

Machine Words in Isabelle/HOL

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Abstract

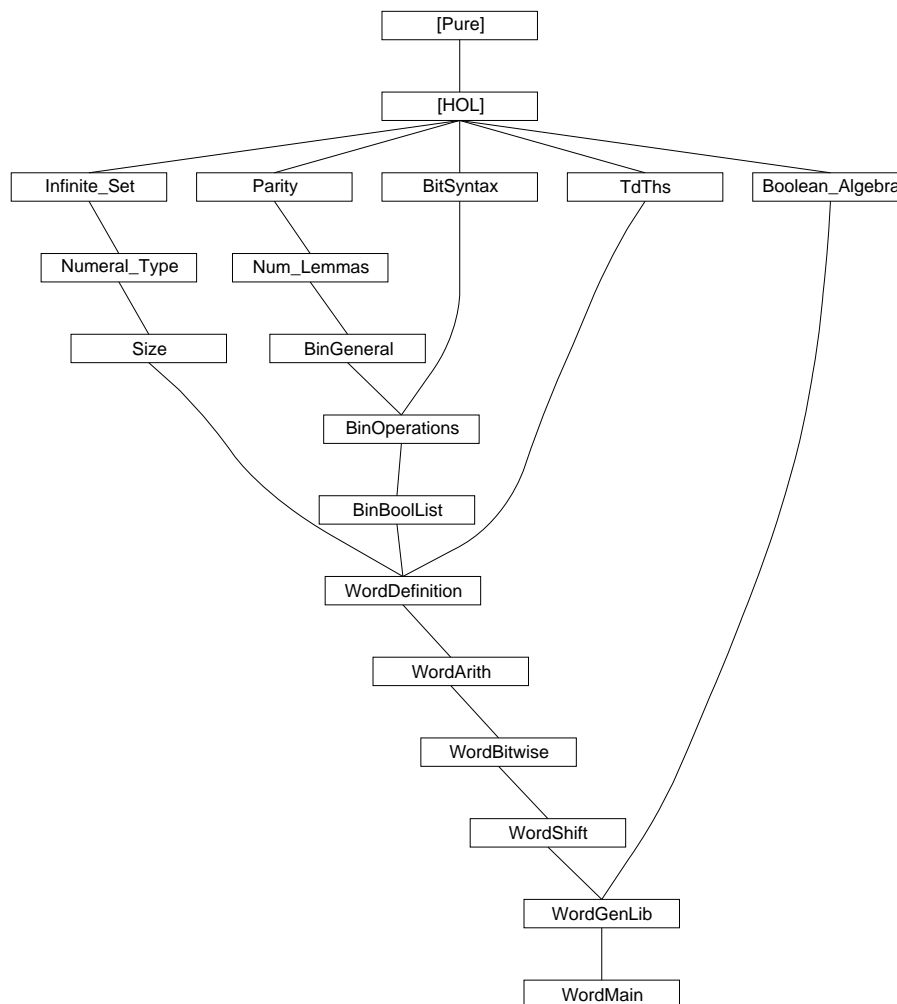
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
  imports Infinite-Set
begin
```

1.1 Preliminary lemmas

```
lemma inj-Inl [simp]: inj-on Inl A
  <proof>
```

```
lemma inj-Inr [simp]: inj-on Inr A
  <proof>
```

```
lemma inj-Some [simp]: inj-on Some A
  <proof>
```

```
lemma card-Plus:
  [| finite A; finite B |] ==> card (A <+> B) = card A + card B
  <proof>
```

```
lemma (in type-definition) univ:
  UNIV = Abs ‘ A
  <proof>
```

```
lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  <proof>
```

1.2 Cardinalities of types

```
syntax -type-card :: type => nat ((1CARD/(1'(-))))
```

```
translations CARD(t) => card (UNIV::t set)
```

```
<ML>
```

```
lemma card-unit: CARD(unit) = 1
  <proof>
```

```
lemma card-bool: CARD(bool) = 2
  <proof>
```

```
lemma card-prod: CARD('a::finite × 'b::finite) = CARD('a) * CARD('b)
  <proof>
```

```
lemma card-sum: CARD('a::finite + 'b::finite) = CARD('a) + CARD('b)
  <proof>
```

```
lemma card-option: CARD('a::finite option) = Suc CARD('a)
  <proof>
```

lemma *card-set*: $CARD('a::finite\ set) = 2 \wedge CARD('a)$
 $\langle proof \rangle$

1.3 Numeral Types

typedef (**open**) *num0* = *UNIV* :: *nat set* $\langle proof \rangle$
typedef (**open**) *num1* = *UNIV* :: *unit set* $\langle proof \rangle$
typedef (**open**) *'a bit0* = *UNIV* :: (*bool* * *'a*) *set* $\langle proof \rangle$
typedef (**open**) *'a bit1* = *UNIV* :: (*bool* * *'a*) *option set* $\langle proof \rangle$

instance *num1* :: *finite*
 $\langle proof \rangle$

instance *bit0* :: (*finite*) *finite*
 $\langle proof \rangle$

instance *bit1* :: (*finite*) *finite*
 $\langle proof \rangle$

lemma *card-num1*: $CARD(num1) = 1$
 $\langle proof \rangle$

lemma *card-bit0*: $CARD('a::finite\ bit0) = 2 * CARD('a)$
 $\langle proof \rangle$

lemma *card-bit1*: $CARD('a::finite\ bit1) = Suc\ (2 * CARD('a))$
 $\langle proof \rangle$

lemma *card-num0*: $CARD\ (num0) = 0$
 $\langle proof \rangle$

lemmas *card-univ-simps* [*simp*] =
card-unit
card-bool
card-prod
card-sum
card-option
card-set
card-num1
card-bit0
card-bit1
card-num0

1.4 Syntax

syntax
-NumeralType :: *num-const* => *type* (-)
-NumeralType0 :: *type* (0)
-NumeralType1 :: *type* (1)

translations

-*NumeralType1* == (*type*) *num1*
 -*NumeralType0* == (*type*) *num0*

⟨*ML*⟩

1.5 Classes with at least 1 and 2

Class *finite* already captures “at least 1”

lemma *zero-less-card-finite* [*simp*]:
 $0 < \text{CARD}('a::\text{finite})$
 ⟨*proof*⟩

lemma *one-le-card-finite* [*simp*]:
 $\text{Suc } 0 \leq \text{CARD}('a::\text{finite})$
 ⟨*proof*⟩

Class for cardinality “at least 2”

class *card2* = *finite* +
assumes *two-le-card*: $2 \leq \text{CARD}('a)$

lemma *one-less-card*: $\text{Suc } 0 < \text{CARD}('a::\text{card2})$
 ⟨*proof*⟩

instance *bit0* :: (*finite*) *card2*
 ⟨*proof*⟩

instance *bit1* :: (*finite*) *card2*
 ⟨*proof*⟩

1.6 Examples

term *TYPE*(10)

lemma $\text{CARD}(0) = 0$ ⟨*proof*⟩
lemma $\text{CARD}(17) = 17$ ⟨*proof*⟩

end

2 Size: The size class

theory *Size*
imports *Numeral-Type*
begin

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in `Numeral_Type`.

```
axclass len0 < type
```

```
consts
```

```
  len-of :: ('a :: len0 itself) ==> nat
```

Some theorems are only true on words with length greater 0.

