hom-syms = wi-homs* [*symmetric*]

lemma *word-sub-def*: $a - b == a + - (b :: 'a :: \text{len0 word})$

unfolding *word-sub-wi diff-def*

by (*simp only* : *word-uint.Rep-inverse wi-hom-syms*)

lemmas *word-diff-minus = word-sub-def* [THEN *meta-eq-to-obj-eq*, *standard*]

lemma *word-of-int-sub-hom*:

$(\text{word-of-int } a) - \text{word-of-int } b = \text{word-of-int } (a - b)$

unfolding *word-sub-def diff-def* **by** (*simp only* : *wi-homs*)

lemmas *new-word-of-int-homs =*

word-of-int-sub-hom wi-homs word-0-wi word-1-wi

lemmas *new-word-of-int-hom-syms = new-word-of-int-homs* [*symmetric*, *standard*]

lemmas *word-of-int-hom-syms =*

new-word-of-int-hom-syms [*unfolded succ-def pred-def*]

lemmas *word-of-int-homs =*

new-word-of-int-homs [*unfolded succ-def pred-def*]

lemmas *word-of-int-add-hom = word-of-int-homs* (2)

lemmas *word-of-int-mult-hom = word-of-int-homs* (3)

lemmas *word-of-int-minus-hom = word-of-int-homs* (4)

lemmas *word-of-int-succ-hom = word-of-int-homs* (5)

lemmas *word-of-int-pred-hom = word-of-int-homs* (6)

lemmas *word-of-int-0-hom = word-of-int-homs* (7)

lemmas *word-of-int-1-hom = word-of-int-homs* (8)

lemmas *word-arith-alt*s =

word-sub-wi [*unfolded succ-def pred-def*, THEN *meta-eq-to-obj-eq*, *standard*]

word-arith-wis [*unfolded succ-def pred-def*, *standard*]

lemmas *word-sub-alt* = *word-arith-alt*s (1)
lemmas *word-add-alt* = *word-arith-alt*s (2)
lemmas *word-mult-alt* = *word-arith-alt*s (3)
lemmas *word-minus-alt* = *word-arith-alt*s (4)
lemmas *word-succ-alt* = *word-arith-alt*s (5)
lemmas *word-pred-alt* = *word-arith-alt*s (6)
lemmas *word-0-alt* = *word-arith-alt*s (7)
lemmas *word-1-alt* = *word-arith-alt*s (8)

11.1 Transferring goals from words to ints

lemma *word-ths*:

shows

word-succ-p1: *word-succ* $a = a + 1$ **and**
word-pred-m1: *word-pred* $a = a - 1$ **and**
word-pred-succ: *word-pred* (*word-succ* a) = a **and**
word-succ-pred: *word-succ* (*word-pred* a) = a **and**
word-mult-succ: *word-succ* $a * b = b + a * b$
by (*rule* *word-uint.Abs-cases* [of b],
rule *word-uint.Abs-cases* [of a],
simp *add*: *pred-def succ-def add-commute mult-commute*
*ring-distrib*s *new-word-of-int-homs*)+

lemmas *uint-cong* = *arg-cong* [**where** $f = \text{uint}$]

lemmas *uint-word-ariths* =

*word-arith-alt*s [*THEN* *trans* [*OF* *uint-cong int-word-uint*], *standard*]

lemmas *uint-word-arith-bintrs* = *uint-word-ariths* [*folded bintrunc-mod2p*]

lemmas *sint-word-ariths* = *uint-word-arith-bintrs*

[*THEN* *uint-sint* [*symmetric*, *THEN* *trans*],
unfolded uint-sint bintr-arith1s bintr-ariths
len-gt-0 [*THEN* *bin-sbin-eq-iff'*] *word-sbin.norm-Rep*, *standard*]

lemmas *uint-div-alt* = *word-div-def*

[*THEN* *meta-eq-to-obj-eq* [*THEN* *trans* [*OF* *uint-cong int-word-uint*]], *standard*]

lemmas *uint-mod-alt* = *word-mod-def*

[*THEN* *meta-eq-to-obj-eq* [*THEN* *trans* [*OF* *uint-cong int-word-uint*]], *standard*]

lemma *word-pred-0-n1*: *word-pred* $0 = \text{word-of-int } - 1$

unfolding *word-pred-def number-of-eq*

by (*simp* *add* : *pred-def word-no-wi*)

lemma *word-pred-0-Min*: *word-pred* $0 = \text{word-of-int Numeral.Min}$

by (*simp* *add*: *word-pred-0-n1 number-of-eq*)

lemma *word-m1-Min*: $- 1 = \text{word-of-int Numeral.Min}$

unfolding *Min-def* **by** (*simp only: word-of-int-hom-syms*)

lemma *succ-pred-no* [*simp*]:

word-succ (*number-of bin*) = *number-of* (*Numeral.succ bin*) &

word-pred (*number-of bin*) = *number-of* (*Numeral.pred bin*)

unfolding *word-number-of-def* **by** (*simp add : new-word-of-int-homs*)

lemma *word-sp-01* [*simp*] :

word-succ -1 = 0 & *word-succ 0 = 1* & *word-pred 0 = -1* & *word-pred 1 = 0*

by (*unfold word-0-no word-1-no*) *auto*

lemma *word-of-int-Ex*:

$\exists y. x = \text{word-of-int } y$

by (*rule-tac x=uint x in exI*) *simp*

lemma *word-arith-egs*:

fixes *a* :: '*a*::len0 *word*

fixes *b* :: '*a*::len0 *word*

shows

word-add-0: $0 + a = a$ **and**

word-add-0-right: $a + 0 = a$ **and**

word-mult-1: $1 * a = a$ **and**

word-mult-1-right: $a * 1 = a$ **and**

word-add-commute: $a + b = b + a$ **and**

word-add-assoc: $a + b + c = a + (b + c)$ **and**

word-add-left-commute: $a + (b + c) = b + (a + c)$ **and**

word-mult-commute: $a * b = b * a$ **and**

word-mult-assoc: $a * b * c = a * (b * c)$ **and**

word-mult-left-commute: $a * (b * c) = b * (a * c)$ **and**

word-left-distrib: $(a + b) * c = a * c + b * c$ **and**

word-right-distrib: $a * (b + c) = a * b + a * c$ **and**

word-left-minus: $- a + a = 0$ **and**

word-diff-0-right: $a - 0 = a$ **and**

word-diff-self: $a - a = 0$

using *word-of-int-Ex* [*of a*]

word-of-int-Ex [*of b*]

word-of-int-Ex [*of c*]

by (*auto simp: word-of-int-hom-syms* [*symmetric*])

zadd-0-right add-commute add-assoc add-left-commute

mult-commute mult-assoc mult-left-commute

plus-times.left-distrib plus-times.right-distrib)

lemmas *word-add-ac = word-add-commute word-add-assoc word-add-left-commute*

lemmas *word-mult-ac = word-mult-commute word-mult-assoc word-mult-left-commute*

lemmas *word-plus-ac0 = word-add-0 word-add-0-right word-add-ac*

lemmas *word-times-ac1 = word-mult-1 word-mult-1-right word-mult-ac*

11.2 Order on fixed-length words

lemma *word-order-trans*: $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** *auto*

lemma *word-order-reft*: $z \leq (z :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** *auto*

lemma *word-order-antisym*: $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** (*auto intro!*: *word-uint.Rep-eqD*)

lemma *word-order-linear*:
 $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** *auto*

lemma *word-zero-le* [*simp*] :
 $0 \leq (y :: 'a :: \text{len0 word})$
unfolding *word-le-def* **by** *auto*

instance *word* :: (*len0*) *semigroup-add*
by *intro-classes* (*simp add*: *word-add-assoc*)

instance *word* :: (*len0*) *linorder*
by *intro-classes* (*auto simp*: *word-less-def word-le-def*)

instance *word* :: (*len0*) *ring*
by *intro-classes*
(*auto simp*: *word-arith-eqs word-diff-minus*
word-diff-self [*unfolded word-diff-minus*])

lemma *word-m1-ge* [*simp*] : *word-pred* 0 $\geq y$
unfolding *word-le-def*
by (*simp only* : *word-pred-0-n1 word-uint.eq-norm m1mod2k*) *auto*

lemmas *word-n1-ge* [*simp*] = *word-m1-ge* [*simplified word-sp-01*]

lemmas *word-not-simps* [*simp*] =
word-zero-le [*THEN leD*] *word-m1-ge* [*THEN leD*] *word-n1-ge* [*THEN leD*]

lemma *word-gt-0*: $0 < y = (0 \sim (y :: 'a :: \text{len0 word}))$
unfolding *word-less-def* **by** *auto*

lemmas *word-gt-0-no* [*simp*] = *word-gt-0* [*of number-of y, standard*]

lemma *word-sless-alt*: $(a <_s b) == (\text{sint } a < \text{sint } b)$
unfolding *word-sle-def word-sless-def*
by (*auto simp add* : *less-eq-less.less-le*)

lemma *word-le-nat-alt*: $(a \leq b) = (\text{unat } a \leq \text{unat } b)$
unfolding *unat-def word-le-def*

by (rule nat-le-eq-zle [symmetric]) simp

lemma word-less-nat-alt: $(a < b) = (\text{unat } a < \text{unat } b)$
unfolding unat-def word-less-alt
by (rule nat-less-eq-zless [symmetric]) simp

lemma wi-less:
 $(\text{word-of-int } n < (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} < m \bmod 2 \wedge \text{len-of TYPE('a)})$
unfolding word-less-alt **by** (simp add: word-uint.eq-norm)

lemma wi-le:
 $(\text{word-of-int } n \leq (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} \leq m \bmod 2 \wedge \text{len-of TYPE('a)})$
unfolding word-le-def **by** (simp add: word-uint.eq-norm)

lemma udvd-nat-alt: $a \text{ udvd } b = (\exists X n \geq 0. \text{unat } b = n * \text{unat } a)$
apply (unfold udvd-def)
apply safe
apply (simp add: unat-def nat-mult-distrib)
apply (simp add: uint-nat int-mult)
apply (rule exI)
apply safe
prefer 2
apply (erule notE)
apply (rule refl)
apply force
done

lemma udvd-iff-dvd: $x \text{ udvd } y \iff \text{unat } x \text{ dvd } \text{unat } y$
unfolding dvd-def udvd-nat-alt **by** force

lemmas unat-mono = word-less-nat-alt [THEN iffD1, standard]

lemma word-zero-neq-one: $0 < \text{len-of TYPE('a)} \implies (0 :: 'a \text{ word}) \approx 1$
unfolding word-arith-wis
by (auto simp add: word-ubin.norm-eq-iff [symmetric] gr0-conv-Suc)

lemmas lenw1-zero-neq-one = len-gt-0 [THEN word-zero-neq-one]

lemma no-no [simp]: $\text{number-of } (\text{number-of } b) = \text{number-of } b$
by (simp add: number-of-eq)

lemma unat-minus-one: $x \approx 0 \implies \text{unat } (x - 1) = \text{unat } x - 1$
apply (unfold unat-def)
apply (simp only: int-word-uint word-arith-alt5 rdmods)
apply (subgoal-tac uint x >= 1)
prefer 2

```

apply (drule contrapos-nn)
apply (erule word-uint.Rep-inverse' [symmetric])
apply (insert uint-ge-0 [of x])[1]
apply arith
apply (rule box-equals)
apply (rule nat-diff-distrib)
prefer 2
apply assumption
apply simp
apply (subst mod-pos-pos-trivial)
apply arith
apply (insert uint-lt2p [of x])[1]
apply arith
apply (rule refl)
apply simp
done

```

lemma *measure-unat*: $p \sim = 0 \implies \text{unat } (p - 1) < \text{unat } p$
by (*simp add: unat-minus-one*) (*simp add: unat-0-iff* [symmetric])

lemmas *uint-add-ge0* [*simp*] =
add-nonneg-nonneg [*OF uint-ge-0 uint-ge-0, standard*]
lemmas *uint-mult-ge0* [*simp*] =
mult-nonneg-nonneg [*OF uint-ge-0 uint-ge-0, standard*]

lemma *uint-sub-lt2p* [*simp*]:
 $\text{uint } (x :: 'a :: \text{len0 word}) - \text{uint } (y :: 'b :: \text{len0 word}) <$
 $2 \wedge \text{len-of TYPE('a)}$
using *uint-ge-0* [of *y*] *uint-lt2p* [of *x*] **by** *arith*

11.3 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:
 $(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x + y :: 'a :: \text{len0 word}) = \text{uint } x + \text{uint } y)$
by (*unfold uint-word-ariths*) (*auto intro!: trans* [*OF - int-mod-lem*])

lemma *uint-mult-lem*:
 $(\text{uint } x * \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x * y :: 'a :: \text{len0 word}) = \text{uint } x * \text{uint } y)$
by (*unfold uint-word-ariths*) (*auto intro!: trans* [*OF - int-mod-lem*])

lemma *uint-sub-lem*:
 $(\text{uint } x \geq \text{uint } y) = (\text{uint } (x - y) = \text{uint } x - \text{uint } y)$
by (*unfold uint-word-ariths*) (*auto intro!: trans* [*OF - int-mod-lem*])

lemma *uint-add-le*: $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$
unfolding *uint-word-ariths* **by** (*auto simp: mod-add-if-z*)

lemma *uint-sub-ge*: $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$
unfolding *uint-word-ariths* **by** (*auto simp: mod-sub-if-z*)

lemmas *uint-sub-if'* =
trans [OF uint-word-ariths(1) mod-sub-if-z, simplified, standard]

lemmas *uint-plus-if'* =
trans [OF uint-word-ariths(2) mod-add-if-z, simplified, standard]

11.4 Definition of `uint_arith`

lemma *word-of-int-inverse*:
 $\text{word-of-int } r = a \implies 0 \leq r \implies r < 2^{\text{len-of TYPE('a)}} \implies$
 $\text{uint } (a::'a::\text{len0 word}) = r$
apply (*erule word-uint.Abs-inverse' [rotated]*)
apply (*simp add: uints-num*)
done

lemma *uint-split*:
fixes $x::'a::\text{len0 word}$
shows $P (\text{uint } x) =$
 $(\text{ALL } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \implies P \ i)$
apply (*fold word-int-case-def*)
apply (*auto dest!: word-of-int-inverse simp: int-word-uint int-mod-eq'*
split: word-int-split)
done

lemma *uint-split-asm*:
fixes $x::'a::\text{len0 word}$
shows $P (\text{uint } x) =$
 $(\sim (\text{EX } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ i))$
by (*auto dest!: word-of-int-inverse*
simp: int-word-uint int-mod-eq'
split: uint-split)

lemmas *uint-splits* = *uint-split uint-split-asm*

lemmas *uint-arith-simps* =
word-le-def word-less-alt
word-uint.Rep-inject [symmetric]
uint-sub-if' uint-plus-if'

lemma *power-False-cong*: $\text{False} \implies a \wedge b = c \wedge d$
by *auto*

ML \ll
fun uint-arith-ss-of ss =

```

ss addsimps @{\thms uint-arith-simps}
delsimps @{\thms word-uint.Rep-inject}
addsplits @{\thms split-if-asm}
addcongs @{\thms power-False-cong}

fun uint-arith-tacs ctxt =
  let fun arith-tac' n t = arith-tac ctxt n t handle COOPER => Seq.empty
  in
    [ CLASET' clarify-tac 1,
      SIMPSET' (full-simp-tac o uint-arith-ss-of) 1,
      ALLGOALS (full-simp-tac (HOL-ss addsplits @{\thms uint-splits}
                             addcongs @{\thms power-False-cong})),
      rewrite-goals-tac @{\thms word-size},
      ALLGOALS (fn n => REPEAT (resolve-tac [allI, impI] n) THEN
                          REPEAT (etac conjE n) THEN
                          REPEAT (dtac @{\thm word-of-int-inverse} n
                                THEN atac n
                                THEN atac n)),
      TRYALL arith-tac' ]
  end

fun uint-arith-tac ctxt = SELECT-GOAL (EVERY (uint-arith-tacs ctxt))
  >>

method-setup uint-arith =
  Method.ctxt-args (fn ctxt => Method.SIMPLE-METHOD (uint-arith-tac ctxt 1))

  solving word arithmetic via integers and arith

11.5 More on overflows and monotonicity

lemma no-plus-overflow-uint-size:
  ((x :: 'a :: len0 word) <= x + y) = (uint x + uint y < 2 ^ size x)
  unfolding word-size by uint-arith

lemmas no-olen-add = no-plus-overflow-uint-size [unfolded word-size]

lemma no-ulen-sub: ((x :: 'a :: len0 word) >= x - y) = (uint y <= uint x)
  by uint-arith

lemma no-olen-add':
  fixes x :: 'a::len0 word
  shows (x ≤ y + x) = (uint y + uint x < 2 ^ len-of TYPE('a))
  by (simp add: word-add-ac add-ac no-olen-add)

lemmas olen-add-quiv = trans [OF no-olen-add no-olen-add' [symmetric], stan-
  dard]

lemmas uint-plus-simple-iff = trans [OF no-olen-add uint-add-lem, standard]

```

lemmas *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1, standard*]

lemmas *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem, standard*]

lemmas *uint-minus-simple-alt* = *uint-sub-lem* [*folded word-le-def*]

lemmas *word-sub-le-iff* = *no-ulen-sub* [*folded word-le-def*]

lemmas *word-sub-le* = *word-sub-le-iff* [*THEN iffD2, standard*]

lemma *word-less-sub1*:

$(x :: 'a :: \text{len } \text{word}) \sim = 0 \implies (1 < x) = (0 < x - 1)$

by *uint-arith*

lemma *word-le-sub1*:

$(x :: 'a :: \text{len } \text{word}) \sim = 0 \implies (1 \leq x) = (0 \leq x - 1)$

by *uint-arith*

lemma *sub-wrap-lt*:

$((x :: 'a :: \text{len } 0 \text{ word}) < x - z) = (x < z)$

by *uint-arith*

lemma *sub-wrap*:

$((x :: 'a :: \text{len } 0 \text{ word}) \leq x - z) = (z = 0 \mid x < z)$

by *uint-arith*

lemma *plus-minus-not-NULL-ab*:

$(x :: 'a :: \text{len } 0 \text{ word}) \leq ab - c \implies c \leq ab \implies c \sim = 0 \implies x + c \sim = 0$

by *uint-arith*

lemma *plus-minus-no-overflow-ab*:

$(x :: 'a :: \text{len } 0 \text{ word}) \leq ab - c \implies c \leq ab \implies x \leq x + c$

by *uint-arith*

lemma *le-minus'*:

$(a :: 'a :: \text{len } 0 \text{ word}) + c \leq b \implies a \leq a + c \implies c \leq b - a$

by *uint-arith*

lemma *le-plus'*:

$(a :: 'a :: \text{len } 0 \text{ word}) \leq b \implies c \leq b - a \implies a + c \leq b$

by *uint-arith*

lemmas *le-plus* = *le-plus'* [*rotated*]

lemmas *le-minus* = *leD* [*THEN thin-rl, THEN le-minus', standard*]

lemma *word-plus-mono-right*:

$(y :: 'a :: \text{len } 0 \text{ word}) \leq z \implies x \leq x + z \implies x + y \leq x + z$

by *uint-arith*

lemma *word-less-minus-cancel*:

$y - x < z - x \implies x \leq z \implies (y :: 'a :: \text{len } 0 \text{ word}) < z$

by *uint-arith*

lemma *word-less-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) < z \implies x \leq y \implies y - x < z - x$
by *uint-arith*

lemma *word-less-minus-mono*:

$a < c \implies d < b \implies a - b < a \implies c - d < c$
 $\implies a - b < c - (d :: 'a :: \text{len word})$
by *uint-arith*

lemma *word-le-minus-cancel*:

$y - x \leq z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) \leq z$
by *uint-arith*

lemma *word-le-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) \leq z \implies x \leq y \implies y - x \leq z - x$
by *uint-arith*

lemma *word-le-minus-mono*:

$a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c$
 $\implies a - b \leq c - (d :: 'a :: \text{len word})$
by *uint-arith*

lemma *plus-le-left-cancel-wrap*:

$(x :: 'a :: \text{len0 word}) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$
by *uint-arith*

lemma *plus-le-left-cancel-nowrap*:

$(x :: 'a :: \text{len0 word}) \leq x + y' \implies x \leq x + y \implies$
 $(x + y' < x + y) = (y' < y)$
by *uint-arith*

lemma *word-plus-mono-right2*:

$(a :: 'a :: \text{len0 word}) \leq a + b \implies c \leq b \implies a \leq a + c$
by *uint-arith*

lemma *word-less-add-right*:

$(x :: 'a :: \text{len0 word}) < y - z \implies z \leq y \implies x + z < y$
by *uint-arith*

lemma *word-less-sub-right*:

$(x :: 'a :: \text{len0 word}) < y + z \implies y \leq x \implies x - y < z$
by *uint-arith*

lemma *word-le-plus-either*:

$(x :: 'a :: \text{len0 word}) \leq y \mid x \leq z \implies y \leq y + z \implies x \leq y + z$
by *uint-arith*

lemma *word-less-nowrapI*:

$(x :: 'a :: \text{len0 word}) < z - k \implies k <= z \implies 0 < k \implies x < x + k$
by *uint-arith*

lemma *inc-le*: $(i :: 'a :: \text{len word}) < m \implies i + 1 <= m$

by *uint-arith*

lemma *inc-i*:

$(1 :: 'a :: \text{len word}) <= i \implies i < m \implies 1 <= (i + 1) \ \& \ i + 1 <= m$
by *uint-arith*

lemma *udvd-incr-lem*:

$up < uq \implies up = ua + n * \text{uint } K \implies$
 $uq = ua + n' * \text{uint } K \implies up + \text{uint } K <= uq$
apply *clarsimp*
apply (*drule less-le-mult*)
apply *safe*
done

lemma *udvd-incr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p + K <= q$
apply (*unfold word-less-alt word-le-def*)
apply (*drule (2) udvd-incr-lem*)
apply (*erule uint-add-le [THEN order-trans]*)
done

lemma *udvd-decr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p <= q - K$
apply (*unfold word-less-alt word-le-def*)
apply (*drule (2) udvd-incr-lem*)
apply (*drule le-diff-eq [THEN iffD2]*)
apply (*erule order-trans*)
apply (*rule uint-sub-ge*)
done

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [**where** *ua=0, simplified*]

lemmas *udvd-incr0* = *udvd-incr'* [**where** *ua=0, simplified*]

lemmas *udvd-decr0* = *udvd-decr'* [**where** *ua=0, simplified*]

lemma *udvd-minus-le'*:

$xy < k \implies z \text{ udvd } xy \implies z \text{ udvd } k \implies xy <= k - z$
apply (*unfold udvd-def*)
apply *clarify*
apply (*erule (2) udvd-decr0*)
done

lemma *udvd-incr2-K*:

```

  p < a + s ==> a <= a + s ==> K udvd s ==> K udvd p - a ==> a <=
p ==>
  0 < K ==> p <= p + K & p + K <= a + s
apply (unfold udvd-def)
apply clarify
apply (simp add: uint-arith-simps split: split-if-asm)
prefer 2
apply (insert uint-range' [of s])[1]
apply arith
apply (drule add-commute [THEN xtr1])
apply (simp add: diff-less-eq [symmetric])
apply (drule less-le-mult)
apply arith
apply simp
done

```

lemma *word-succ-rbl*:

```

to-bl w = bl ==> to-bl (word-succ w) = (rev (rbl-succ (rev bl)))
apply (unfold word-succ-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-succ)
done

```

lemma *word-pred-rbl*:

```

to-bl w = bl ==> to-bl (word-pred w) = (rev (rbl-pred (rev bl)))
apply (unfold word-pred-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-pred)
done

```

lemma *word-add-rbl*:

```

to-bl v = vbl ==> to-bl w = wbl ==>
  to-bl (v + w) = (rev (rbl-add (rev vbl) (rev wbl)))
apply (unfold word-add-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-add)
done

```

lemma *word-mult-rbl*:

```

to-bl v = vbl ==> to-bl w = wbl ==>
  to-bl (v * w) = (rev (rbl-mult (rev vbl) (rev wbl)))
apply (unfold word-mult-def)
apply clarify
apply (simp add: to-bl-of-bin)
apply (simp add: to-bl-def rbl-mult)

```

done

lemma *rtb-rbl-ariths*:

$rev (to-bl w) = ys \implies rev (to-bl (word-succ w)) = rbl-succ ys$

$rev (to-bl w) = ys \implies rev (to-bl (word-pred w)) = rbl-pred ys$

$[| rev (to-bl v) = ys; rev (to-bl w) = xs |]$
 $\implies rev (to-bl (v * w)) = rbl-mult ys xs$

$[| rev (to-bl v) = ys; rev (to-bl w) = xs |]$
 $\implies rev (to-bl (v + w)) = rbl-add ys xs$

by (*auto simp: rev-swap [symmetric] word-succ-rbl*
word-pred-rbl word-mult-rbl word-add-rbl)

11.6 Arithmetic type class instantiations

instance *word* :: (*len0*) *comm-monoid-add* ..

instance *word* :: (*len0*) *comm-monoid-mult*

apply (*intro-classes*)

apply (*simp add: word-mult-commute*)

apply (*simp add: word-mult-1*)

done

instance *word* :: (*len0*) *comm-semiring*

by (*intro-classes*) (*simp add : word-left-distrib*)

instance *word* :: (*len0*) *ab-group-add* ..

instance *word* :: (*len0*) *comm-ring* ..

instance *word* :: (*len*) *comm-semiring-1*

by (*intro-classes*) (*simp add: lenw1-zero-neq-one*)

instance *word* :: (*len*) *comm-ring-1* ..

instance *word* :: (*len0*) *comm-semiring-0* ..

instance *word* :: (*len0*) *order* ..

instance *word* :: (*len*) *recpower*

by (*intro-classes*) (*simp-all add: word-pow*)

lemma *zero-bintrunc*:

iszero (number-of x :: 'a :: len word) =

(bintrunc (len-of TYPE('a)) x = Numeral.Pls)

apply (*unfold iszero-def word-0-wi word-no-wi*)

```

apply (rule word-ubin.norm-eq-iff [symmetric, THEN trans])
apply (simp add : Pls-def [symmetric])
done

lemmas word-le-0-iff [simp] =
  word-zero-le [THEN leD, THEN linorder-antisym-conv1]

lemma word-of-nat: of-nat n = word-of-int (int n)
  by (induct n) (auto simp add : word-of-int-hom-syms)

lemma word-of-int: of-int = word-of-int
  apply (rule ext)
  apply (unfold of-int-def)
  apply (rule contentsI)
  apply safe
  apply (simp-all add: word-of-nat word-of-int-homs)
  defer
  apply (rule Rep-Integ-ne [THEN nonemptyE])
  apply (rule beXI)
  prefer 2
  apply assumption
  apply (auto simp add: RI-eq-diff)
done

lemma word-of-int-nat:
  0 <= x ==> word-of-int x = of-nat (nat x)
  by (simp add: of-nat-nat word-of-int)

lemma word-number-of-eq:
  number-of w = (of-int w :: 'a :: len word)
  unfolding word-number-of-def word-of-int by auto

instance word :: (len) number-ring
  by (intro-classes) (simp add : word-number-of-eq)

lemma iszero-word-no [simp] :
  iszero (number-of bin :: 'a :: len word) =
    iszero (number-of (bintrunc (len-of TYPE('a)) bin) :: int)
  apply (simp add: zero-bintrunc number-of-is-id)
  apply (unfold iszero-def Pls-def)
  apply (rule refl)
done

```

11.7 Word and nat

```

lemma td-ext-unat':
  n = len-of TYPE ('a :: len) ==>
  td-ext (unat :: 'a word => nat) of-nat
  (unats n) (%i. i mod 2 ^ n)

```

```

apply (unfold td-ext-def' unat-def word-of-nat unats-uints)
apply (auto intro!: imageI simp add : word-of-int-hom-syms)
apply (erule word-uint.Abs-inverse [THEN arg-cong])
apply (simp add: int-word-uint nat-mod-distrib nat-power-eq)
done

```

```

lemmas td-ext-unat = refl [THEN td-ext-unat]
lemmas unat-of-nat = td-ext-unat [THEN td-ext.eq-norm, standard]

```

```

interpretation word-unat:
  td-ext [unat::'a::len word ==> nat
          of-nat
          unats (len-of TYPE('a::len))
          %i. i mod 2 ^ len-of TYPE('a::len)]
by (rule td-ext-unat)

```

```

lemmas td-unat = word-unat.td-thm

```

```

lemmas unat-lt2p [iff] = word-unat.Rep [unfolded unats-def mem-Collect-eq]

```

```

lemma unat-le: y <= unat (z :: 'a :: len word) ==> y : unats (len-of TYPE
('a))
apply (unfold unats-def)
apply clarsimp
apply (rule xtrans, rule unat-lt2p, assumption)
done

```

```

lemma word-nchotomy:
  ALL w. EX n. (w :: 'a :: len word) = of-nat n & n < 2 ^ len-of TYPE ('a)
apply (rule allI)
apply (rule word-unat.Abs-cases)
apply (unfold unats-def)
apply auto
done

```

```

lemma of-nat-eq:
  fixes w :: 'a::len word
  shows (of-nat n = w) = (∃ q. n = unat w + q * 2 ^ len-of TYPE('a))
apply (rule trans)
apply (rule word-unat.inverse-norm)
apply (rule iffI)
apply (rule mod-eqD)
apply simp
apply clarsimp
done

```

```

lemma of-nat-eq-size:
  (of-nat n = w) = (EX q. n = unat w + q * 2 ^ size w)
unfolding word-size by (rule of-nat-eq)

```

lemma *of-nat-0*:

$(\text{of-nat } m = (0 :: 'a :: \text{len word})) = (\exists q. m = q * 2 \wedge \text{len-of TYPE}('a))$
by (*simp add: of-nat-eq*)

lemmas *of-nat-2p = mult-1* [*symmetric, THEN iffD2 [OF of-nat-0 exI]*]

lemma *of-nat-gt-0*: $\text{of-nat } k \sim = 0 \implies 0 < k$

by (*cases k*) *auto*

lemma *of-nat-neq-0*:

$0 < k \implies k < 2 \wedge \text{len-of TYPE} ('a :: \text{len}) \implies \text{of-nat } k \sim = (0 :: 'a \text{ word})$
by (*clarsimp simp add : of-nat-0*)

lemma *Abs-fnat-hom-add*:

$\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$
by *simp*

lemma *Abs-fnat-hom-mult*:

$\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$
by (*simp add: word-of-nat word-of-int-mult-hom zmult-int*)

lemma *Abs-fnat-hom-Suc*:

$\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$
by (*simp add: word-of-nat word-of-int-succ-hom add-ac*)

lemma *Abs-fnat-hom-0*: $(0 :: 'a :: \text{len word}) = \text{of-nat } 0$

by (*simp add: word-of-nat word-0-wi*)

lemma *Abs-fnat-hom-1*: $(1 :: 'a :: \text{len word}) = \text{of-nat } (\text{Suc } 0)$

by (*simp add: word-of-nat word-1-wi*)

lemmas *Abs-fnat-homs =*

Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc
Abs-fnat-hom-0 Abs-fnat-hom-1

lemma *word-arith-nat-add*:

$a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$
by *simp*

lemma *word-arith-nat-mult*:

$a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$
by (*simp add: Abs-fnat-hom-mult [symmetric]*)

lemma *word-arith-nat-Suc*:

$\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$
by (*subst Abs-fnat-hom-Suc [symmetric]*) *simp*

lemma *word-arith-nat-div*:

$a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div } \text{unat } b)$
by (*simp add: word-div-def word-of-nat zdiv-int uint-nat*)

lemma *word-arith-nat-mod*:
 $a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod } \text{unat } b)$
by (*simp add: word-mod-def word-of-nat zmod-int uint-nat*)

lemmas *word-arith-nat-defs* =
word-arith-nat-add word-arith-nat-mult
word-arith-nat-Suc Abs-fnat-hom-0
Abs-fnat-hom-1 word-arith-nat-div
word-arith-nat-mod

lemmas *unat-cong* = *arg-cong* [**where** $f = \text{unat}$]

lemmas *unat-word-ariths* = *word-arith-nat-defs*
[*THEN trans [OF unat-cong unat-of-nat], standard*]

lemmas *word-sub-less-iff* = *word-sub-le-iff*
[*simplified linorder-not-less [symmetric], simplified*]

lemma *unat-add-lem*:
 $(\text{unat } x + \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{unat } (x + y :: 'a :: \text{len word}) = \text{unat } x + \text{unat } y)$
unfolding *unat-word-ariths*
by (*auto intro!: trans [OF - nat-mod-lem]*)

lemma *unat-mult-lem*:
 $(\text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{unat } (x * y :: 'a :: \text{len word}) = \text{unat } x * \text{unat } y)$
unfolding *unat-word-ariths*
by (*auto intro!: trans [OF - nat-mod-lem]*)

lemmas *unat-plus-if'* =
trans [OF unat-word-ariths(1) mod-nat-add, simplified, standard]

lemma *le-no-overflow*:
 $x \leq b \implies a \leq a + b \implies x \leq a + (b :: 'a :: \text{len0 word})$
apply (*erule order-trans*)
apply (*erule olen-add-eqv [THEN iffD1]*)
done

lemmas *un-ui-le* = *trans*
[*OF word-le-nat-alt [symmetric]*
word-le-def [THEN meta-eq-to-obj-eq],
standard]

lemma *unat-sub-if-size*:
 $\text{unat } (x - y) = (\text{if } \text{unat } y \leq \text{unat } x$

```

    then unat x - unat y
    else unat x + 2 ^ size x - unat y)
apply (unfold word-size)
apply (simp add: un-ui-le)
apply (auto simp add: unat-def uint-sub-if')
apply (rule nat-diff-distrib)
  prefer 3
apply (simp add: group-simps)
apply (rule nat-diff-distrib [THEN trans])
  prefer 3
apply (subst nat-add-distrib)
  prefer 3
apply (simp add: nat-power-eq)
apply auto
apply uint-arith
done

```

lemmas unat-sub-if' = unat-sub-if-size [unfolded word-size]

```

lemma unat-div: unat ((x :: 'a :: len word) div y) = unat x div unat y
apply (simp add : unat-word-ariths)
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq'])
apply (rule div-le-dividend)
done

```

```

lemma unat-mod: unat ((x :: 'a :: len word) mod y) = unat x mod unat y
apply (clarsimp simp add : unat-word-ariths)
apply (cases unat y)
  prefer 2
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq'])
apply (rule mod-le-divisor)
apply auto
done

```

```

lemma uint-div: uint ((x :: 'a :: len word) div y) = uint x div uint y
  unfolding uint-nat by (simp add : unat-div zdiv-int)

```

```

lemma uint-mod: uint ((x :: 'a :: len word) mod y) = uint x mod uint y
  unfolding uint-nat by (simp add : unat-mod zmod-int)

```

11.8 Definition of unat.arith tactic

```

lemma unat-split:
  fixes x::'a::len word
  shows P (unat x) =
    (ALL n. of-nat n = x & n < 2^len-of TYPE('a) --> P n)
  by (auto simp: unat-of-nat)

```

```

lemma unat-split-asm:

```

fixes $x :: 'a :: \text{len word}$
shows $P (\text{unat } x) =$
 $(\sim (EX n. \text{of-nat } n = x \ \& \ n < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ n))$
by $(\text{auto simp: unat-of-nat})$

lemmas $\text{of-nat-inverse} =$
 $\text{word-unat.Abs-inverse'}$ [*rotated, unfolded unats-def, simplified*]

lemmas $\text{unat-splits} = \text{unat-split unat-split-asm}$

lemmas $\text{unat-arith-simps} =$
 $\text{word-le-nat-alt word-less-nat-alt}$
 $\text{word-unat.Rep-inject}$ [*symmetric*]
 $\text{unat-sub-if' unat-plus-if' unat-div unat-mod}$

ML $\langle\langle$

```

fun unat-arith-ss-of ss =
  ss addsimps @{thms unat-arith-simps}
  delsimps @{thms word-unat.Rep-inject}
  addsplits @{thms split-if-asm}
  addcongs @{thms power-False-cong}

fun unat-arith-tacs ctxt =
  let fun arith-tac' n t = arith-tac ctxt n t handle COOPER => Seq.empty
  in
    [ CLASET' clarify-tac 1,
      SIMPSET' (full-simp-tac o unat-arith-ss-of) 1,
      ALLGOALS (full-simp-tac (HOL-ss addsplits @{thms unat-splits}
                             addcongs @{thms power-False-cong})),
      rewrite-goals-tac @{thms word-size},
      ALLGOALS (fn n => REPEAT (resolve-tac [allI, impI] n) THEN
                    REPEAT (etac conjE n) THEN
                    REPEAT (dtac @{thm of-nat-inverse} n THEN atac n)),
      TRYALL arith-tac' ]
  end

```

$\text{fun unat-arith-tac ctxt} = \text{SELECT-GOAL (EVERY (unat-arith-tacs ctxt))}$
 $\rangle\rangle$

method-setup $\text{unat-arith} =$
 $\text{Method.ctxt-args (fn ctxt => Method.SIMPLE-METHOD (unat-arith-tac ctxt 1))}$
solving word arithmetic via natural numbers and arith

lemma $\text{no-plus-overflow-unat-size:}$
 $((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2^{\text{size } x})$
unfolding word-size **by** unat-arith

lemma *unat-sub*: $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$
by *unat-arith*

lemmas *no-olen-add-nat* = *no-plus-overflow-unat-size* [*unfolded word-size*]

lemmas *unat-plus-simple* = *trans* [*OF no-olen-add-nat unat-add-lem, standard*]

lemma *word-div-mult*:

$(0 :: 'a :: \text{len word}) < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies$
 $x * y \text{ div } y = x$
apply *unat-arith*
apply *clarsimp*
apply (*subst unat-mult-lem [THEN iffD1]*)
apply *auto*
done

lemma *div-lt'*: $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies$

$\text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$
apply *unat-arith*
apply *clarsimp*
apply (*drule mult-le-mono1*)
apply (*erule order-le-less-trans*)
apply (*rule xtr7 [OF unat-lt2p div-mult-le]*)
done

lemmas *div-lt''* = *order-less-imp-le* [*THEN div-lt'*]

lemma *div-lt-mult*: $(i :: 'a :: \text{len word}) < k \text{ div } x \implies 0 < x \implies i * x < k$

apply (*frule div-lt'' [THEN unat-mult-lem [THEN iffD1]]*)
apply (*simp add: unat-arith-simps*)
apply (*drule (1) mult-less-mono1*)
apply (*erule order-less-le-trans*)
apply (*rule div-mult-le*)
done

lemma *div-le-mult*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$
apply (*frule div-lt' [THEN unat-mult-lem [THEN iffD1]]*)
apply (*simp add: unat-arith-simps*)
apply (*drule mult-le-mono1*)
apply (*erule order-trans*)
apply (*rule div-mult-le*)
done

lemma *div-lt-uint'*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2^{\text{len-of TYPE('a)}}$
apply (*unfold uint-nat*)
apply (*drule div-lt'*)
apply (*simp add: zmult-int zless-nat-eq-int-zless [symmetric]*)

```

    nat-power-eq)
done

lemmas div-lt-uint'' = order-less-imp-le [THEN div-lt-uint']

lemma word-le-exists':
  (x :: 'a :: len0 word) <= y ==>
    (EX z. y = x + z & uint x + uint z < 2 ^ len-of TYPE('a))
  apply (rule exI)
  apply (rule conjI)
  apply (rule zadd-diff-inverse)
  apply uint-arith
done

lemmas plus-minus-not-NULL = order-less-imp-le [THEN plus-minus-not-NULL-ab]

lemmas plus-minus-no-overflow =
  order-less-imp-le [THEN plus-minus-no-overflow-ab]

lemmas mcs = word-less-minus-cancel word-less-minus-mono-left
  word-le-minus-cancel word-le-minus-mono-left

lemmas word-l-diffs = mcs [where y = w + x, unfolded add-diff-cancel, standard]
lemmas word-diff-ls = mcs [where z = w + x, unfolded add-diff-cancel, standard]
lemmas word-plus-mcs = word-diff-ls
  [where y = v + x, unfolded add-diff-cancel, standard]

lemmas le-unat-uoI = unat-le [THEN word-unat.Abs-inverse]

lemmas thd = refl [THEN [2] split-div-lemma [THEN iffD2], THEN conjunct1]

lemma thd1:
  a div b * b ≤ (a::nat)
  using gt-or-eq-0 [of b]
  apply (rule disjE)
  apply (erule xtr4 [OF thd mult-commute])
  apply clarsimp
done

lemmas uno-simps [THEN le-unat-uoI, standard] =
  mod-le-divisor div-le-dividend thd1

lemma word-mod-div-equality:
  (n div b) * b + (n mod b) = (n :: 'a :: len word)
  apply (unfold word-less-nat-alt word-arith-nat-defs)
  apply (cut-tac y=unat b in gt-or-eq-0)
  apply (erule disjE)
  apply (simp add: mod-div-equality uno-simps)
  apply simp

