

Isabelle/FOL — First-Order Logic

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1 Intuitionistic first-order logic

theory *IFOL*

imports *Pure*

uses

~~/src/Provers/splitter.ML

~~/src/Provers/hypsubst.ML

~~/src/Tools/IsaPlanner/zipper.ML

~~/src/Tools/IsaPlanner/isand.ML

~~/src/Tools/IsaPlanner/rw-tools.ML

~~/src/Tools/IsaPlanner/rw-inst.ML

~~/src/Provers/eqsubst.ML

~~/src/Provers/quantifier1.ML

~~/src/Provers/project-rule.ML

(fologic.ML)

(hypsubstdata.ML)

(intprover.ML)

begin

1.1 Syntax and axiomatic basis

global

classes *term*

defaultsort *term*

typeddecl *o*

judgment

Trueprop :: $o \Rightarrow prop$ ($((-) \ 5)$)

consts

True :: *o*

False :: *o*

op = :: $['a, 'a] \Rightarrow o$ (**infixl** = 50)

Not :: $o \Rightarrow o$ (\sim - [40] 40)

op & :: $[o, o] \Rightarrow o$ (**infixr** & 35)

op | :: $[o, o] \Rightarrow o$ (**infixr** | 30)

op --> :: $[o, o] \Rightarrow o$ (**infixr** --> 25)

op <-> :: $[o, o] \Rightarrow o$ (**infixr** <-> 25)

All :: $('a \Rightarrow o) \Rightarrow o$ (**binder** *ALL* 10)

Ex :: $('a \Rightarrow o) \Rightarrow o$ (**binder** *EX* 10)

Ex1 :: $('a \Rightarrow o) \Rightarrow o$ (**binder** *EX!* 10)

abbreviation

not-equal :: $['a, 'a] \Rightarrow o$ (**infixl** $\sim =$ 50) **where**

$x \sim = y == \sim (x = y)$

notation (*xsymbols*)

not-equal (**infixl** \neq 50)

notation (*HTML output*)

not-equal (**infixl** \neq 50)

notation (*xsymbols*)

Not (\neg - [40] 40) **and**

op & (**infixr** \wedge 35) **and**

op | (**infixr** \vee 30) **and**

All (**binder** \forall 10) **and**

Ex (**binder** \exists 10) **and**

Ex1 (**binder** $\exists!$ 10) **and**

$op \dashrightarrow$ (**infixr** \longrightarrow 25) **and**
 $op \longleftrightarrow$ (**infixr** \longleftrightarrow 25)

notation (*HTML output*)

Not (\neg - [40] 40) **and**
 $op \ \&$ (**infixr** \wedge 35) **and**
 $op \ |$ (**infixr** \vee 30) **and**
 All (**binder** \forall 10) **and**
 Ex (**binder** \exists 10) **and**
 $Ex1$ (**binder** $\exists!$ 10)

local

finalconsts

$False$ All Ex
 $op =$
 $op \ \&$
 $op \ |$
 $op \dashrightarrow$

axioms

$refl:$ $a=a$

$conjI:$ $[| P; Q |] \implies P \& Q$
 $conjunct1:$ $P \& Q \implies P$
 $conjunct2:$ $P \& Q \implies Q$

$disjI1:$ $P \implies P | Q$
 $disjI2:$ $Q \implies P | Q$
 $disjE:$ $[| P | Q; P \implies R; Q \implies R |] \implies R$

$impI:$ $(P \implies Q) \implies P \dashrightarrow Q$
 $mp:$ $[| P \dashrightarrow Q; P |] \implies Q$

$FalseE:$ $False \implies P$

$allI:$ $(!!x. P(x)) \implies (ALL\ x. P(x))$
 $spec:$ $(ALL\ x. P(x)) \implies P(x)$