```
axclass len < len0
```

```
  len-gt-0 [iff]: 0 < len-of TYPE ('a :: len0)
```

```
instance num0 :: len0 <proof>
```

```
instance num1 :: len0 <proof>
```

```
instance bit0 :: (len0) len0 <proof>
```

```
instance bit1 :: (len0) len0 <proof>
```

```
defs (overloaded)
```

```
  len-num0: len-of (x::num0 itself) == 0
```

```
  len-num1: len-of (x::num1 itself) == 1
```

```
  len-bit0: len-of (x::'a::len0 bit0 itself) == 2 * len-of TYPE ('a)
```

```
  len-bit1: len-of (x::'a::len0 bit1 itself) == 2 * len-of TYPE ('a) + 1
```

```
lemmas len-of-numeral-defs [simp] = len-num0 len-num1 len-bit0 len-bit1
```

```
instance num1 :: len <proof>
```

```
instance bit0 :: (len) len <proof>
```

```
instance bit1 :: (len0) len <proof>
```

```
lemma len-of TYPE(17) = 17 <proof>
```

```
lemma len-of TYPE(0) = 0 <proof>
```

```
lemma len-of TYPE('a::len0) = x  
  <proof>
```

```
end
```

3 Num-Lemmas: Useful Numerical Lemmas

```
theory Num-Lemmas imports Parity begin
```

```
lemma contentsI: y = {x} ==> contents y = x  
  <proof>
```

```
lemma prod-case-split: prod-case = split  
  <proof>
```

```
lemmas split-split = prod.split [unfolded prod-case-split]
```

lemmas *split-split-asm* = *prod.split-asm* [*unfolded prod-case-split*]
lemmas *split.splits* = *split-split split-split-asm*

lemmas *funpow-0* = *funpow.simps*(1)
lemmas *funpow-Suc* = *funpow.simps*(2)

lemma *nonemptyE*: $S \sim = \{\} \implies (!x. x : S \implies R) \implies R$
 ⟨*proof*⟩

lemma *gt-or-eq-0*: $0 < y \vee 0 = (y::nat)$ ⟨*proof*⟩

constdefs

mod-alt :: $'a \implies 'a \implies 'a :: \text{Divides.div}$
mod-alt $n\ m == n \text{ mod } m$

— alternative way of defining *bin-last*, *bin-rest*

bin-rl :: $int \implies int * bit$
bin-rl $w == \text{SOME } (r, l). w = r \text{ BIT } l$

declare *iszero-0* [*iff*]

lemmas *xtr1* = *xtrans*(1)
lemmas *xtr2* = *xtrans*(2)
lemmas *xtr3* = *xtrans*(3)
lemmas *xtr4* = *xtrans*(4)
lemmas *xtr5* = *xtrans*(5)
lemmas *xtr6* = *xtrans*(6)
lemmas *xtr7* = *xtrans*(7)
lemmas *xtr8* = *xtrans*(8)

lemma *Min-ne-Pls* [*iff*]:
Numeral.Min $\sim = \text{Numerals.Pl}$
 ⟨*proof*⟩

lemmas *Pls-ne-Min* [*iff*] = *Min-ne-Pls* [*symmetric*]

lemmas *PlsMin-defs* [*intro!*] =
Pls-def Min-def Pls-def [*symmetric*] *Min-def* [*symmetric*]

lemmas *PlsMin-simps* [*simp*] = *PlsMin-defs* [*THEN Eq-TrueI*]

lemma *number-of-False-cong*:
 $\text{False} \implies \text{number-of } x = \text{number-of } y$
 ⟨*proof*⟩

lemmas *nat-simps* = *diff-add-inverse2 diff-add-inverse*
lemmas *nat-iffs* = *le-add1 le-add2*

lemma *sum-imp-diff*: $j = k + i \implies j - i = (k :: nat)$

$\langle \text{proof} \rangle$

lemma nobm1:

$0 < (\text{number-of } w :: \text{nat}) ==>$
 $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Numeral.pred } w)$
 $\langle \text{proof} \rangle$

lemma of-int-power:

$\text{of-int } (a \wedge n) = (\text{of-int } a \wedge n :: 'a :: \{\text{recpower}, \text{comm-ring-1}\})$
 $\langle \text{proof} \rangle$

lemma zless2: $0 < (2 :: \text{int})$

$\langle \text{proof} \rangle$

lemmas $\text{zless2p } [\text{simp}] = \text{zless2 } [\text{THEN zero-less-power}]$

lemmas $\text{zle2p } [\text{simp}] = \text{zless2p } [\text{THEN order-less-imp-le}]$

lemmas $\text{pos-mod-sign2} = \text{zless2 } [\text{THEN pos-mod-sign } [\text{where } b = 2 :: \text{int}]]$

lemmas $\text{pos-mod-bound2} = \text{zless2 } [\text{THEN pos-mod-bound } [\text{where } b = 2 :: \text{int}]]$

— the inverse(s) of *number-of*

lemma nmod2: $n \bmod (2 :: \text{int}) = 0 \mid n \bmod 2 = 1$

$\langle \text{proof} \rangle$

lemma emep1:

$\text{even } n ==> \text{even } d ==> 0 \leq d ==> (n + 1) \bmod (d :: \text{int}) = (n \bmod d) + 1$
 $\langle \text{proof} \rangle$

lemmas $\text{eme1p} = \text{emep1 } [\text{simplified add-commute}]$

lemma le-diff-eq': $(a \leq c - b) = (b + a \leq (c :: \text{int}))$

$\langle \text{proof} \rangle$

lemma less-diff-eq': $(a < c - b) = (b + a < (c :: \text{int}))$

$\langle \text{proof} \rangle$

lemma diff-le-eq': $(a - b \leq c) = (a \leq b + (c :: \text{int}))$

$\langle \text{proof} \rangle$

lemma diff-less-eq': $(a - b < c) = (a < b + (c :: \text{int}))$

$\langle \text{proof} \rangle$

lemmas $\text{m1mod2k} = \text{zless2p } [\text{THEN zmod-minus1}]$

lemmas $\text{m1mod22k} = \text{mult-pos-pos } [\text{OF zless2 zless2p}, \text{ THEN zmod-minus1}]$

lemmas $\text{p1mod22k}' = \text{zless2p } [\text{THEN order-less-imp-le}, \text{ THEN pos-zmod-mult-2}]$

lemmas $\text{z1pmod2}' = \text{zero-le-one } [\text{THEN pos-zmod-mult-2}, \text{ simplified}]$

lemmas $\text{z1pdiv2}' = \text{zero-le-one } [\text{THEN pos-zdiv-mult-2}, \text{ simplified}]$

lemma *p1mod22k*:

$(2 * b + 1) \text{ mod } (2 * 2^n) = 2 * (b \text{ mod } 2^n) + (1::int)$
 $\langle \text{proof} \rangle$

lemma *z1pmod2*:

$(2 * b + 1) \text{ mod } 2 = (1::int)$
 $\langle \text{proof} \rangle$

lemma *z1pdiv2*:

$(2 * b + 1) \text{ div } 2 = (b::int)$
 $\langle \text{proof} \rangle$

lemmas *zdiv-le-dividend* = *xtr3* [*OF zdiv-1* [*symmetric*] *zdiv-mono2*,
simplified int-one-le-iff-zero-less, simplified, standard]

lemma *BIT-eq*: $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$
 $\langle \text{proof} \rangle$

lemmas *BIT-eqE* [*elim!*] = *BIT-eq* [*THEN conjE, standard*]

lemma *BIT-eq-iff* [*simp*]:

$(u \text{ BIT } b = v \text{ BIT } c) = (u = v \ \wedge \ b = c)$
 $\langle \text{proof} \rangle$

lemmas *BIT-eqI* [*intro!*] = *conjI* [*THEN BIT-eq-iff* [*THEN iffD2*]]

lemma *less-Bits*:

$(v \text{ BIT } b < w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *le-Bits*:

$(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ (b \sim = \text{bit.B1} \mid c \sim = \text{bit.B0}))$
 $\langle \text{proof} \rangle$

lemma *neB1E* [*elim!*]:

assumes *ne*: $y \neq \text{bit.B1}$
assumes *y*: $y = \text{bit.B0} \implies P$
shows *P*
 $\langle \text{proof} \rangle$

lemma *bin-ex-rl*: $EX \ w \ b. \ w \text{ BIT } b = \text{bin}$
 $\langle \text{proof} \rangle$

lemma *bin-exhaust*:

assumes *Q*: $\bigwedge x \ b. \ \text{bin} = x \text{ BIT } b \implies Q$
shows *Q*

$\langle proof \rangle$

lemma *bin-rl-char*: $(bin\text{-}rl\ w = (r, l)) = (r\ BIT\ l = w)$
 $\langle proof \rangle$

lemmas *bin-rl-simps* [*THEN bin-rl-char* [*THEN iffD2*], *standard*, *simp*] =
Pls-0-eq Min-1-eq refl

lemma *bin-abs-lem*:
 $bin = (w\ BIT\ b) ==> \sim\ bin = Numeral.Min\ --> \sim\ bin = Numeral.Pls\ -->$
 $nat\ (abs\ w) < nat\ (abs\ bin)$
 $\langle proof \rangle$

lemma *bin-induct*:
assumes *PPls*: $P\ Numeral.Pls$
and *PMin*: $P\ Numeral.Min$
and *PBit*: $!!bin\ bit. P\ bin ==> P\ (bin\ BIT\ bit)$
shows $P\ bin$
 $\langle proof \rangle$

lemma *no-no* [*simp*]: *number-of* (*number-of* *i*) = *i*
 $\langle proof \rangle$

lemma *Bit-B0*:
 $k\ BIT\ bit.B0 = k + k$
 $\langle proof \rangle$

lemma *Bit-B1*:
 $k\ BIT\ bit.B1 = k + k + 1$
 $\langle proof \rangle$

lemma *Bit-B0-2t*: $k\ BIT\ bit.B0 = 2 * k$
 $\langle proof \rangle$

lemma *Bit-B1-2t*: $k\ BIT\ bit.B1 = 2 * k + 1$
 $\langle proof \rangle$

lemma *B-mod-2'*:
 $X = 2 ==> (w\ BIT\ bit.B1)\ mod\ X = 1 \ \&\ (w\ BIT\ bit.B0)\ mod\ X = 0$
 $\langle proof \rangle$

lemmas *B1-mod-2* [*simp*] = *B-mod-2'* [*OF refl*, *THEN conjunct1*, *standard*]

lemmas *B0-mod-2* [*simp*] = *B-mod-2'* [*OF refl*, *THEN conjunct2*, *standard*]

lemma *axbbyy*:
 $a + m + m = b + n + n ==> (a = 0 \mid a = 1) ==> (b = 0 \mid b = 1) ==>$
 $a = b \ \&\ m = (n :: int)$
 $\langle proof \rangle$

lemma *axxmod2*:

$(1 + x + x) \bmod 2 = (1 :: \text{int}) \ \& \ (0 + x + x) \bmod 2 = (0 :: \text{int})$
 $\langle \text{proof} \rangle$

lemma *axxdiv2*:

$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int})$
 $\langle \text{proof} \rangle$

lemmas *iszero-minus* = *trans* [THEN *trans*,

OF iszero-def neg-equal-0-iff-equal iszero-def [symmetric], *standard*]

lemmas *zadd-diff-inverse* = *trans* [*OF diff-add-cancel* [symmetric] *add-commute*,
standard]

lemmas *add-diff-cancel2* = *add-commute* [THEN *diff-eq-eq* [THEN *iffD2*], *standard*]

lemma *zmod-uminus*: $-(a :: \text{int}) \bmod b \bmod b = -a \bmod b$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-distrib*: $((a :: \text{int}) - b) \bmod c = (a \bmod c - b \bmod c) \bmod c$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-right-eq*: $((a :: \text{int}) - b) \bmod c = (a - b \bmod c) \bmod c$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-left-eq*: $((a :: \text{int}) - b) \bmod c = (a \bmod c - b) \bmod c$
 $\langle \text{proof} \rangle$

lemma *zmod-zsub-self* [simp]:

$((b :: \text{int}) - a) \bmod a = b \bmod a$
 $\langle \text{proof} \rangle$

lemma *zmod-zmult1-eq-rev*:

$b * a \bmod c = b \bmod c * a \bmod (c :: \text{int})$
 $\langle \text{proof} \rangle$

lemmas *rdmods* [symmetric] = *zmod-uminus* [symmetric]

zmod-zsub-left-eq zmod-zsub-right-eq zmod-zadd-left-eq
zmod-zadd-right-eq zmod-zmult1-eq zmod-zmult1-eq-rev

lemma *mod-plus-right*:

$((a + x) \bmod m = (b + x) \bmod m) = (a \bmod m = b \bmod (m :: \text{nat}))$
 $\langle \text{proof} \rangle$

lemma *nat-minus-mod*: $(n - n \bmod m) \bmod m = (0 :: \text{nat})$
 $\langle \text{proof} \rangle$

lemmas *nat-minus-mod-plus-right* = *trans* [*OF nat-minus-mod mod-0* [symmetric],

THEN mod-plus-right [*THEN iffD2*], *standard*, *simplified*]

lemmas *push-mods'* = *zmod-zadd1-eq* [*standard*]
zmod-zmult-distrib [*standard*] *zmod-zsub-distrib* [*standard*]
zmod-uminus [*symmetric*, *standard*]

lemmas *push-mods* = *push-mods'* [*THEN eq-reflection*, *standard*]
lemmas *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN eq-reflection*, *standard*]
lemmas *mod-simps* =
zmod-zmult-self1 [*THEN eq-reflection*] *zmod-zmult-self2* [*THEN eq-reflection*]
mod-mod-trivial [*THEN eq-reflection*]

lemma *nat-mod-eq*:
 $!!b. b < n \implies a \bmod n = b \bmod n \implies a \bmod n = (b :: nat)$
<proof>

lemmas *nat-mod-eq'* = *refl* [*THEN* [2] *nat-mod-eq*]

lemma *nat-mod-lem*:
 $(0 :: nat) < n \implies b < n \implies (b \bmod n = b)$
<proof>

lemma *mod-nat-add*:
 $(x :: nat) < z \implies y < z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
<proof>

lemma *mod-nat-sub*:
 $(x :: nat) < z \implies (x - y) \bmod z = x - y$
<proof>

lemma *int-mod-lem*:
 $(0 :: int) < n \implies (0 \leq b \ \& \ b < n) \implies (b \bmod n = b)$
<proof>

lemma *int-mod-eq*:
 $(0 :: int) \leq b \implies b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$
<proof>

lemmas *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

lemma *int-mod-le*: $0 \leq a \implies 0 < (n :: int) \implies a \bmod n \leq a$
<proof>

lemma *int-mod-le'*: $0 \leq b - n \implies 0 < (n :: int) \implies b \bmod n \leq b - n$
<proof>

lemma *int-mod-ge*: $a < n \implies 0 < (n :: int) \implies a \leq a \bmod n$

$\langle \text{proof} \rangle$

lemma *int-mod-ge'*: $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$
 $\langle \text{proof} \rangle$

lemma *mod-add-if-z*:
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
 $\langle \text{proof} \rangle$

lemma *mod-sub-if-z*:
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
 $\langle \text{proof} \rangle$

lemmas *zmde* = *zmod-zdiv-equality* [THEN *diff-eq-eq* [THEN *iffD2*], *symmetric*]
lemmas *mcl* = *mult-cancel-left* [THEN *iffD1*, THEN *make-pos-rule*]

lemma *zdiv-mult-self*: $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$
 $\langle \text{proof} \rangle$

lemma *eqne*: $\text{equiv } A \ r \implies X : A // r \implies X \sim = \{\}$
 $\langle \text{proof} \rangle$

lemmas *Rep-Integ-ne* = *Integ.Rep-Integ*
[THEN *equiv-intrel* [THEN *eqne*, *simplified Integ-def* [symmetric]], *standard*]

lemmas *riq* = *Integ.Rep-Integ* [*simplified Integ-def*]
lemmas *intrel-refl* = *refl* [THEN *equiv-intrel-iff* [THEN *iffD1*], *standard*]
lemmas *Rep-Integ-equiv* = *quotient-eq-iff*
[OF *equiv-intrel riq riq*, *simplified Integ.Rep-Integ-inject*, *standard*]
lemmas *Rep-Integ-same* =
Rep-Integ-equiv [THEN *intrel-refl* [THEN *rev-iffD2*], *standard*]

lemma *RI-int*: $(a, 0) : \text{Rep-Integ } (int \ a)$
 $\langle \text{proof} \rangle$

lemmas *RI-intrel* [*simp*] = *UNIV-I* [THEN *quotientI*,
THEN *Integ.Abs-Integ-inverse* [*simplified Integ-def*], *standard*]

lemma *RI-minus*: $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (-x)$
 $\langle \text{proof} \rangle$

lemma *RI-add*:
 $(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$
 $(a + c, b + d) : \text{Rep-Integ } (x + y)$
 $\langle \text{proof} \rangle$

lemma *mem-same*: $a : S \implies a = b \implies b : S$
 ⟨proof⟩

lemma *RI-eq-diff'*: $(a, b) : \text{Rep-Integ } (\text{int } a - \text{int } b)$
 ⟨proof⟩

lemma *RI-eq-diff*: $((a, b) : \text{Rep-Integ } x) = (\text{int } a - \text{int } b = x)$
 ⟨proof⟩

lemma *mod-power-lem*:
 $a > 1 \implies a \wedge^n \text{ mod } a \wedge^m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int}) \wedge^n)$
 ⟨proof⟩

lemma *min-pm* [simp]: $\min a b + (a - b) = (a :: \text{nat})$
 ⟨proof⟩

lemmas *min-pm1* [simp] = trans [OF add-commute min-pm]

lemma *rev-min-pm* [simp]: $\min b a + (a - b) = (a :: \text{nat})$
 ⟨proof⟩

lemmas *rev-min-pm1* [simp] = trans [OF add-commute rev-min-pm]

lemma *pl-pl-rels*:
 $a + b = c + d \implies$
 $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$
 ⟨proof⟩

lemmas *pl-pl-rels'* = add-commute [THEN [2] trans, THEN pl-pl-rels]

lemma *minus-eq*: $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$
 ⟨proof⟩

lemma *pl-pl-mm*: $(a :: \text{nat}) + b = c + d \implies a - c = d - b$
 ⟨proof⟩

lemmas *pl-pl-mm'* = add-commute [THEN [2] trans, THEN pl-pl-mm]

lemma *min-minus* [simp]: $\min m (m - k) = (m - k :: \text{nat})$
 ⟨proof⟩

lemmas *min-minus'* [simp] = trans [OF min-max.inf-commute min-minus]

lemma *nat-no-eq-iff*:
 $(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$
 $(\text{number-of } b = (\text{number-of } c :: \text{nat})) = (b = c)$
 ⟨proof⟩

lemmas *dme* = *box-equals* [*OF* *div-mod-equality* *add-0-right* *add-0-right*]
lemmas *dtle* = *xtr3* [*OF* *dme* [*symmetric*] *le-add1*]
lemmas *th2* = *order-trans* [*OF* *order-refl* [*THEN* [2] *mult-le-mono*] *dtle*]

lemma *td-gal*:

$0 < c \implies (a \geq b * c) = (a \text{ div } c \geq (b :: \text{nat}))$
 $\langle \text{proof} \rangle$

lemmas *td-gal-lt* = *td-gal* [*simplified le-def*, *simplified*]

lemma *div-mult-le*: $(a :: \text{nat}) \text{ div } b * b \leq a$
 $\langle \text{proof} \rangle$

lemmas *sdl* = *split-div-lemma* [*THEN iffD1*, *symmetric*]

lemma *given-quot*: $f > (0 :: \text{nat}) \implies (f * l + (f - 1)) \text{ div } f = l$
 $\langle \text{proof} \rangle$

lemma *given-quot-alt*: $f > (0 :: \text{nat}) \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$
 $\langle \text{proof} \rangle$

lemma *diff-mod-le*: $(a :: \text{nat}) < d \implies b \text{ dvd } d \implies a - a \text{ mod } b \leq d - b$
 $\langle \text{proof} \rangle$

lemma *less-le-mult'*:

$w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * (c :: \text{int})$
 $\langle \text{proof} \rangle$

lemmas *less-le-mult* = *less-le-mult'* [*simplified left-distrib*, *simplified*]

lemmas *less-le-mult-minus* = *iffD2* [*OF* *le-diff-eq* *less-le-mult*,
simplified left-diff-distrib, *standard*]

lemma *lrlem'*:

assumes *d*: $(i :: \text{nat}) \leq j \vee m < j'$
assumes *R1*: $i * k \leq j * k \implies R$
assumes *R2*: $\text{Suc } m * k' \leq j' * k' \implies R$
shows *R* $\langle \text{proof} \rangle$

lemma *lrlem*: $(0 :: \text{nat}) < sc \implies$

$(sc - n + (n + lb * n) \leq m * n) = (sc + lb * n \leq m * n)$
 $\langle \text{proof} \rangle$

lemma *gen-minus*: $0 < n \implies f \ n = f \ (\text{Suc } (n - 1))$
 $\langle \text{proof} \rangle$

lemma *mpl-lem*: $j \leq (i :: \text{nat}) \implies k < j \implies i - j + k < i$
 $\langle \text{proof} \rangle$

lemma *nonneg-mod-div*:

$0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b :: \text{int}) \ \& \ 0 \leq a \operatorname{div} b$
 $\langle \text{proof} \rangle$

end

4 BinGeneral: Basic Definitions for Binary Integers

theory *BinGeneral* **imports** *Num-Lemmas*

begin

4.1 Recursion combinator for binary integers

lemma *brlem*: $(\text{bin} = \text{Numeral.Min}) = (- \text{bin} + \text{Numeral.pred } 0 = 0)$
 $\langle \text{proof} \rangle$

function

$\text{bin-rec}' :: \text{int} * 'a * 'a * (\text{int} \Rightarrow \text{bit} \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'a$

where

$\text{bin-rec}' (\text{bin}, f1, f2, f3) = (\text{if } \text{bin} = \text{Numeral.Plus} \text{ then } f1$

$\text{else if } \text{bin} = \text{Numeral.Min} \text{ then } f2$

$\text{else case bin-rl bin of } (w, b) \Rightarrow f3 \ w \ b \ (\text{bin-rec}' (w, f1, f2, f3)))$

$\langle \text{proof} \rangle$

termination

$\langle \text{proof} \rangle$

constdefs

$\text{bin-rec} :: 'a \Rightarrow 'a \Rightarrow (\text{int} \Rightarrow \text{bit} \Rightarrow 'a \Rightarrow 'a) \Rightarrow \text{int} \Rightarrow 'a$

$\text{bin-rec } f1 \ f2 \ f3 \ \text{bin} == \text{bin-rec}' (\text{bin}, f1, f2, f3)$

lemma *bin-rec-PM*:

$f = \text{bin-rec } f1 \ f2 \ f3 \implies f \ \text{Numeral.Plus} = f1 \ \& \ f \ \text{Numeral.Min} = f2$