```

done

lemma *word-div-mult-le*: $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$
apply (*unfold word-le-nat-alt word-arith-nat-defs*)
apply (*cut-tac y=unat b in gt-or-eq-0*)
apply (*erule disjE*)
apply (*simp add: div-mult-le uno-simps*)
apply *simp*
done

lemma *word-mod-less-divisor*: $0 < n \implies m \text{ mod } n < (n :: 'a :: \text{len word})$
apply (*simp only: word-less-nat-alt word-arith-nat-defs*)
apply (*clarsimp simp add : uno-simps*)
done

lemma *word-of-int-power-hom*:
 $\text{word-of-int } a \wedge n = (\text{word-of-int } (a \wedge n) :: 'a :: \text{len word})$
by (*induct n*) (*simp-all add : word-of-int-hom-syms power-Suc*)

lemma *word-arith-power-alt*:
 $a \wedge n = (\text{word-of-int } (\text{uint } a \wedge n) :: 'a :: \text{len word})$
by (*simp add : word-of-int-power-hom [symmetric]*)

lemma *of-bl-length-less*:
 $\text{length } x = k \implies k < \text{len-of TYPE('a)} \implies (\text{of-bl } x :: 'a :: \text{len word}) < 2 \wedge k$
apply (*unfold of-bl-no [unfolded word-number-of-def]*
word-less-alt word-number-of-alt)
apply *safe*
apply (*simp (no-asm) add: word-of-int-power-hom word-uint.eq-norm*
del: word-of-int-bin)
apply (*simp add: mod-pos-pos-trivial*)
apply (*subst mod-pos-pos-trivial*)
apply (*rule bl-to-bin-ge0*)
apply (*rule order-less-trans*)
apply (*rule bl-to-bin-lt2p*)
apply *simp*
apply (*rule bl-to-bin-lt2p*)
done

11.9 Cardinality, finiteness of set of words

lemmas *card-lessThan' = card-lessThan* [*unfolded lessThan-def*]

lemmas *card-eq = word-unat.Abs-inj-on* [*THEN card-image,*
unfolded word-unat.image, unfolded unats-def, standard]

lemmas *card-word = trans* [*OF card-eq card-lessThan', standard*]

lemma *finite-word-UNIV*: *finite* (*UNIV :: 'a :: len word set*)

```

apply (rule contrapos-np)
prefer 2
apply (erule card-infinite)
apply (simp add: card-word)
done

lemma card-word-size:
  card (UNIV :: 'a :: len word set) = (2 ^ size (x :: 'a word))
unfolding word-size by (rule card-word)

end

```

12 WordBitwise: Bitwise Operations on Words

```

theory WordBitwise imports WordArith begin

```

```

lemmas bin-log-bintrs = bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or

```

```

lemmas wils1 = bin-log-bintrs [THEN word-ubin.norm-eq-iff [THEN iffD1],
  folded word-ubin.eq-norm, THEN eq-reflection, standard]

```

```

lemmas word-log-binary-defs =
  word-and-def word-or-def word-xor-def

```

```

lemmas word-no-log-defs [simp] =
  word-not-def [where a=number-of a,
    unfolded word-no-wi wils1, folded word-no-wi, standard]
  word-log-binary-defs [where a=number-of a and b=number-of b,
    unfolded word-no-wi wils1, folded word-no-wi, standard]

```

```

lemmas word-wi-log-defs = word-no-log-defs [unfolded word-no-wi]

```

```

lemma uint-or: uint (x OR y) = (uint x) OR (uint y)
by (simp add: word-or-def word-no-wi [symmetric] number-of-is-id
  bin-trunc-ao(2) [symmetric])

```

```

lemma uint-and: uint (x AND y) = (uint x) AND (uint y)
by (simp add: word-and-def number-of-is-id word-no-wi [symmetric]
  bin-trunc-ao(1) [symmetric])

```

```

lemma word-ops-nth-size:
  n < size (x::'a::len0 word) ==>
  (x OR y) !! n = (x !! n | y !! n) &

```

```

(x AND y) !! n = (x !! n & y !! n) &
(x XOR y) !! n = (x !! n ~ = y !! n) &
(NOT x) !! n = (~ x !! n)
unfolding word-size word-no-wi word-test-bit-def word-log-defs
by (clarsimp simp add : word-ubin.eq-norm nth-bintr bin-nth-ops)

```

```

lemma word-ao-nth:
fixes x :: 'a::len0 word
shows (x OR y) !! n = (x !! n | y !! n) &
      (x AND y) !! n = (x !! n & y !! n)
apply (cases n < size x)
apply (drule-tac y = y in word-ops-nth-size)
apply simp
apply (simp add : test-bit-bin word-size)
done

```

```

lemmas bwsimps =
  word-of-int-homs(2)
  word-0-wi-Pls
  word-m1-wi-Min
  word-wi-log-defs

```

```

lemma word-bw-assocs:
fixes x :: 'a::len0 word
shows
  (x AND y) AND z = x AND y AND z
  (x OR y) OR z = x OR y OR z
  (x XOR y) XOR z = x XOR y XOR z
using word-of-int-Ex [where x=x]
      word-of-int-Ex [where x=y]
      word-of-int-Ex [where x=z]
by (auto simp: bwsimps bbw-assocs)

```

```

lemma word-bw-comms:
fixes x :: 'a::len0 word
shows
  x AND y = y AND x
  x OR y = y OR x
  x XOR y = y XOR x
using word-of-int-Ex [where x=x]
      word-of-int-Ex [where x=y]
by (auto simp: bwsimps bin-ops-comm)

```

```

lemma word-bw-lcs:
fixes x :: 'a::len0 word
shows
  y AND x AND z = x AND y AND z

```

$y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
using *word-of-int-Ex* [**where** $x=x$]
 word-of-int-Ex [**where** $x=y$]
 word-of-int-Ex [**where** $x=z$]
by (*auto simp: bwsimps*)

lemma *word-log-esimps* [*simp*]:
fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$
 $0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$
using *word-of-int-Ex* [**where** $x=x$]
by (*auto simp: bwsimps*)

lemma *word-not-dist*:
fixes $x :: 'a::len0 \text{ word}$
shows
 $\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$
 $\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$
using *word-of-int-Ex* [**where** $x=x$]
 word-of-int-Ex [**where** $x=y$]
by (*auto simp: bwsimps bbw-not-dist*)

lemma *word-bw-same*:
fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } x = x$
 $x \text{ OR } x = x$
 $x \text{ XOR } x = 0$
using *word-of-int-Ex* [**where** $x=x$]
by (*auto simp: bwsimps*)

lemma *word-ao-absorbs* [*simp*]:
fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } (y \text{ OR } x) = x$
 $x \text{ OR } y \text{ AND } x = x$
 $x \text{ AND } (x \text{ OR } y) = x$

$y \text{ AND } x \text{ OR } x = x$
 $(y \text{ OR } x) \text{ AND } x = x$
 $x \text{ OR } x \text{ AND } y = x$
 $(x \text{ OR } y) \text{ AND } x = x$
 $x \text{ AND } y \text{ OR } x = x$
using *word-of-int-Ex* [**where** $x=x$]
 word-of-int-Ex [**where** $x=y$]
by (*auto simp: bwsimps*)

lemma *word-not-not* [*simp*]:
 $\text{NOT NOT } (x::'a::\text{len0 word}) = x$
using *word-of-int-Ex* [**where** $x=x$]
by (*auto simp: bwsimps*)

lemma *word-ao-dist*:
fixes $x :: 'a::\text{len0 word}$
shows $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$
using *word-of-int-Ex* [**where** $x=x$]
 word-of-int-Ex [**where** $x=y$]
 word-of-int-Ex [**where** $x=z$]
by (*auto simp: bwsimps bbw-ao-dist simp del: bin-ops-comm*)

lemma *word-oa-dist*:
fixes $x :: 'a::\text{len0 word}$
shows $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
using *word-of-int-Ex* [**where** $x=x$]
 word-of-int-Ex [**where** $x=y$]
 word-of-int-Ex [**where** $x=z$]
by (*auto simp: bwsimps bbw-oa-dist simp del: bin-ops-comm*)

lemma *word-add-not* [*simp*]:
fixes $x :: 'a::\text{len0 word}$
shows $x + \text{NOT } x = -1$
using *word-of-int-Ex* [**where** $x=x$]
by (*auto simp: bwsimps bin-add-not*)

lemma *word-plus-and-or* [*simp*]:
fixes $x :: 'a::\text{len0 word}$
shows $(x \text{ AND } y) + (x \text{ OR } y) = x + y$
using *word-of-int-Ex* [**where** $x=x$]
 word-of-int-Ex [**where** $x=y$]
by (*auto simp: bwsimps plus-and-or*)

lemma *lea0*:
fixes $x :: 'a::\text{len0 word}$
shows $(w = (x \text{ OR } y)) ==> (y = (w \text{ AND } y))$ **by** *auto*

lemma *lea0*:
fixes $x' :: 'a::\text{len0 word}$
shows $(w' = (x' \text{ AND } y')) ==> (x' = (x' \text{ OR } w'))$ **by** *auto*

```

lemmas word-ao-equiv = leao [COMP leoa [COMP iffI]]

lemma le-word-or2: x <= x OR (y::'a::len0 word)
  unfolding word-le-def uint-or
  by (auto intro: le-int-or)

lemmas le-word-or1 = xtr3 [OF word-bw-comms (2) le-word-or2, standard]
lemmas word-and-le1 =
  xtr3 [OF word-ao-absorbs (4) [symmetric] le-word-or2, standard]
lemmas word-and-le2 =
  xtr3 [OF word-ao-absorbs (8) [symmetric] le-word-or2, standard]

lemma bl-word-not: to-bl (NOT w) = map Not (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-not-bin number-of-is-id word-no-wi [symmetric])

lemma bl-word-xor: to-bl (v XOR w) = app2 op ~ = (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs bl-xor-bin
  by (simp add: number-of-is-id word-no-wi [symmetric])

lemma bl-word-or: to-bl (v OR w) = app2 op | (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-or-bin number-of-is-id word-no-wi [symmetric])

lemma bl-word-and: to-bl (v AND w) = app2 op & (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-and-bin number-of-is-id word-no-wi [symmetric])

lemma word-lsb-alt: lsb (w::'a::len0 word) = test-bit w 0
  by (auto simp: word-test-bit-def word-lsb-def)

lemma word-lsb-1-0: lsb (1::'a::len word) & ~ lsb (0::'b::len0 word)
  unfolding word-lsb-def word-1-no word-0-no by auto

lemma word-lsb-last: lsb (w::'a::len word) = last (to-bl w)
  apply (unfold word-lsb-def uint-bl bin-to-bl-def)
  apply (rule-tac bin=uint w in bin-exhaust)
  apply (cases size w)
  apply auto
  apply (auto simp add: bin-to-bl-aux-alt)
  done

lemma word-lsb-int: lsb w = (uint w mod 2 = 1)
  unfolding word-lsb-def bin-last-mod by auto

lemma word-msb-sint: msb w = (sint w < 0)
  unfolding word-msb-def
  by (simp add : sign-Min-lt-0 number-of-is-id)

```

lemma *word-msb-no'*:

$w = \text{number-of bin} \implies \text{msb } (w::'a::\text{len word}) = \text{bin-nth bin } (\text{size } w - 1)$
unfolding *word-msb-def word-number-of-def*
by (*clarsimp simp add: word-sbin.eq-norm word-size bin-sign-lem*)

lemmas *word-msb-no = refl [THEN word-msb-no', unfolded word-size]*

lemma *word-msb-nth'*: $\text{msb } (w::'a::\text{len word}) = \text{bin-nth } (\text{uint } w) (\text{size } w - 1)$

apply (*unfold word-size*)
apply (*rule trans [OF - word-msb-no]*)
apply (*simp add : word-number-of-def*)
done

lemmas *word-msb-nth = word-msb-nth' [unfolded word-size]*

lemma *word-msb-alt*: $\text{msb } (w::'a::\text{len word}) = \text{hd } (\text{to-bl } w)$

apply (*unfold word-msb-nth uint-bl*)
apply (*subst hd-conv-nth*)
apply (*rule length-greater-0-conv [THEN iffD1]*)
apply *simp*
apply (*simp add : nth-bin-to-bl word-size*)
done

lemma *word-set-nth*:

$\text{set-bit } w \ n \ (\text{test-bit } w \ n) = (w::'a::\text{len0 word})$
unfolding *word-test-bit-def word-set-bit-def* **by** *auto*

lemma *bin-nth-uint'*:

$\text{bin-nth } (\text{uint } w) \ n = (\text{rev } (\text{bin-to-bl } (\text{size } w) (\text{uint } w))) \ ! \ n \ \& \ n < \text{size } w$
apply (*unfold word-size*)
apply (*safe elim!: bin-nth-uint-imp*)
apply (*frule bin-nth-uint-imp*)
apply (*fast dest!: bin-nth-bl*)
done

lemmas *bin-nth-uint = bin-nth-uint' [unfolded word-size]*

lemma *test-bit-bl*: $w \ ! \ n = (\text{rev } (\text{to-bl } w)) \ ! \ n \ \& \ n < \text{size } w$

unfolding *to-bl-def word-test-bit-def word-size*
by (*rule bin-nth-uint*)

lemma *to-bl-nth*: $n < \text{size } w \implies \text{to-bl } w \ ! \ n = w \ ! \ (\text{size } w - \text{Suc } n)$

apply (*unfold test-bit-bl*)
apply *clarsimp*
apply (*rule trans*)
apply (*rule nth-rev-alt*)
apply (*auto simp add: word-size*)
done

lemma *test-bit-set*:

fixes $w :: 'a::len0$ word
shows $(set\text{-}bit\ w\ n\ x) !! n = (n < size\ w \ \&\ x)$
unfolding *word-size word-test-bit-def word-set-bit-def*
by $(clarsimp\ simp\ add : word\text{-}ubin.eq\text{-}norm\ nth\text{-}bintr)$

lemma *test-bit-set-gen*:

fixes $w :: 'a::len0$ word
shows $test\text{-}bit\ (set\text{-}bit\ w\ n\ x)\ m =$
 $(if\ m = n\ then\ n < size\ w \ \&\ x\ else\ test\text{-}bit\ w\ m)$
apply $(unfold\ word\text{-}size\ word\text{-}test\text{-}bit\text{-}def\ word\text{-}set\text{-}bit\text{-}def)$
apply $(clarsimp\ simp\ add : word\text{-}ubin.eq\text{-}norm\ nth\text{-}bintr\ bin\text{-}nth\text{-}sc\text{-}gen)$
apply $(auto\ elim! : test\text{-}bit\text{-}size\ [unfolded\ word\text{-}size]$
 $simp\ add : word\text{-}test\text{-}bit\text{-}def\ [symmetric])$
done

lemma *of-bl-rep-False*: $of\text{-}bl\ (replicate\ n\ False\ @\ bs) = of\text{-}bl\ bs$

unfolding *of-bl-def bl-to-bin-rep-F* **by** *auto*

lemma *msb-nth'*:

fixes $w :: 'a::len$ word
shows $msb\ w = w !! (size\ w - 1)$
unfolding *word-msb-nth' word-test-bit-def* **by** *simp*

lemmas $msb\text{-}nth = msb\text{-}nth'\ [unfolded\ word\text{-}size]$

lemmas $msb0 = len\text{-}gt\ 0\ [THEN\ diff\text{-}Suc\text{-}less,\ THEN$

$word\text{-}ops\text{-}nth\text{-}size\ [unfolded\ word\text{-}size],\ standard]$

lemmas $msb1 = msb0\ [where\ i = 0]$

lemmas $word\text{-}ops\text{-}msb = msb1\ [unfolded\ msb\text{-}nth\ [symmetric,\ unfolded\ One\text{-}nat\text{-}def]]$

lemmas $lsb0 = len\text{-}gt\ 0\ [THEN\ word\text{-}ops\text{-}nth\text{-}size\ [unfolded\ word\text{-}size],\ standard]$

lemmas $word\text{-}ops\text{-}lsb = lsb0\ [unfolded\ word\text{-}lsb\text{-}alt]$

lemma *td-ext-nth'*:

$n = size\ (w :: 'a::len0\ word) ==> ofn = set\text{-}bits ==> [w,\ ofn\ g] = l ==>$

$td\text{-}ext\ test\text{-}bit\ ofn\ \{f.\ ALL\ i.\ f\ i \text{---}> i < n\}\ (\%h\ i.\ h\ i \ \&\ i < n)$

apply $(unfold\ word\text{-}size\ td\text{-}ext\text{-}def')$

apply *safe*

apply $(rule\text{-}tac\ [3]\ ext)$

apply $(rule\text{-}tac\ [4]\ ext)$

apply $(unfold\ word\text{-}size\ of\text{-}nth\text{-}def\ test\text{-}bit\text{-}bl)$

apply *safe*

defer

apply $(clarsimp\ simp : word\text{-}bl.Abs\text{-}inverse)+$

apply $(rule\ word\text{-}bl.Rep\text{-}inverse')$

apply $(rule\ sym\ [THEN\ trans])$

apply $(rule\ bl\text{-}of\text{-}nth\text{-}nth)$

```

apply simp
apply (rule bl-of-nth-inj)
apply (clarsimp simp add : test-bit-bl word-size)
done

```

lemmas *td-ext-nth = td-ext-nth'* [*OF refl refl refl, unfolded word-size*]

interpretation *test-bit*:

```

td-ext [op !! :: 'a::len0 word => nat => bool
  set-bits
  {f. ∀ i. f i → i < len-of TYPE('a::len0)}
  (λh i. h i ∧ i < len-of TYPE('a::len0))]
by (rule td-ext-nth)

```

declare *test-bit.Rep'* [*simp del*]

declare *test-bit.Rep'* [*rule del*]

lemmas *td-nth = test-bit.td-thm*

lemma *word-set-set-same*:

```

fixes w :: 'a::len0 word
shows set-bit (set-bit w n x) n y = set-bit w n y
by (rule word-eqI) (simp add : test-bit-set-gen word-size)

```

lemma *word-set-set-diff*:

```

fixes w :: 'a::len0 word
assumes m ~ n
shows set-bit (set-bit w m x) n y = set-bit (set-bit w n y) m x
by (rule word-eqI) (clarsimp simp add : test-bit-set-gen word-size prems)

```

lemma *test-bit-no'*:

```

fixes w :: 'a::len0 word
shows w = number-of bin ==> test-bit w n = (n < size w & bin-nth bin n)
unfolding word-test-bit-def word-number-of-def word-size
by (simp add : nth-bintr [symmetric] word-ubin.eq-norm)

```

lemmas *test-bit-no =*

```

refl [THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard]

```

lemma *nth-0: ~ (0::'a::len0 word) !! n*

```

unfolding test-bit-no word-0-no by auto

```

lemma *nth-sint*:

```

fixes w :: 'a::len word
defines l ≡ len-of TYPE ('a)
shows bin-nth (sint w) n = (if n < l - 1 then w !! n else w !! (l - 1))
unfolding sint-uint l-def
by (clarsimp simp add: nth-sbintr word-test-bit-def [symmetric])