$exI:$ $P(x) \implies (EX\ x. P(x))$
 $exE:$ $[| EX\ x. P(x); !!x. P(x) \implies R |] \implies R$

eq-reflection: $(x=y) \implies (x==y)$
iff-reflection: $(P<->Q) \implies (P==Q)$

lemmas *strip* = *impI allI*

Thanks to Stephan Merz

theorem *subst*:

assumes *eq*: $a = b$ **and** $p: P(a)$

shows $P(b)$

$\langle proof \rangle$

defs

True-def: $True == False \dashv\dashv False$

not-def: $\sim P == P \dashv\dashv False$

iff-def: $P<->Q == (P \dashv\dashv Q) \ \& \ (Q \dashv\dashv P)$

ex1-def: $Ex1(P) == EX\ x. P(x) \ \& \ (ALL\ y. P(y) \dashv\dashv y=x)$

1.2 Lemmas and proof tools

lemma *TrueI*: *True*

$\langle proof \rangle$

lemma *conjE*:

assumes *major*: $P \ \& \ Q$

and $r: [| P; Q |] \implies R$

shows R

$\langle proof \rangle$

lemma *impE*:

assumes *major*: $P \dashv\dashv Q$

and P

and $r: Q \implies R$

shows R

$\langle proof \rangle$

lemma *allE*:

assumes *major*: $ALL\ x. P(x)$

and $r: P(x) ==> R$
shows R
 $\langle proof \rangle$

lemma *all-dupE*:
assumes *major*: $ALL\ x.\ P(x)$
and $r: [P(x); ALL\ x.\ P(x)] ==> R$
shows R
 $\langle proof \rangle$

lemma *notI*: $(P ==> False) ==> \sim P$
 $\langle proof \rangle$

lemma *notE*: $[\sim P; P] ==> R$
 $\langle proof \rangle$

lemma *rev-notE*: $[P; \sim P] ==> R$
 $\langle proof \rangle$

lemma *not-to-imp*:
assumes $\sim P$
and $r: P --> False ==> Q$
shows Q
 $\langle proof \rangle$

lemma *rev-mp*: $[P; P --> Q] ==> Q$
 $\langle proof \rangle$

lemma *contrapos*:
assumes *major*: $\sim Q$
and *minor*: $P ==> Q$
shows $\sim P$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *iffI*: $[| P ==> Q; Q ==> P |] ==> P <-> Q$
 $\langle proof \rangle$

lemma *iffE*:
assumes *major*: $P <-> Q$
and *r*: $P --> Q ==> Q --> P ==> R$
shows *R*
 $\langle proof \rangle$

lemma *iffD1*: $[| P <-> Q; P |] ==> Q$
 $\langle proof \rangle$

lemma *iffD2*: $[| P <-> Q; Q |] ==> P$
 $\langle proof \rangle$

lemma *rev-iffD1*: $[| P; P <-> Q |] ==> Q$
 $\langle proof \rangle$

lemma *rev-iffD2*: $[| Q; P <-> Q |] ==> P$
 $\langle proof \rangle$

lemma *iff-refl*: $P <-> P$
 $\langle proof \rangle$

lemma *iff-sym*: $Q <-> P ==> P <-> Q$
 $\langle proof \rangle$

lemma *iff-trans*: $[| P <-> Q; Q <-> R |] ==> P <-> R$
 $\langle proof \rangle$

lemma *exII*:
 $P(a) ==> (!x. P(x) ==> x=a) ==> EX! x. P(x)$
 $\langle proof \rangle$

lemma *ex-exII*:
 $EX x. P(x) ==> (!x y. [| P(x); P(y) |] ==> x=y) ==> EX! x. P(x)$
 $\langle proof \rangle$

lemma *exIE*:
 $EX! x. P(x) ==> (!x. [| P(x); ALL y. P(y) --> y=x |] ==> R) ==> R$

$\langle proof \rangle$

$\langle ML \rangle$

lemma *conj-cong*:

assumes $P \leftrightarrow P'$
and $P' \implies Q \leftrightarrow Q'$
shows $(P \& Q) \leftrightarrow (P' \& Q')$
 $\langle proof \rangle$

lemma *conj-cong2*:

assumes $P \leftrightarrow P'$
and $P' \implies Q \leftrightarrow Q'$
shows $(Q \& P) \leftrightarrow (Q' \& P')$
 $\langle proof \rangle$

lemma *disj-cong*:

assumes $P \leftrightarrow P'$ and $Q \leftrightarrow Q'$
shows $(P | Q) \leftrightarrow (P' | Q')$
 $\langle proof \rangle$

lemma *imp-cong*:

assumes $P \leftrightarrow P'$
and $P' \implies Q \leftrightarrow Q'$
shows $(P \rightarrow Q) \leftrightarrow (P' \rightarrow Q')$
 $\langle proof \rangle$

lemma *iff-cong*: $[[P \leftrightarrow P'; Q \leftrightarrow Q']] \implies (P \leftrightarrow Q) \leftrightarrow (P' \leftrightarrow Q')$
 $\langle proof \rangle$

lemma *not-cong*: $P \leftrightarrow P' \implies \sim P \leftrightarrow \sim P'$
 $\langle proof \rangle$

lemma *all-cong*:

assumes $!!x. P(x) \leftrightarrow Q(x)$
shows $(ALL x. P(x)) \leftrightarrow (ALL x. Q(x))$
 $\langle proof \rangle$

lemma *ex-cong*:

assumes $!!x. P(x) \leftrightarrow Q(x)$
shows $(EX x. P(x)) \leftrightarrow (EX x. Q(x))$
 $\langle proof \rangle$

lemma *ex1-cong*:

assumes $!!x. P(x) <-> Q(x)$
shows $(EX! x. P(x)) <-> (EX! x. Q(x))$
 $\langle proof \rangle$

lemma *sym*: $a=b ==> b=a$
 $\langle proof \rangle$

lemma *trans*: $[| a=b; b=c |] ==> a=c$
 $\langle proof \rangle$

lemma *not-sym*: $b \sim = a ==> a \sim = b$
 $\langle proof \rangle$

lemma *def-imp-iff*: $(A == B) ==> A <-> B$
 $\langle proof \rangle$

lemma *meta-eq-to-obj-eq*: $(A == B) ==> A = B$
 $\langle proof \rangle$

lemma *meta-eq-to-iff*: $x==y ==> x<->y$
 $\langle proof \rangle$

lemma *ssubst*: $[| b = a; P(a) |] ==> P(b)$
 $\langle proof \rangle$

lemma *ex1-equalsE*:
 $[| EX! x. P(x); P(a); P(b) |] ==> a=b$
 $\langle proof \rangle$

lemma *subst-context*: $[| a=b |] ==> t(a)=t(b)$
 $\langle proof \rangle$

lemma *subst-context2*: $[| a=b; c=d |] ==> t(a,c)=t(b,d)$
 $\langle proof \rangle$

lemma *subst-context3*: $[| a=b; c=d; e=f |] ==> t(a,c,e)=t(b,d,f)$
 $\langle proof \rangle$

lemma *box-equals*: $[| a=b; a=c; b=d |] ==> c=d$

$\langle proof \rangle$

lemma *simp-equals*: $[\mid a=c; \ b=d; \ c=d \mid] ==> a=b$
 $\langle proof \rangle$

lemma *pred1-cong*: $a=a' ==> P(a) <-> P(a')$
 $\langle proof \rangle$

lemma *pred2-cong*: $[\mid a=a'; \ b=b' \mid] ==> P(a,b) <-> P(a',b')$
 $\langle proof \rangle$

lemma *pred3-cong*: $[\mid a=a'; \ b=b'; \ c=c' \mid] ==> P(a,b,c) <-> P(a',b',c')$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *eq-cong*: $[\mid a = a'; \ b = b' \mid] ==> a = b <-> a' = b'$
 $\langle proof \rangle$

lemma *conj-impE*:
 assumes *major*: $(P \& Q) --> S$
 and *r*: $P --> (Q --> S) ==> R$
 shows *R*
 $\langle proof \rangle$

lemma *disj-impE*:
 assumes *major*: $(P \mid Q) --> S$
 and *r*: $[\mid P --> S; \ Q --> S \mid] ==> R$
 shows *R*
 $\langle proof \rangle$

lemma *imp-impE*:
 assumes *major*: $(P --> Q) --> S$
 and *r1*: $[\mid P; \ Q --> S \mid] ==> Q$
 and *r2*: $S ==> R$
 shows *R*
 $\langle proof \rangle$