$\langle \text{proof} \rangle$

lemmas *bin-rec-Plus* = *refl [THEN bin-rec-PM, THEN conjunct1, standard]*

lemmas *bin-rec-Min* = *refl [THEN bin-rec-PM, THEN conjunct2, standard]*

lemma *bin-rec-Bit*:

$f = \text{bin-rec } f1 \ f2 \ f3 \implies f3 \ \text{Numeral.Plus} \ \text{bit.B0 } f1 = f1 \implies$

$f3 \ \text{Numeral.Min} \ \text{bit.B1 } f2 = f2 \implies f \ (w \ \text{BIT } b) = f3 \ w \ b \ (f \ w)$

$\langle \text{proof} \rangle$

lemmas *bin-rec-simps* = *refl [THEN bin-rec-Bit] bin-rec-Plus bin-rec-Min*

4.2 Destructors for binary integers

consts

— corresponding operations analysing bins
 $bin_last :: int \Rightarrow bit$
 $bin_rest :: int \Rightarrow int$
 $bin_sign :: int \Rightarrow int$
 $bin_nth :: int \Rightarrow nat \Rightarrow bool$

primrec

$Z : bin_nth\ w\ 0 = (bin_last\ w = bit.B1)$
 $Suc : bin_nth\ w\ (Suc\ n) = bin_nth\ (bin_rest\ w)\ n$

defs

$bin_rest_def : bin_rest\ w == fst\ (bin_rl\ w)$
 $bin_last_def : bin_last\ w == snd\ (bin_rl\ w)$
 $bin_sign_def : bin_sign == bin_rec\ Numeral.Pls\ Numeral.Min\ (\%w\ b\ s.\ s)$

lemma bin_rl : $bin_rl\ w = (bin_rest\ w, bin_last\ w)$
 $\langle proof \rangle$

lemmas bin_rl_simp $[simp] = iffD1\ [OF\ bin_rl_char\ bin_rl]$

lemma bin_rest_sims $[simp]$:

$bin_rest\ Numeral.Pls = Numeral.Pls$
 $bin_rest\ Numeral.Min = Numeral.Min$
 $bin_rest\ (w\ BIT\ b) = w$
 $\langle proof \rangle$

lemma bin_last_sims $[simp]$:

$bin_last\ Numeral.Pls = bit.B0$
 $bin_last\ Numeral.Min = bit.B1$
 $bin_last\ (w\ BIT\ b) = b$
 $\langle proof \rangle$

lemma bin_sign_sims $[simp]$:

$bin_sign\ Numeral.Pls = Numeral.Pls$
 $bin_sign\ Numeral.Min = Numeral.Min$
 $bin_sign\ (w\ BIT\ b) = bin_sign\ w$
 $\langle proof \rangle$

lemma $bin_r_l_extras$ $[simp]$:

$bin_last\ 0 = bit.B0$
 $bin_last\ (-\ 1) = bit.B1$
 $bin_last\ -1 = bit.B1$
 $bin_last\ 1 = bit.B1$
 $bin_rest\ 1 = 0$
 $bin_rest\ 0 = 0$
 $bin_rest\ (-\ 1) = -\ 1$
 $bin_rest\ -1 = -1$

$\langle \text{proof} \rangle$

lemma *bin-last-mod*:

bin-last $w = (\text{if } w \bmod 2 = 0 \text{ then } \text{bit.B0} \text{ else } \text{bit.B1})$

$\langle \text{proof} \rangle$

lemma *bin-rest-div*:

bin-rest $w = w \text{ div } 2$

$\langle \text{proof} \rangle$

lemma *Bit-div2* [simp]: $(w \text{ BIT } b) \text{ div } 2 = w$

$\langle \text{proof} \rangle$

lemma *bin-nth-lem* [rule-format]:

ALL y . *bin-nth* $x = \text{bin-nth } y \dashrightarrow x = y$

$\langle \text{proof} \rangle$

lemma *bin-nth-eq-iff*: $(\text{bin-nth } x = \text{bin-nth } y) = (x = y)$

$\langle \text{proof} \rangle$

lemmas *bin-eqI* = *ext* [THEN *bin-nth-eq-iff* [THEN *iffD1*], *standard*]

lemma *bin-nth-Pls* [simp]: $\sim \text{bin-nth Numeral.Pls } n$

$\langle \text{proof} \rangle$

lemma *bin-nth-Min* [simp]: *bin-nth Numeral.Min* n

$\langle \text{proof} \rangle$

lemma *bin-nth-0-BIT*: *bin-nth* $(w \text{ BIT } b) 0 = (b = \text{bit.B1})$

$\langle \text{proof} \rangle$

lemma *bin-nth-Suc-BIT*: *bin-nth* $(w \text{ BIT } b) (\text{Suc } n) = \text{bin-nth } w n$

$\langle \text{proof} \rangle$

lemma *bin-nth-minus* [simp]: $0 < n \implies \text{bin-nth } (w \text{ BIT } b) n = \text{bin-nth } w (n - 1)$

$\langle \text{proof} \rangle$

lemmas *bin-nth-0* = *bin-nth.simps*(1)

lemmas *bin-nth-Suc* = *bin-nth.simps*(2)

lemmas *bin-nth-simps* =

bin-nth-0 bin-nth-Suc bin-nth-Pls bin-nth-Min bin-nth-minus

lemma *bin-sign-rest* [simp]:

bin-sign $(\text{bin-rest } w) = (\text{bin-sign } w)$

$\langle \text{proof} \rangle$

4.3 Truncating binary integers

consts

bintrunc :: nat => int => int

primrec

Z : *bintrunc* 0 *bin* = *Numeral.Pls*

Suc : *bintrunc* (*Suc* *n*) *bin* = *bintrunc* *n* (*bin-rest* *bin*) *BIT* (*bin-last* *bin*)

consts

sbintrunc :: nat => int => int

primrec

Z : *sbintrunc* 0 *bin* =

(case *bin-last* *bin* of *bit.B1* => *Numeral.Min* | *bit.B0* => *Numeral.Pls*)

Suc : *sbintrunc* (*Suc* *n*) *bin* = *sbintrunc* *n* (*bin-rest* *bin*) *BIT* (*bin-last* *bin*)

lemma *sign-bintr*:

!!*w*. *bin-sign* (*bintrunc* *n* *w*) = *Numeral.Pls*

⟨*proof*⟩

lemma *bintrunc-mod2p*:

!!*w*. *bintrunc* *n* *w* = (*w mod* 2 ^ *n* :: int)

⟨*proof*⟩

lemma *sbintrunc-mod2p*:

!!*w*. *sbintrunc* *n* *w* = ((*w* + 2 ^ *n*) mod 2 ^ (*Suc* *n*) - 2 ^ *n* :: int)

⟨*proof*⟩

4.4 Simplifications for (s)bintrunc

lemma *bit-bool*:

(*b* = (*b'* = *bit.B1*)) = (*b'* = (if *b* then *bit.B1* else *bit.B0*))

⟨*proof*⟩

lemmas *bit-bool1* [*simp*] = *refl* [*THEN* *bit-bool* [*THEN* *iffD1*], *symmetric*]

lemma *bin-sign-lem*:

!!*bin*. (*bin-sign* (*sbintrunc* *n* *bin*) = *Numeral.Min*) = *bin-nth* *bin* *n*

⟨*proof*⟩

lemma *nth-bintr*:

!!*w m*. *bin-nth* (*bintrunc* *m* *w*) *n* = (*n* < *m* & *bin-nth* *w* *n*)

⟨*proof*⟩

lemma *nth-sbintr*:

!!*w m*. *bin-nth* (*sbintrunc* *m* *w*) *n* =

(if *n* < *m* then *bin-nth* *w* *n* else *bin-nth* *w* *m*)

⟨*proof*⟩

lemma *bin-nth-Bit*:

bin-nth (*w BIT b*) *n* = (*n* = 0 & *b* = *bit.B1* | (EX *m*. *n* = *Suc* *m* & *bin-nth* *w*

$m))$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-l*:
 $n \leq m \implies (\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } n w)$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-sbintrunc-l*:
 $n \leq m \implies (\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } n w)$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-ge*:
 $n \leq m \implies (\text{bintrunc } n (\text{bintrunc } m w) = \text{bintrunc } n w)$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-min* [simp]:
 $\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } (\min m n) w$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-sbintrunc-min* [simp]:
 $\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } (\min m n) w$
 $\langle \text{proof} \rangle$

lemmas *bintrunc-Pls* =
 bintrunc.Suc [where $\text{bin} = \text{Numeral.Pls}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *bintrunc-Min* [simp] =
 bintrunc.Suc [where $\text{bin} = \text{Numeral.Min}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *bintrunc-BIT* [simp] =
 bintrunc.Suc [where $\text{bin} = w \text{ BIT } b$, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

lemmas *sbintrunc-Suc-Pls* =
 sbintrunc.Suc [where $\text{bin} = \text{Numeral.Pls}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *sbintrunc-Suc-Min* =
 sbintrunc.Suc [where $\text{bin} = \text{Numeral.Min}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *sbintrunc-Suc-BIT* [simp] =
 sbintrunc.Suc [where $\text{bin} = w \text{ BIT } b$, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls* *sbintrunc-Suc-Min* *sbintrunc-Suc-BIT*

lemmas *sbintrunc-Pls* =
sbintrunc.Z [**where** *bin*=*Numeral.Pls*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-Min* =
sbintrunc.Z [**where** *bin*=*Numeral.Min*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-BIT-B0* [*simp*] =
sbintrunc.Z [**where** *bin*=*w BIT bit.B0*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-BIT-B1* [*simp*] =
sbintrunc.Z [**where** *bin*=*w BIT bit.B1*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls *sbintrunc-Min* *sbintrunc-0-BIT-B0* *sbintrunc-0-BIT-B1*

lemmas *bintrunc-simps* = *bintrunc.Z* *bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps* *sbintrunc-Sucs*

lemma *bintrunc-minus*:
 $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$
<proof>

lemma *sbintrunc-minus*:
 $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$
<proof>

lemmas *bintrunc-minus-simps* =
bintrunc-Sucs [*THEN* [2] *bintrunc-minus* [*symmetric, THEN trans*], *standard*]

lemmas *sbintrunc-minus-simps* =
sbintrunc-Sucs [*THEN* [2] *sbintrunc-minus* [*symmetric, THEN trans*], *standard*]

lemma *bintrunc-n-Pls* [*simp*]:
 $\text{bintrunc } n \ \text{Numeral.Pls} = \text{Numeral.Pls}$
<proof>

lemma *sbintrunc-n-PM* [*simp*]:
 $\text{sbintrunc } n \ \text{Numeral.Pls} = \text{Numeral.Pls}$
 $\text{sbintrunc } n \ \text{Numeral.Min} = \text{Numeral.Min}$
<proof>

lemmas *thobini1* = *arg-cong* [**where** *f* = $\%w. \ w \ \text{BIT } b$, *standard*]

lemmas *bintrunc-BIT-I* = *trans* [*OF* *bintrunc-BIT* *thobini1*]

lemmas *bintrunc-Min-I* = *trans* [*OF bintrunc-Min thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps* [*THEN thobini1* [*THEN* [2] *trans*], *standard*]

lemmas *bintrunc-Pls-minus-I* = *bmsts*(1)

lemmas *bintrunc-Min-minus-I* = *bmsts*(2)

lemmas *bintrunc-BIT-minus-I* = *bmsts*(3)

lemma *bintrunc-0-Min*: *bintrunc 0 Numeral.Min* = *Numeral.Pls*
 ⟨*proof*⟩

lemma *bintrunc-0-BIT*: *bintrunc 0 (w BIT b)* = *Numeral.Pls*
 ⟨*proof*⟩

lemma *bintrunc-Suc-lem*:

bintrunc (Suc n) x = y ==> m = Suc n ==> bintrunc m x = y
 ⟨*proof*⟩

lemmas *bintrunc-Suc-Ialts* =

bintrunc-Min-I bintrunc-BIT-I [*THEN bintrunc-Suc-lem, standard*]

lemmas *sbintrunc-BIT-I* = *trans* [*OF sbintrunc-Suc-BIT thobini1*]

lemmas *sbintrunc-Suc-Is* =

sbintrunc-Sucs [*THEN thobini1* [*THEN* [2] *trans*], *standard*]

lemmas *sbintrunc-Suc-minus-Is* =

sbintrunc-minus-simps [*THEN thobini1* [*THEN* [2] *trans*], *standard*]

lemma *sbintrunc-Suc-lem*:

sbintrunc (Suc n) x = y ==> m = Suc n ==> sbintrunc m x = y
 ⟨*proof*⟩

lemmas *sbintrunc-Suc-Ialts* =

sbintrunc-Suc-Is [*THEN sbintrunc-Suc-lem, standard*]

lemma *sbintrunc-bintrunc-lt*:

m > n ==> sbintrunc n (bintrunc m w) = sbintrunc n w
 ⟨*proof*⟩

lemma *bintrunc-sbintrunc-le*:

m <= Suc n ==> bintrunc m (sbintrunc n w) = bintrunc m w
 ⟨*proof*⟩

lemmas *bintrunc-sbintrunc [simp]* = *order-refl* [*THEN bintrunc-sbintrunc-le*]

lemmas *sbintrunc-bintrunc [simp]* = *lessI* [*THEN sbintrunc-bintrunc-lt*]

lemmas *bintrunc-bintrunc [simp]* = *order-refl* [*THEN bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc [simp]* = *order-refl* [*THEN sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc' [simp]*:

$0 < n \implies \text{bintrunc } n (\text{sbintrunc } (n - 1) w) = \text{bintrunc } n w$
 ⟨proof⟩

lemma *sbintrunc-bintrunc'* [simp]:

$0 < n \implies \text{sbintrunc } (n - 1) (\text{bintrunc } n w) = \text{sbintrunc } (n - 1) w$
 ⟨proof⟩

lemma *bin-sbin-eq-iff*:

$\text{bintrunc } (\text{Suc } n) x = \text{bintrunc } (\text{Suc } n) y <->$
 $\text{sbintrunc } n x = \text{sbintrunc } n y$
 ⟨proof⟩

lemma *bin-sbin-eq-iff'*:

$0 < n \implies \text{bintrunc } n x = \text{bintrunc } n y <->$
 $\text{sbintrunc } (n - 1) x = \text{sbintrunc } (n - 1) y$
 ⟨proof⟩

lemmas *bintrunc-sbintruncS0* [simp] = *bintrunc-sbintrunc'* [unfolded One-nat-def]

lemmas *sbintrunc-bintruncS0* [simp] = *sbintrunc-bintrunc'* [unfolded One-nat-def]

lemmas *bintrunc-bintrunc-l'* = *le-add1* [THEN *bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc-l'* = *le-add1* [THEN *sbintrunc-sbintrunc-l*]

lemmas *nat-non0-gr* =

trans [OF *iszero-def* [THEN *Not-eq-iff* [THEN *iffD2*]] *refl, standard*]

lemmas *bintrunc-pred-simps* [simp] =

bintrunc-minus-simps [of *number-of bin, simplified nobm1, standard*]

lemmas *sbintrunc-pred-simps* [simp] =

sbintrunc-minus-simps [of *number-of bin, simplified nobm1, standard*]

lemma *no-bintr-alt*:

$\text{number-of } (\text{bintrunc } n w) = w \bmod 2^n$
 ⟨proof⟩

lemma *no-bintr-alt1*: $\text{bintrunc } n = (\%w. w \bmod 2^n :: \text{int})$

⟨proof⟩

lemma *range-bintrunc*: $\text{range } (\text{bintrunc } n) = \{i. 0 \leq i \ \& \ i < 2^n\}$

⟨proof⟩

lemma *no-bintr*:

$\text{number-of } (\text{bintrunc } n w) = (\text{number-of } w \bmod 2^n :: \text{int})$
 ⟨proof⟩

lemma *no-sbintr-alt2*:

$sbintrunc\ n = (\%w. (w + 2^{\wedge} n) \bmod 2^{\wedge} Suc\ n - 2^{\wedge} n :: int)$
 $\langle proof \rangle$

lemma *no-sbintr*:

$number-of\ (sbintrunc\ n\ w) =$
 $((number-of\ w + 2^{\wedge} n) \bmod 2^{\wedge} Suc\ n - 2^{\wedge} n :: int)$
 $\langle proof \rangle$

lemma *range-sbintrunc*:

$range\ (sbintrunc\ n) = \{i. - (2^{\wedge} n) \leq i \ \& \ i < 2^{\wedge} n\}$
 $\langle proof \rangle$

lemma *sb-inc-lem*:

$(a::int) + 2^{\wedge} k < 0 \implies a + 2^{\wedge} k + 2^{\wedge} (Suc\ k) \leq (a + 2^{\wedge} k) \bmod 2^{\wedge} (Suc\ k)$
 $\langle proof \rangle$

lemma *sb-inc-lem'*:

$(a::int) < - (2^{\wedge} k) \implies a + 2^{\wedge} k + 2^{\wedge} (Suc\ k) \leq (a + 2^{\wedge} k) \bmod 2^{\wedge} (Suc\ k)$
 $\langle proof \rangle$

lemma *sbintrunc-inc*:

$x < - (2^{\wedge} n) \implies x + 2^{\wedge} (Suc\ n) \leq sbintrunc\ n\ x$
 $\langle proof \rangle$

lemma *sb-dec-lem*:

$(0::int) \leq - (2^{\wedge} k) + a \implies (a + 2^{\wedge} k) \bmod (2 * 2^{\wedge} k) \leq - (2^{\wedge} k) + a$
 $\langle proof \rangle$

lemma *sb-dec-lem'*:

$(2::int)^{\wedge} k \leq a \implies (a + 2^{\wedge} k) \bmod (2 * 2^{\wedge} k) \leq - (2^{\wedge} k) + a$
 $\langle proof \rangle$

lemma *sbintrunc-dec*:

$x >= (2^{\wedge} n) \implies x - 2^{\wedge} (Suc\ n) >= sbintrunc\ n\ x$
 $\langle proof \rangle$

lemmas *zmod-uminus'* = *zmod-uminus* [where *b=c*, standard]

lemmas *zpower-zmod'* = *zpower-zmod* [where *m=c* and *y=k*, standard]

lemmas *brdmod1s'* [symmetric] =

zmod-zadd-left-eq *zmod-zadd-right-eq*
zmod-zsub-left-eq *zmod-zsub-right-eq*
zmod-zmult1-eq *zmod-zmult1-eq-rev*

lemmas *brdmods'* [symmetric] =

zpower-zmod' [symmetric]
trans [OF *zmod-zadd-left-eq* *zmod-zadd-right-eq*]
trans [OF *zmod-zsub-left-eq* *zmod-zsub-right-eq*]
trans [OF *zmod-zmult1-eq* *zmod-zmult1-eq-rev*]

zmod-uminus' [symmetric]
zmod-zadd-left-eq [where $b = 1$]
zmod-zsub-left-eq [where $b = 1$]

lemmas *bintr-arith1s* =
brdmod1s' [where $c=2^n$, folded pred-def succ-def bintrunc-mod2p, standard]

lemmas *bintr-ariths* =
brdmods' [where $c=2^n$, folded pred-def succ-def bintrunc-mod2p, standard]

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [OF *zless2p*, standard]

lemma *bintr-ge0*: $(0 :: \text{int}) \leq \text{number-of } (\text{bintrunc } n \ w)$
 ⟨proof⟩

lemma *bintr-lt2p*: $\text{number-of } (\text{bintrunc } n \ w) < (2^n :: \text{int})$
 ⟨proof⟩

lemma *bintr-Min*:
 $\text{number-of } (\text{bintrunc } n \ \text{Numeral.Min}) = (2^n :: \text{int}) - 1$
 ⟨proof⟩

lemma *sbintr-ge*: $(-(2^n :: \text{int})) \leq \text{number-of } (\text{sbintrunc } n \ w)$
 ⟨proof⟩

lemma *sbintr-lt*: $\text{number-of } (\text{sbintrunc } n \ w) < (2^n :: \text{int})$
 ⟨proof⟩

lemma *bintrunc-Suc*:
 $\text{bintrunc } (\text{Suc } n) \ \text{bin} = \text{bintrunc } n \ (\text{bin-rest } \text{bin}) \ \text{BIT } \text{bin-last } \text{bin}$
 ⟨proof⟩

lemma *sign-Pls-ge-0*:
 $(\text{bin-sign } \text{bin} = \text{Numeral.Pls}) = (\text{number-of } \text{bin} \geq (0 :: \text{int}))$
 ⟨proof⟩

lemma *sign-Min-lt-0*:
 $(\text{bin-sign } \text{bin} = \text{Numeral.Min}) = (\text{number-of } \text{bin} < (0 :: \text{int}))$
 ⟨proof⟩

lemmas *sign-Min-neg* = *trans* [OF *sign-Min-lt-0 neg-def* [symmetric]]

lemma *bin-rest-trunc*:
 $!!\text{bin}. (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bintrunc } (n - 1) \ (\text{bin-rest } \text{bin})$
 ⟨proof⟩

lemma *bin-rest-power-trunc* [rule-format] :
 $(\text{bin-rest } ^k) (\text{bintrunc } n \ \text{bin}) =$
 $\text{bintrunc } (n - k) \ ((\text{bin-rest } ^k) \ \text{bin})$
 ⟨proof⟩

lemma *bin-rest-trunc-i*:

$\text{bintrunc } n \ (\text{bin-rest } \text{bin}) = \text{bin-rest } (\text{bintrunc } (\text{Suc } n) \ \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bin-rest-strunc*:

$!!\text{bin. } \text{bin-rest } (\text{sbintrunc } (\text{Suc } n) \ \text{bin}) = \text{sbintrunc } n \ (\text{bin-rest } \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bintrunc-rest [simp]*:

$!!\text{bin. } \text{bintrunc } n \ (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bin-rest } (\text{bintrunc } n \ \text{bin})$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-rest [simp]*:

$!!\text{bin. } \text{sbintrunc } n \ (\text{bin-rest } (\text{sbintrunc } n \ \text{bin})) = \text{bin-rest } (\text{sbintrunc } n \ \text{bin})$
 $\langle \text{proof} \rangle$

lemma *bintrunc-rest'*:

$\text{bintrunc } n \ o \ \text{bin-rest } o \ \text{bintrunc } n = \text{bin-rest } o \ \text{bintrunc } n$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-rest'*:

$\text{sbintrunc } n \ o \ \text{bin-rest } o \ \text{sbintrunc } n = \text{bin-rest } o \ \text{sbintrunc } n$
 $\langle \text{proof} \rangle$

lemma *rco-lem*:

$f \ o \ g \ o \ f = g \ o \ f \implies f \ o \ (g \ o \ f) \ ^n = g \ ^n \ o \ f$
 $\langle \text{proof} \rangle$

lemma *rco-alt*: $(f \ o \ g) \ ^n \ o \ f = f \ o \ (g \ o \ f) \ ^n$

$\langle \text{proof} \rangle$

lemmas *rco-bintr = bintrunc-rest'*

$[\text{THEN } \text{rco-lem } [\text{THEN } \text{fun-cong}], \text{unfolded } \text{o-def}]$

lemmas *rco-sbintr = sbintrunc-rest'*

$[\text{THEN } \text{rco-lem } [\text{THEN } \text{fun-cong}], \text{unfolded } \text{o-def}]$

4.5 Splitting and concatenation

consts

$\text{bin-split} :: \text{nat} \Rightarrow \text{int} \Rightarrow \text{int} * \text{int}$

primrec

$Z : \text{bin-split } 0 \ w = (w, \text{Numeral.Pls})$

$\text{Suc} : \text{bin-split } (\text{Suc } n) \ w = (\text{let } (w1, w2) = \text{bin-split } n \ (\text{bin-rest } w) \\ \text{in } (w1, w2 \ \text{BIT } \text{bin-last } w))$

consts

$\text{bin-cat} :: \text{int} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int}$