```

lemma *word-lsb-no*:

lsb (number-of bin :: 'a :: len word) = (bin-last bin = bit.B1)

unfolding *word-lsb-alt test-bit-no* **by** *auto*

lemma *word-set-no*:

set-bit (number-of bin::'a::len0 word) n b =

number-of (bin-sc n (if b then bit.B1 else bit.B0) bin)

apply (*unfold word-set-bit-def word-number-of-def [symmetric]*)

apply (*rule word-eqI*)

apply (*clarsimp simp: word-size bin-nth-sc-gen number-of-is-id
test-bit-no nth-bintr*)

done

lemmas *setBit-no = setBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],
simplified if-simps, THEN eq-reflection, standard]*

lemmas *clearBit-no = clearBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],
simplified if-simps, THEN eq-reflection, standard]*

lemma *to-bl-n1*:

to-bl (-1::'a::len0 word) = replicate (len-of TYPE ('a)) True

apply (*rule word-bl.Abs-inverse'*)

apply *simp*

apply (*rule word-eqI*)

apply (*clarsimp simp add: word-size test-bit-no*)

apply (*auto simp add: word-bl.Abs-inverse test-bit-bl word-size*)

done

lemma *word-msb-n1: msb (-1::'a::len word)*

unfolding *word-msb-alt word-msb-alt to-bl-n1* **by** *simp*

declare *word-set-set-same [simp] word-set-nth [simp]*

test-bit-no [simp] word-set-no [simp] nth-0 [simp]

setBit-no [simp] clearBit-no [simp]

word-lsb-no [simp] word-msb-no [simp] word-msb-n1 [simp] word-lsb-1-0 [simp]

lemma *word-set-nth-iff*:

(set-bit w n b = w) = (w !! n = b | n >= size (w::'a::len0 word))

apply (*rule iffI*)

apply (*rule disjCI*)

apply (*drule word-eqD*)

apply (*erule sym [THEN trans]*)

apply (*simp add: test-bit-set*)

apply (*erule disjE*)

apply *clarsimp*

apply (*rule word-eqI*)

apply (*clarsimp simp add : test-bit-set-gen*)

apply (*drule test-bit-size*)

apply *force*

done

```

lemma test-bit-2p':
  w = word-of-int (2 ^ n) ==>
    w !! m = (m = n & m < size (w :: 'a :: len word))
  unfolding word-test-bit-def word-size
  by (auto simp add: word-ubin.eq-norm nth-bintr nth-2p-bin)

lemmas test-bit-2p = refl [THEN test-bit-2p', unfolded word-size]

lemmas nth-w2p = test-bit-2p [unfolded of-int-number-of-eq
  word-of-int [symmetric] of-int-power]

lemma uint-2p:
  (0::'a::len word) < 2 ^ n ==> uint (2 ^ n::'a::len word) = 2 ^ n
  apply (unfold word-arith-power-alt)
  apply (case-tac len-of TYPE ('a))
  apply clarsimp
  apply (case-tac nat)
  apply clarsimp
  apply (case-tac n)
  apply (clarsimp simp add : word-1-wi [symmetric])
  apply (clarsimp simp add : word-0-wi [symmetric])
  apply (drule word-gt-0 [THEN iffD1])
  apply (safe intro!: word-eqI bin-nth-lem ext)
  apply (auto simp add: test-bit-2p nth-2p-bin word-test-bit-def [symmetric])
  done

lemma word-of-int-2p: (word-of-int (2 ^ n) :: 'a :: len word) = 2 ^ n
  apply (unfold word-arith-power-alt)
  apply (case-tac len-of TYPE ('a))
  apply clarsimp
  apply (case-tac nat)
  apply (rule word-ubin.norm-eq-iff [THEN iffD1])
  apply (rule box-equals)
  apply (rule-tac [2] bintr-ariths (1))+
  apply (clarsimp simp add : number-of-is-id)
  apply simp
  done

lemma bang-is-le: x !! m ==> 2 ^ m <= (x :: 'a :: len word)
  apply (rule xtr3)
  apply (rule-tac [2] y = x in le-word-or2)
  apply (rule word-eqI)
  apply (auto simp add: word-ao-nth nth-w2p word-size)
  done

lemma word-clr-le:
  fixes w :: 'a::len0 word
  shows w >= set-bit w n False

```

```

apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
apply simp
apply (rule order-trans)
  apply (rule bintr-bin-clr-le)
apply simp
done

lemma word-set-ge:
  fixes w :: 'a::len word
  shows w <= set-bit w n True
  apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
  apply simp
  apply (rule order-trans [OF - bintr-bin-set-ge])
  apply simp
  done

end

```

13 WordShift: Shifting, Rotating, and Splitting Words

theory WordShift **imports** WordBitwise **begin**

13.1 Bit shifting

```

lemma shiffl1-number [simp] :
  shiffl1 (number-of w) = number-of (w BIT bit.B0)
  apply (unfold shiffl1-def word-number-of-def)
  apply (simp add: word-ubin.norm-eq-iff [symmetric] word-ubin.eq-norm)
  apply (subst refl [THEN bintrunc-BIT-I, symmetric])
  apply (subst bintrunc-bintrunc-min)
  apply simp
  done

```

```

lemma shiffl1-0 [simp] : shiffl1 0 = 0
  unfolding word-0-no shiffl1-number by auto

```

```

lemmas shiffl1-def-u = shiffl1-def [folded word-number-of-def]

```

```

lemma shiffl1-def-s: shiffl1 w = number-of (sint w BIT bit.B0)
  by (rule trans [OF - shiffl1-number]) simp

```

```

lemma shiftr1-0 [simp] : shiftr1 0 = 0
  unfolding shiftr1-def
  by simp (simp add: word-0-wi)

```

```

lemma sshiftr1-0 [simp] : sshiftr1 0 = 0
  apply (unfold sshiftr1-def)
  apply simp
  apply (simp add : word-0-wi)
  done

lemma sshiftr1-n1 [simp] : sshiftr1 -1 = -1
  unfolding sshiftr1-def by auto

lemma shiftr1-0 [simp] : (0::'a::len0 word) << n = 0
  unfolding shiftr1-def by (induct n) auto

lemma shiftr1-0 [simp] : (0::'a::len0 word) >> n = 0
  unfolding shiftr1-def by (induct n) auto

lemma sshiftr1-0 [simp] : 0 >>> n = 0
  unfolding sshiftr1-def by (induct n) auto

lemma sshiftr1-n1 [simp] : -1 >>> n = -1
  unfolding sshiftr1-def by (induct n) auto

lemma nth-shiftr1: shiftr1 w !! n = (n < size w & n > 0 & w !! (n - 1))
  apply (unfold shiftr1-def word-test-bit-def)
  apply (simp add: nth-bintr word-ubin.eq-norm word-size)
  apply (cases n)
  apply auto
  done

lemma nth-shiftr1' [rule-format]:
  ALL n. ((w::'a::len0 word) << m) !! n = (n < size w & n >= m & w !! (n -
m))
  apply (unfold shiftr1-def)
  apply (induct m)
  apply (force elim!: test-bit-size)
  apply (clarsimp simp add : nth-shiftr1 word-size)
  apply arith
  done

lemmas nth-shiftr1 = nth-shiftr1' [unfolded word-size]

lemma nth-shiftr1: shiftr1 w !! n = w !! Suc n
  apply (unfold shiftr1-def word-test-bit-def)
  apply (simp add: nth-bintr word-ubin.eq-norm)
  apply safe
  apply (drule bin-nth.Suc [THEN iffD2, THEN bin-nth-uint-imp])
  apply simp
  done

lemma nth-shiftr:

```

```

 $\wedge n. ((w::'a::len0 \text{ word}) \gg m) !! n = w !! (n + m)$ 
apply (unfold shiftr-def)
apply (induct m)
apply (auto simp add : nth-shiftr1)
done

```

```

lemma uint-shiftr1: uint (shiftr1 w) = bin-rest (uint w)
apply (unfold shiftr1-def word-ubin.eq-norm bin-rest-trunc-i)
apply (subst bintr-uint [symmetric, OF order-refl])
apply (simp only : bintrunc-bintrunc-l)
apply simp
done

```

```

lemma nth-sshiftr1:
  sshiftr1 w !! n = (if n = size w - 1 then w !! n else w !! Suc n)
apply (unfold sshiftr1-def word-test-bit-def)
apply (simp add: nth-bintr word-ubin.eq-norm
            bin-nth.Suc [symmetric] word-size
            del: bin-nth.simps)
apply (simp add: nth-bintr uint-sint del : bin-nth.simps)
apply (auto simp add: bin-nth-sint)
done

```

```

lemma nth-sshiftr [rule-format] :
  ALL n. sshiftr w m !! n = (n < size w &
    (if n + m >= size w then w !! (size w - 1) else w !! (n + m)))
apply (unfold sshiftr-def)
apply (induct-tac m)
apply (simp add: test-bit-bl)
apply (clarsimp simp add: nth-sshiftr1 word-size)
apply safe
  apply arith
  apply arith
apply (erule thin-rl)
apply (case-tac n)
  apply safe
  apply simp
apply simp
apply (erule thin-rl)
apply (case-tac n)
  apply safe
  apply simp
apply simp
apply arith+
done

```

```

lemma shiftr1-div-2: uint (shiftr1 w) = uint w div 2

```

```

apply (unfold shiftr1-def bin-rest-div)
apply (rule word-uint.Abs-inverse)
apply (simp add: uints-num pos-imp-zdiv-nonneg-iff)
apply (rule xtr7)
prefer 2
apply (rule zdiv-le-dividend)
apply auto
done

```

```

lemma sshiftr1-div-2: sint (sshiftr1 w) = sint w div 2
apply (unfold sshiftr1-def bin-rest-div [symmetric])
apply (simp add: word-sbin.eq-norm)
apply (rule trans)
defer
apply (subst word-sbin.norm-Rep [symmetric])
apply (rule refl)
apply (subst word-sbin.norm-Rep [symmetric])
apply (unfold One-nat-def)
apply (rule sbintrunc-rest)
done

```

```

lemma shiftr-div-2n: uint (shiftr w n) = uint w div 2 ^ n
apply (unfold shiftr-def)
apply (induct n)
apply simp
apply (simp add: shiftr1-div-2 mult-commute
                zdiv-zmult2-eq [symmetric])
done

```

```

lemma sshiftr-div-2n: sint (sshiftr w n) = sint w div 2 ^ n
apply (unfold sshiftr-def)
apply (induct n)
apply simp
apply (simp add: sshiftr1-div-2 mult-commute
                zdiv-zmult2-eq [symmetric])
done

```

13.1.1 shift functions in terms of lists of bools

```

lemmas bshiftr1-no-bin [simp] =
  bshiftr1-def [where w=number-of w, unfolded to-bl-no-bin, standard]

```

```

lemma bshiftr1-bl: to-bl (bshiftr1 b w) = b # butlast (to-bl w)
unfolding bshiftr1-def by (rule word-bl.Abs-inverse) simp

```

```

lemma shiftr1-of-bl: shiftr1 (of-bl bl) = of-bl (bl @ [False])
unfolding uint-bl of-bl-no
by (simp add: bl-to-bin-aux-append bl-to-bin-def)

```

lemma *shiffl1-bl*: $\text{shiffl1 } (w :: 'a :: \text{len0 word}) = \text{of-bl } (\text{to-bl } w \text{ @ } [\text{False}])$

proof –

have $\text{shiffl1 } w = \text{shiffl1 } (\text{of-bl } (\text{to-bl } w))$ **by** *simp*

also have $\dots = \text{of-bl } (\text{to-bl } w \text{ @ } [\text{False}])$ **by** (*rule shiffl1-of-bl*)

finally show *?thesis* .

qed

lemma *bl-shiffl1*:

$\text{to-bl } (\text{shiffl1 } (w :: 'a :: \text{len word})) = \text{tl } (\text{to-bl } w) \text{ @ } [\text{False}]$

apply (*simp add: shiffl1-bl word-rep-drop drop-Suc drop-Cons'*)

apply (*fast intro!: Suc-leI*)

done

lemma *shiftr1-bl*: $\text{shiftr1 } w = \text{of-bl } (\text{butlast } (\text{to-bl } w))$

apply (*unfold shiftr1-def wint-bl of-bl-def*)

apply (*simp add: butlast-rest-bin word-size*)

apply (*simp add: bin-rest-trunc [symmetric, unfolded One-nat-def]*)

done

lemma *bl-shiftr1*:

$\text{to-bl } (\text{shiftr1 } (w :: 'a :: \text{len word})) = \text{False} \# \text{butlast } (\text{to-bl } w)$

unfolding *shiftr1-bl*

by (*simp add : word-rep-drop len-gt-0 [THEN Suc-leI]*)

lemma *shiffl1-rev*:

$\text{shiffl1 } (w :: 'a :: \text{len word}) = \text{word-reverse } (\text{shiftr1 } (\text{word-reverse } w))$

apply (*unfold word-reverse-def*)

apply (*rule word-bl.Rep-inverse' [symmetric]*)

apply (*simp add: bl-shiffl1 bl-shiftr1 word-bl.Abs-inverse*)

apply (*cases to-bl w*)

apply *auto*

done

lemma *shiffl-rev*:

$\text{shiffl } (w :: 'a :: \text{len word}) \ n = \text{word-reverse } (\text{shiftr } (\text{word-reverse } w) \ n)$

apply (*unfold shiffl-def shiftr-def*)

apply (*induct n*)

apply (*auto simp add : shiffl-rev*)

done

lemmas *rev-shiffl* =

shiffl-rev [**where** $w = \text{word-reverse } w$, *simplified, standard*]

lemmas *shiftr-rev* = *rev-shiffl* [*THEN word-rev-gal'*, *standard*]

lemmas *rev-shiftr* = *shiffl-rev* [*THEN word-rev-gal'*, *standard*]

lemma *bl-sshiftr1*:

```

to-bl (sshiftr1 (w :: 'a :: len word)) = hd (to-bl w) # butlast (to-bl w)
apply (unfold sshiftr1-def uint-bl word-size)
apply (simp add: butlast-rest-bin word-ubin.eq-norm)
apply (simp add: sint-uint)
apply (rule nth-equalityI)
  apply clarsimp
  apply clarsimp
apply (case-tac i)
  apply (simp-all add: hd-conv-nth length-0-conv [symmetric]
    nth-bin-to-bl bin-nth.Suc [symmetric]
    nth-sbintr
    del: bin-nth.Suc)
  apply force
apply (rule impI)
apply (rule-tac f = bin-nth (uint w) in arg-cong)
apply simp
done

```

lemma drop-shiftr:

```

drop n (to-bl ((w :: 'a :: len word) >> n)) = take (size w - n) (to-bl w)
apply (unfold shiftr-def)
apply (induct n)
  prefer 2
  apply (simp add: drop-Suc bl-shiftr1 butlast-drop [symmetric])
  apply (rule butlast-take [THEN trans])
apply (auto simp: word-size)
done

```

lemma drop-sshiftr:

```

drop n (to-bl ((w :: 'a :: len word) >>> n)) = take (size w - n) (to-bl w)
apply (unfold sshiftr-def)
apply (induct n)
  prefer 2
  apply (simp add: drop-Suc bl-sshiftr1 butlast-drop [symmetric])
  apply (rule butlast-take [THEN trans])
apply (auto simp: word-size)
done

```

lemma take-shiftr [rule-format] :

```

n <= size (w :: 'a :: len word) --> take n (to-bl (w >> n)) =
  replicate n False
apply (unfold shiftr-def)
apply (induct n)
  prefer 2
  apply (simp add: bl-shiftr1)
  apply (rule impI)
  apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

```

```

lemma take-sshiftr' [rule-format] :
  n <= size (w :: 'a :: len word) --> hd (to-bl (w >>> n)) = hd (to-bl w) &
  take n (to-bl (w >>> n)) = replicate n (hd (to-bl w))
apply (unfold sshiftr-def)
apply (induct n)
prefer 2
apply (simp add: bl-sshiftr1)
apply (rule impI)
apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

```

```

lemmas hd-sshiftr = take-sshiftr' [THEN conjunct1, standard]
lemmas take-sshiftr = take-sshiftr' [THEN conjunct2, standard]

```

```

lemma atd-lem: take n xs = t ==> drop n xs = d ==> xs = t @ d
by (auto intro: append-take-drop-id [symmetric])

```

```

lemmas bl-shiftr = atd-lem [OF take-shiftr drop-shiftr]
lemmas bl-sshiftr = atd-lem [OF take-sshiftr drop-sshiftr]

```

```

lemma shiftl-of-bl: of-bl bl << n = of-bl (bl @ replicate n False)
unfolding shiftl-def
by (induct n) (auto simp: shiftl1-of-bl replicate-app-Cons-same)

```

```

lemma shiftl-bl:
  (w::'a::len0 word) << (n::nat) = of-bl (to-bl w @ replicate n False)
proof –
  have w << n = of-bl (to-bl w) << n by simp
  also have ... = of-bl (to-bl w @ replicate n False) by (rule shiftl-of-bl)
  finally show ?thesis .
qed

```

```

lemmas shiftl-number [simp] = shiftl-def [where w=number-of w, standard]

```

```

lemma bl-shiftl:
  to-bl (w << n) = drop n (to-bl w) @ replicate (min (size w) n) False
by (simp add: shiftl-bl word-rep-drop word-size min-def)

```

```

lemma shiftl-zero-size:
  fixes x :: 'a::len0 word
  shows size x <= n ==> x << n = 0
apply (unfold word-size)
apply (rule word-eqI)
apply (clarsimp simp add: shiftl-bl word-size test-bit-of-bl nth-append)
done

```

```

lemma shiffl1-2t: shiffl1 (w :: 'a :: len word) = 2 * w
  apply (simp add: shiffl1-def-u)
  apply (simp only: double-number-of-BIT [symmetric])
  apply simp
  done

lemma shiffl1-p: shiffl1 (w :: 'a :: len word) = w + w
  apply (simp add: shiffl1-def-u)
  apply (simp only: double-number-of-BIT [symmetric])
  apply simp
  done

lemma shiffl-t2n: shiffl (w :: 'a :: len word) n = 2 ^ n * w
  unfolding shiffl-def
  by (induct n) (auto simp: shiffl1-2t power-Suc)

lemma shiftr1-bintr [simp]:
  (shiftr1 (number-of w) :: 'a :: len0 word) =
    number-of (bin-rest (bintrunc (len-of TYPE ('a)) w))
  unfolding shiftr1-def word-number-of-def
  by (simp add : word-ubin.eq-norm)

lemma sshiftr1-sbintr [simp] :
  (sshiftr1 (number-of w) :: 'a :: len word) =
    number-of (bin-rest (sbintrunc (len-of TYPE ('a) - 1) w))
  unfolding sshiftr1-def word-number-of-def
  by (simp add : word-sbin.eq-norm)

lemma shiftr-no':
  w = number-of bin ==>
  (w::'a::len0 word) >> n = number-of ((bin-rest ^ n) (bintrunc (size w) bin))
  apply clarsimp
  apply (rule word-eqI)
  apply (auto simp: nth-shiftr nth-rest-power-bin nth-bintr word-size)
  done

lemma sshiftr-no':
  w = number-of bin ==> w >>> n = number-of ((bin-rest ^ n)
    (sbintrunc (size w - 1) bin))
  apply clarsimp
  apply (rule word-eqI)
  apply (auto simp: nth-sshiftr nth-rest-power-bin nth-sbintr word-size)
  apply (subgoal-tac na + n = len-of TYPE('a) - Suc 0, simp, simp)+
  done

lemmas sshiftr-no [simp] =
  shiftr-no' [where w = number-of w, OF refl, unfolded word-size, standard]