lemma *not-impE*:
 $\sim P \multimap S \implies (P \implies \text{False}) \implies (S \implies R) \implies R$
 $\langle \text{proof} \rangle$

lemma *iff-impE*:
assumes *major*: $(P \multimap Q) \multimap S$
and *r1*: $\llbracket P; Q \multimap S \rrbracket \implies Q$
and *r2*: $\llbracket Q; P \multimap S \rrbracket \implies P$
and *r3*: $S \implies R$
shows *R*
 $\langle \text{proof} \rangle$

lemma *all-impE*:
assumes *major*: $(\text{ALL } x. P(x)) \multimap S$
and *r1*: $\llbracket x. P(x) \rrbracket$
and *r2*: $S \implies R$
shows *R*
 $\langle \text{proof} \rangle$

lemma *ex-impE*:
assumes *major*: $(\text{EX } x. P(x)) \multimap S$
and *r*: $P(x) \multimap S \implies R$
shows *R*
 $\langle \text{proof} \rangle$

lemma *disj-imp-disj*:
 $P \mid Q \implies (P \implies R) \implies (Q \implies S) \implies R \mid S$
 $\langle \text{proof} \rangle$

$\langle \text{ML} \rangle$

lemma *thin-refl*: $\llbracket x=x; \text{PROP } W \rrbracket \implies \text{PROP } W \langle \text{proof} \rangle$

$\langle \text{ML} \rangle$

1.3 Intuitionistic Reasoning

lemma *impE'*:
assumes *1*: $P \multimap Q$
and *2*: $Q \implies R$
and *3*: $P \multimap Q \implies P$
shows *R*
 $\langle \text{proof} \rangle$

lemma *allE'*:
 assumes 1: $\text{ALL } x. P(x)$
 and 2: $P(x) \implies \text{ALL } x. P(x) \implies Q$
 shows Q
 $\langle \text{proof} \rangle$

lemma *notE'*:
 assumes 1: $\sim P$
 and 2: $\sim P \implies P$
 shows R
 $\langle \text{proof} \rangle$

lemmas [*Pure.elim!*] = *disjE iffE FalseE conjE exE*
 and [*Pure.intro!*] = *iffI conjI impI TrueI notI allI refl*
 and [*Pure.elim 2*] = *allE notE' impE'*
 and [*Pure.intro*] = *exI disjI2 disjI1*

$\langle \text{ML} \rangle$

lemma *iff-not-sym*: $\sim (Q \longleftrightarrow P) \implies \sim (P \longleftrightarrow Q)$
 $\langle \text{proof} \rangle$

lemmas [*sym*] = *sym iff-sym not-sym iff-not-sym*
 and [*Pure.elim?*] = *iffD1 iffD2 impE*

lemma *eq-commute*: $a=b \longleftrightarrow b=a$
 $\langle \text{proof} \rangle$

1.4 Atomizing meta-level rules

lemma *atomize-all* [*atomize*]: $(!!x. P(x)) \implies \text{Trueprop } (\text{ALL } x. P(x))$
 $\langle \text{proof} \rangle$

lemma *atomize-imp* [*atomize*]: $(A \implies B) \implies \text{Trueprop } (A \longrightarrow B)$
 $\langle \text{proof} \rangle$

lemma *atomize-eq* [*atomize*]: $(x == y) \implies \text{Trueprop } (x = y)$
 $\langle \text{proof} \rangle$

lemma *atomize-iff* [*atomize*]: $(A == B) \implies \text{Trueprop } (A \longleftrightarrow B)$
 $\langle \text{proof} \rangle$

lemma *atomize-conj* [*atomize*]:
 includes *meta-conjunction-syntax*
 shows $(A \ \&\& \ B) \implies \text{Trueprop } (A \ \& \ B)$
 $\langle \text{proof} \rangle$

lemmas $[symmetric, rulify] = atomize-all\ atomize-imp$
and $[symmetric, defn] = atomize-all\ atomize-imp\ atomize-eq\ atomize-iff$