primrec

$Z : \text{bin-cat } w \ 0 \ v = w$
 $\text{Suc} : \text{bin-cat } w \ (\text{Suc } n) \ v = \text{bin-cat } w \ n \ (\text{bin-rest } v) \ \text{BIT } \text{bin-last } v$

4.6 Miscellaneous lemmas

lemmas *funpow-minus-simp* =
 $\text{trans } [\text{OF } \text{gen-minus } [\mathbf{where} \ f = \text{power } f] \ \text{funpow-Suc}, \text{ standard}]$

lemmas *funpow-pred-simp* [simp] =
 $\text{funpow-minus-simp } [\text{of number-of bin}, \text{ simplified nobm1}, \text{ standard}]$

lemmas *replicate-minus-simp* =
 $\text{trans } [\text{OF } \text{gen-minus } [\mathbf{where} \ f = \%n. \text{replicate } n \ x] \ \text{replicate.replicate-Suc},$
 $\text{standard}]$

lemmas *replicate-pred-simp* [simp] =
 $\text{replicate-minus-simp } [\text{of number-of bin}, \text{ simplified nobm1}, \text{ standard}]$

lemmas *power-Suc-no* [simp] = *power-Suc* [of number-of a, standard]

lemmas *power-minus-simp* =
 $\text{trans } [\text{OF } \text{gen-minus } [\mathbf{where} \ f = \text{power } f] \ \text{power-Suc}, \text{ standard}]$

lemmas *power-pred-simp* =
 $\text{power-minus-simp } [\text{of number-of bin}, \text{ simplified nobm1}, \text{ standard}]$

lemmas *power-pred-simp-no* [simp] = *power-pred-simp* [where $f = \text{number-of } f$, standard]

lemma *list-exhaust-size-gt0*:
assumes $y : \bigwedge a \ \text{list}. \ y = a \ \# \ \text{list} \implies P$
shows $0 < \text{length } y \implies P$
 $\langle \text{proof} \rangle$

lemma *list-exhaust-size-eq0*:
assumes $y : y = [] \implies P$
shows $\text{length } y = 0 \implies P$
 $\langle \text{proof} \rangle$

lemma *size-Cons-lem-eq*:
 $y = xa \ \# \ \text{list} \implies \text{size } y = \text{Suc } k \implies \text{size } \text{list} = k$
 $\langle \text{proof} \rangle$

lemma *size-Cons-lem-eq-bin*:
 $y = xa \ \# \ \text{list} \implies \text{size } y = \text{number-of } (\text{Numeral.succ } k) \implies$
 $\text{size } \text{list} = \text{number-of } k$
 $\langle \text{proof} \rangle$

lemmas *ls-splits* =
 $\text{prod.split split-split prod.split-asm split-split-asm split-if-asm}$

lemma *not-B1-is-B0*: $y \neq \text{bit}.B1 \implies y = \text{bit}.B0$
 ⟨*proof*⟩

lemma *B1-ass-B0*:
assumes $y: y = \text{bit}.B0 \implies y = \text{bit}.B1$
shows $y = \text{bit}.B1$
 ⟨*proof*⟩

lemmas *n2s-ths* [*THEN eq-reflection*] = *add-2-eq-Suc add-2-eq-Suc'*

lemmas *s2n-ths* = *n2s-ths* [*symmetric*]

end

5 BitSyntax: Syntactic class for bitwise operations

theory *BitSyntax*
imports *Main*
begin

class *bit* = *type* +
fixes *bitNOT* :: $'a \Rightarrow 'a$
and *bitAND* :: $'a \Rightarrow 'a \Rightarrow 'a$
and *bitOR* :: $'a \Rightarrow 'a \Rightarrow 'a$
and *bitXOR* :: $'a \Rightarrow 'a \Rightarrow 'a$

We want the bitwise operations to bind slightly weaker than + and −, but ~ to bind slightly stronger than *.

notation
bitNOT (*NOT* - [70] 71) **and**
bitAND (**infixr** *AND* 64) **and**
bitOR (**infixr** *OR* 59) **and**
bitXOR (**infixr** *XOR* 59)

Testing and shifting operations.

consts
test-bit :: $'a::\text{bit} \Rightarrow \text{nat} \Rightarrow \text{bool}$ (**infixl** !! 100)
lsb :: $'a::\text{bit} \Rightarrow \text{bool}$
msb :: $'a::\text{bit} \Rightarrow \text{bool}$
set-bit :: $'a::\text{bit} \Rightarrow \text{nat} \Rightarrow \text{bool} \Rightarrow 'a$
set-bits :: $(\text{nat} \Rightarrow \text{bool}) \Rightarrow 'a::\text{bit}$ (**binder** *BITS* 10)
shiffl :: $'a::\text{bit} \Rightarrow \text{nat} \Rightarrow 'a$ (**infixl** << 55)
shiftr :: $'a::\text{bit} \Rightarrow \text{nat} \Rightarrow 'a$ (**infixl** >> 55)

5.1 Bitwise operations on type *bit*

instance *bit* :: *bit*

NOT-bit-def: $NOT\ x \equiv \text{case } x \text{ of } bit.B0 \Rightarrow bit.B1 \mid bit.B1 \Rightarrow bit.B0$

AND-bit-def: $x\ AND\ y \equiv \text{case } x \text{ of } bit.B0 \Rightarrow bit.B0 \mid bit.B1 \Rightarrow y$

OR-bit-def: $x\ OR\ y :: bit \equiv NOT\ (NOT\ x\ AND\ NOT\ y)$

XOR-bit-def: $x\ XOR\ y :: bit \equiv (x\ AND\ NOT\ y)\ OR\ (NOT\ x\ AND\ y)$

<proof>

lemma *bit-simps* [*simp*]:

$NOT\ bit.B0 = bit.B1$

$NOT\ bit.B1 = bit.B0$

$bit.B0\ AND\ y = bit.B0$

$bit.B1\ AND\ y = y$

$bit.B0\ OR\ y = y$

$bit.B1\ OR\ y = bit.B1$

$bit.B0\ XOR\ y = y$

$bit.B1\ XOR\ y = NOT\ y$

<proof>

lemma *bit-extra-simps* [*simp*]:

$x\ AND\ bit.B0 = bit.B0$

$x\ AND\ bit.B1 = x$

$x\ OR\ bit.B1 = bit.B1$

$x\ OR\ bit.B0 = x$

$x\ XOR\ bit.B1 = NOT\ x$

$x\ XOR\ bit.B0 = x$

<proof>

lemma *bit-ops-comm*:

$(x::bit)\ AND\ y = y\ AND\ x$

$(x::bit)\ OR\ y = y\ OR\ x$

$(x::bit)\ XOR\ y = y\ XOR\ x$

<proof>

lemma *bit-ops-same* [*simp*]:

$(x::bit)\ AND\ x = x$

$(x::bit)\ OR\ x = x$

$(x::bit)\ XOR\ x = bit.B0$

<proof>

lemma *bit-not-not* [*simp*]: $NOT\ (NOT\ (x::bit)) = x$

<proof>

end

6 BinOperations: Bitwise Operations on Binary Integers

theory *BinOperations* **imports** *BinGeneral BitSyntax*

begin

6.1 Logical operations

bit-wise logical operations on the int type

instance *int* :: *bit*

int-not-def: $\text{bitNOT} \equiv \text{bin-rec Numeral.Min Numeral.Pls}$
 $(\lambda w \ b \ s. \ s \ \text{BIT} \ (\text{NOT} \ b))$
int-and-def: $\text{bitAND} \equiv \text{bin-rec} \ (\lambda x. \ \text{Numeral.Pls}) \ (\lambda y. \ y)$
 $(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT} \ (b \ \text{AND} \ \text{bin-last } y))$
int-or-def: $\text{bitOR} \equiv \text{bin-rec} \ (\lambda x. \ x) \ (\lambda y. \ \text{Numeral.Min})$
 $(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT} \ (b \ \text{OR} \ \text{bin-last } y))$
int-xor-def: $\text{bitXOR} \equiv \text{bin-rec} \ (\lambda x. \ x) \ \text{bitNOT}$
 $(\lambda w \ b \ s \ y. \ s \ (\text{bin-rest } y) \ \text{BIT} \ (b \ \text{XOR} \ \text{bin-last } y))$
 $\langle \text{proof} \rangle$

lemma *int-not-simps* [*simp*]:

$\text{NOT Numeral.Pls} = \text{Numeral.Min}$
 $\text{NOT Numeral.Min} = \text{Numeral.Pls}$
 $\text{NOT} \ (w \ \text{BIT} \ b) = (\text{NOT} \ w) \ \text{BIT} \ (\text{NOT} \ b)$
 $\langle \text{proof} \rangle$

lemma *int-xor-Pls* [*simp*]:

$\text{Numeral.Pls XOR } x = x$
 $\langle \text{proof} \rangle$

lemma *int-xor-Min* [*simp*]:

$\text{Numeral.Min XOR } x = \text{NOT } x$
 $\langle \text{proof} \rangle$

lemma *int-xor-Bits* [*simp*]:

$(x \ \text{BIT} \ b) \ \text{XOR} \ (y \ \text{BIT} \ c) = (x \ \text{XOR} \ y) \ \text{BIT} \ (b \ \text{XOR} \ c)$
 $\langle \text{proof} \rangle$

lemma *int-xor-x-simps'*:

$w \ \text{XOR} \ (\text{Numeral.Pls} \ \text{BIT} \ \text{bit.B0}) = w$
 $w \ \text{XOR} \ (\text{Numeral.Min} \ \text{BIT} \ \text{bit.B1}) = \text{NOT } w$
 $\langle \text{proof} \rangle$

lemmas *int-xor-extra-simps* [*simp*] = *int-xor-x-simps'* [*simplified arith-simps*]

lemma *int-or-Pls* [*simp*]:

$\text{Numeral.Pls OR } x = x$

$\langle \text{proof} \rangle$

lemma *int-or-Min* [simp]:
 $\text{Numeral.Min OR } x = \text{Numeral.Min}$
 $\langle \text{proof} \rangle$

lemma *int-or-Bits* [simp]:
 $(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \text{ OR } c)$
 $\langle \text{proof} \rangle$

lemma *int-or-x-simps'*:
 $w \text{ OR } (\text{Numeral.Pls BIT bit.B0}) = w$
 $w \text{ OR } (\text{Numeral.Min BIT bit.B1}) = \text{Numeral.Min}$
 $\langle \text{proof} \rangle$

lemmas *int-or-extra-simps* [simp] = *int-or-x-simps'* [simplified arith-simps]

lemma *int-and-Pls* [simp]:
 $\text{Numeral.Pls AND } x = \text{Numeral.Pls}$
 $\langle \text{proof} \rangle$

lemma *int-and-Min* [simp]:
 $\text{Numeral.Min AND } x = x$
 $\langle \text{proof} \rangle$

lemma *int-and-Bits* [simp]:
 $(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \text{ AND } c)$
 $\langle \text{proof} \rangle$

lemma *int-and-x-simps'*:
 $w \text{ AND } (\text{Numeral.Pls BIT bit.B0}) = \text{Numeral.Pls}$
 $w \text{ AND } (\text{Numeral.Min BIT bit.B1}) = w$
 $\langle \text{proof} \rangle$

lemmas *int-and-extra-simps* [simp] = *int-and-x-simps'* [simplified arith-simps]

lemma *bin-ops-comm*:
shows
int-and-comm: $!!y::\text{int}. x \text{ AND } y = y \text{ AND } x$ **and**
int-or-comm: $!!y::\text{int}. x \text{ OR } y = y \text{ OR } x$ **and**
int-xor-comm: $!!y::\text{int}. x \text{ XOR } y = y \text{ XOR } x$
 $\langle \text{proof} \rangle$

lemma *bin-ops-same* [simp]:
 $(x::\text{int}) \text{ AND } x = x$
 $(x::\text{int}) \text{ OR } x = x$
 $(x::\text{int}) \text{ XOR } x = \text{Numeral.Pls}$

$\langle \text{proof} \rangle$

lemma *int-not-not* [simp]: $\text{NOT } (\text{NOT } (x::\text{int})) = x$
 $\langle \text{proof} \rangle$

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min

lemma *bbw-ao-absorb*:
 $!!y::\text{int}. x \text{ AND } (y \text{ OR } x) = x \ \& \ x \text{ OR } (y \text{ AND } x) = x$
 $\langle \text{proof} \rangle$

lemma *bbw-ao-absorbs-other*:
 $x \text{ AND } (x \text{ OR } y) = x \wedge (y \text{ AND } x) \text{ OR } x = (x::\text{int})$
 $(y \text{ OR } x) \text{ AND } x = x \wedge x \text{ OR } (x \text{ AND } y) = (x::\text{int})$
 $(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::\text{int})$
 $\langle \text{proof} \rangle$

lemmas *bbw-ao-absorbs* [simp] = *bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:
 $!!y::\text{int}. (\text{NOT } x) \text{ XOR } y = \text{NOT } (x \text{ XOR } y) \ \&$
 $x \text{ XOR } (\text{NOT } y) = \text{NOT } (x \text{ XOR } y)$
 $\langle \text{proof} \rangle$

lemma *bbw-assocs'*:
 $!!y \ z::\text{int}. (x \text{ AND } y) \text{ AND } z = x \text{ AND } (y \text{ AND } z) \ \&$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } (y \text{ OR } z) \ \&$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } (y \text{ XOR } z)$
 $\langle \text{proof} \rangle$

lemma *int-and-assoc*:
 $(x \text{ AND } y) \text{ AND } (z::\text{int}) = x \text{ AND } (y \text{ AND } z)$
 $\langle \text{proof} \rangle$

lemma *int-or-assoc*:
 $(x \text{ OR } y) \text{ OR } (z::\text{int}) = x \text{ OR } (y \text{ OR } z)$
 $\langle \text{proof} \rangle$

lemma *int-xor-assoc*:
 $(x \text{ XOR } y) \text{ XOR } (z::\text{int}) = x \text{ XOR } (y \text{ XOR } z)$
 $\langle \text{proof} \rangle$

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [simp]:

$(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$
 $(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$
 $(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$
 $\langle proof \rangle$

lemma *bbw-not-dist*:

$!!y::int. \text{ NOT } (x \text{ OR } y) = (\text{ NOT } x) \text{ AND } (\text{ NOT } y)$
 $!!y::int. \text{ NOT } (x \text{ AND } y) = (\text{ NOT } x) \text{ OR } (\text{ NOT } y)$
 $\langle proof \rangle$

lemma *bbw-ao-dist*:

$!!y \ z::int. (x \text{ AND } y) \text{ OR } z =$
 $(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
 $\langle proof \rangle$

lemma *bbw-a-o-dist*:

$!!y \ z::int. (x \text{ OR } y) \text{ AND } z =$
 $(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$
 $\langle proof \rangle$

lemma *plus-and-or* [rule-format]:

$ALL \ y::int. (x \text{ AND } y) + (x \text{ OR } y) = x + y$
 $\langle proof \rangle$

lemma *le-int-or*:

$!!x. \text{ bin-sign } y = \text{ Numeral.Pls } ==> x \leq x \text{ OR } y$
 $\langle proof \rangle$

lemmas *int-and-le* =

xtr3 [OF *bbw-ao-absorbs* (2) [THEN *conjunct2*, *symmetric*] *le-int-or*]

lemma *bin-nth-ops*:

$!!x \ y. \text{ bin-nth } (x \text{ AND } y) \ n = (\text{ bin-nth } x \ n \ \& \ \text{ bin-nth } y \ n)$
 $!!x \ y. \text{ bin-nth } (x \text{ OR } y) \ n = (\text{ bin-nth } x \ n \ | \ \text{ bin-nth } y \ n)$
 $!!x \ y. \text{ bin-nth } (x \text{ XOR } y) \ n = (\text{ bin-nth } x \ n \ \sim = \text{ bin-nth } y \ n)$
 $!!x. \text{ bin-nth } (\text{ NOT } x) \ n = (\sim \text{ bin-nth } x \ n)$
 $\langle proof \rangle$

lemma *bin-add-not*: $x + \text{ NOT } x = \text{ Numeral.Min}$

$\langle proof \rangle$

lemma *bin-trunc-ao*:

$!!x \ y. (\text{ bintrunc } n \ x) \text{ AND } (\text{ bintrunc } n \ y) = \text{ bintrunc } n \ (x \text{ AND } y)$
 $!!x \ y. (\text{ bintrunc } n \ x) \text{ OR } (\text{ bintrunc } n \ y) = \text{ bintrunc } n \ (x \text{ OR } y)$

$\langle \text{proof} \rangle$

lemma *bin-trunc-xor*:

$!!x\ y. \text{bintrunc } n (\text{bintrunc } n\ x\ \text{XOR } \text{bintrunc } n\ y) =$
 $\text{bintrunc } n\ (x\ \text{XOR } y)$

$\langle \text{proof} \rangle$

lemma *bin-trunc-not*:

$!!x. \text{bintrunc } n (\text{NOT } (\text{bintrunc } n\ x)) = \text{bintrunc } n (\text{NOT } x)$

$\langle \text{proof} \rangle$

lemma *bintr-bintr-i*:

$x = \text{bintrunc } n\ y \implies \text{bintrunc } n\ x = \text{bintrunc } n\ y$

$\langle \text{proof} \rangle$

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

6.2 Setting and clearing bits

consts

bin-sc :: *nat* => *bit* => *int* => *int*

primrec

Z : *bin-sc* 0 *b w* = *bin-rest w BIT b*

Suc :

bin-sc (*Suc n*) *b w* = *bin-sc n b (bin-rest w) BIT bin-last w*

lemma *bin-nth-sc* [simp]:

$!!w. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ n = (b = \text{bit.B1})$

$\langle \text{proof} \rangle$

lemma *bin-sc-sc-same* [simp]:

$!!w. \text{bin-sc } n\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ c\ w$

$\langle \text{proof} \rangle$

lemma *bin-sc-sc-diff*:

$!!w\ m. m \sim n \implies$

$\text{bin-sc } m\ c\ (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ b\ (\text{bin-sc } m\ c\ w)$

$\langle \text{proof} \rangle$

lemma *bin-nth-sc-gen*:

$!!w\ m. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ m = (\text{if } m = n \text{ then } b = \text{bit.B1} \text{ else } \text{bin-nth } w\ m)$

$\langle \text{proof} \rangle$

lemma *bin-sc-nth* [simp]:

$!!w. (bin-sc\ n\ (If\ (bin-nth\ w\ n)\ bit.B1\ bit.B0)\ w) = w$
 $\langle proof \rangle$

lemma *bin-sign-sc* [simp]:
 $!!w. bin-sign\ (bin-sc\ n\ b\ w) = bin-sign\ w$
 $\langle proof \rangle$

lemma *bin-sc-bintr* [simp]:
 $!!w\ m. bintrunc\ m\ (bin-sc\ n\ x\ (bintrunc\ m\ (w))) = bintrunc\ m\ (bin-sc\ n\ x\ w)$
 $\langle proof \rangle$

lemma *bin-clr-le*:
 $!!w. bin-sc\ n\ bit.B0\ w \leq w$
 $\langle proof \rangle$

lemma *bin-set-ge*:
 $!!w. bin-sc\ n\ bit.B1\ w \geq w$
 $\langle proof \rangle$

lemma *bintr-bin-clr-le*:
 $!!w\ m. bintrunc\ n\ (bin-sc\ m\ bit.B0\ w) \leq bintrunc\ n\ w$
 $\langle proof \rangle$

lemma *bintr-bin-set-ge*:
 $!!w\ m. bintrunc\ n\ (bin-sc\ m\ bit.B1\ w) \geq bintrunc\ n\ w$
 $\langle proof \rangle$

lemma *bin-sc-FP* [simp]: $bin-sc\ n\ bit.B0\ Numeral.Pls = Numeral.Pls$
 $\langle proof \rangle$

lemma *bin-sc-TM* [simp]: $bin-sc\ n\ bit.B1\ Numeral.Min = Numeral.Min$
 $\langle proof \rangle$

lemmas *bin-sc-simps* = *bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP*

lemma *bin-sc-minus*:
 $0 < n \implies bin-sc\ (Suc\ (n - 1))\ b\ w = bin-sc\ n\ b\ w$
 $\langle proof \rangle$

lemmas *bin-sc-Suc-minus* =
 $trans\ [OF\ bin-sc-minus\ [symmetric]\ bin-sc.Suc,\ standard]$

lemmas *bin-sc-Suc-pred* [simp] =
 $bin-sc-Suc-minus\ [of\ number-of\ bin,\ simplified\ nobm1,\ standard]$

6.3 Operations on lists of booleans

consts
 $bin-to-bl :: nat \Rightarrow int \Rightarrow bool\ list$

$bin\text{-}to\text{-}bl\text{-}aux :: nat \Rightarrow int \Rightarrow bool\ list \Rightarrow bool\ list$
 $bl\text{-}to\text{-}bin :: bool\ list \Rightarrow int$
 $bl\text{-}to\text{-}bin\text{-}aux :: int \Rightarrow bool\ list \Rightarrow int$

$bl\text{-}of\text{-}nth :: nat \Rightarrow (nat \Rightarrow bool) \Rightarrow bool\ list$

primrec

$Nil : bl\text{-}to\text{-}bin\text{-}aux\ w\ [] = w$
 $Cons : bl\text{-}to\text{-}bin\text{-}aux\ w\ (b \# bs) =$
 $\quad bl\text{-}to\text{-}bin\text{-}aux\ (w\ BIT\ (if\ b\ then\ bit.B1\ else\ bit.B0))\ bs$

primrec

$Z : bin\text{-}to\text{-}bl\text{-}aux\ 0\ w\ bl = bl$
 $Suc : bin\text{-}to\text{-}bl\text{-}aux\ (Suc\ n)\ w\ bl =$
 $\quad bin\text{-}to\text{-}bl\text{-}aux\ n\ (bin\text{-}rest\ w)\ ((bin\text{-}last\ w = bit.B1) \# bl)$

defs

$bin\text{-}to\text{-}bl\text{-}def : bin\text{-}to\text{-}bl\ n\ w == bin\text{-}to\text{-}bl\text{-}aux\ n\ w\ []$
 $bl\text{-}to\text{-}bin\text{-}def : bl\text{-}to\text{-}bin\ bs == bl\text{-}to\text{-}bin\text{-}aux\ Numeral.Pls\ bs$

primrec

$Suc : bl\text{-}of\text{-}nth\ (Suc\ n)\ f = f\ n \# bl\text{-}of\text{-}nth\ n\ f$
 $Z : bl\text{-}of\text{-}nth\ 0\ f = []$

consts

$takefill :: 'a \Rightarrow nat \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $app2 :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow 'c\ list$

— takefill - like take but if argument list too short,
 — extends result to get requested length

primrec

$Z : takefill\ fill\ 0\ xs = []$
 $Suc : takefill\ fill\ (Suc\ n)\ xs =$
 $\quad case\ xs\ of\ [] \Rightarrow fill \# takefill\ fill\ n\ xs$
 $\quad | y \# ys \Rightarrow y \# takefill\ fill\ n\ ys$

defs

$app2\text{-}def : app2\ f\ as\ bs == map\ (split\ f)\ (zip\ as\ bs)$

6.4 Splitting and concatenation

— rcat and rsplit

consts

$bin\text{-}rcat :: nat \Rightarrow int\ list \Rightarrow int$
 $bin\text{-}rsplit\text{-}aux :: nat * int\ list * nat * int \Rightarrow int\ list$
 $bin\text{-}rsplit :: nat \Rightarrow (nat * int) \Rightarrow int\ list$
 $bin\text{-}rsplitl\text{-}aux :: nat * int\ list * nat * int \Rightarrow int\ list$
 $bin\text{-}rsplitl :: nat \Rightarrow (nat * int) \Rightarrow int\ list$