```

lemmas *shiftr-no* [*simp*] =
shiftr-no' [**where** *w* = *number-of w*, *OF refl*, *unfolded word-size*, *standard*]

lemma *shiftr1-bl-of'*:
 $us = \text{shiftr1 } (of\text{-bl } bl) \implies \text{length } bl \leq \text{size } us \implies$
 $us = of\text{-bl } (\text{butlast } bl)$
by (*clarsimp simp: shiftr1-def of-bl-def word-size butlast-rest-bl2bin*
word-ubin.eq-norm trunc-bl2bin)

lemmas *shiftr1-bl-of* = *refl* [*THEN shiftr1-bl-of'*, *unfolded word-size*]

lemma *shiftr-bl-of'* [*rule-format*]:
 $us = of\text{-bl } bl \gg n \implies \text{length } bl \leq \text{size } us \implies$
 $us = of\text{-bl } (\text{take } (\text{length } bl - n) \text{ } bl)$
apply (*unfold shiftr-def*)
apply *hypsubst*
apply (*unfold word-size*)
apply (*induct n*)
apply *clarsimp*
apply *clarsimp*
apply (*subst shiftr1-bl-of*)
apply *simp*
apply (*simp add: butlast-take*)
done

lemmas *shiftr-bl-of* = *refl* [*THEN shiftr-bl-of'*, *unfolded word-size*]

lemmas *shiftr-bl* = *word-bl.Rep'* [*THEN eq-imp-le*, *THEN shiftr-bl-of*,
simplified word-size, *simplified*, *THEN eq-reflection*, *standard*]

lemma *msb-shift'*: $msb (w::'a::len \text{ word}) \leftrightarrow (w \gg (\text{size } w - 1)) \approx 0$
apply (*unfold shiftr-bl word-msb-alt*)
apply (*simp add: word-size Suc-le-eq take-Suc*)
apply (*cases hd (to-bl w)*)
apply (*auto simp: word-1-bl word-0-bl*
of-bl-rep-False [**where** *n=1 and bs=[]*, *simplified*])
done

lemmas *msb-shift* = *msb-shift'* [*unfolded word-size*]

lemma *align-lem-or* [*rule-format*]:
 $ALL \ x \ m. \ \text{length } x = n + m \implies \text{length } y = n + m \implies$
 $\text{drop } m \ x = \text{replicate } n \ \text{False} \implies \text{take } m \ y = \text{replicate } m \ \text{False} \implies$
 $\text{app2 } op \ | \ x \ y = \text{take } m \ x \ @ \ \text{drop } m \ y$
apply (*induct-tac y*)
apply *force*
apply *clarsimp*
apply (*case-tac x, force*)
apply (*case-tac m, auto*)

```

apply (drule sym)
apply auto
apply (induct-tac list, auto)
done

```

```

lemma align-lem-and [rule-format] :
  ALL x m. length x = n + m --> length y = n + m -->
    drop m x = replicate n False --> take m y = replicate m False -->
    app2 op & x y = replicate (n + m) False
apply (induct-tac y)
apply force
apply clarsimp
apply (case-tac x, force)
apply (case-tac m, auto)
apply (drule sym)
apply auto
apply (induct-tac list, auto)
done

```

```

lemma aligned-bl-add-size':
  size x - n = m ==> n <= size x ==> drop m (to-bl x) = replicate n False
==>
  take m (to-bl y) = replicate m False ==>
  to-bl (x + y) = take m (to-bl x) @ drop m (to-bl y)
apply (subgoal-tac x AND y = 0)
prefer 2
apply (rule word-bl.Rep-eqD)
apply (simp add: bl-word-and to-bl-0)
apply (rule align-lem-and [THEN trans])
  apply (simp-all add: word-size)[5]
apply (rule-tac f = %n. replicate n False in arg-cong)
apply simp
apply (subst word-plus-and-or [symmetric])
apply (simp add : bl-word-or)
apply (rule align-lem-or)
  apply (simp-all add: word-size)
done

```

```

lemmas aligned-bl-add-size = refl [THEN aligned-bl-add-size']

```

13.1.2 Mask

```

lemma nth-mask': m = mask n ==> test-bit m i = (i < n & i < size m)
apply (unfold mask-def test-bit-bl)
apply (simp only: word-1-bl [symmetric] shiftl-of-bl)
apply (clarsimp simp add: word-size)
apply (simp only: of-bl-no mask-lem number-of-succ add-diff-cancel2)
apply (fold of-bl-no)
apply (simp add: word-1-bl)

```

```

apply (rule test-bit-of-bl [THEN trans, unfolded test-bit-bl word-size])
apply auto
done

lemmas nth-mask [simp] = refl [THEN nth-mask']

lemma mask-bl: mask n = of-bl (replicate n True)
  by (auto simp add : test-bit-of-bl word-size intro: word-eqI)

lemma mask-bin: mask n = number-of (bintrunc n Numeral.Min)
  by (auto simp add: nth-bintr word-size intro: word-eqI)

lemma and-mask-bintr: w AND mask n = number-of (bintrunc n (uint w))
  apply (rule word-eqI)
  apply (simp add: nth-bintr word-size word-ops-nth-size)
  apply (auto simp add: test-bit-bin)
done

lemma and-mask-no: number-of i AND mask n = number-of (bintrunc n i)
  by (auto simp add : nth-bintr word-size word-ops-nth-size intro: word-eqI)

lemmas and-mask-wi = and-mask-no [unfolded word-number-of-def]

lemma bl-and-mask:
  to-bl (w AND mask n :: 'a :: len word) =
    replicate (len-of TYPE('a) - n) False @
    drop (len-of TYPE('a) - n) (to-bl w)
  apply (rule nth-equalityI)
  apply simp
  apply (clarsimp simp add: to-bl-nth word-size)
  apply (simp add: word-size word-ops-nth-size)
  apply (auto simp add: word-size test-bit-bl nth-append nth-rev)
done

lemmas and-mask-mod-2p =
  and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]

lemma and-mask-lt-2p: uint (w AND mask n) < 2 ^ n
  apply (simp add : and-mask-bintr no-bintr-alt)
  apply (rule xtr8)
  prefer 2
  apply (rule pos-mod-bound)
  apply auto
done

lemmas eq-mod-iff = trans [symmetric, OF int-mod-lem eq-sym-conv]

lemma mask-eq-iff: (w AND mask n) = w <-> uint w < 2 ^ n
  apply (simp add: and-mask-bintr word-number-of-def)

```

```

apply (simp add: word-ubin.inverse-norm)
apply (simp add: eq-mod-iff bintrunc-mod2p min-def)
apply (fast intro!: lt2p-lem)
done

```

```

lemma and-mask-dvd:  $2^n \text{ dvd uint } w = (w \text{ AND mask } n = 0)$ 
apply (simp add: zdvd-iff-zmod-eq-0 and-mask-mod-2p)
apply (simp add: word-uint.norm-eq-iff [symmetric] word-of-int-homs)
apply (subst word-uint.norm-Rep [symmetric])
apply (simp only: bintrunc-bintrunc-min bintrunc-mod2p [symmetric] min-def)
apply auto
done

```

```

lemma and-mask-dvd-nat:  $2^n \text{ dvd unat } w = (w \text{ AND mask } n = 0)$ 
apply (unfold unat-def)
apply (rule trans [OF - and-mask-dvd])
apply (unfold dvd-def)
apply auto
apply (drule uint-ge-0 [THEN nat-int.Abs-inverse' [simplified], symmetric])
apply (simp add : int-mult int-power)
apply (simp add : nat-mult-distrib nat-power-eq)
done

```

```

lemma word-2p-lem:
   $n < \text{size } w \implies w < 2^n = (\text{uint } (w :: 'a :: \text{len word}) < 2^n)$ 
apply (unfold word-size word-less-alt word-number-of-alt)
apply (clarsimp simp add: word-of-int-power-hom word-uint.eq-norm
  int-mod-eq'
  simp del: word-of-int-bin)
done

```

```

lemma less-mask-eq:  $x < 2^n \implies x \text{ AND mask } n = (x :: 'a :: \text{len word})$ 
apply (unfold word-less-alt word-number-of-alt)
apply (clarsimp simp add: and-mask-mod-2p word-of-int-power-hom
  word-uint.eq-norm
  simp del: word-of-int-bin)
apply (drule xtr8 [rotated])
apply (rule int-mod-le)
apply (auto simp add : mod-pos-pos-trivial)
done

```

```

lemmas mask-eq-iff-w2p =
  trans [OF mask-eq-iff word-2p-lem [symmetric], standard]

```

```

lemmas and-mask-less' =
  iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]

```

```

lemma and-mask-less-size:  $n < \text{size } x \implies x \text{ AND mask } n < 2^n$ 
unfolding word-size by (erule and-mask-less')

```

lemma *word-mod-2p-is-mask'*:

$c = 2 \wedge n \implies c > 0 \implies x \text{ mod } c = (x :: 'a :: \text{len word}) \text{ AND mask } n$
by (*clarsimp simp add: word-mod-def uint-2p and-mask-mod-2p*)

lemmas *word-mod-2p-is-mask* = refl [THEN *word-mod-2p-is-mask'*]

lemma *mask-egs*:

$(a \text{ AND mask } n) + b \text{ AND mask } n = a + b \text{ AND mask } n$
 $a + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - b \text{ AND mask } n = a - b \text{ AND mask } n$
 $a - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $a * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $(b \text{ AND mask } n) * a \text{ AND mask } n = b * a \text{ AND mask } n$
 $(a \text{ AND mask } n) + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $(a \text{ AND mask } n) * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $-(a \text{ AND mask } n) \text{ AND mask } n = -a \text{ AND mask } n$
 $\text{word-succ } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-succ } a \text{ AND mask } n$
 $\text{word-pred } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-pred } a \text{ AND mask } n$
using *word-of-int-Ex [where x=a] word-of-int-Ex [where x=b]*
by (*auto simp: and-mask-wi bintr-ariths bintr-arith1s new-word-of-int-homs*)

lemma *mask-power-eq*:

$(x \text{ AND mask } n) \wedge k \text{ AND mask } n = x \wedge k \text{ AND mask } n$
using *word-of-int-Ex [where x=x]*
by (*clarsimp simp: and-mask-wi word-of-int-power-hom bintr-ariths*)

13.1.3 Recast

lemmas *revcast-def'* = *revcast-def* [*simplified*]

lemmas *revcast-def''* = *revcast-def'* [*simplified word-size*]

lemmas *revcast-no-def* [*simp*] =

revcast-def' [**where** *w=number-of w, unfolded word-size, standard*]

lemma *to-bl-revcast*:

$\text{to-bl } (\text{revcast } w :: 'a :: \text{len0 word}) =$
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$
apply (*unfold revcast-def' word-size*)
apply (*rule word-bl.Abs-inverse*)
apply *simp*
done

lemma *revcast-rev-ucast'*:

$cs = [rc, uc] \implies rc = \text{revcast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$
 $rc = \text{word-reverse } uc$
apply (*unfold ucast-def revcast-def' Let-def word-reverse-def*)
apply (*clarsimp simp add : to-bl-of-bin takefill-bintrunc*)
apply (*simp add : word-bl.Abs-inverse word-size*)

done

lemmas *revcast-rev-ucast = revcast-rev-ucast'* [*OF refl refl refl*]

lemmas *revcast-ucast = revcast-rev-ucast*
 [**where** *w = word-reverse w, simplified word-rev-rev, standard*]

lemmas *ucast-revcast = revcast-rev-ucast* [*THEN word-rev-gal', standard*]

lemmas *ucast-rev-revcast = revcast-ucast* [*THEN word-rev-gal', standard*]

— linking revcast and cast via shift

lemmas *wsst-TYs = source-size target-size word-size*

lemma *revcast-down-uu'*:

rc = revcast ==> source-size rc = target-size rc + n ==>

rc (w :: 'a :: len word) = ucast (w >> n)

apply (*simp add: revcast-def'*)

apply (*rule word-bl.Rep-inverse'*)

apply (*rule trans, rule ucast-down-drop*)

prefer 2

apply (*rule trans, rule drop-shiftr*)

apply (*auto simp: takefill-alt wsst-TYs*)

done

lemma *revcast-down-us'*:

rc = revcast ==> source-size rc = target-size rc + n ==>

rc (w :: 'a :: len word) = ucast (w >>> n)

apply (*simp add: revcast-def'*)

apply (*rule word-bl.Rep-inverse'*)

apply (*rule trans, rule ucast-down-drop*)

prefer 2

apply (*rule trans, rule drop-sshiftr*)

apply (*auto simp: takefill-alt wsst-TYs*)

done

lemma *revcast-down-su'*:

rc = revcast ==> source-size rc = target-size rc + n ==>

rc (w :: 'a :: len word) = scast (w >> n)

apply (*simp add: revcast-def'*)

apply (*rule word-bl.Rep-inverse'*)

apply (*rule trans, rule scast-down-drop*)

prefer 2

apply (*rule trans, rule drop-shiftr*)

apply (*auto simp: takefill-alt wsst-TYs*)

done

lemma *revcast-down-ss'*:

```

rc = revcast ==> source-size rc = target-size rc + n ==>
  rc (w :: 'a :: len word) = scast (w >>> n)
apply (simp add: revcast-def')
apply (rule word-bl.Rep-inverse')
apply (rule trans, rule scast-down-drop)
prefer 2
apply (rule trans, rule drop-sshiftr)
apply (auto simp: takefill-alt wsst-TYs)
done

```

```

lemmas revcast-down-uu = refl [THEN revcast-down-uu']
lemmas revcast-down-us = refl [THEN revcast-down-us']
lemmas revcast-down-su = refl [THEN revcast-down-su']
lemmas revcast-down-ss = refl [THEN revcast-down-ss']

```

```

lemma cast-down-rev:
  uc = ucast ==> source-size uc = target-size uc + n ==>
    uc w = revcast ((w :: 'a :: len word) << n)
apply (unfold shiftl-rev)
apply clarify
apply (simp add: revcast-rev-ucast)
apply (rule word-rev-gal')
apply (rule trans [OF - revcast-rev-ucast])
apply (rule revcast-down-uu [symmetric])
apply (auto simp add: wsst-TYs)
done

```

```

lemma revcast-up':
  rc = revcast ==> source-size rc + n = target-size rc ==>
    rc w = (ucast w :: 'a :: len word) << n
apply (simp add: revcast-def')
apply (rule word-bl.Rep-inverse')
apply (simp add: takefill-alt)
apply (rule bl-shiftl [THEN trans])
apply (subst ucast-up-app)
apply (auto simp add: wsst-TYs)
apply (drule sym)
apply (simp add: min-def)
done

```

```

lemmas revcast-up = refl [THEN revcast-up']

```

```

lemmas rc1 = revcast-up [THEN
  revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

```

```

lemmas rc2 = revcast-down-uu [THEN
  revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

```

```

lemmas ucast-up =
  rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]

```

lemmas *ucast-down* =
rc2 [*simplified rev-shiftr revcast-ucast* [*symmetric*]]

13.1.4 Slices

lemmas *slice1-no-bin* [*simp*] =
slice1-def [**where** *w=number-of w, unfolded to-bl-no-bin, standard*]

lemmas *slice-no-bin* [*simp*] =
trans [*OF slice-def* [*THEN meta-eq-to-obj-eq*]
slice1-no-bin [*THEN meta-eq-to-obj-eq*],
unfolded word-size, standard]

lemma *slice1-0* [*simp*] : *slice1 n 0 = 0*
unfolding *slice1-def* **by** (*simp add : to-bl-0*)

lemma *slice-0* [*simp*] : *slice n 0 = 0*
unfolding *slice-def* **by** *auto*

lemma *slice-take'*: *slice n w = of-bl (take (size w - n) (to-bl w))*
unfolding *slice-def' slice1-def*
by (*simp add : takefill-alt word-size*)

lemmas *slice-take = slice-take'* [*unfolded word-size*]

— *shiftr* to a word of the same size is just *slice*, *slice* is just *shiftr* then *ucast*

lemmas *shiftr-slice = trans*
[*OF shiftr-bl* [*THEN meta-eq-to-obj-eq*] *slice-take* [*symmetric*], *standard*]

lemma *slice-shiftr*: *slice n w = ucast (w >> n)*
apply (*unfold slice-take shiftr-bl*)
apply (*rule ucast-of-bl-up* [*symmetric*])
apply (*simp add: word-size*)
done

lemma *nth-slice*:
(*slice n w :: 'a :: len0 word*) !! *m* =
(*w* !! (*m + n*) & *m < len-of TYPE ('a)*)
unfolding *slice-shiftr*
by (*simp add : nth-ucast nth-shiftr*)

lemma *slice1-down-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>
to-bl sl = takefill False fs (drop k (to-bl w))
unfolding *slice1-def word-size of-bl-def uint-bl*
by (*clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop drop-takefill*)

lemma *slice1-up-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>

```

    to-bl sl = takefill False fs (replicate k False @ (to-bl w))
  apply (unfold slice1-def word-size of-bl-def uint-bl)
  apply (clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop
         takefill-append [symmetric])
  apply (rule-tac f = %k. takefill False (len-of TYPE('a))
        (replicate k False @ bin-to-bl (len-of TYPE('b)) (uint w)) in arg-cong)
  apply arith
  done

```

```

lemmas sd1 = slice1-down-alt' [OF refl refl, unfolded word-size]
lemmas su1 = slice1-up-alt' [OF refl refl, unfolded word-size]
lemmas slice1-down-alt = le-add-diff-inverse [THEN sd1]
lemmas slice1-up-alt =
  le-add-diff-inverse [symmetric, THEN su1]
  le-add-diff-inverse2 [symmetric, THEN su1]

```

```

lemma ucast-slice1: ucast w = slice1 (size w) w
  unfolding slice1-def ucast-bl
  by (simp add : takefill-same' word-size)

```

```

lemma ucast-slice: ucast w = slice 0 w
  unfolding slice-def by (simp add : ucast-slice1)

```

```

lemmas slice-id = trans [OF ucast-slice [symmetric] ucast-id]

```

```

lemma revcast-slice1':
  rc = revcast w ==> slice1 (size rc) w = rc
  unfolding slice1-def revcast-def' by (simp add : word-size)

```

```

lemmas revcast-slice1 = refl [THEN revcast-slice1']

```

```

lemma slice1-tf-tf':
  to-bl (slice1 n w :: 'a :: len0 word) =
    rev (takefill False (len-of TYPE('a)) (rev (takefill False n (to-bl w))))
  unfolding slice1-def by (rule word-rev-tf)

```

```

lemmas slice1-tf-tf = slice1-tf-tf'
  [THEN word-bl.Rep-inverse', symmetric, standard]

```

```

lemma rev-slice1:
  n + k = len-of TYPE('a) + len-of TYPE('b) ==>
  slice1 n (word-reverse w :: 'b :: len0 word) =
  word-reverse (slice1 k w :: 'a :: len0 word)
  apply (unfold word-reverse-def slice1-tf-tf)
  apply (rule word-bl.Rep-inverse')
  apply (rule rev-swap [THEN iffD1])
  apply (rule trans [symmetric])
  apply (rule tf-rev)
  apply (simp add: word-bl.Abs-inverse)

```

```

apply (simp add: word-bl.Abs-inverse)
done

```

lemma *rev-slice'*:

```

res = slice n (word-reverse w) ==> n + k + size res = size w ==>
  res = word-reverse (slice k w)
apply (unfold slice-def word-size)
apply clarify
apply (rule rev-slice1)
apply arith
done

```

lemmas *rev-slice = refl [THEN rev-slice', unfolded word-size]*

lemmas *sym-notr =*

```

not-iff [THEN iffD2, THEN not-sym, THEN not-iff [THEN iffD1]]

```

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:

```

(sint (x :: 'a :: len word) + sint y = sint (x + y)) =
  (((x+y) XOR x) AND ((x+y) XOR y)) >> (size x - 1) = 0)
apply (unfold word-size)
apply (cases len-of TYPE('a), simp)
apply (subst msb-shift [THEN sym-notr])
apply (simp add: word-ops-msb)
apply (simp add: word-msb-sint)
apply safe
  apply simp-all
apply (unfold sint-word-ariths)
apply (unfold word-sbin.set-iff-norm [symmetric] sints-num)
apply safe
  apply (insert sint-range' [where x=x])
  apply (insert sint-range' [where x=y])
  defer
  apply (simp (no-asm), arith)
  apply (simp (no-asm), arith)
  defer
  defer
  apply (simp (no-asm), arith)
  apply (simp (no-asm), arith)
apply (rule notI [THEN notnotD],
  drule leI not-leE,
  drule sbintrunc-inc sbintrunc-dec,
  simp)+
done

```

13.2 Split and cat

lemmas *word-split-bin'* = *word-split-def* [*THEN meta-eq-to-obj-eq, standard*]

lemmas *word-cat-bin'* = *word-cat-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-rsplit-no*:

(*word-rsplit* (*number-of bin* :: '*b* :: *len0 word*) :: '*a word list*) =
 map *number-of* (*bin-rsplit* (*len-of TYPE*('a :: *len*))
 (*len-of TYPE*('b), *bintrunc* (*len-of TYPE*('b)) *bin*))

apply (*unfold word-rsplit-def word-no-wi*)

apply (*simp add: word-ubin.eq-norm*)

done

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*

[*unfolded bin-rsplittl-def bin-rsplit-l [symmetric]*]

lemma *test-bit-cat*:

wc = *word-cat a b* ==> *wc !! n* = (*n < size wc* &
 (*if n < size b then b !! n else a !! (n - size b)*))

apply (*unfold word-cat-bin' test-bit-bin*)

apply (*auto simp add : word-ubin.eq-norm nth-bintr bin-nth-cat word-size*)

apply (*erule bin-nth-uint-imp*)

done

lemma *word-cat-bl*: *word-cat a b* = *of-bl (to-bl a @ to-bl b)*

apply (*unfold of-bl-def to-bl-def word-cat-bin'*)

apply (*simp add: bl-to-bin-app-cat*)

done

lemma *of-bl-append*:

(*of-bl* (*xs @ ys*) :: '*a :: len word*) = *of-bl xs* * $2^{(\text{length } ys)}$ + *of-bl ys*

apply (*unfold of-bl-def*)

apply (*simp add: bl-to-bin-app-cat bin-cat-num*)

apply (*simp add: word-of-int-power-hom [symmetric] new-word-of-int-hom-syms*)

done

lemma *of-bl-False* [*simp*]:

of-bl (False#xs) = *of-bl xs*

by (*rule word-eqI*)

(*auto simp add: test-bit-of-bl nth-append*)

lemma *of-bl-True*:

(*of-bl (True#xs)::'a::len word*) = $2^{\text{length } xs}$ + *of-bl xs*

by (*subst of-bl-append [where xs=[True], simplified]*)

(*simp add: word-1-bl*)

lemma *of-bl-Cons*:

of-bl (x#xs) = *of-bool x* * $2^{\text{length } xs}$ + *of-bl xs*

by (*cases x*) (*simp-all add: of-bl-True*)

```

lemma split-uint-lem: bin-split n (uint (w :: 'a :: len0 word)) = (a, b) ==>
  a = bintrunc (len-of TYPE('a) - n) a & b = bintrunc (len-of TYPE('a)) b
apply (frule word-ubin.norm-Rep [THEN ssubst])
apply (drule bin-split-trunc1)
apply (drule sym [THEN trans])
apply assumption
apply safe
done

```

```

lemma word-split-bl':
  std = size c - size b ==> (word-split c = (a, b)) ==>
    (a = of-bl (take std (to-bl c)) & b = of-bl (drop std (to-bl c)))
apply (unfold word-split-bin')
apply safe
  defer
    apply (clarsimp split: prod.splits)
    apply (drule word-ubin.norm-Rep [THEN ssubst])
    apply (drule split-bintrunc)
    apply (simp add : of-bl-def bl2bin-drop word-size
      word-ubin.norm-eq-iff [symmetric] min-def del : word-ubin.norm-Rep)
    apply (clarsimp split: prod.splits)
    apply (frule split-uint-lem [THEN conjunct1])
    apply (unfold word-size)
    apply (cases len-of TYPE('a) >= len-of TYPE('b))
    defer
      apply (simp add: word-0-bl word-0-wi-Pls)
      apply (simp add : of-bl-def to-bl-def)
      apply (subst bin-split-take1 [symmetric])
      prefer 2
      apply assumption
      apply simp
    apply (erule thin-rl)
    apply (erule arg-cong [THEN trans])
    apply (simp add : word-ubin.norm-eq-iff [symmetric])
done

```

```

lemma word-split-bl: std = size c - size b ==>
  (a = of-bl (take std (to-bl c)) & b = of-bl (drop std (to-bl c))) <->
  word-split c = (a, b)
apply (rule iffI)
defer
  apply (erule (1) word-split-bl')
apply (case-tac word-split c)
apply (auto simp add : word-size)
apply (frule word-split-bl' [rotated])
apply (auto simp add : word-size)
done

```

```

lemma word-split-bl-eq:

```

```

(word-split (c::'a::len word) :: ('c :: len0 word * 'd :: len0 word)) =
  (of-bl (take (len-of TYPE('a::len) - len-of TYPE('d::len0)) (to-bl c)),
   of-bl (drop (len-of TYPE('a) - len-of TYPE('d)) (to-bl c)))
apply (rule word-split-bl [THEN iffD1])
apply (unfold word-size)
apply (rule refl conjI)+
done

```

— keep quantifiers for use in simplification

```

lemma test-bit-split':
  word-split c = (a, b) --> (ALL n m. b !! n = (n < size b & c !! n) &
    a !! m = (m < size a & c !! (m + size b)))
apply (unfold word-split-bin' test-bit-bin)
apply (clarify)
apply (clarsimp simp: word-ubin.eq-norm nth-bintr word-size split: prod.splits)
apply (drule bin-nth-split)
apply safe
  apply (simp-all add: add-commute)
  apply (erule bin-nth-wint-imp)+
done

```

```

lemmas test-bit-split =
  test-bit-split' [THEN mp, simplified all-simps, standard]

```

```

lemma test-bit-split-eq: word-split c = (a, b) <->
  ((ALL n::nat. b !! n = (n < size b & c !! n)) &
   (ALL m::nat. a !! m = (m < size a & c !! (m + size b))))
apply (rule-tac iffI)
apply (rule-tac conjI)
  apply (erule test-bit-split [THEN conjunct1])
  apply (erule test-bit-split [THEN conjunct2])
apply (case-tac word-split c)
apply (frule test-bit-split)
apply (erule trans)
apply (fastsimp intro ! : word-eqI simp add : word-size)
done

```

— this odd result is analogous to ucast_id, result to the length given by the result type

```

lemma word-cat-id: word-cat a b = b
  unfolding word-cat-bin' by (simp add: word-ubin.inverse-norm)

```

— limited hom result

```

lemma word-cat-hom:
  len-of TYPE('a::len0) <= len-of TYPE('b::len0) + len-of TYPE('c::len0)
  ==>
  (word-cat (word-of-int w :: 'b word) (b :: 'c word) :: 'a word) =
  word-of-int (bin-cat w (size b) (uint b))

```

```

apply (unfold word-cat-def word-size)
apply (clarsimp simp add : word-ubin.norm-eq-iff [symmetric]
      word-ubin.eq-norm bintr-cat min-def)
apply arith
done

```

```

lemma word-cat-split-alt:
  size w <= size u + size v ==> word-split w = (u, v) ==> word-cat u v = w
apply (rule word-eqI)
apply (drule test-bit-split)
apply (clarsimp simp add : test-bit-cat word-size)
apply safe
apply arith
done

```

```

lemmas word-cat-split-size =
  sym [THEN [2] word-cat-split-alt [symmetric], standard]

```

13.2.1 Split and slice

```

lemma split-slices:
  word-split w = (u, v) ==> u = slice (size v) w & v = slice 0 w
apply (drule test-bit-split)
apply (rule conjI)
apply (rule word-eqI, clarsimp simp: nth-slice word-size)+
done

```

```

lemma slice-cat1':
  wc = word-cat a b ==> size wc >= size a + size b ==> slice (size b) wc = a
apply safe
apply (rule word-eqI)
apply (simp add: nth-slice test-bit-cat word-size)
done

```

```

lemmas slice-cat1 = refl [THEN slice-cat1']
lemmas slice-cat2 = trans [OF slice-id word-cat-id]

```

```

lemma cat-slices:
  a = slice n c ==> b = slice 0 c ==> n = size b ==>
  size a + size b >= size c ==> word-cat a b = c
apply safe
apply (rule word-eqI)
apply (simp add: nth-slice test-bit-cat word-size)
apply safe
apply arith
done

```

```

lemma word-split-cat-alt:
  w = word-cat u v ==> size u + size v <= size w ==> word-split w = (u, v)

```

```

apply (case-tac word-split ?w)
apply (rule trans, assumption)
apply (drule test-bit-split)
apply safe
apply (rule word-eqI, clarsimp simp: test-bit-cat word-size)+
done

```

```

lemmas word-cat-bl-no-bin [simp] =
  word-cat-bl [where a=number-of a
    and b=number-of b,
    unfolded to-bl-no-bin, standard]

```

```

lemmas word-split-bl-no-bin [simp] =
  word-split-bl-eq [where c=number-of c, unfolded to-bl-no-bin, standard]

```

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

```

lemma word-rsplit-same: word-rsplit w = [w]
unfolding word-rsplit-def by (simp add : bin-rsplit-all)

```

```

lemma word-rsplit-empty-iff-size:
  (word-rsplit w = []) = (size w = 0)
unfolding word-rsplit-def bin-rsplit-def word-size
by (simp add: bin-rsplit-aux-simp-alt Let-def split: split-split)

```

```

lemma test-bit-rsplit:
  sw = word-rsplit w ==> m < size (hd sw :: 'a :: len word) ==>
    k < length sw ==> (rev sw ! k) !! m = (w !! (k * size (hd sw) + m))
apply (unfold word-rsplit-def word-test-bit-def)
apply (rule trans)
apply (rule-tac f = %x. bin-nth x m in arg-cong)
apply (rule nth-map [symmetric])
apply simp
apply (rule bin-nth-rsplit)
apply simp-all
apply (simp add : word-size rev-map map-compose [symmetric])
apply (rule trans)
defer
apply (rule map-ident [THEN fun-cong])
apply (rule refl [THEN map-cong])
apply (simp add : word-ubin.eq-norm)
apply (erule bin-rsplit-size-sign [OF len-gt-0 refl])
done

```

```

lemma word-rcat-bl: word-rcat wl == of-bl (concat (map to-bl wl))
unfolding word-rcat-def to-bl-def' of-bl-def
by (clarsimp simp add : bin-rcat-bl map-compose)

```

```

lemma size-rcat-lem':
  size (concat (map to-bl wl)) = length wl * size (hd wl)
  unfolding word-size by (induct wl) auto

lemmas size-rcat-lem = size-rcat-lem' [unfolded word-size]

lemmas td-gal-lt-len = len-gt-0 [THEN td-gal-lt, standard]

lemma nth-rcat-lem' [rule-format] :
  sw = size (hd wl :: 'a :: len word) ==> (ALL n. n < size wl * sw -->
    rev (concat (map to-bl wl)) ! n =
    rev (to-bl (rev wl ! (n div sw))) ! (n mod sw))
  apply (unfold word-size)
  apply (induct wl)
  apply clarsimp
  apply (clarsimp simp add : nth-append size-rcat-lem)
  apply (simp (no-asm-use) only: mult-Suc [symmetric]
    td-gal-lt-len less-Suc-eq-le mod-div-equality')
  apply clarsimp
  done

lemmas nth-rcat-lem = refl [THEN nth-rcat-lem', unfolded word-size]

lemma test-bit-rcat:
  sw = size (hd wl :: 'a :: len word) ==> rc = word-rcat wl ==> rc !! n =
    (n < size rc & n div sw < size wl & (rev wl) ! (n div sw) !! (n mod sw))
  apply (unfold word-rcat-bl word-size)
  apply (clarsimp simp add :
    test-bit-of-bl size-rcat-lem word-size td-gal-lt-len)
  apply safe
  apply (auto simp add :
    test-bit-bl word-size td-gal-lt-len [THEN iffD2, THEN nth-rcat-lem])
  done

lemma foldl-eq-foldr [rule-format] :
  ALL x. foldl op + x xs = foldr op + (x # xs) (0 :: 'a :: comm-monoid-add)
  by (induct xs) (auto simp add : add-assoc)

lemmas test-bit-cong = arg-cong [where f = test-bit, THEN fun-cong]

lemmas test-bit-rsplit-alt =
  trans [OF nth-rev-alt [THEN test-bit-cong]
    test-bit-rsplit [OF refl asm-rl diff-Suc-less]]

— lazy way of expressing that u and v, and su and sv, have same types
lemma word-rsplit-len-indep':
  [u,v] = p ==> [su,sv] = q ==> word-rsplit u = su ==>
    word-rsplit v = sv ==> length su = length sv

```

```

apply (unfold word-rsplit-def)
apply (auto simp add : bin-rsplit-len-indep)
done

```

lemmas word-rsplit-len-indep = word-rsplit-len-indep' [OF refl refl refl refl]

```

lemma length-word-rsplit-size:
  n = len-of TYPE ('a :: len) ==>
    (length (word-rsplit w :: 'a word list) <= m) = (size w <= m * n)
apply (unfold word-rsplit-def word-size)
apply (clarsimp simp add : bin-rsplit-len-le)
done

```

lemmas length-word-rsplit-lt-size =
length-word-rsplit-size [unfolded Not-eq-iff linorder-not-less [symmetric]]

```

lemma length-word-rsplit-exp-size:
  n = len-of TYPE ('a :: len) ==>
    length (word-rsplit w :: 'a word list) = (size w + n - 1) div n
unfolding word-rsplit-def by (clarsimp simp add : word-size bin-rsplit-len)

```

```

lemma length-word-rsplit-even-size:
  n = len-of TYPE ('a :: len) ==> size w = m * n ==>
    length (word-rsplit w :: 'a word list) = m
by (clarsimp simp add : length-word-rsplit-exp-size given-quot-alt)

```

lemmas length-word-rsplit-exp-size' = refl [THEN length-word-rsplit-exp-size]

lemmas ttle = iffD2 [OF split-div-lemma refl, THEN conjunct1]
lemmas dtle = xtr4 [OF ttle mult-commute]

```

lemma word-rcat-rsplit: word-rcat (word-rsplit w) = w
apply (rule word-eqI)
apply (clarsimp simp add : test-bit-rcat word-size)
apply (subst refl [THEN test-bit-rsplit])
  apply (simp-all add: word-size
    refl [THEN length-word-rsplit-size [simplified le-def, simplified]])
apply safe
apply (erule xtr7, rule len-gt-0 [THEN dtle])+
done

```

```

lemma size-word-rsplit-rcat-size':
  word-rcat (ws :: 'a :: len word list) = frcw ==>
    size frcw = length ws * len-of TYPE ('a) ==>
    size (hd [word-rsplit frcw, ws]) = size ws
apply (clarsimp simp add : word-size length-word-rsplit-exp-size')
apply (fast intro: given-quot-alt)
done

```

lemmas *size-word-rsplit-rcat-size* =
size-word-rsplit-rcat-size' [*simplified*]

lemma *msrevs*:

fixes *n::nat*
shows $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$
and $(k * n + m) \text{ mod } n = m \text{ mod } n$
by (*auto simp: add-commute*)

lemma *word-rsplit-rcat-size'*:

word-rcat (*ws* :: '*a* :: len word list) = *frcw* ==>
size frcw = *length ws* * *len-of TYPE ('a)* ==> *word-rsplit frcw* = *ws*
apply (*frule size-word-rsplit-rcat-size, assumption*)
apply (*clarsimp simp add : word-size*)
apply (*rule nth-equalityI, assumption*)
apply *clarsimp*
apply (*rule word-eqI*)
apply (*rule trans*)
apply (*rule test-bit-rsplit-alt*)
apply (*clarsimp simp: word-size*)+
apply (*rule trans*)
apply (*rule test-bit-rcat [OF refl refl]*)
apply (*simp add : word-size msrevs*)
apply (*subst nth-rev*)
apply *arith*
apply (*simp add : le0 [THEN [2] xtr7, THEN diff-Suc-less]*)
apply *safe*
apply (*simp add : diff-mult-distrib*)
apply (*rule mpl-lem*)
apply (*cases size ws*)
apply *simp-all*
done

lemmas *word-rsplit-rcat-size* = *refl [THEN word-rsplit-rcat-size']*

13.3 Rotation

lemmas *rotater-0'* [*simp*] = *rotater-def* [**where** $n = 0$, *simplified*]

lemmas *word-rot-defs* = *word-roti-def word-rotr-def word-rotl-def*

lemma *rotate-eq-mod*:

$m \text{ mod } \text{length } xs = n \text{ mod } \text{length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$
apply (*rule box-equals*)
defer
apply (*rule rotate-conv-mod [symmetric]*)+
apply *simp*
done

lemmas *rotate-eqs* [*standard*] =
trans [*OF rotate0* [*THEN fun-cong*] *id-apply*]
rotate-rotate [*symmetric*]
rotate-id
rotate-conv-mod
rotate-eq-mod

13.3.1 Rotation of list to right

lemma *rotate1-rl'*: *rotater1* (*l* @ [*a*]) = *a* # *l*
unfolding *rotater1-def* **by** (*cases l*) *auto*

lemma *rotate1-rl* [*simp*] : *rotater1* (*rotate1 l*) = *l*
apply (*unfold rotater1-def*)
apply (*cases l*)
apply (*case-tac* [2] *list*)
apply *auto*
done

lemma *rotate1-lr* [*simp*] : *rotate1* (*rotater1 l*) = *l*
unfolding *rotater1-def* **by** (*cases l*) *auto*

lemma *rotater1-rev'*: *rotater1* (*rev xs*) = *rev* (*rotate1 xs*)
apply (*cases xs*)
apply (*simp add* : *rotater1-def*)
apply (*simp add* : *rotate1-rl'*)
done

lemma *rotater-rev'*: *rotater n* (*rev xs*) = *rev* (*rotate n xs*)
unfolding *rotater-def* **by** (*induct n*) (*auto intro: rotater1-rev'*)

lemmas *rotater-rev* = *rotater-rev'* [**where** *xs* = *rev ys*, *simplified*, *standard*]

lemma *rotater-drop-take*:
rotater n xs =
drop (*length xs* - *n mod length xs*) *xs* @
take (*length xs* - *n mod length xs*) *xs*
by (*clarsimp simp add* : *rotater-rev rotate-drop-take rev-take rev-drop*)

lemma *rotater-Suc* [*simp*] :
rotater (*Suc n*) *xs* = *rotater1* (*rotater n xs*)
unfolding *rotater-def* **by** *auto*

lemma *rotate-inv-plus* [*rule-format*] :
ALL k. k = *m* + *n* \longrightarrow *rotater k* (*rotate n xs*) = *rotater m xs* &
rotate k (*rotater n xs*) = *rotate m xs* &
rotater n (*rotate k xs*) = *rotate m xs* &
rotate n (*rotater k xs*) = *rotater m xs*

unfolding *rotater-def rotate-def*
by (*induct n*) (*auto intro: funpow-swap1 [THEN trans]*)

lemmas *rotate-inv-rel = le-add-diff-inverse2 [symmetric, THEN rotate-inv-plus]*

lemmas *rotate-inv-eq = order-refl [THEN rotate-inv-rel, simplified]*

lemmas *rotate-lr [simp] = rotate-inv-eq [THEN conjunct1, standard]*
lemmas *rotate-rl [simp] =*
rotate-inv-eq [THEN conjunct2, THEN conjunct1, standard]

lemma *rotate-gal: (rotater n xs = ys) = (rotate n ys = xs)*
by *auto*

lemma *rotate-gal': (ys = rotater n xs) = (xs = rotate n ys)*
by *auto*

lemma *length-rotater [simp]:*
length (rotater n xs) = length xs
by (*simp add : rotater-rev*)

lemmas *rrs0 = rotate-eqs [THEN restrict-to-left,*
simplified rotate-gal [symmetric] rotate-gal' [symmetric], standard]
lemmas *rrs1 = rrs0 [THEN refl [THEN rev-iffD1]]*
lemmas *rotater-eqs = rrs1 [simplified length-rotater, standard]*
lemmas *rotater-0 = rotater-eqs (1)*
lemmas *rotater-add = rotater-eqs (2)*

13.3.2 map, app2, commuting with rotate(r)

lemma *last-map: xs $\sim = [] \implies \text{last (map f xs)} = f (\text{last xs})$*
by (*induct xs*) *auto*

lemma *butlast-map:*
xs $\sim = [] \implies \text{butlast (map f xs)} = \text{map f (butlast xs)}$
by (*induct xs*) *auto*

lemma *rotater1-map: rotater1 (map f xs) = map f (rotater1 xs)*
unfolding *rotater1-def*
by (*cases xs*) (*auto simp add: last-map butlast-map*)

lemma *rotater-map:*
rotater n (map f xs) = map f (rotater n xs)
unfolding *rotater-def*
by (*induct n*) (*auto simp add : rotater1-map*)

lemma *but-last-zip [rule-format] :*
ALL ys. length xs = length ys \implies xs $\sim = [] \implies$
last (zip xs ys) = (last xs, last ys) &

```

  butlast (zip xs ys) = zip (butlast xs) (butlast ys)
apply (induct xs)
apply auto
  apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
done

```

```

lemma but-last-app2 [rule-format] :
  ALL ys. length xs = length ys --> xs ~ = [] -->
  last (app2 f xs ys) = f (last xs) (last ys) &
  butlast (app2 f xs ys) = app2 f (butlast xs) (butlast ys)
apply (induct xs)
apply auto
  apply (unfold app2-def)
  apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
done