1.5 Calculational rules

lemma *forw-subst*: $a = b ==> P(b) ==> P(a)$
 $\langle proof \rangle$

lemma *back-subst*: $P(a) ==> a = b ==> P(b)$
 $\langle proof \rangle$

Note that this list of rules is in reverse order of priorities.

lemmas *basic-trans-rules* $[trans] =$
forw-subst
back-subst
rev-mp
mp
trans

1.6 “Let” declarations

nonterminals *letbinds letbind*

constdefs
 $Let :: ['a::\{\}, 'a ==> 'b] ==> ('b::\{\})$
 $Let(s, f) == f(s)$

syntax
 $-bind \quad :: [pttrn, 'a] ==> letbind \quad ((2- = / -) 10)$
 $\quad \quad :: letbind ==> letbinds \quad (-)$
 $-binds \quad :: [letbind, letbinds] ==> letbinds \quad (-; / -)$
 $-Let \quad \quad :: [letbinds, 'a] ==> 'a \quad ((let (-) / in (-)) 10)$

translations
 $-Let(-binds(b, bs), e) == -Let(b, -Let(bs, e))$
 $let\ x = a\ in\ e \quad == Let(a, \%x. e)$

lemma *LetI*:
assumes $!!x. x=t ==> P(u(x))$
shows $P(let\ x=t\ in\ u(x))$
 $\langle proof \rangle$

1.7 ML bindings

$\langle ML \rangle$

end

2 Classical first-order logic

```
theory FOL
imports IFOL
uses
  ~~ /src/Provers/classical.ML
  ~~ /src/Provers/blast.ML
  ~~ /src/Provers/clasimp.ML
  ~~ /src/Tools/induct.ML
  (cladata.ML)
  (blastdata.ML)
  (simpdata.ML)
begin
```

2.1 The classical axiom

```
axioms
  classical: ( $\sim P \implies P$ )  $\implies P$ 
```

2.2 Lemmas and proof tools

```
lemma ccontr: ( $\neg P \implies False$ )  $\implies P$ 
  <proof>
```

```
lemma disjCI: ( $\sim Q \implies P$ )  $\implies P \mid Q$ 
  <proof>
```

```
lemma ex-classical:
  assumes r:  $\sim (EX\ x. P(x)) \implies P(a)$ 
  shows  $EX\ x. P(x)$ 
  <proof>
```

```
lemma exCI:
  assumes r:  $ALL\ x. \sim P(x) \implies P(a)$ 
  shows  $EX\ x. P(x)$ 
  <proof>
```

```
lemma excluded-middle:  $\sim P \mid P$ 
  <proof>
```

<ML>

```
lemma case-split-thm:
  assumes r1:  $P \implies Q$ 
  and r2:  $\sim P \implies Q$ 
```

shows Q
 $\langle proof \rangle$

lemmas $case-split = case-split-thm$ [$case-names$ $True$ $False$]

$\langle ML \rangle$

lemma $impCE$:
 assumes $major$: $P \dashrightarrow Q$
 and $r1$: $\sim P \implies R$
 and $r2$: $Q \implies R$
 shows R
 $\langle proof \rangle$

lemma $impCE'$:
 assumes $major$: $P \dashrightarrow Q$
 and $r1$: $Q \implies R$
 and $r2$: $\sim P \implies R$
 shows R
 $\langle proof \rangle$

lemma $notnotD$: $\sim\sim P \implies P$
 $\langle proof \rangle$

lemma $contrapos2$: $[[Q; \sim P \implies \sim Q]] \implies P$
 $\langle proof \rangle$

lemma $iffCE$:
 assumes $major$: $P <-> Q$
 and $r1$: $[[P; Q]] \implies R$
 and $r2$: $[[\sim P; \sim Q]] \implies R$
 shows R
 $\langle proof \rangle$

lemma $alt-ex1E$:
 assumes $major$: $EX! x. P(x)$

and $r: !!x. [| P(x); ALL y y'. P(y) \& P(y') \dashv\vdash y=y' |] \implies R$
shows R
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *ex1-functional*: $[| EX! z. P(a,z); P(a,b); P(a,c) |] \implies b = c$
 $\langle proof \rangle$