```

recdef bin-rsplit-aux measure (fst o snd o snd)
  bin-rsplit-aux (n, bs, (m, c)) =
    (if m = 0 | n = 0 then bs else
     let (a, b) = bin-split n c
     in bin-rsplit-aux (n, b # bs, (m - n, a)))

recdef bin-rsplittl-aux measure (fst o snd o snd)
  bin-rsplittl-aux (n, bs, (m, c)) =
    (if m = 0 | n = 0 then bs else
     let (a, b) = bin-split (min m n) c
     in bin-rsplittl-aux (n, b # bs, (m - n, a)))

defs
  bin-rcat-def : bin-rcat n bs == foldl (%u v. bin-cat u n v) Numeral.Pls bs
  bin-rsplit-def : bin-rsplit n w == bin-rsplit-aux (n, [], w)
  bin-rsplittl-def : bin-rsplittl n w == bin-rsplittl-aux (n, [], w)

declare bin-rsplit-aux.simps [simp del]
declare bin-rsplittl-aux.simps [simp del]

lemma bin-sign-cat:
  !!y. bin-sign (bin-cat x n y) = bin-sign x
  <proof>

lemma bin-cat-Suc-Bit:
  bin-cat w (Suc n) (v BIT b) = bin-cat w n v BIT b
  <proof>

lemma bin-nth-cat:
  !!n y. bin-nth (bin-cat x k y) n =
    (if n < k then bin-nth y n else bin-nth x (n - k))
  <proof>

lemma bin-nth-split:
  !!b c. bin-split n c = (a, b) ==>
    (ALL k. bin-nth a k = bin-nth c (n + k)) &
    (ALL k. bin-nth b k = (k < n & bin-nth c k))
  <proof>

lemma bin-cat-assoc:
  !!z. bin-cat (bin-cat x m y) n z = bin-cat x (m + n) (bin-cat y n z)
  <proof>