```

```

lemma rotater1-zip:
  length xs = length ys ==>
  rotater1 (zip xs ys) = zip (rotater1 xs) (rotater1 ys)
apply (unfold rotater1-def)
apply (cases xs)
apply auto
  apply ((case-tac ys, auto simp: neq-Nil-conv but-last-zip)[1])+
done

```

```

lemma rotater1-app2:
  length xs = length ys ==>
  rotater1 (app2 f xs ys) = app2 f (rotater1 xs) (rotater1 ys)
unfolding app2-def by (simp add: rotater1-map rotater1-zip)

```

```

lemmas lrth =
  box-equals [OF asm-rl length-rotater [symmetric]
    length-rotater [symmetric],
    THEN rotater1-app2]

```

```

lemma rotater-app2:
  length xs = length ys ==>
  rotater n (app2 f xs ys) = app2 f (rotater n xs) (rotater n ys)
by (induct n) (auto intro!: lrth)

```

```

lemma rotate1-app2:
  length xs = length ys ==>
  rotate1 (app2 f xs ys) = app2 f (rotate1 xs) (rotate1 ys)
apply (unfold app2-def)
apply (cases xs)
  apply (cases ys, auto simp add : rotate1-def)+
done

```

```

lemmas lth = box-equals [OF asm-rl length-rotate [symmetric]

```

length-rotate [*symmetric*], *THEN rotate1-app2*]

lemma *rotate-app2*:

length xs = length ys ==>

rotate n (app2 f xs ys) = app2 f (rotate n xs) (rotate n ys)

by (*induct n*) (*auto intro! lth*)

— corresponding equalities for word rotation

lemma *to-bl-rotl*:

to-bl (word-rotl n w) = rotate n (to-bl w)

by (*simp add: word-bl.Abs-inverse' word-rotl-def*)

lemmas *blrs0 = rotate-egs* [*THEN to-bl-rotl* [*THEN trans*]]

lemmas *word-rotl-egs =*

blrs0 [*simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotl* [*symmetric*]]

lemma *to-bl-rotr*:

to-bl (word-rotr n w) = rotater n (to-bl w)

by (*simp add: word-bl.Abs-inverse' word-rotr-def*)

lemmas *brrs0 = rotater-egs* [*THEN to-bl-rotr* [*THEN trans*]]

lemmas *word-rotr-egs =*

brrs0 [*simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotr* [*symmetric*]]

declare *word-rotr-egs* (1) [*simp*]

declare *word-rotl-egs* (1) [*simp*]

lemma

word-rot-rl [*simp*]:

word-rotl k (word-rotr k v) = v **and**

word-rot-lr [*simp*]:

word-rotr k (word-rotl k v) = v

by (*auto simp add: to-bl-rotr to-bl-rotl word-bl.Rep-inject* [*symmetric*])

lemma

word-rot-gal:

(word-rotr n v = w) = (word-rotl n w = v) **and**

word-rot-gal':

(w = word-rotr n v) = (v = word-rotl n w)

by (*auto simp: to-bl-rotr to-bl-rotl word-bl.Rep-inject* [*symmetric*]
dest: sym)

lemma *word-rotr-rev*:

word-rotr n w = word-reverse (word-rotl n (word-reverse w))

by (*simp add: word-bl.Rep-inject* [*symmetric*] *to-bl-word-rev*)

to-bl-rotr to-bl-rotl rotater-rev)

lemma *word-roti-0* [simp]: *word-roti 0 w = w*
by (*unfold word-rot-defs*) *auto*

lemmas *abl-cong = arg-cong* [where *f = of-bl*]

lemma *word-roti-add*:

word-roti (m + n) w = word-roti m (word-roti n w)

proof –

have *rotater-eq-lem*:

$\bigwedge m n xs. m = n \implies \text{rotater } m \text{ } xs = \text{rotater } n \text{ } xs$
by *auto*

have *rotate-eq-lem*:

$\bigwedge m n xs. m = n \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$
by *auto*

note *rpts* [symmetric, standard] =

rotate-inv-plus [THEN *conjunct1*]

rotate-inv-plus [THEN *conjunct2*, THEN *conjunct1*]

rotate-inv-plus [THEN *conjunct2*, THEN *conjunct2*, THEN *conjunct1*]

rotate-inv-plus [THEN *conjunct2*, THEN *conjunct2*, THEN *conjunct2*]

note *rrp = trans* [symmetric, OF *rotate-rotate rotate-eq-lem*]

note *rrrp = trans* [symmetric, OF *rotater-add* [symmetric] *rotater-eq-lem*]

show *?thesis*

apply (*unfold word-rot-defs*)

apply (*simp only: split: split-if*)

apply (*safe intro!: abl-cong*)

apply (*simp-all only: to-bl-rotl* [THEN *word-bl.Rep-inverse*']

to-bl-rotl

to-bl-rotr [THEN *word-bl.Rep-inverse*']

to-bl-rotr)

apply (*rule rrp rrrp rpts*,

simp add: nat-add-distrib [symmetric]

nat-diff-distrib [symmetric])+

done

qed

lemma *word-roti-conv-mod'*: *word-roti n w = word-roti (n mod int (size w)) w*

apply (*unfold word-rot-defs*)

apply (*cut-tac y=size w in gt-or-eq-0*)

apply (*erule disjE*)

apply *simp-all*

apply (*safe intro!: abl-cong*)

apply (*rule rotater-eqs*)

apply (*simp add: word-size nat-mod-distrib*)

```

apply (simp add: rotater-add [symmetric] rotate-gal [symmetric])
apply (rule rotater-eqs)
apply (simp add: word-size nat-mod-distrib)
apply (rule int-eq-0-conv [THEN iffD1])
apply (simp only: zmod-int zadd-int [symmetric])
apply (simp add: rmods)
done

```

lemmas *word-roti-conv-mod = word-roti-conv-mod'* [unfolded word-size]

13.3.3 Word rotation commutes with bit-wise operations

locale *word-rotate*

context *word-rotate*

begin

lemmas *word-rot-defs' = to-bl-rotl to-bl-rotr*

lemmas *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

lemmas *lbl-lbl = trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

lemmas *ths-app2* [OF *lbl-lbl*] = *rotate-app2 rotater-app2*

lemmas *ths-map* [where *xs = to-bl v*] = *rotate-map rotater-map*

lemmas *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-app2 ths-map*

lemma *word-rot-logs*:

```

word-rotl n (NOT v) = NOT word-rotl n v
word-rotr n (NOT v) = NOT word-rotr n v
word-rotl n (x AND y) = word-rotl n x AND word-rotl n y
word-rotr n (x AND y) = word-rotr n x AND word-rotr n y
word-rotl n (x OR y) = word-rotl n x OR word-rotl n y
word-rotr n (x OR y) = word-rotr n x OR word-rotr n y
word-rotl n (x XOR y) = word-rotl n x XOR word-rotl n y
word-rotr n (x XOR y) = word-rotr n x XOR word-rotr n y
by (rule word-bl.Rep-eqD,
      rule word-rot-defs' [THEN trans],
      simp only: blwl-syms [symmetric],
      rule th1s [THEN trans],
      rule refl)+

```

end

lemmas *word-rot-logs = word-rotate.word-rot-logs*

lemmas *bl-word-rotl-dt = trans* [OF *to-bl-rotl rotate-drop-take*,
simplified word-bl.Rep', *standard*]

```

lemmas bl-word-rotr-dt = trans [OF to-bl-rotr rotater-drop-take,
  simplified word-bl.Rep', standard]

lemma bl-word-roti-dt':
  n = nat ((- i) mod int (size (w :: 'a :: len word))) ==>
    to-bl (word-roti i w) = drop n (to-bl w) @ take n (to-bl w)
apply (unfold word-roti-def)
apply (simp add: bl-word-rotl-dt bl-word-rotr-dt word-size)
apply safe
apply (simp add: zmod-zminus1-eq-if)
apply safe
apply (simp add: nat-mult-distrib)
apply (simp add: nat-diff-distrib [OF pos-mod-sign pos-mod-conj
  [THEN conjunct2, THEN order-less-imp-le]]
  nat-mod-distrib)
apply (simp add: nat-mod-distrib)
done

lemmas bl-word-roti-dt = bl-word-roti-dt' [unfolded word-size]

lemmas word-rotl-dt = bl-word-rotl-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-rotr-dt = bl-word-rotr-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-roti-dt = bl-word-roti-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]

lemma word-rotx-0 [simp] : word-rotr i 0 = 0 & word-rotl i 0 = 0
  by (simp add : word-rotr-dt word-rotl-dt to-bl-0 replicate-add [symmetric])

lemma word-roti-0' [simp] : word-roti n 0 = 0
  unfolding word-roti-def by auto

lemmas word-rotr-dt-no-bin' [simp] =
  word-rotr-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

lemmas word-rotl-dt-no-bin' [simp] =
  word-rotl-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

declare word-roti-def [simp]

end

```

14 Boolean-Algebra: Boolean Algebras

theory *Boolean-Algebra*

```

imports Main
begin

locale boolean =
  fixes conj :: 'a ⇒ 'a ⇒ 'a (infixr  $\sqcap$  70)
  fixes disj :: 'a ⇒ 'a ⇒ 'a (infixr  $\sqcup$  65)
  fixes compl :: 'a ⇒ 'a ( $\sim$  - [81] 80)
  fixes zero :: 'a (0)
  fixes one  :: 'a (1)
  assumes conj-assoc:  $(x \sqcap y) \sqcap z = x \sqcap (y \sqcap z)$ 
  assumes disj-assoc:  $(x \sqcup y) \sqcup z = x \sqcup (y \sqcup z)$ 
  assumes conj-commute:  $x \sqcap y = y \sqcap x$ 
  assumes disj-commute:  $x \sqcup y = y \sqcup x$ 
  assumes conj-disj-distrib:  $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$ 
  assumes disj-conj-distrib:  $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$ 
  assumes conj-one-right [simp]:  $x \sqcap \mathbf{1} = x$ 
  assumes disj-zero-right [simp]:  $x \sqcup \mathbf{0} = x$ 
  assumes conj-cancel-right [simp]:  $x \sqcap \sim x = \mathbf{0}$ 
  assumes disj-cancel-right [simp]:  $x \sqcup \sim x = \mathbf{1}$ 
begin

lemmas disj-ac =
  disj-assoc disj-commute
  mk-left-commute [where 'a = 'a, of disj, OF disj-assoc disj-commute]

lemmas conj-ac =
  conj-assoc conj-commute
  mk-left-commute [where 'a = 'a, of conj, OF conj-assoc conj-commute]

lemma dual: boolean disj conj compl one zero
apply (rule boolean.intro)
apply (rule disj-assoc)
apply (rule conj-assoc)
apply (rule disj-commute)
apply (rule conj-commute)
apply (rule disj-conj-distrib)
apply (rule conj-disj-distrib)
apply (rule disj-zero-right)
apply (rule conj-one-right)
apply (rule disj-cancel-right)
apply (rule conj-cancel-right)
done

14.1 Complement

lemma complement-unique:
  assumes 1:  $a \sqcap x = \mathbf{0}$ 
  assumes 2:  $a \sqcup x = \mathbf{1}$ 
  assumes 3:  $a \sqcap y = \mathbf{0}$ 

```

assumes $\not\downarrow$: $a \sqcup y = \mathbf{1}$
shows $x = y$
proof –
have $(a \sqcap x) \sqcup (x \sqcap y) = (a \sqcap y) \sqcup (x \sqcap y)$ **using** 1 3 **by** *simp*
hence $(x \sqcap a) \sqcup (x \sqcap y) = (y \sqcap a) \sqcup (y \sqcap x)$ **using** *conj-commute* **by** *simp*
hence $x \sqcap (a \sqcup y) = y \sqcap (a \sqcup x)$ **using** *conj-disj-distrib* **by** *simp*
hence $x \sqcap \mathbf{1} = y \sqcap \mathbf{1}$ **using** 2 4 **by** *simp*
thus $x = y$ **using** *conj-one-right* **by** *simp*
qed

lemma *compl-unique*: $\llbracket x \sqcap y = \mathbf{0}; x \sqcup y = \mathbf{1} \rrbracket \implies \sim x = y$
by (*rule complement-unique* [*OF conj-cancel-right disj-cancel-right*])

lemma *double-compl* [*simp*]: $\sim(\sim x) = x$
proof (*rule compl-unique*)
from *conj-cancel-right* **show** $\sim x \sqcap x = \mathbf{0}$ **by** (*simp only: conj-commute*)
from *disj-cancel-right* **show** $\sim x \sqcup x = \mathbf{1}$ **by** (*simp only: disj-commute*)
qed

lemma *compl-eq-compl-iff* [*simp*]: $(\sim x = \sim y) = (x = y)$
by (*rule inj-eq* [*OF inj-on-inverseI*], *rule double-compl*)

14.2 Conjunction

lemma *conj-absorb* [*simp*]: $x \sqcap x = x$
proof –
have $x \sqcap x = (x \sqcap x) \sqcup \mathbf{0}$ **using** *disj-zero-right* **by** *simp*
also have $\dots = (x \sqcap x) \sqcup (x \sqcap \sim x)$ **using** *conj-cancel-right* **by** *simp*
also have $\dots = x \sqcap (x \sqcup \sim x)$ **using** *conj-disj-distrib* **by** (*simp only:*)
also have $\dots = x \sqcap \mathbf{1}$ **using** *disj-cancel-right* **by** *simp*
also have $\dots = x$ **using** *conj-one-right* **by** *simp*
finally show *?thesis* .
qed

lemma *conj-zero-right* [*simp*]: $x \sqcap \mathbf{0} = \mathbf{0}$
proof –
have $x \sqcap \mathbf{0} = x \sqcap (x \sqcap \sim x)$ **using** *conj-cancel-right* **by** *simp*
also have $\dots = (x \sqcap x) \sqcap \sim x$ **using** *conj-assoc* **by** (*simp only:*)
also have $\dots = x \sqcap \sim x$ **using** *conj-absorb* **by** *simp*
also have $\dots = \mathbf{0}$ **using** *conj-cancel-right* **by** *simp*
finally show *?thesis* .
qed

lemma *compl-one* [*simp*]: $\sim \mathbf{1} = \mathbf{0}$
by (*rule compl-unique* [*OF conj-zero-right disj-zero-right*])

lemma *conj-zero-left* [*simp*]: $\mathbf{0} \sqcap x = \mathbf{0}$
by (*subst conj-commute*) (*rule conj-zero-right*)

lemma *conj-one-left* [*simp*]: $\mathbf{1} \sqcap x = x$
by (*subst conj-commute*) (*rule conj-one-right*)

lemma *conj-cancel-left* [*simp*]: $\sim x \sqcap x = \mathbf{0}$
by (*subst conj-commute*) (*rule conj-cancel-right*)

lemma *conj-left-absorb* [*simp*]: $x \sqcap (x \sqcap y) = x \sqcap y$
by (*simp only: conj-assoc [symmetric] conj-absorb*)

lemma *conj-disj-distrib2*:
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$
by (*simp only: conj-commute conj-disj-distrib*)

lemmas *conj-disj-distrib* =
conj-disj-distrib conj-disj-distrib2

14.3 Disjunction

lemma *disj-absorb* [*simp*]: $x \sqcup x = x$
by (*rule boolean.conj-absorb [OF dual]*)

lemma *disj-one-right* [*simp*]: $x \sqcup \mathbf{1} = \mathbf{1}$
by (*rule boolean.conj-zero-right [OF dual]*)

lemma *compl-zero* [*simp*]: $\sim \mathbf{0} = \mathbf{1}$
by (*rule boolean.compl-one [OF dual]*)

lemma *disj-zero-left* [*simp*]: $\mathbf{0} \sqcup x = x$
by (*rule boolean.conj-one-left [OF dual]*)

lemma *disj-one-left* [*simp*]: $\mathbf{1} \sqcup x = \mathbf{1}$
by (*rule boolean.conj-zero-left [OF dual]*)

lemma *disj-cancel-left* [*simp*]: $\sim x \sqcup x = \mathbf{1}$
by (*rule boolean.conj-cancel-left [OF dual]*)

lemma *disj-left-absorb* [*simp*]: $x \sqcup (x \sqcup y) = x \sqcup y$
by (*rule boolean.conj-left-absorb [OF dual]*)

lemma *disj-conj-distrib2*:
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$
by (*rule boolean.conj-disj-distrib2 [OF dual]*)

lemmas *disj-conj-distrib* =
disj-conj-distrib disj-conj-distrib2

14.4 De Morgan’s Laws

lemma *de-Morgan-conj* [*simp*]: $\sim (x \sqcap y) = \sim x \sqcup \sim y$
proof (*rule compl-unique*)

```

have  $(x \sqcap y) \sqcap (\sim x \sqcup \sim y) = ((x \sqcap y) \sqcap \sim x) \sqcup ((x \sqcap y) \sqcap \sim y)$ 
  by (rule conj-disj-distrib)
also have ... =  $(y \sqcap (x \sqcap \sim x)) \sqcup (x \sqcap (y \sqcap \sim y))$ 
  by (simp only: conj-ac)
finally show  $(x \sqcap y) \sqcap (\sim x \sqcup \sim y) = \mathbf{0}$ 
  by (simp only: conj-cancel-right conj-zero-right disj-zero-right)
next
have  $(x \sqcap y) \sqcup (\sim x \sqcup \sim y) = (x \sqcup (\sim x \sqcup \sim y)) \sqcap (y \sqcup (\sim x \sqcup \sim y))$ 
  by (rule disj-conj-distrib2)
also have ... =  $(\sim y \sqcup (x \sqcup \sim x)) \sqcap (\sim x \sqcup (y \sqcup \sim y))$ 
  by (simp only: disj-ac)
finally show  $(x \sqcap y) \sqcup (\sim x \sqcup \sim y) = \mathbf{1}$ 
  by (simp only: disj-cancel-right disj-one-right conj-one-right)
qed

```

```

lemma de-Morgan-disj [simp]:  $\sim (x \sqcup y) = \sim x \sqcap \sim y$ 
by (rule boolean.de-Morgan-conj [OF dual])

```

end

14.5 Symmetric Difference

```

locale boolean-xor = boolean +
  fixes xor :: 'a => 'a => 'a (infixr  $\oplus$  65)
  assumes xor-def:  $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$ 
begin

```

```

lemma xor-def2:
   $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$ 
by (simp only: xor-def conj-disj-distribs
      disj-ac conj-ac conj-cancel-right disj-zero-left)

```

```

lemma xor-commute:  $x \oplus y = y \oplus x$ 
by (simp only: xor-def conj-commute disj-commute)

```

```

lemma xor-assoc:  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$ 

```

proof –

```

let ?t =  $(x \sqcap y \sqcap z) \sqcup (x \sqcap \sim y \sqcap \sim z) \sqcup$ 
           $(\sim x \sqcap y \sqcap \sim z) \sqcup (\sim x \sqcap \sim y \sqcap z)$ 

```

```

have ?t  $\sqcup (z \sqcap x \sqcap \sim x) \sqcup (z \sqcap y \sqcap \sim y) =$ 
  ?t  $\sqcup (x \sqcap y \sqcap \sim y) \sqcup (x \sqcap z \sqcap \sim z)$ 

```

```

by (simp only: conj-cancel-right conj-zero-right)

```

```

thus  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$ 

```

```

apply (simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl)

```

```

apply (simp only: conj-disj-distribs conj-ac disj-ac)