lemma *True-implies-equals*: $(True \implies PROP P) == PROP P$
 $\langle proof \rangle$

lemma *uncurry*: $P \dashv\vdash Q \dashv\vdash R \implies P \& Q \dashv\vdash R$
 $\langle proof \rangle$

lemma *iff-allI*: $(!!x. P(x) \dashv\vdash Q(x)) \implies (ALL x. P(x)) \dashv\vdash (ALL x. Q(x))$
 $\langle proof \rangle$

lemma *iff-exI*: $(!!x. P(x) \dashv\vdash Q(x)) \implies (EX x. P(x)) \dashv\vdash (EX x. Q(x))$
 $\langle proof \rangle$

lemma *all-comm*: $(ALL x y. P(x,y)) \dashv\vdash (ALL y x. P(x,y))$ $\langle proof \rangle$

lemma *ex-comm*: $(EX x y. P(x,y)) \dashv\vdash (EX y x. P(x,y))$ $\langle proof \rangle$

$\langle ML \rangle$

2.3 Other simple lemmas

lemma *[simp]*: $((P \dashv\vdash R) \dashv\vdash (Q \dashv\vdash R)) \dashv\vdash ((P \dashv\vdash Q) \mid R)$
 $\langle proof \rangle$

lemma *[simp]*: $((P \dashv\vdash Q) \dashv\vdash (P \dashv\vdash R)) \dashv\vdash (P \dashv\vdash (Q \dashv\vdash R))$
 $\langle proof \rangle$

lemma *not-disj-iff-imp*: $\sim P \mid Q \dashv\vdash (P \dashv\vdash Q)$
 $\langle proof \rangle$

lemma *conj-mono*: $[| P1 \dashv\vdash Q1; P2 \dashv\vdash Q2 |] \implies (P1 \& P2) \dashv\vdash (Q1 \& Q2)$
 $\langle proof \rangle$

lemma *disj-mono*: $[| P1 \dashv\vdash Q1; P2 \dashv\vdash Q2 |] \implies (P1 \mid P2) \dashv\vdash (Q1 \mid Q2)$
 $\langle proof \rangle$

lemma *imp-mono*: $[| Q1 \dashv\vdash P1; P2 \dashv\vdash Q2 |] \implies (P1 \dashv\vdash P2) \dashv\vdash (Q1 \dashv\vdash Q2)$

<proof>

lemma *imp-refl*: $P \dashv\dashv P$
<proof>

lemma *ex-mono*: $(!!x. P(x) \dashv\dashv Q(x)) \implies (EX x. P(x)) \dashv\dashv (EX x. Q(x))$
<proof>

lemma *all-mono*: $(!!x. P(x) \dashv\dashv Q(x)) \implies (ALL x. P(x)) \dashv\dashv (ALL x. Q(x))$
<proof>

2.4 Proof by cases and induction

Proper handling of non-atomic rule statements.

constdefs

induct-forall **where** $induct-forall(P) == \forall x. P(x)$
induct-implies **where** $induct-implies(A, B) == A \longrightarrow B$
induct-equal **where** $induct-equal(x, y) == x = y$
induct-conj **where** $induct-conj(A, B) == A \wedge B$

lemma *induct-forall-eq*: $(!!x. P(x)) == Trueprop(induct-forall(\lambda x. P(x)))$
<proof>

lemma *induct-implies-eq*: $(A \implies B) == Trueprop(induct-implies(A, B))$
<proof>

lemma *induct-equal-eq*: $(x == y) == Trueprop(induct-equal(x, y))$
<proof>

lemma *induct-conj-eq*:

includes *meta-conjunction-syntax*
shows $(A \ \&\& \ B) == Trueprop(induct-conj(A, B))$
<proof>

lemmas *induct-atomize* = *induct-forall-eq induct-implies-eq induct-equal-eq induct-conj-eq*

lemmas *induct-rulify* [*symmetric, standard*] = *induct-atomize*

lemmas *induct-rulify-fallback* =

induct-forall-def induct-implies-def induct-equal-def induct-conj-def

hide *const induct-forall induct-implies induct-equal induct-conj*

Method setup.

<ML>

declare *case-split* [*cases type: o*]

end