lemma bin-cat-assoc-sym: !!z m.
  bin-cat x m (bin-cat y n z) = bin-cat (bin-cat x (m - n) y) (min m n) z
  <proof>

```

lemma *bin-cat-Pls* [simp]:

!!w. *bin-cat Numeral.Pls* n w = *bintrunc* n w

⟨proof⟩

lemma *bintr-cat1*:

!!b. *bintrunc* (k + n) (*bin-cat* a n b) = *bin-cat* (*bintrunc* k a) n b

⟨proof⟩

lemma *bintr-cat*: *bintrunc* m (*bin-cat* a n b) =

bin-cat (*bintrunc* (m - n) a) n (*bintrunc* (min m n) b)

⟨proof⟩

lemma *bintr-cat-same* [simp]:

bintrunc n (*bin-cat* a n b) = *bintrunc* n b

⟨proof⟩

lemma *cat-bintr* [simp]:

!!b. *bin-cat* a n (*bintrunc* n b) = *bin-cat* a n b

⟨proof⟩

lemma *split-bintrunc*:

!!b c. *bin-split* n c = (a, b) ==> b = *bintrunc* n c

⟨proof⟩

lemma *bin-cat-split*:

!!v w. *bin-split* n w = (u, v) ==> w = *bin-cat* u n v

⟨proof⟩

lemma *bin-split-cat*:

!!w. *bin-split* n (*bin-cat* v n w) = (v, *bintrunc* n w)

⟨proof⟩

lemma *bin-split-Pls* [simp]:

bin-split n *Numeral.Pls* = (*Numeral.Pls*, *Numeral.Pls*)

⟨proof⟩

lemma *bin-split-Min* [simp]:

bin-split n *Numeral.Min* = (*Numeral.Min*, *bintrunc* n *Numeral.Min*)

⟨proof⟩

lemma *bin-split-trunc*:

!!m b c. *bin-split* (min m n) c = (a, b) ==>

bin-split n (*bintrunc* m c) = (*bintrunc* (m - n) a, b)

⟨proof⟩

lemma *bin-split-trunc1*:

!!m b c. *bin-split* n c = (a, b) ==>

bin-split n (*bintrunc* m c) = (*bintrunc* (m - n) a, *bintrunc* m b)

⟨proof⟩

lemma *bin-cat-num*:

!!*b*. *bin-cat* *a* *n* *b* = *a* * $2^{\wedge} n$ + *bintrunc* *n* *b*
 ⟨*proof*⟩

lemma *bin-split-num*:

!!*b*. *bin-split* *n* *b* = (*b* *div* $2^{\wedge} n$, *b* *mod* $2^{\wedge} n$)
 ⟨*proof*⟩

6.5 Miscellaneous lemmas

lemma *nth-2p-bin*:

!!*m*. *bin-nth* ($2^{\wedge} n$) *m* = (*m* = *n*)
 ⟨*proof*⟩

lemma *ex-eq-or*:

(*EX* *m*. *n* = *Suc* *m* & (*m* = *k* | *P* *m*)) = (*n* = *Suc* *k* | (*EX* *m*. *n* = *Suc* *m* & *P* *m*))
 ⟨*proof*⟩

end

7 BinBoolList: Bool lists and integers

theory *BinBoolList* **imports** *BinOperations* **begin**

7.1 Arithmetic in terms of bool lists

consts

rbl-succ :: *bool list* => *bool list*
rbl-pred :: *bool list* => *bool list*
rbl-add :: *bool list* => *bool list* => *bool list*
rbl-mult :: *bool list* => *bool list* => *bool list*

primrec

Nil: *rbl-succ* *Nil* = *Nil*
Cons: *rbl-succ* (*x* # *xs*) = (if *x* then *False* # *rbl-succ* *xs* else *True* # *xs*)

primrec

Nil : *rbl-pred* *Nil* = *Nil*
Cons : *rbl-pred* (*x* # *xs*) = (if *x* then *False* # *xs* else *True* # *rbl-pred* *xs*)

primrec

Nil : *rbl-add* *Nil* *x* = *Nil*
Cons : *rbl-add* (*y* # *ys*) *x* = (let *ws* = *rbl-add* *ys* (tl *x*) in

$$(y \sim = hd\ x) \# (if\ hd\ x \ \&\ y\ then\ rbl\ succ\ ws\ else\ ws))$$
primrec

$$Nil : rbl\mult\ Nil\ x = Nil$$

$$Cons : rbl\mult\ (y \# ys)\ x = (let\ ws = False \# rbl\mult\ ys\ x\ in \\ if\ y\ then\ rbl\add\ ws\ x\ else\ ws)$$

lemma *tl-take*: $tl\ (take\ n\ l) = take\ (n - 1)\ (tl\ l)$
 $\langle proof \rangle$

lemma *take-butlast* [rule-format] :
 $ALL\ n.\ n < length\ l \longrightarrow take\ n\ (butlast\ l) = take\ n\ l$
 $\langle proof \rangle$

lemma *butlast-take* [rule-format] :
 $ALL\ n.\ n \leq length\ l \longrightarrow butlast\ (take\ n\ l) = take\ (n - 1)\ l$
 $\langle proof \rangle$

lemma *butlast-drop* [rule-format] :
 $ALL\ n.\ butlast\ (drop\ n\ l) = drop\ n\ (butlast\ l)$
 $\langle proof \rangle$

lemma *butlast-power*:
 $(butlast\ ^\ n)\ bl = take\ (length\ bl - n)\ bl$
 $\langle proof \rangle$

lemma *bin-to-bl-aux-Pls-minus-simp*:
 $0 < n \implies bin\to\bl\aux\ n\ Numeral.Pls\ bl = \\ bin\to\bl\aux\ (n - 1)\ Numeral.Pls\ (False \# bl)$
 $\langle proof \rangle$

lemma *bin-to-bl-aux-Min-minus-simp*:
 $0 < n \implies bin\to\bl\aux\ n\ Numeral.Min\ bl = \\ bin\to\bl\aux\ (n - 1)\ Numeral.Min\ (True \# bl)$
 $\langle proof \rangle$

lemma *bin-to-bl-aux-Bit-minus-simp*:
 $0 < n \implies bin\to\bl\aux\ n\ (w\ BIT\ b)\ bl = \\ bin\to\bl\aux\ (n - 1)\ w\ ((b = bit.B1) \# bl)$
 $\langle proof \rangle$

declare *bin-to-bl-aux-Pls-minus-simp* [simp]
 $bin\to\bl\aux\ Min\ minus\ simp$ [simp]
 $bin\to\bl\aux\ Bit\ minus\ simp$ [simp]

lemma *bl-to-bin-aux-append* [rule-format] :
 $ALL\ w.\ bl\to\bin\aux\ w\ (bs\ @\ cs) = bl\to\bin\aux\ (bl\to\bin\aux\ w\ bs)\ cs$

$\langle \text{proof} \rangle$

lemma *bin-to-bl-aux-append* [rule-format] :

ALL w bs. bin-to-bl-aux n w bs @ cs = bin-to-bl-aux n w (bs @ cs)

$\langle \text{proof} \rangle$

lemma *bl-to-bin-append*:

bl-to-bin (bs @ cs) = bl-to-bin-aux (bl-to-bin bs) cs

$\langle \text{proof} \rangle$

lemma *bin-to-bl-aux-alt*:

bin-to-bl-aux n w bs = bin-to-bl n w @ bs

$\langle \text{proof} \rangle$

lemma *bin-to-bl-0*: *bin-to-bl 0 bs = []*

$\langle \text{proof} \rangle$

lemma *size-bin-to-bl-aux* [rule-format] :

ALL w bs. size (bin-to-bl-aux n w bs) = n + length bs

$\langle \text{proof} \rangle$

lemma *size-bin-to-bl*: *size (bin-to-bl n w) = n*

$\langle \text{proof} \rangle$

lemma *bin-bl-bin'* [rule-format] :

ALL w bs. bl-to-bin (bin-to-bl-aux n w bs) =

bl-to-bin-aux (bintrunc n w) bs

$\langle \text{proof} \rangle$

lemma *bin-bl-bin*: *bl-to-bin (bin-to-bl n w) = bintrunc n w*

$\langle \text{proof} \rangle$

lemma *bl-bin-bl'* [rule-format] :

ALL w n. bin-to-bl (n + length bs) (bl-to-bin-aux w bs) =

bin-to-bl-aux n w bs

$\langle \text{proof} \rangle$

lemma *bl-bin-bl*: *bin-to-bl (length bs) (bl-to-bin bs) = bs*

$\langle \text{proof} \rangle$

declare

bin-to-bl-0 [simp]

size-bin-to-bl [simp]

bin-bl-bin [simp]

bl-bin-bl [simp]

lemma *bl-to-bin-inj*:

bl-to-bin bs = bl-to-bin cs ==> length bs = length cs ==> bs = cs

$\langle \text{proof} \rangle$

lemma *bl-to-bin-False*: *bl-to-bin (False # bl) = bl-to-bin bl*
 ⟨proof⟩

lemma *bl-to-bin-Nil*: *bl-to-bin [] = Numeral.Pls*
 ⟨proof⟩

lemma *bin-to-bl-Pls-aux* [rule-format] :
ALL bl. bin-to-bl-aux n Numeral.Pls bl = replicate n False @ bl
 ⟨proof⟩

lemma *bin-to-bl-Pls*: *bin-to-bl n Numeral.Pls = replicate n False*
 ⟨proof⟩

lemma *bin-to-bl-Min-aux* [rule-format] :
ALL bl. bin-to-bl-aux n Numeral.Min bl = replicate n True @ bl
 ⟨proof⟩

lemma *bin-to-bl-Min*: *bin-to-bl n Numeral.Min = replicate n True*
 ⟨proof⟩

lemma *bl-to-bin-rep-F*:
bl-to-bin (replicate n False @ bl) = bl-to-bin bl
 ⟨proof⟩

lemma *bin-to-bl-trunc*:
n <= m ==> bin-to-bl n (bintrunc m w) = bin-to-bl n w
 ⟨proof⟩

declare
bin-to-bl-trunc [simp]
bl-to-bin-False [simp]
bl-to-bin-Nil [simp]

lemma *bin-to-bl-aux-bintr* [rule-format] :
ALL m bin bl. bin-to-bl-aux n (bintrunc m bin) bl =
replicate (n - m) False @ bin-to-bl-aux (min n m) bin bl
 ⟨proof⟩

lemmas *bin-to-bl-bintr* =
bin-to-bl-aux-bintr [where *bl* = [], folded *bin-to-bl-def*]

lemma *bl-to-bin-rep-False*: *bl-to-bin (replicate n False) = Numeral.Pls*
 ⟨proof⟩

lemma *len-bin-to-bl-aux* [rule-format] :
ALL w bs. length (bin-to-bl-aux n w bs) = n + length bs
 ⟨proof⟩

lemma *len-bin-to-bl* [simp]: $\text{length } (\text{bin-to-bl } n \ w) = n$
 ⟨proof⟩

lemma *sign-bl-bin'* [rule-format] :
 ALL w . $\text{bin-sign } (\text{bl-to-bin-aux } w \ bs) = \text{bin-sign } w$
 ⟨proof⟩

lemma *sign-bl-bin*: $\text{bin-sign } (\text{bl-to-bin } bs) = \text{Numeral.Pls}$
 ⟨proof⟩

lemma *bl-sbin-sign-aux* [rule-format] :
 ALL $w \ bs$. $\text{hd } (\text{bin-to-bl-aux } (\text{Suc } n) \ w \ bs) =$
 $(\text{bin-sign } (\text{sbintrunc } n \ w) = \text{Numeral.Min})$
 ⟨proof⟩

lemma *bl-sbin-sign*:
 $\text{hd } (\text{bin-to-bl } (\text{Suc } n) \ w) = (\text{bin-sign } (\text{sbintrunc } n \ w) = \text{Numeral.Min})$
 ⟨proof⟩

lemma *bin-nth-of-bl-aux* [rule-format] :
 ALL w . $\text{bin-nth } (\text{bl-to-bin-aux } w \ bl) \ n =$
 $(n < \text{size } bl \ \& \ \text{rev } bl \ ! \ n \mid n >= \text{length } bl \ \& \ \text{bin-nth } w \ (n - \text{size } bl))$
 ⟨proof⟩

lemma *bin-nth-of-bl*: $\text{bin-nth } (\text{bl-to-bin } bl) \ n = (n < \text{length } bl \ \& \ \text{rev } bl \ ! \ n)$
 ⟨proof⟩

lemma *bin-nth-bl* [rule-format] : ALL $m \ w$. $n < m \ \longrightarrow$
 $\text{bin-nth } w \ n = \text{nth } (\text{rev } (\text{bin-to-bl } m \ w)) \ n$
 ⟨proof⟩

lemma *nth-rev* [rule-format] :
 $n < \text{length } xs \ \longrightarrow \ \text{rev } xs \ ! \ n = xs \ ! \ (\text{length } xs - 1 - n)$
 ⟨proof⟩

lemmas *nth-rev-alt* = *nth-rev* [where $xs = \text{rev } ys$, simplified, standard]

lemma *nth-bin-to-bl-aux* [rule-format] :
 ALL $w \ n \ bl$. $n < m + \text{length } bl \ \longrightarrow \ (\text{bin-to-bl-aux } m \ w \ bl) \ ! \ n =$
 $(\text{if } n < m \ \text{then } \text{bin-nth } w \ (m - 1 - n) \ \text{else } bl \ ! \ (n - m))$
 ⟨proof⟩

lemma *nth-bin-to-bl*: $n < m \ \Longrightarrow \ (\text{bin-to-bl } m \ w) \ ! \ n = \text{bin-nth } w \ (m - \text{Suc } n)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p-aux* [rule-format] :
 ALL w . $\text{bl-to-bin-aux } w \ bs < (w + 1) * (2 \wedge \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-lt2p*: $bl\text{-}to\text{-}bin\ bs < (2 \wedge length\ bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge2p-aux* [rule-format] :
 $ALL\ w.\ bl\text{-}to\text{-}bin\text{-}aux\ w\ bs \geq w * (2 \wedge length\ bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge0*: $bl\text{-}to\text{-}bin\ bs \geq 0$
 ⟨proof⟩

lemma *butlast-rest-bin*:
 $butlast\ (bin\text{-}to\text{-}bl\ n\ w) = bin\text{-}to\text{-}bl\ (n - 1)\ (bin\text{-}rest\ w)$
 ⟨proof⟩

lemmas *butlast-bin-rest = butlast-rest-bin*
 [where $w = bl\text{-}to\text{-}bin\ bl$ and $n = length\ bl$, simplified, standard]

lemma *butlast-rest-bl2bin-aux* [rule-format] :
 $ALL\ w.\ bl \sim = [] \dashrightarrow$
 $bl\text{-}to\text{-}bin\text{-}aux\ w\ (butlast\ bl) = bin\text{-}rest\ (bl\text{-}to\text{-}bin\text{-}aux\ w\ bl)$
 ⟨proof⟩

lemma *butlast-rest-bl2bin*:
 $bl\text{-}to\text{-}bin\ (butlast\ bl) = bin\text{-}rest\ (bl\text{-}to\text{-}bin\ bl)$
 ⟨proof⟩

lemma *trunc-bl2bin-aux* [rule-format] :
 $ALL\ w.\ bintrunc\ m\ (bl\text{-}to\text{-}bin\text{-}aux\ w\ bl) =$
 $bl\text{-}to\text{-}bin\text{-}aux\ (bintrunc\ (m - length\ bl)\ w)\ (drop\ (length\ bl - m)\ bl)$
 ⟨proof⟩

lemma *trunc-bl2bin*:
 $bintrunc\ m\ (bl\text{-}to\text{-}bin\ bl) = bl\text{-}to\text{-}bin\ (drop\ (length\ bl - m)\ bl)$
 ⟨proof⟩

lemmas *trunc-bl2bin-len* [simp] =
 $trunc\text{-}bl2bin\ [of\ length\ bl\ bl,\ simplified,\ standard]$

lemma *bl2bin-drop*:
 $bl\text{-}to\text{-}bin\ (drop\ k\ bl) = bintrunc\ (length\ bl - k)\ (bl\text{-}to\text{-}bin\ bl)$
 ⟨proof⟩

lemma *nth-rest-power-bin* [rule-format] :
 $ALL\ n.\ bin\text{-}nth\ ((bin\text{-}rest\ \wedge k)\ w)\ n = bin\text{-}nth\ w\ (n + k)$
 ⟨proof⟩

lemma *take-rest-power-bin*:
 $m \leq n \implies take\ m\ (bin\text{-}to\text{-}bl\ n\ w) = bin\text{-}to\text{-}bl\ m\ ((bin\text{-}rest\ \wedge (n - m))\ w)$
 ⟨proof⟩

lemma *hd-butlast*: $\text{size } xs > 1 \implies \text{hd } (\text{butlast } xs) = \text{hd } xs$
 ⟨proof⟩

lemma *last-bin-last'* [rule-format] :
 $\text{ALL } w. \text{size } xs > 0 \dashrightarrow \text{last } xs = (\text{bin-last } (\text{bl-to-bin-aux } w \ xs) = \text{bit.B1})$
 ⟨proof⟩

lemma *last-bin-last*:
 $\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last } (\text{bl-to-bin } xs) = \text{bit.B1})$
 ⟨proof⟩

lemma *bin-last-last*:
 $\text{bin-last } w = (\text{if last } (\text{bin-to-bl } (\text{Suc } n) \ w) \text{ then bit.B1 else bit.B0})$
 ⟨proof⟩

lemma *app2-Nil* [simp]: $\text{app2 } f \ [] \ ys = []$
 ⟨proof⟩

lemma *app2-Cons* [simp]:
 $\text{app2 } f \ (x \# \ xs) \ (y \# \ ys) = f \ x \ y \# \ \text{app2 } f \ xs \ ys$
 ⟨proof⟩

lemma *bl-xor-aux-bin* [rule-format] : $\text{ALL } v \ w \ bs \ cs.$
 $\text{app2 } (\%x \ y. \ x \ \sim = \ y) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ \text{XOR } w) \ (\text{app2 } (\%x \ y. \ x \ \sim = \ y) \ bs \ cs)$
 ⟨proof⟩

lemma *bl-or-aux-bin* [rule-format] : $\text{ALL } v \ w \ bs \ cs.$
 $\text{app2 } (\text{op } |) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ \text{OR } w) \ (\text{app2 } (\text{op } |) \ bs \ cs)$
 ⟨proof⟩

lemma *bl-and-aux-bin* [rule-format] : $\text{ALL } v \ w \ bs \ cs.$
 $\text{app2 } (\text{op } \&) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ \text{AND } w) \ (\text{app2 } (\text{op } \&) \ bs \ cs)$
 ⟨proof⟩

lemma *bl-not-aux-bin* [rule-format] :
 $\text{ALL } w \ cs. \text{map Not } (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (\text{NOT } w) \ (\text{map Not } cs)$
 ⟨proof⟩

lemmas *bl-not-bin* = *bl-not-aux-bin*
 [where $cs = []$, unfolded *bin-to-bl-def* [symmetric] map.simps]

lemmas *bl-and-bin* = *bl-and-aux-bin* [where $bs=[]$ and $cs=[]$,

unfolded app2-Nil, folded bin-to-bl-def]

lemmas *bl-or-bin* = *bl-or-aux-bin* [**where** *bs*=[] **and** *cs*=[],
unfolded app2-Nil, folded bin-to-bl-def]

lemmas *bl-xor-bin* = *bl-xor-aux-bin* [**where** *bs*=[] **and** *cs*=[],
unfolded app2-Nil, folded bin-to-bl-def]

lemma *drop-bin2bl-aux* [*rule-format*] :
ALL m bin bs. drop m (bin-to-bl-aux n bin bs) =
bin-to-bl-aux (n - m) bin (drop (m - n) bs)
<proof>

lemma *drop-bin2bl*: *drop m (bin-to-bl n bin) = bin-to-bl (n - m) bin*
<proof>

lemma *take-bin2bl-lem1* [*rule-format*] :
ALL w bs. take m (bin-to-bl-aux m w bs) = bin-to-bl m w
<proof>

lemma *take-bin2bl-lem* [*rule-format*] :
ALL w bs. take m (bin-to-bl-aux (m + n) w bs) =
take m (bin-to-bl (m + n) w)
<proof>

lemma *bin-split-take* [*rule-format*] :
ALL b c. bin-split n c = (a, b) -->
bin-to-bl m a = take m (bin-to-bl (m + n) c)
<proof>

lemma *bin-split-take1*:
k = m + n ==> bin-split n c = (a, b) ==>
bin-to-bl m a = take m (bin-to-bl k c)
<proof>

lemma *nth-takefill* [*rule-format*] : *ALL m l. m < n -->*
takefill fill n l ! m = (if m < length l then l ! m else fill)
<proof>

lemma *takefill-alt* [*rule-format*] :
ALL l. takefill fill n l = take n l @ replicate (n - length l) fill
<proof>

lemma *takefill-replicate* [*simp*]:
takefill fill n (replicate m fill) = replicate n fill
<proof>

lemma *takefill-le'* [*rule-format*] :
ALL l n. n = m + k --> takefill x m (takefill x n l) = takefill x m l

$\langle \text{proof} \rangle$

lemma *length-takefill* [simp]: $\text{length } (\text{takefill fill } n \ l) = n$
 $\langle \text{proof} \rangle$

lemma *take-takefill'*:
 $!!w \ n. \ n = k + m \implies \text{take } k \ (\text{takefill fill } n \ w) = \text{takefill fill } k \ w$
 $\langle \text{proof} \rangle$

lemma *drop-takefill*:
 $!!w. \ \text{drop } k \ (\text{takefill fill } (m + k) \ w) = \text{takefill fill } m \ (\text{drop } k \ w)$
 $\langle \text{proof} \rangle$

lemma *takefill-le* [simp]:
 $m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
 $\langle \text{proof} \rangle$

lemma *take-takefill* [simp]:
 $m \leq n \implies \text{take } m \ (\text{takefill fill } n \ w) = \text{takefill fill } m \ w$
 $\langle \text{proof} \rangle$

lemma *takefill-append*:
 $\text{takefill fill } (m + \text{length } xs) \ (xs \ @ \ w) = xs \ @ \ (\text{takefill fill } m \ w)$
 $\langle \text{proof} \rangle$

lemma *takefill-same'*:
 $l = \text{length } xs \implies \text{takefill fill } l \ xs = xs$
 $\langle \text{proof} \rangle$

lemmas *takefill-same* [simp] = *takefill-same'* [OF refl]

lemma *takefill-bintrunc*:
 $\text{takefill False } n \ bl = \text{rev } (\text{bin-to-bl } n \ (\text{bl-to-bin } (\text{rev } bl)))$
 $\langle \text{proof} \rangle$

lemma *bl-bin-bl-rtf*:
 $\text{bin-to-bl } n \ (\text{bl-to-bin } bl) = \text{rev } (\text{takefill False } n \ (\text{rev } bl))$
 $\langle \text{proof} \rangle$

lemmas *bl-bin-bl-rep-drop* =
 bl-bin-bl-rtf [simplified takefill-alt,
simplified, simplified rev-take, simplified]

lemma *tf-rev*:
 $n + k = m + \text{length } bl \implies \text{takefill } x \ m \ (\text{rev } (\text{takefill } y \ n \ bl)) =$
 $\text{rev } (\text{takefill } y \ m \ (\text{rev } (\text{takefill } x \ k \ (\text{rev } bl))))$
 $\langle \text{proof} \rangle$

lemma *takefill-minus*:

$0 < n ==> \text{takefill fill (Suc (n - 1)) w} = \text{takefill fill n w}$
 $\langle \text{proof} \rangle$

lemmas *takefill-Suc-cases* =
list.cases [THEN takefill.Suc [THEN trans], standard]

lemmas *takefill-Suc-Nil* = *takefill-Suc-cases (1)*
lemmas *takefill-Suc-Cons* = *takefill-Suc-cases (2)*

lemmas *takefill-minus-simps* = *takefill-Suc-cases [THEN [2]*
takefill-minus [symmetric, THEN trans], standard]

lemmas *takefill-pred-simps* [*simp*] =
takefill-minus-simps [where n=number-of bin, simplified nobm1, standard]

lemma *bl-to-bin-aux-cat*:
 $!!nv\ v. \text{bl-to-bin-aux (bin-cat w nv v) bs} =$
 $\text{bin-cat w (nv + length bs) (bl-to-bin-aux v bs)}$
 $\langle \text{proof} \rangle$

lemma *bin-to-bl-aux-cat*:
 $!!w\ bs. \text{bin-to-bl-aux (nv + nw) (bin-cat v nw w) bs} =$
 $\text{bin-to-bl-aux nv v (bin-to-bl-aux nw w bs)}$
 $\langle \text{proof} \rangle$

lemmas *bl-to-bin-aux-alt* =
bl-to-bin-aux-cat [where nv = 0 and v = Numeral.Pls,
simplified bl-to-bin-def [symmetric], simplified]

lemmas *bin-to-bl-cat* =
bin-to-bl-aux-cat [where bs = [], folded bin-to-bl-def]

lemmas *bl-to-bin-aux-app-cat* =
trans [OF bl-to-bin-aux-append bl-to-bin-aux-alt]

lemmas *bin-to-bl-aux-cat-app* =
trans [OF bin-to-bl-aux-cat bin-to-bl-aux-alt]

lemmas *bl-to-bin-app-cat* = *bl-to-bin-aux-app-cat*
 $[\text{where } w = \text{Numeral.Pls, folded bl-to-bin-def}]$

lemmas *bin-to-bl-cat-app* = *bin-to-bl-aux-cat-app*
 $[\text{where } bs = [], \text{folded bin-to-bl-def}]$

lemma *bl-to-bin-app-cat-alt*:
 $\text{bin-cat (bl-to-bin cs) n w} = \text{bl-to-bin (cs @ bin-to-bl n w)}$

$\langle proof \rangle$

lemma *mask-lem*: $(bl\text{-}to\text{-}bin\ (True \# replicate\ n\ False)) =$
 $Numeral.succ\ (bl\text{-}to\text{-}bin\ (replicate\ n\ True))$
 $\langle proof \rangle$

lemma *length-bl-of-nth* [simp]: $length\ (bl\text{-}of\text{-}nth\ n\ f) = n$
 $\langle proof \rangle$

lemma *nth-bl-of-nth* [simp]:
 $m < n \implies rev\ (bl\text{-}of\text{-}nth\ n\ f) ! m = f\ m$
 $\langle proof \rangle$

lemma *bl-of-nth-inj*:
 $(!!k. k < n \implies f\ k = g\ k) \implies bl\text{-}of\text{-}nth\ n\ f = bl\text{-}of\text{-}nth\ n\ g$
 $\langle proof \rangle$

lemma *bl-of-nth-nth-le* [rule-format] : *ALL xs.*
 $length\ xs \geq n \longrightarrow bl\text{-}of\text{-}nth\ n\ (nth\ (rev\ xs)) = drop\ (length\ xs - n)\ xs$
 $\langle proof \rangle$

lemmas *bl-of-nth-nth* [simp] = *order-refl* [THEN *bl-of-nth-nth-le*, *simplified*]

lemma *size-rbl-pred*: $length\ (rbl\text{-}pred\ bl) = length\ bl$
 $\langle proof \rangle$

lemma *size-rbl-succ*: $length\ (rbl\text{-}succ\ bl) = length\ bl$
 $\langle proof \rangle$

lemma *size-rbl-add*:
 $!!cl. length\ (rbl\text{-}add\ bl\ cl) = length\ bl$
 $\langle proof \rangle$

lemma *size-rbl-mult*:
 $!!cl. length\ (rbl\text{-}mult\ bl\ cl) = length\ bl$
 $\langle proof \rangle$

lemmas *rbl-sizes* [simp] =
 $size\text{-}rbl\text{-}pred\ size\text{-}rbl\text{-}succ\ size\text{-}rbl\text{-}add\ size\text{-}rbl\text{-}mult$

lemmas *rbl-Nils* =
 $rbl\text{-}pred.Nil\ rbl\text{-}succ.Nil\ rbl\text{-}add.Nil\ rbl\text{-}mult.Nil$

lemma *rbl-pred*:
 $!!bin. rbl\text{-}pred\ (rev\ (bin\text{-}to\text{-}bl\ n\ bin)) = rev\ (bin\text{-}to\text{-}bl\ n\ (Numeral.pred\ bin))$
 $\langle proof \rangle$

lemma *rbl-succ*:

$!!bin. rbl_succ (rev (bin_to_bl\ n\ bin)) = rev (bin_to_bl\ n\ (Numeral.succ\ bin))$
 $\langle proof \rangle$

lemma *rbl-add*:

$!!bina\ binb. rbl_add (rev (bin_to_bl\ n\ bina)) (rev (bin_to_bl\ n\ binb)) =$
 $rev (bin_to_bl\ n\ (bina + binb))$
 $\langle proof \rangle$

lemma *rbl-add-app2*:

$!!blb. length\ blb \geq length\ bla ==>$
 $rbl_add\ bla\ (blb\ @\ blc) = rbl_add\ bla\ blb$
 $\langle proof \rangle$

lemma *rbl-add-take2*:

$!!blb. length\ blb \geq length\ bla ==>$
 $rbl_add\ bla\ (take\ (length\ bla)\ blb) = rbl_add\ bla\ blb$
 $\langle proof \rangle$

lemma *rbl-add-long*:

$m \geq n ==> rbl_add (rev (bin_to_bl\ n\ bina)) (rev (bin_to_bl\ m\ binb)) =$
 $rev (bin_to_bl\ n\ (bina + binb))$
 $\langle proof \rangle$

lemma *rbl-mult-app2*:

$!!blb. length\ blb \geq length\ bla ==>$
 $rbl_mult\ bla\ (blb\ @\ blc) = rbl_mult\ bla\ blb$
 $\langle proof \rangle$

lemma *rbl-mult-take2*:

$length\ blb \geq length\ bla ==>$
 $rbl_mult\ bla\ (take\ (length\ bla)\ blb) = rbl_mult\ bla\ blb$
 $\langle proof \rangle$

lemma *rbl-mult-gt1*:

$m \geq length\ bl ==> rbl_mult\ bl\ (rev (bin_to_bl\ m\ binb)) =$
 $rbl_mult\ bl\ (rev (bin_to_bl\ (length\ bl)\ binb))$
 $\langle proof \rangle$

lemma *rbl-mult-gt*:

$m > n ==> rbl_mult (rev (bin_to_bl\ n\ bina)) (rev (bin_to_bl\ m\ binb)) =$
 $rbl_mult (rev (bin_to_bl\ n\ bina)) (rev (bin_to_bl\ n\ binb))$
 $\langle proof \rangle$

lemmas *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

lemma *rbbl-Cons*:

$b \# rev (bin_to_bl\ n\ x) = rev (bin_to_bl\ (Suc\ n)\ (x\ BIT\ If\ b\ bit.B1\ bit.B0))$
 $\langle proof \rangle$

lemma *rbl-mult: !!bina binb.*

*rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =
rev (bin-to-bl n (bina * binb))*
<proof>

lemma *rbl-add-split:*

*P (rbl-add (y # ys) (x # xs)) =
(ALL ws. length ws = length ys --> ws = rbl-add ys xs -->
(y --> ((x --> P (False # rbl-succ ws)) & (~ x --> P (True # ws))))*
&
(~ y --> P (x # ws)))
<proof>

lemma *rbl-mult-split:*

*P (rbl-mult (y # ys) xs) =
(ALL ws. length ws = Suc (length ys) --> ws = False # rbl-mult ys xs -->
(y --> P (rbl-add ws xs)) & (~ y --> P ws))*
<proof>

lemma *and-len: xs = ys ==> xs = ys & length xs = length ys*

<proof>

lemma *size-if: size (if p then xs else ys) = (if p then size xs else size ys)*

<proof>

lemma *tl-if: tl (if p then xs else ys) = (if p then tl xs else tl ys)*

<proof>

lemma *hd-if: hd (if p then xs else ys) = (if p then hd xs else hd ys)*

<proof>

lemma *if-Not-x: (if p then ~ x else x) = (p = (~ x))*

<proof>

lemma *if-x-Not: (if p then x else ~ x) = (p = x)*

<proof>

lemma *if-same-and: (If p x y & If p u v) = (if p then x & u else y & v)*

<proof>

lemma *if-same-eq: (If p x y = (If p u v)) = (if p then x = (u) else y = (v))*

<proof>

lemma *if-same-eq-not:*

(If p x y = (~ If p u v)) = (if p then x = (~u) else y = (~v))
<proof>

lemma *if-Cons*: $(\text{if } p \text{ then } x \# xs \text{ else } y \# ys) = \text{If } p \ x \ y \# \text{If } p \ xs \ ys$
 ⟨proof⟩

lemma *if-single*:
 $(\text{if } xc \text{ then } [xab] \text{ else } [an]) = [\text{if } xc \text{ then } xab \text{ else } an]$
 ⟨proof⟩

lemma *if-bool-simps*:
 $\text{If } p \ \text{True} \ y = (p \mid y) \ \& \ \text{If } p \ \text{False} \ y = (\sim p \ \& \ y) \ \&$
 $\text{If } p \ y \ \text{True} = (p \dashrightarrow y) \ \& \ \text{If } p \ y \ \text{False} = (p \ \& \ y)$
 ⟨proof⟩

lemmas *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

lemmas *segr* = *eq-reflection* [where $x = \text{size } w, \text{ standard}$]

lemmas *tl-Nil* = *tl.simps* (1)
lemmas *tl-Cons* = *tl.simps* (2)

7.2 Repeated splitting or concatenation

lemma *sclem*:
 $\text{size } (\text{concat } (\text{map } (\text{bin-to-bl } n) \ xs)) = \text{length } xs * n$
 ⟨proof⟩

lemma *bin-cat-foldl-lem* [rule-format] :
 $\text{ALL } x. \text{foldl } (\%u. \text{bin-cat } u \ n) \ x \ xs =$
 $\text{bin-cat } x \ (\text{size } xs * n) \ (\text{foldl } (\%u. \text{bin-cat } u \ n) \ y \ xs)$
 ⟨proof⟩

lemma *bin-rcat-bl*:
 $(\text{bin-rcat } n \ wl) = \text{bl-to-bin } (\text{concat } (\text{map } (\text{bin-to-bl } n) \ wl))$
 ⟨proof⟩

lemmas *bin-rsplit-aux-simps* = *bin-rsplit-aux.simps bin-rsplitl-aux.simps*
lemmas *rsplit-aux-simps* = *bin-rsplit-aux-simps*

lemmas *th-if-simp1* = *split-if* [where $P = op = l,$
 $\text{THEN } \text{iffD1}, \text{ THEN } \text{conjunct1}, \text{ THEN } mp, \text{ standard}$]
lemmas *th-if-simp2* = *split-if* [where $P = op = l,$
 $\text{THEN } \text{iffD1}, \text{ THEN } \text{conjunct2}, \text{ THEN } mp, \text{ standard}$]

lemmas *rsplit-aux-simp1s* = *rsplit-aux-simps* [THEN *th-if-simp1*]

lemmas *rsplit-aux-simp2ls* = *rsplit-aux-simps* [THEN *th-if-simp2*]

lemmas *bin-rsplit-aux-simp2s* [simp] = *rsplit-aux-simp2ls* [unfolded *Let-def*]
lemmas *rbocl* = *bin-rsplit-aux-simp2s* (2)

lemmas *rsplit-aux-0-simps* [*simp*] =
rsplit-aux-simp1s [*OF disjI1*] *rsplit-aux-simp1s* [*OF disjI2*]

lemma *bin-rsplit-aux-append*:
 $\text{bin-rsplit-aux } (n, \text{bs} @ \text{cs}, m, c) = \text{bin-rsplit-aux } (n, \text{bs}, m, c) @ \text{cs}$
 ⟨*proof*⟩

lemma *bin-rsplittl-aux-append*:
 $\text{bin-rsplittl-aux } (n, \text{bs} @ \text{cs}, m, c) = \text{bin-rsplittl-aux } (n, \text{bs}, m, c) @ \text{cs}$
 ⟨*proof*⟩

lemmas *rsplit-aux-apps* [where *bs* = []] =
bin-rsplit-aux-append *bin-rsplittl-aux-append*

lemmas *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplittl-def*

lemmas *rsplit-aux-alt*s = *rsplit-aux-apps*
 [*unfolded append-Nil* *rsplit-def-auxs* [*symmetric*]]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) \text{ } w = \text{bin-split } n \text{ } w$
 ⟨*proof*⟩

lemmas *bin-split-minus-simp* =
bin-split.Suc [*THEN* [2] *bin-split-minus* [*symmetric*, *THEN trans*], *standard*]

lemma *bin-split-pred-simp* [*simp*]:
 $(0::\text{nat}) < \text{number-of bin} \implies$
 $\text{bin-split } (\text{number-of bin}) \text{ } w =$
 $(\text{let } (w1, w2) = \text{bin-split } (\text{number-of } (\text{Numeral.pred bin})) \text{ } (\text{bin-rest } w)$
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$
 ⟨*proof*⟩

declare *bin-split-pred-simp* [*simp*]

lemma *bin-rsplit-aux-simp-alt*:
 $\text{bin-rsplit-aux } (n, \text{bs}, m, c) =$
 $(\text{if } m = 0 \vee n = 0$
 $\text{then } \text{bs}$
 $\text{else let } (a, b) = \text{bin-split } n \text{ } c \text{ in bin-rsplit } n \text{ } (m - n, a) @ b \# \text{bs})$
 ⟨*proof*⟩

lemmas *bin-rsplit-simp-alt* =
trans [*OF bin-rsplit-def* [*THEN meta-eq-to-obj-eq*]
bin-rsplit-aux-simp-alt, *standard*]

lemmas *bthrs* = *bin-rsplit-simp-alt* [*THEN* [2] *trans*]

lemma *bin-rsplit-size-sign'* [*rule-format*] :
 $n > 0 \implies (\text{ALL } nw \text{ } w. \text{rev } sw = \text{bin-rsplit } n \text{ } (nw, w) \dashrightarrow$

(*ALL* *v*: *set sw*. *bintrunc n v* = *v*)
 ⟨*proof*⟩

lemmas *bin-rsplit-size-sign* = *bin-rsplit-size-sign'* [*OF asm-rl*
rev-rev-ident [*THEN trans*] *set-rev* [*THEN equalityD2* [*THEN subsetD*]],
standard]

lemma *bin-nth-rsplit* [*rule-format*] :
 $n > 0 \implies m < n \implies (\text{ALL } w \ k \ nw. \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \dashrightarrow$
 $k < \text{size } sw \dashrightarrow \text{bin-nth } (sw ! k) \ m = \text{bin-nth } w \ (k * n + m))$
 ⟨*proof*⟩

lemma *bin-rsplit-all*:
 $0 < nw \implies nw \leq n \implies \text{bin-rsplit } n \ (nw, w) = [\text{bintrunc } n \ w]$
 ⟨*proof*⟩

lemma *bin-rsplit-l* [*rule-format*] :
 $\text{ALL } bin. \text{bin-rsplittl } n \ (m, bin) = \text{bin-rsplit } n \ (m, \text{bintrunc } m \ bin)$
 ⟨*proof*⟩

lemma *bin-rsplit-rcat* [*rule-format*] :
 $n > 0 \dashrightarrow \text{bin-rsplit } n \ (n * \text{size } ws, \text{bin-rcat } n \ ws) = \text{map } (\text{bintrunc } n) \ ws$
 ⟨*proof*⟩

lemma *bin-rsplit-aux-len-le* [*rule-format*] :
 $\text{ALL } ws \ m. \ n \neq 0 \dashrightarrow ws = \text{bin-rsplit-aux } (n, bs, nw, w) \dashrightarrow$
 $(\text{length } ws \leq m) = (nw + \text{length } bs * n \leq m * n)$
 ⟨*proof*⟩

lemma *bin-rsplit-len-le*:
 $n \neq 0 \dashrightarrow ws = \text{bin-rsplit } n \ (nw, w) \dashrightarrow (\text{length } ws \leq m) = (nw \leq m * n)$
 ⟨*proof*⟩

lemma *bin-rsplit-aux-len* [*rule-format*] :
 $n \neq 0 \dashrightarrow \text{length } (\text{bin-rsplit-aux } (n, cs, nw, w)) =$
 $(nw + n - 1) \text{ div } n + \text{length } cs$
 ⟨*proof*⟩

lemma *bin-rsplit-len*:
 $n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, w)) = (nw + n - 1) \text{ div } n$
 ⟨*proof*⟩

lemma *bin-rsplit-aux-len-indep* [*rule-format*] :
 $n \neq 0 \implies (\text{ALL } v \ bs. \text{length } bs = \text{length } cs \dashrightarrow$
 $\text{length } (\text{bin-rsplit-aux } (n, bs, nw, v)) =$
 $\text{length } (\text{bin-rsplit-aux } (n, cs, nw, w)))$
 ⟨*proof*⟩

lemma *bin-rsplit-len-indep*:
 $n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, v)) = \text{length } (\text{bin-rsplit } n \ (nw, w))$
 $\langle \text{proof} \rangle$

end

8 TdThs: Type Definition Theorems

theory *TdThs* **imports** *Main* **begin**

9 More lemmas about normal type definitions

lemma
 $tdD1$: *type-definition* $Rep\ Abs\ A \implies \forall x. Rep\ x \in A$ **and**
 $tdD2$: *type-definition* $Rep\ Abs\ A \implies \forall x. Abs\ (Rep\ x) = x$ **and**
 $tdD3$: *type-definition* $Rep\ Abs\ A \implies \forall y. y \in A \longrightarrow Rep\ (Abs\ y) = y$
 $\langle \text{proof} \rangle$

lemma *td-nat-int*:
type-definition $int\ nat\ (Collect\ (op\ \leq\ 0))$
 $\langle \text{proof} \rangle$

context *type-definition*
begin

lemmas $Rep' \ [iff] = Rep \ [simplified]$

declare $Rep\text{-}inverse \ [simp] \ Rep\text{-}inject \ [simp]$

lemma *Abs-eqD*: $Abs\ x = Abs\ y \implies x \in A \implies y \in A \implies x = y$
 $\langle \text{proof} \rangle$

lemma *Abs-inverse'*:
 $r : A \implies Abs\ r = a \implies Rep\ a = r$
 $\langle \text{proof} \rangle$

lemma *Rep-comp-inverse*:
 $Rep\ o\ f = g \implies Abs\ o\ g = f$
 $\langle \text{proof} \rangle$

lemma *Rep-eqD* $[elim!]$: $Rep\ x = Rep\ y \implies x = y$
 $\langle \text{proof} \rangle$

lemma *Rep-inverse'*: $Rep\ a = r \implies Abs\ r = a$
 $\langle \text{proof} \rangle$

lemma *comp-Abs-inverse*:

$$f \circ \text{Abs} = g \implies g \circ \text{Rep} = f$$

<proof>

lemma *set-Rep*:

$$A = \text{range } \text{Rep}$$

<proof>

lemma *set-Rep-Abs*: $A = \text{range } (\text{Rep} \circ \text{Abs})$

<proof>

lemma *Abs-inj-on*: *inj-on* Abs A

<proof>

lemma *image*: $\text{Abs} \text{ ` } A = \text{UNIV}$

<proof>

lemmas *td-thm* = *type-definition-axioms*

lemma *fns1*:

$$\text{Rep} \circ \text{fa} = \text{fr} \circ \text{Rep} \mid \text{fa} \circ \text{Abs} = \text{Abs} \circ \text{fr} \implies \text{Abs} \circ \text{fr} \circ \text{Rep} = \text{fa}$$

<proof>

lemmas *fns1a* = *disjI1* [THEN *fns1*]

lemmas *fns1b* = *disjI2* [THEN *fns1*]

lemma *fns4*:

$$\text{Rep} \circ \text{fa} \circ \text{Abs} = \text{fr} \implies$$

$$\text{Rep} \circ \text{fa} = \text{fr} \circ \text{Rep} \ \& \ \text{fa} \circ \text{Abs} = \text{Abs} \circ \text{fr}$$

<proof>

end

interpretation *nat-int*: *type-definition* [*int nat Collect* (*op* <= 0)]

<proof>

declare *Nat.induct* [*case-names* 0 *Suc*, *induct type*]

declare *Nat.exhaust* [*case-names* 0 *Suc*, *cases type*]

9.1 Extended form of type definition predicate

lemma *td-conds*:

$$\text{norm} \circ \text{norm} = \text{norm} \implies (\text{fr} \circ \text{norm} = \text{norm} \circ \text{fr}) =$$

$$(\text{norm} \circ \text{fr} \circ \text{norm} = \text{fr} \circ \text{norm} \ \& \ \text{norm} \circ \text{fr} \circ \text{norm} = \text{norm} \circ \text{fr})$$

<proof>

lemma *fn-comm-power*:

$$\text{fa} \circ \text{tr} = \text{tr} \circ \text{fr} \implies \text{fa} \circ \text{tr} \circ \text{tr} \circ \text{tr} = \text{tr} \circ \text{fr} \circ \text{tr} \circ \text{tr}$$

<proof>

lemmas *fn-comm-power'* =
ext [*THEN fn-comm-power*, *THEN fun-cong*, *unfolded o-def*, *standard*]

locale *td-ext* = *type-definition* +
fixes *norm*
assumes *eq-norm*: $\bigwedge x. \text{Rep } (\text{Abs } x) = \text{norm } x$
begin

lemma *Abs-norm* [*simp*]:
 $\text{Abs } (\text{norm } x) = \text{Abs } x$
 $\langle \text{proof} \rangle$

lemma *td-th*:
 $g \circ \text{Abs} = f \implies f (\text{Rep } x) = g x$
 $\langle \text{proof} \rangle$

lemma *eq-norm'*: $\text{Rep } o \text{ Abs} = \text{norm}$
 $\langle \text{proof} \rangle$

lemma *norm-Rep* [*simp*]: $\text{norm } (\text{Rep } x) = \text{Rep } x$
 $\langle \text{proof} \rangle$

lemmas *td* = *td-thm*

lemma *set-iff-norm*: $w : A \longleftrightarrow w = \text{norm } w$
 $\langle \text{proof} \rangle$

lemma *inverse-norm*:
 $(\text{Abs } n = w) = (\text{Rep } w = \text{norm } n)$
 $\langle \text{proof} \rangle$

lemma *norm-eq-iff*:
 $(\text{norm } x = \text{norm } y) = (\text{Abs } x = \text{Abs } y)$
 $\langle \text{proof} \rangle$

lemma *norm-comps*:
 $\text{Abs } o \text{ norm} = \text{Abs}$
 $\text{norm } o \text{ Rep} = \text{Rep}$
 $\text{norm } o \text{ norm} = \text{norm}$
 $\langle \text{proof} \rangle$

lemmas *norm-norm* [*simp*] = *norm-comps*

lemma *fns5*:
 $\text{Rep } o \text{ fa } o \text{ Abs} = \text{fr} \implies$
 $\text{fr } o \text{ norm} = \text{fr} \ \& \ \text{norm } o \text{ fr} = \text{fr}$
 $\langle \text{proof} \rangle$

lemma *fns2*:
 $Abs \circ fr \circ Rep = fa ==>$
 $(norm \circ fr \circ norm = fr \circ norm) = (Rep \circ fa = fr \circ Rep)$
 $\langle proof \rangle$

lemma *fns3*:
 $Abs \circ fr \circ Rep = fa ==>$
 $(norm \circ fr \circ norm = norm \circ fr) = (fa \circ Abs = Abs \circ fr)$
 $\langle proof \rangle$

lemma *fns*:
 $fr \circ norm = norm \circ fr ==>$
 $(fa \circ Abs = Abs \circ fr) = (Rep \circ fa = fr \circ Rep)$
 $\langle proof \rangle$

lemma *range-norm*:
 $range (Rep \circ Abs) = A$
 $\langle proof \rangle$

end

lemmas *td-ext-def'* =
 $td-ext-def [unfolded\ type-definition-def\ td-ext-axioms-def]$

end

10 WordDefinition: Definition of Word Type

theory *WordDefinition* **imports** *Size BinBoolList TdThs* **begin**

typedef (**open** *word*) '*a word*'
 $= \{(0::int) ..< 2^{len-of\ TYPE('a::len0)}\} \langle proof \rangle$

instance *word* :: (*len0*) *number* $\langle proof \rangle$
instance *word* :: (*type*) *minus* $\langle proof \rangle$
instance *word* :: (*type*) *plus* $\langle proof \rangle$
instance *word* :: (*type*) *one* $\langle proof \rangle$
instance *word* :: (*type*) *zero* $\langle proof \rangle$
instance *word* :: (*type*) *times* $\langle proof \rangle$
instance *word* :: (*type*) *Divides.div* $\langle proof \rangle$
instance *word* :: (*type*) *power* $\langle proof \rangle$
instance *word* :: (*type*) *ord* $\langle proof \rangle$
instance *word* :: (*type*) *size* $\langle proof \rangle$
instance *word* :: (*type*) *inverse* $\langle proof \rangle$
instance *word* :: (*type*) *bit* $\langle proof \rangle$

10.1 Type conversions and casting

constdefs

— representation of words using unsigned or signed bins, only difference in these is the type class