```

```

done

```

qed

```

lemmas xor-ac =

```

xor-assoc xor-commute
mk-left-commute [where 'a = 'a, of xor, OF xor-assoc xor-commute]

lemma *xor-zero-right* [simp]: $x \oplus \mathbf{0} = x$
by (simp only: xor-def compl-zero conj-one-right conj-zero-right disj-zero-right)

lemma *xor-zero-left* [simp]: $\mathbf{0} \oplus x = x$
by (subst xor-commute) (rule xor-zero-right)

lemma *xor-one-right* [simp]: $x \oplus \mathbf{1} = \sim x$
by (simp only: xor-def compl-one conj-zero-right conj-one-right disj-zero-left)

lemma *xor-one-left* [simp]: $\mathbf{1} \oplus x = \sim x$
by (subst xor-commute) (rule xor-one-right)

lemma *xor-self* [simp]: $x \oplus x = \mathbf{0}$
by (simp only: xor-def conj-cancel-right conj-cancel-left disj-zero-right)

lemma *xor-left-self* [simp]: $x \oplus (x \oplus y) = y$
by (simp only: xor-assoc [symmetric] xor-self xor-zero-left)

lemma *xor-compl-left*: $\sim x \oplus y = \sim (x \oplus y)$
apply (simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl)
apply (simp only: conj-disj-distrib)
apply (simp only: conj-cancel-right conj-cancel-left)
apply (simp only: disj-zero-left disj-zero-right)
apply (simp only: disj-ac conj-ac)
done

lemma *xor-compl-right*: $x \oplus \sim y = \sim (x \oplus y)$
apply (simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl)
apply (simp only: conj-disj-distrib)
apply (simp only: conj-cancel-right conj-cancel-left)
apply (simp only: disj-zero-left disj-zero-right)
apply (simp only: disj-ac conj-ac)
done

lemma *xor-cancel-right* [simp]: $x \oplus \sim x = \mathbf{1}$
by (simp only: xor-compl-right xor-self compl-zero)

lemma *xor-cancel-left* [simp]: $\sim x \oplus x = \mathbf{1}$
by (subst xor-commute) (rule xor-cancel-right)

lemma *conj-xor-distrib*: $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

proof –

have $(x \sqcap y \sqcap \sim z) \sqcup (x \sqcap \sim y \sqcap z) =$
 $(y \sqcap x \sqcap \sim z) \sqcup (z \sqcap x \sqcap \sim y) \sqcup (x \sqcap y \sqcap \sim z) \sqcup (x \sqcap \sim y \sqcap z)$
by (simp only: conj-cancel-right conj-zero-right disj-zero-left)
thus $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

by (*simp* (*no-asm-use*) *only*:
xor-def de-Morgan-disj de-Morgan-conj double-compl
conj-disj-distrib conj-ac disj-ac)
qed

lemma *conj-xor-distrib2*:
 $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$
proof –
have $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$
by (*rule conj-xor-distrib*)
thus $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$
by (*simp only: conj-commute*)
qed

lemmas *conj-xor-distrib* =
conj-xor-distrib conj-xor-distrib2

end

end

15 WordGenLib: Miscellaneous Library for Words

theory *WordGenLib* **imports** *WordShift Boolean-Algebra*
begin

declare *of-nat-2p* [*simp*]

lemma *word-int-cases*:
 $\llbracket \bigwedge n. \llbracket (x :: 'a :: \text{len } 0 \text{ word}) = \text{word-of-int } n; 0 \leq n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$
 $\implies P$
by (*cases x rule: word-uint.Abs-cases*) (*simp add: uints-num*)

lemma *word-nat-cases* [*cases type: word*]:
 $\llbracket \bigwedge n. \llbracket (x :: 'a :: \text{len } \text{word}) = \text{of-nat } n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$
 $\implies P$
by (*cases x rule: word-unat.Abs-cases*) (*simp add: unats-def*)

lemma *max-word-eq*:
 $(\text{max-word} :: 'a :: \text{len } \text{word}) = 2^{\text{len-of TYPE('a)}} - 1$
by (*simp add: max-word-def word-of-int-hom-syms word-of-int-2p*)

lemma *max-word-max* [*simp,intro!*]:
 $n \leq \text{max-word}$
by (*cases n rule: word-int-cases*)
(simp add: max-word-def word-le-def int-word-uint int-mod-eq')

lemma *word-of-int-2p-len*:
 $word-of-int (2 \wedge len-of TYPE('a)) = (0::'a::len0 word)$
by (*subst word-uint.Abs-norm [symmetric]*)
(simp add: word-of-int-hom-syms)

lemma *word-pow-0*:
 $(2::'a::len word) \wedge len-of TYPE('a) = 0$
proof –
have $word-of-int (2 \wedge len-of TYPE('a)) = (0::'a word)$
by (*rule word-of-int-2p-len*)
thus *?thesis* **by** (*simp add: word-of-int-2p*)
qed

lemma *max-word-wrap*: $x + 1 = 0 \implies x = max-word$
apply (*simp add: max-word-eq*)
apply *uint-arith*
apply *auto*
apply (*simp add: word-pow-0*)
done

lemma *max-word-minus*:
 $max-word = (-1::'a::len word)$
proof –
have $-1 + 1 = (0::'a word)$ **by** *simp*
thus *?thesis* **by** (*rule max-word-wrap [symmetric]*)
qed

lemma *max-word-bl [simp]*:
 $to-bl (max-word::'a::len word) = replicate (len-of TYPE('a)) True$
by (*subst max-word-minus to-bl-n1*) *+ simp*

lemma *max-test-bit [simp]*:
 $(max-word::'a::len word) !! n = (n < len-of TYPE('a))$
by (*auto simp add: test-bit-bl word-size*)

lemma *word-and-max [simp]*:
 $x AND max-word = x$
by (*rule word-eqI*) (*simp add: word-ops-nth-size word-size*)

lemma *word-or-max [simp]*:
 $x OR max-word = max-word$
by (*rule word-eqI*) (*simp add: word-ops-nth-size word-size*)

lemma *word-ao-dist2*:
 $x AND (y OR z) = x AND y OR x AND (z::'a::len0 word)$
by (*rule word-eqI*) (*auto simp add: word-ops-nth-size word-size*)

lemma *word-oa-dist2*:
 $x OR y AND z = (x OR y) AND (x OR (z::'a::len0 word))$

by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-and-not [simp]:

$x \text{ AND } \text{NOT } x = (0::'a::\text{len } 0 \text{ word})$

by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-or-not [simp]:

$x \text{ OR } \text{NOT } x = \text{max-word}$

by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-boolean:

boolean (op AND) (op OR) bitNOT 0 max-word

apply (rule boolean.intro)

apply (rule word-bw-assocs)

apply (rule word-bw-assocs)

apply (rule word-bw-comms)

apply (rule word-bw-comms)

apply (rule word-ao-dist2)

apply (rule word-oa-dist2)

apply (rule word-and-max)

apply (rule word-log-esimps)

apply (rule word-and-not)

apply (rule word-or-not)

done

interpretation word-bool-alg:

boolean [op AND op OR bitNOT 0 max-word]

by (rule word-boolean)

lemma word-xor-and-or:

$x \text{ XOR } y = x \text{ AND } \text{NOT } y \text{ OR } \text{NOT } x \text{ AND } (y::'a::\text{len } 0 \text{ word})$

by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

interpretation word-bool-alg:

boolean-xor [op AND op OR bitNOT 0 max-word op XOR]

apply (rule boolean-xor.intro)

apply (rule word-boolean)

apply (rule boolean-xor-axioms.intro)

apply (rule word-xor-and-or)

done

lemma shiftr-0 [iff]:

$(x::'a::\text{len } 0 \text{ word}) \gg 0 = x$

by (simp add: shiftr-bl)

lemma shiffl-0 [simp]:

$(x :: 'a :: \text{len } \text{word}) \ll 0 = x$

by (simp add: shiffl-t2n)

```

lemma shiftr1-1 [simp]:
  (1::a::len word) << n = 2^n
  by (simp add: shiftr1-t2n)

lemma uint-lt-0 [simp]:
  uint x < 0 = False
  by (simp add: linorder-not-less)

lemma shiftr1-1 [simp]:
  shiftr1 (1::a::len word) = 0
  by (simp add: shiftr1-def word-0-alt)

lemma shiftr-1 [simp]:
  (1::a::len word) >> n = (if n = 0 then 1 else 0)
  by (induct n) (auto simp: shiftr-def)

lemma word-less-1 [simp]:
  ((x::a::len word) < 1) = (x = 0)
  by (simp add: word-less-nat-alt unat-0-iff)

lemma to-bl-mask:
  to-bl (mask n :: a::len word) =
    replicate (len-of TYPE(a) - n) False @
    replicate (min (len-of TYPE(a)) n) True
  by (simp add: mask-bl word-rep-drop min-def)

lemma map-replicate-True:
  n = length xs ==>
    map (λ(x,y). x & y) (zip xs (replicate n True)) = xs
  by (induct xs arbitrary: n) auto

lemma map-replicate-False:
  n = length xs ==> map (λ(x,y). x & y)
    (zip xs (replicate n False)) = replicate n False
  by (induct xs arbitrary: n) auto

lemma bl-and-mask:
  fixes w :: a::len word
  fixes n
  defines n' ≡ len-of TYPE(a) - n
  shows to-bl (w AND mask n) = replicate n' False @ drop n' (to-bl w)
proof -
  note [simp] = map-replicate-True map-replicate-False
  have to-bl (w AND mask n) =
    app2 op & (to-bl w) (to-bl (mask n::a::len word))
  by (simp add: bl-word-and)
  also
  have to-bl w = take n' (to-bl w) @ drop n' (to-bl w) by simp
  also

```

```

have app2 op & ... (to-bl (mask n::'a::len word)) =
  replicate n' False @ drop n' (to-bl w)
unfolding to-bl-mask n'-def app2-def
by (subst zip-append) auto
finally
show ?thesis .
qed

```

```

lemma drop-rev-takefill:
  length xs ≤ n ==>
    drop (n - length xs) (rev (takefill False n (rev xs))) = xs
by (simp add: takefill-alt rev-take)

```

```

lemma map-nth-0 [simp]:
  map (op !! (0::'a::len0 word)) xs = replicate (length xs) False
by (induct xs) auto

```

```

lemma uint-plus-if-size:
  uint (x + y) =
    (if uint x + uint y < 2^size x then
      uint x + uint y
    else
      uint x + uint y - 2^size x)
by (simp add: word-arith-alts int-word-uint mod-add-if-z
  word-size)

```

```

lemma unat-plus-if-size:
  unat (x + (y::'a::len word)) =
    (if unat x + unat y < 2^size x then
      unat x + unat y
    else
      unat x + unat y - 2^size x)
apply (subst word-arith-nat-defs)
apply (subst unat-of-nat)
apply (simp add: mod-nat-add word-size)
done

```

```

lemma word-neq-0-conv [simp]:
  fixes w :: 'a :: len word
  shows (w ≠ 0) = (0 < w)
proof -
  have 0 ≤ w by (rule word-zero-le)
  thus ?thesis by (auto simp add: word-less-def)
qed

```

```

lemma max-lt:
  unat (max a b div c) = unat (max a b) div unat (c:: 'a :: len word)
apply (subst word-arith-nat-defs)
apply (subst word-unat.eq-norm)

```

```

apply (subst mod-if)
apply clarsimp
apply (erule notE)
apply (insert div-le-dividend [of unat (max a b) unat c])
apply (erule order-le-less-trans)
apply (insert unat-lt2p [of max a b])
apply simp
done

```

lemma *uint-sub-if-size*:

```

uint (x - y) =
  (if uint y ≤ uint x then
    uint x - uint y
  else
    uint x - uint y + 2size x)
by (simp add: word-arith-alts int-word-uint mod-sub-if-z
      word-size)

```

lemma *unat-sub-simple*:

```

x ≤ y ==> unat (y - x) = unat y - unat x
by (simp add: unat-def uint-sub-if-size word-le-def nat-diff-distrib)

```

lemmas *unat-sub = unat-sub-simple*

lemma *word-less-sub1*:

```

fixes x :: 'a :: len word
shows x ≠ 0 ==> 1 < x = (0 < x - 1)
by (simp add: unat-sub-if-size word-less-nat-alt)

```

lemma *word-le-sub1*:

```

fixes x :: 'a :: len word
shows x ≠ 0 ==> 1 ≤ x = (0 ≤ x - 1)
by (simp add: unat-sub-if-size order-le-less word-less-nat-alt)

```

lemmas *word-less-sub1-numberof [simp] =*
word-less-sub1 [of number-of w, standard]

lemmas *word-le-sub1-numberof [simp] =*
word-le-sub1 [of number-of w, standard]

lemma *word-of-int-minus*:

```

word-of-int (2len-of TYPE('a) - i) = (word-of-int (-i)::'a::len word)

```

proof –

```

have x: 2len-of TYPE('a) - i = -i + 2len-of TYPE('a) by simp
show ?thesis
  apply (subst x)
  apply (subst word-uint.Abs-norm [symmetric], subst zmod-zadd-self2)
  apply simp
done

```

qed

lemmas *word-of-int-inj* =
word-uint.Abs-inject [*unfolded uints-num, simplified*]

lemma *word-le-less-eq*:
 $(x :: 'z :: \text{len } \text{word}) \leq y = (x = y \vee x < y)$
by (*auto simp add: word-less-def*)

lemma *mod-plus-cong*:
assumes 1: $(b :: \text{int}) = b'$
and 2: $x \bmod b' = x' \bmod b'$
and 3: $y \bmod b' = y' \bmod b'$
and 4: $x' + y' = z'$
shows $(x + y) \bmod b = z' \bmod b'$
proof –
from 1 2[*symmetric*] 3[*symmetric*] **have** $(x + y) \bmod b = (x' \bmod b' + y' \bmod b') \bmod b'$
by (*simp add: zmod-zadd1-eq[symmetric]*)
also have $\dots = (x' + y') \bmod b'$
by (*simp add: zmod-zadd1-eq[symmetric]*)
finally show ?thesis **by** (*simp add: 4*)
qed

lemma *mod-minus-cong*:
assumes 1: $(b :: \text{int}) = b'$
and 2: $x \bmod b' = x' \bmod b'$
and 3: $y \bmod b' = y' \bmod b'$
and 4: $x' - y' = z'$
shows $(x - y) \bmod b = z' \bmod b'$
using *assms*
apply (*subst zmod-zsub-left-eq*)
apply (*subst zmod-zsub-right-eq*)
apply (*simp add: zmod-zsub-left-eq [symmetric] zmod-zsub-right-eq [symmetric]*)
done

lemma *word-induct-less*:
 $\llbracket P (0 :: 'a :: \text{len } \text{word}); \bigwedge n. \llbracket n < m; P n \rrbracket \implies P (1 + n) \rrbracket \implies P m$
apply (*cases m*)
apply *atomize*
apply (*erule rev-mp*)
apply (*rule-tac x=m in spec*)
apply (*induct-tac n*)
apply *simp*
apply *clarsimp*
apply (*erule impE*)
apply *clarsimp*
apply (*erule-tac x=n in allE*)
apply (*erule impE*)
apply (*simp add: unat-arith-simps*)

```

  apply (clarsimp simp: unat-of-nat)
  apply simp
  apply (erule-tac x=of-nat na in allE)
  apply (erule impE)
  apply (simp add: unat-arith-simps)
  apply (clarsimp simp: unat-of-nat)
  apply simp
done

```

lemma *word-induct*:

```

[[P (0::'a::len word);  $\bigwedge n. P n \implies P (1 + n)$ ]]  $\implies P m$ 
by (erule word-induct-less, simp)

```

lemma *word-induct2* [*induct type*]:

```

[[P 0;  $\bigwedge n. [1 + n \neq 0; P n] \implies P (1 + n)$ ]]  $\implies P (n::'b::len word)$ 
  apply (rule word-induct, simp)
  apply (case-tac 1+n = 0, auto)
done

```

constdefs

```

word-rec :: 'a  $\implies$  ('b::len word  $\implies$  'a  $\implies$  'a)  $\implies$  'b word  $\implies$  'a
word-rec forZero forSuc n  $\equiv$  nat-rec forZero (forSuc  $\circ$  of-nat) (unat n)

```

lemma *word-rec-0*: *word-rec z s 0 = z*

```

by (simp add: word-rec-def)

```

lemma *word-rec-Suc*:

```

1 + n  $\neq$  (0::'a::len word)  $\implies$  word-rec z s (1 + n) = s n (word-rec z s n)
  apply (simp add: word-rec-def unat-word-ariths)
  apply (subst nat-mod-eq')
  apply (cut-tac x=n in unat-lt2p)
  apply (drule Suc-mono)
  apply (simp add: less-Suc-eq-le)
  apply (simp only: order-less-le, simp)
  apply (erule contrapos-pn, simp)
  apply (drule arg-cong[where f=of-nat])
  apply simp
  apply (subst (asm) word-unat.Rep-Abs-A.Rep-inverse[of n])
  apply simp
  apply simp
done

```

lemma *word-rec-Pred*:

```

n  $\neq$  0  $\implies$  word-rec z s n = s (n - 1) (word-rec z s (n - 1))
  apply (rule subst[where t=n and s=1 + (n - 1)])
  apply simp
  apply (subst word-rec-Suc)
  apply simp
  apply simp

```

done

lemma *word-rec-in*:

$f (\text{word-rec } z (\lambda-. f) n) = \text{word-rec } (f z) (\lambda-. f) n$
by (*induct n*) (*simp-all add: word-rec-0 word-rec-Suc*)

lemma *word-rec-in2*:

$f n (\text{word-rec } z f n) = \text{word-rec } (f 0 z) (f \circ \text{op} + 1) n$
by (*induct n*) (*simp-all add: word-rec-0 word-rec-Suc*)

lemma *word-rec-twice*:

$m \leq n \implies \text{word-rec } z f n = \text{word-rec } (\text{word-rec } z f (n - m)) (f \circ \text{op} + (n - m)) m$
apply (*erule rev-mp*)
apply (*rule-tac x=z in spec*)
apply (*rule-tac x=f in spec*)
apply (*induct n*)
apply (*simp add: word-rec-0*)
apply *clarsimp*
apply (*rule-tac t=1 + n - m and s=1 + (n - m) in subst*)
apply *simp*
apply (*case-tac 1 + (n - m) = 0*)
apply (*simp add: word-rec-0*)
apply (*rule-tac f = word-rec ?a ?b in arg-cong*)
apply (*rule-tac t=m and s=m + (1 + (n - m)) in subst*)
apply *simp*
apply (*simp (no-asm-use)*)
apply (*simp add: word-rec-Suc word-rec-in2*)
apply (*erule impE*)
apply *uint-arith*
apply (*drule-tac x=x \circ op + 1 in spec*)
apply (*drule-tac x=x 0 xa in spec*)
apply *simp*
apply (*rule-tac t=\lambda a. x (1 + (n - m + a)) and s=\lambda a. x (1 + (n - m) + a)*
in subst)
apply (*clarsimp simp add: expand-fun-eq*)
apply (*rule-tac t=(1 + (n - m + xb)) and s=1 + (n - m) + xb in subst*)
apply *simp*
apply (*rule refl*)
apply (*rule refl*)
done

lemma *word-rec-id*: $\text{word-rec } z (\lambda-. \text{id}) n = z$

by (*induct n*) (*auto simp add: word-rec-0 word-rec-Suc*)

lemma *word-rec-id-eq*: $\forall m < n. f m = \text{id} \implies \text{word-rec } z f n = z$

apply (*erule rev-mp*)

apply (*induct n*)

apply (*auto simp add: word-rec-0 word-rec-Suc*)

```

apply (drule spec, erule mp)
apply uint-arith
apply (drule-tac x=n in spec, erule impE)
apply uint-arith
apply simp
done

```

lemma word-rec-max:

```

 $\forall m \geq n. m \neq -1 \longrightarrow f m = id \implies \text{word-rec } z f -1 = \text{word-rec } z f n$ 
apply (subst word-rec-twice[where n=-1 and m=-1 - n])
apply simp
apply simp
apply (rule word-rec-id-eq)
apply clarsimp
apply (drule spec, rule mp, erule mp)
apply (rule word-plus-mono-right2[OF - order-less-imp-le])
prefer 2
apply assumption
apply simp
apply (erule contrapos-pn)
apply simp
apply (drule arg-cong[where f= $\lambda x. x - n$ ])
apply simp
done

```

lemma unatSuc:

```

 $1 + n \neq (0::'a::\text{len word}) \implies \text{unat } (1 + n) = \text{Suc } (\text{unat } n)$ 
by unat-arith

```

end

16 WordMain: Main Word Library

```

theory WordMain imports WordGenLib
begin

```

```

lemmas word-no-1 [simp] = word-1-no [symmetric]
lemmas word-no-0 [simp] = word-0-no [symmetric]

```

```

declare word-0-bl [simp]
declare bin-to-bl-def [simp]
declare to-bl-0 [simp]
declare of-bl-True [simp]

```

Examples

```

types word32 = 32 word
types word8 = 8 word
types byte = word8

```

for more see WordExampes.thy
end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.