```
word-of-int :: int => 'a :: len0 word
word-of-int w == Abs-word (bintrunc (len-of TYPE ('a)) w)
```

— uint and sint cast a word to an integer, uint treats the word as unsigned, sint treats the most-significant-bit as a sign bit

```
uint :: 'a :: len0 word => int
uint w == Rep-word w
sint :: 'a :: len word => int
sint-uint: sint w == sbintrunc (len-of TYPE ('a) - 1) (uint w)
unat :: 'a :: len0 word => nat
unat w == nat (uint w)
```

— the sets of integers representing the words

```
uints :: nat => int set
uints n == range (bintrunc n)
sints :: nat => int set
sints n == range (sbintrunc (n - 1))
unats :: nat => nat set
unats n == {i. i < 2 ^ n}
norm-sint :: nat => int => int
norm-sint n w == (w + 2 ^ (n - 1)) mod 2 ^ n - 2 ^ (n - 1)
```

— cast a word to a different length

```
scast :: 'a :: len word => 'b :: len word
scast w == word-of-int (sint w)
ucast :: 'a :: len0 word => 'b :: len0 word
ucast w == word-of-int (uint w)
```

— whether a cast (or other) function is to a longer or shorter length

```
source-size :: ('a :: len0 word => 'b) => nat
source-size c == let arb = arbitrary ; x = c arb in size arb
target-size :: ('a => 'b :: len0 word) => nat
target-size c == size (c arbitrary)
is-up :: ('a :: len0 word => 'b :: len0 word) => bool
is-up c == source-size c <= target-size c
is-down :: ('a :: len0 word => 'b :: len0 word) => bool
is-down c == target-size c <= source-size c
```

constdefs

```
of-bl :: bool list => 'a :: len0 word
of-bl bl == word-of-int (bl-to-bin bl)
to-bl :: 'a :: len0 word => bool list
to-bl w ==
bin-to-bl (len-of TYPE ('a)) (uint w)
```

```
word-reverse :: 'a :: len0 word => 'a word
word-reverse w == of-bl (rev (to-bl w))
```

defs (overloaded)

```
word-size: size (w :: 'a :: len0 word) == len-of TYPE('a)
word-number-of-def: number-of w == word-of-int w
```

constdefs

```
word-int-case :: (int => 'b) => ('a :: len0 word) => 'b
word-int-case f w == f (uint w)
```

syntax

```
of-int :: int => 'a
```

translations

```
case x of of-int y => b == word-int-case (%y. b) x
```

10.2 Arithmetic operations**defs (overloaded)**

```
word-1-wi: (1 :: ('a :: len0) word) == word-of-int 1
word-0-wi: (0 :: ('a :: len0) word) == word-of-int 0
```

```
word-le-def: a <= b == uint a <= uint b
word-less-def: x < y == x <= y & x ~ = (y :: 'a :: len0 word)
```

constdefs

```
word-succ :: 'a :: len0 word => 'a word
word-succ a == word-of-int (Numeral.succ (uint a))
```

```
word-pred :: 'a :: len0 word => 'a word
word-pred a == word-of-int (Numeral.pred (uint a))
```

```
udvd :: 'a::len word => 'a::len word => bool (infixl udvd 50)
a udvd b == EX n>=0. uint b = n * uint a
```

```
word-sle :: 'a :: len word => 'a word => bool ((-/ <=s -) [50, 51] 50)
a <=s b == sint a <= sint b
```

```
word-sless :: 'a :: len word => 'a word => bool ((-/ <s -) [50, 51] 50)
(x <s y) == (x <=s y & x ~ = y)
```

consts

```
word-power :: 'a :: len0 word => nat => 'a word
```

primrec

```
word-power a 0 = 1
word-power a (Suc n) = a * word-power a n
```

defs (overloaded)

```
word-pow: power == word-power
```

```

word-add-def: a + b == word-of-int (uint a + uint b)
word-sub-wi: a - b == word-of-int (uint a - uint b)
word-minus-def: - a == word-of-int (- uint a)
word-mult-def: a * b == word-of-int (uint a * uint b)
word-div-def: a div b == word-of-int (uint a div uint b)
word-mod-def: a mod b == word-of-int (uint a mod uint b)

```

10.3 Bit-wise operations

defs (overloaded)

```

word-and-def:
(a::'a::len0 word) AND b == word-of-int (uint a AND uint b)

```

```

word-or-def:
(a::'a::len0 word) OR b == word-of-int (uint a OR uint b)

```

```

word-xor-def:
(a::'a::len0 word) XOR b == word-of-int (uint a XOR uint b)

```

```

word-not-def:
NOT (a::'a::len0 word) == word-of-int (NOT (uint a))

```

```

word-test-bit-def:
test-bit (a::'a::len0 word) == bin-nth (uint a)

```

```

word-set-bit-def:
set-bit (a::'a::len0 word) n x ==
word-of-int (bin-sc n (If x bit.B1 bit.B0) (uint a))

```

```

word-set-bits-def:
(BITS n. f n)::'a::len0 word == of-bl (bl-of-nth (len-of TYPE ('a)) f)

```

```

word-lsb-def:
lsb (a::'a::len0 word) == bin-last (uint a) = bit.B1

```

```

word-msb-def:
msb (a::'a::len word) == bin-sign (sint a) = Numeral.Min

```

constdefs

```

setBit :: 'a :: len0 word => nat => 'a word
setBit w n == set-bit w n True

```

```

clearBit :: 'a :: len0 word => nat => 'a word
clearBit w n == set-bit w n False

```

10.4 Shift operations

constdefs

```

shifl1 :: 'a :: len0 word => 'a word

```

shiftr1 w == word-of-int (uint w BIT bit.B0)

— shift right as unsigned or as signed, ie logical or arithmetic

shiftr1 :: 'a :: len0 word => 'a word

shiftr1 w == word-of-int (bin-rest (uint w))

sshiftr1 :: 'a :: len word => 'a word

sshiftr1 w == word-of-int (bin-rest (sint w))

bshiftr1 :: bool => 'a :: len word => 'a word

bshiftr1 b w == of-bl (b # butlast (to-bl w))

sshiftr :: 'a :: len word => nat => 'a word (infixl >>> 55)

w >>> n == (sshiftr1 ^ n) w

mask :: nat => 'a::len word

mask n == (1 << n) - 1

revcast :: 'a :: len0 word => 'b :: len0 word

revcast w == of-bl (takefill False (len-of TYPE('b)) (to-bl w))

slice1 :: nat => 'a :: len0 word => 'b :: len0 word

slice1 n w == of-bl (takefill False n (to-bl w))

slice :: nat => 'a :: len0 word => 'b :: len0 word

slice n w == slice1 (size w - n) w

defs (overloaded)

shiftr-def: (w::'a::len0 word) << n == (shiftr1 ^ n) w

shiftr-def: (w::'a::len0 word) >> n == (shiftr1 ^ n) w

10.5 Rotation

constdefs

rotater1 :: 'a list => 'a list

rotater1 ys ==

case ys of [] => [] | x # xs => last xs # butlast xs

rotater :: nat => 'a list => 'a list

rotater n == rotater1 ^ n

word-rotr :: nat => 'a :: len0 word => 'a :: len0 word

word-rotr n w == of-bl (rotater n (to-bl w))

word-rotl :: nat => 'a :: len0 word => 'a :: len0 word

word-rotl n w == of-bl (rotate n (to-bl w))

word-roti :: int => 'a :: len0 word => 'a :: len0 word

word-roti i $w ==$ if $i \geq 0$ then *word-rotr* (nat i) w
 else *word-rotl* (nat $(- i)$) w

10.6 Split and cat operations

constdefs

word-cat $:: 'a :: \text{len0 word} \Rightarrow 'b :: \text{len0 word} \Rightarrow 'c :: \text{len0 word}$
word-cat a $b ==$ *word-of-int* (bin-cat (uint a) (len-of TYPE (' b)) (uint b))

word-split $:: 'a :: \text{len0 word} \Rightarrow ('b :: \text{len0 word}) * ('c :: \text{len0 word})$
word-split $a ==$
 case bin-split (len-of TYPE (' c)) (uint a) of
 (u, v) \Rightarrow (*word-of-int* u , *word-of-int* v)

word-rcat $:: 'a :: \text{len0 word list} \Rightarrow 'b :: \text{len0 word}$
word-rcat $ws ==$
word-of-int (bin-rcat (len-of TYPE (' a)) (map uint ws))

word-rsplit $:: 'a :: \text{len0 word} \Rightarrow 'b :: \text{len word list}$
word-rsplit $w ==$
 map *word-of-int* (bin-rsplit (len-of TYPE (' b)) (len-of TYPE (' a), uint w))

constdefs

— Largest representable machine integer.
max-word $:: 'a::\text{len word}$
max-word \equiv *word-of-int* ($2^{\text{len-of TYPE}('a)} - 1$)

consts

of-bool $:: \text{bool} \Rightarrow 'a::\text{len word}$

primrec

of-bool False = 0
of-bool True = 1

lemmas *of-nth-def* = *word-set-bits-def*

lemmas *word-size-gt-0* [iff] =

xtr1 [OF *word-size* [THEN *meta-eq-to-obj-eq*] *len-gt-0*, *standard*]

lemmas *lens-gt-0* = *word-size-gt-0* *len-gt-0*

lemmas *lens-not-0* [iff] = *lens-gt-0* [THEN *gr-implies-not0*, *standard*]

lemma *uints-num*: $\text{uints } n = \{i. 0 \leq i \wedge i < 2^n\}$
 <proof>

lemma *sints-num*: $\text{sints } n = \{i. -(2^n - 1) \leq i \wedge i < 2^n - 1\}$
 <proof>

lemmas *atLeastLessThan-alt* = *atLeastLessThan-def* [unfolded]

atLeast-def lessThan-def Collect-conj-eq [symmetric]]

lemma *mod-in-reps*: $m > 0 \implies y \bmod m : \{0::\text{int} \ ..< m\}$
 ⟨proof⟩

lemma
Rep-word-0:0 \leq *Rep-word* x **and**
Rep-word-lt: *Rep-word* $(x::'a::\text{len0 word}) < 2 \wedge \text{len-of TYPE}('a)$
 ⟨proof⟩

lemma *Rep-word-mod-same*:
Rep-word $x \bmod 2 \wedge \text{len-of TYPE}('a) = \text{Rep-word } (x::'a::\text{len0 word})$
 ⟨proof⟩

lemma *td-ext-uint*:
td-ext $(\text{uint}::'a::\text{len0 word} \Rightarrow \text{int}) \text{ word-of-int } (\text{uints } (\text{len-of TYPE}('a::\text{len0})))$
 $(\%w::\text{int}. w \bmod 2 \wedge \text{len-of TYPE}('a))$
 ⟨proof⟩

lemmas *int-word-uint* = *td-ext-uint* [THEN *td-ext.eq-norm*, *standard*]

interpretation *word-uint*:
td-ext $[\text{uint}::'a::\text{len0 word} \Rightarrow \text{int}]$
word-of-int
uints $(\text{len-of TYPE}('a::\text{len0}))$
 $\lambda w. w \bmod 2 \wedge \text{len-of TYPE}('a::\text{len0})]$
 ⟨proof⟩

lemmas *td-uint* = *word-uint.td-thm*

lemmas *td-ext-ubin* = *td-ext-uint*
 [simplified *len-gt-0 no-bintr-alt1* [symmetric]]

interpretation *word-ubin*:
td-ext $[\text{uint}::'a::\text{len0 word} \Rightarrow \text{int}]$
word-of-int
uints $(\text{len-of TYPE}('a::\text{len0}))$
bintrunc $(\text{len-of TYPE}('a::\text{len0}))]$
 ⟨proof⟩

lemma *sint-sbintrunc'*:
sint $(\text{word-of-int bin}::'a \text{ word}) =$
 $(\text{sbintrunc } (\text{len-of TYPE}('a::\text{len}) - 1) \text{ bin})$
 ⟨proof⟩

lemma *uint-sint*:
uint $w = \text{bintrunc } (\text{len-of TYPE}('a)) (\text{sint } (w::'a::\text{len word}))$
 ⟨proof⟩

lemma *bintr-uint'*:

$n \geq \text{size } w \implies \text{bintrunc } n \text{ (uint } w) = \text{uint } w$
 ⟨proof⟩

lemma *wi-bintr'*:

$wb = \text{word-of-int bin} \implies n \geq \text{size } wb \implies$
 $\text{word-of-int (bintrunc } n \text{ bin)} = wb$
 ⟨proof⟩

lemmas *bintr-uint = bintr-uint'* [unfolded word-size]

lemmas *wi-bintr = wi-bintr'* [unfolded word-size]

lemma *td-ext-sbin*:

$\text{td-ext (sint :: 'a word} \Rightarrow \text{int) word-of-int (sints (len-of TYPE('a::len)))}$
 $(\text{sbintrunc (len-of TYPE('a) - 1)})$
 ⟨proof⟩

lemmas *td-ext-sint = td-ext-sbin*

[simplified len-gt-0 no-sbintr-alt2 Suc-pred' [symmetric]]

interpretation *word-sint*:

$\text{td-ext [sint :: 'a::len word} \Rightarrow \text{int}$
 word-of-int
 $\text{sints (len-of TYPE('a::len))}$
 $\%w. (w + 2^{(\text{len-of TYPE('a::len) - 1})}) \bmod 2^{\text{len-of TYPE('a::len) - 1}}$
 $2^{(\text{len-of TYPE('a::len) - 1})}]$
 ⟨proof⟩

interpretation *word-sbin*:

$\text{td-ext [sint :: 'a::len word} \Rightarrow \text{int}$
 word-of-int
 $\text{sints (len-of TYPE('a::len))}$
 $\text{sbintrunc (len-of TYPE('a::len) - 1)]}$
 ⟨proof⟩

lemmas *int-word-sint = td-ext-sint* [THEN td-ext.eq-norm, standard]

lemmas *td-sint = word-sint.td*

lemma *word-number-of-alt*: $\text{number-of } b == \text{word-of-int (number-of } b)$

⟨proof⟩

lemma *word-no-wi*: $\text{number-of} = \text{word-of-int}$

⟨proof⟩

lemma *to-bl-def'*:

$(\text{to-bl :: 'a :: len0 word} \Rightarrow \text{bool list}) =$
 $\text{bin-to-bl (len-of TYPE('a)) o uint}$

<proof>

lemmas *word-reverse-no-def* [*simp*] = *word-reverse-def* [*of number-of w, standard*]

lemmas *uints-mod* = *uints-def* [*unfolded no-bintr-alt1*]

lemma *uint-bintrunc*: *uint* (*number-of bin* :: 'a word) =
number-of (*bintrunc* (*len-of TYPE* ('a :: len0)) *bin*)
<proof>

lemma *sint-sbintrunc*: *sint* (*number-of bin* :: 'a word) =
number-of (*sbintrunc* (*len-of TYPE* ('a :: len) - 1) *bin*)
<proof>

lemma *unat-bintrunc*:
unat (*number-of bin* :: 'a :: len0 word) =
number-of (*bintrunc* (*len-of TYPE* ('a)) *bin*)
<proof>

declare

uint-bintrunc [*simp*]
sint-sbintrunc [*simp*]
unat-bintrunc [*simp*]

lemma *size-0-eq*: *size* (*w* :: 'a :: len0 word) = 0 ==> *v* = *w*
<proof>

lemmas *uint-lem* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]

lemmas *sint-lem* = *word-sint.Rep* [*unfolded sints-num mem-Collect-eq*]

lemmas *uint-ge-0* [*iff*] = *uint-lem* [*THEN conjunct1, standard*]

lemmas *uint-lt2p* [*iff*] = *uint-lem* [*THEN conjunct2, standard*]

lemmas *sint-ge* = *sint-lem* [*THEN conjunct1, standard*]

lemmas *sint-lt* = *sint-lem* [*THEN conjunct2, standard*]

lemma *sign-uint-Pls* [*simp*]:
bin-sign (*uint x*) = *Numeral.Pl*
<proof>

lemmas *uint-m2p-neg* = *iffD2* [*OF diff-less-0-iff-less uint-lt2p, standard*]

lemmas *uint-m2p-not-non-neg* =
iffD2 [*OF linorder-not-le uint-m2p-neg, standard*]

lemma *lt2p-lem*:
len-of TYPE ('a) <= *n* ==> *uint* (*w* :: 'a :: len0 word) < 2 ^ *n*
<proof>

lemmas *uint-le-0-iff* [*simp*] =
uint-ge-0 [*THEN leD, THEN linorder-antisym-conv1, standard*]

lemma *uint-nat*: $\text{uint } w == \text{int } (\text{unat } w)$
 ⟨proof⟩

lemma *uint-number-of*:
 $\text{uint } (\text{number-of } b :: 'a :: \text{len0 word}) = \text{number-of } b \bmod 2^{\text{len-of TYPE } ('a)}$
 ⟨proof⟩

lemma *unat-number-of*:
 $\text{bin-sign } b = \text{Numeral.Pls} ==>$
 $\text{unat } (\text{number-of } b :: 'a :: \text{len0 word}) = \text{number-of } b \bmod 2^{\text{len-of TYPE } ('a)}$
 ⟨proof⟩

lemma *sint-number-of*: $\text{sint } (\text{number-of } b :: 'a :: \text{len word}) = (\text{number-of } b + 2^{\text{len-of TYPE } ('a) - 1}) \bmod 2^{\text{len-of TYPE } ('a)}$
 ⟨proof⟩

lemma *word-of-int-bin* [simp] :
 $(\text{word-of-int } (\text{number-of bin}) :: 'a :: \text{len0 word}) = (\text{number-of bin})$
 ⟨proof⟩

lemma *word-int-case-wi*:
 $\text{word-int-case } f (\text{word-of-int } i :: 'b \text{ word}) = f (i \bmod 2^{\text{len-of TYPE } ('b :: \text{len0})})$
 ⟨proof⟩

lemma *word-int-split*:
 $P (\text{word-int-case } f x) =$
 $(\text{ALL } i. x = (\text{word-of-int } i :: 'b :: \text{len0 word}) \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE } ('b)}) \longrightarrow P (f i))$
 ⟨proof⟩

lemma *word-int-split-asm*:
 $P (\text{word-int-case } f x) =$
 $(\sim (\text{EX } n. x = (\text{word-of-int } n :: 'b :: \text{len0 word}) \ \& \ 0 \leq n \ \& \ n < 2^{\text{len-of TYPE } ('b :: \text{len0})}) \ \& \ \sim P (f n)))$
 ⟨proof⟩

lemmas *uint-range'* =
 $\text{word-uint.Rep } [\text{unfolded uints-num mem-Collect-eq, standard}]$
lemmas *sint-range'* = $\text{word-sint.Rep } [\text{unfolded One-nat-def sints-num mem-Collect-eq, standard}]$

lemma *uint-range-size*: $0 \leq \text{uint } w \ \& \ \text{uint } w < 2^{\text{size } w}$
 ⟨proof⟩

lemma *sint-range-size*:
 $-(2^{\text{size } w - \text{Suc } 0}) \leq \text{sint } w \ \& \ \text{sint } w < 2^{\text{size } w - \text{Suc } 0}$

$\langle proof \rangle$

lemmas *sint-above-size* = *sint-range-size*
 [THEN *conjunct2*, THEN [2] *xtr8*, folded *One-nat-def*, *standard*]

lemmas *sint-below-size* = *sint-range-size*
 [THEN *conjunct1*, THEN [2] *order-trans*, folded *One-nat-def*, *standard*]

lemma *test-bit-eq-iff*: (*test-bit* (*u*::'*a*::len0 *word*) = *test-bit* *v*) = (*u* = *v*)
 $\langle proof \rangle$

lemma *test-bit-size* [rule-format] : (*w*::'*a*::len0 *word*) !! *n* \longrightarrow *n* < size *w*
 $\langle proof \rangle$

lemma *word-eqI* [rule-format] :
 fixes *u* :: '*a*::len0 *word*
 shows (ALL *n*. *n* < size *u* \longrightarrow *u* !! *n* = *v* !! *n*) \implies *u* = *v*
 $\langle proof \rangle$

lemmas *word-eqD* = *test-bit-eq-iff* [THEN *iffD2*, THEN *fun-cong*, *standard*]

lemma *test-bit-bin'*: *w* !! *n* = (*n* < size *w* & *bin-nth* (*uint w*) *n*)
 $\langle proof \rangle$

lemmas *test-bit-bin* = *test-bit-bin'* [unfolded *word-size*]

lemma *bin-nth-uint-imp'*: *bin-nth* (*uint w*) *n* \longrightarrow *n* < size *w*
 $\langle proof \rangle$

lemma *bin-nth-sint'*:
n >= size *w* \longrightarrow *bin-nth* (*sint w*) *n* = *bin-nth* (*sint w*) (size *w* - 1)
 $\langle proof \rangle$

lemmas *bin-nth-uint-imp* = *bin-nth-uint-imp'* [rule-format, unfolded *word-size*]

lemmas *bin-nth-sint* = *bin-nth-sint'* [rule-format, unfolded *word-size*]

lemma *td-bl*:
type-definition (*to-bl* :: '*a*::len0 *word* \Rightarrow *bool list*)
 of-bl
 {*bl*. length *bl* = len-of TYPE('a)}
 $\langle proof \rangle$

interpretation *word-bl*:
type-definition [*to-bl* :: '*a*::len0 *word* \Rightarrow *bool list*
 of-bl
 {*bl*. length *bl* = len-of TYPE('a::len0)}]
 $\langle proof \rangle$

lemma *word-size-bl*: $\text{size } w == \text{size } (\text{to-bl } w)$
 ⟨proof⟩

lemma *to-bl-use-of-bl*:
 $(\text{to-bl } w = \text{bl}) = (w = \text{of-bl } \text{bl} \wedge \text{length } \text{bl} = \text{length } (\text{to-bl } w))$
 ⟨proof⟩

lemma *to-bl-word-rev*: $\text{to-bl } (\text{word-reverse } w) = \text{rev } (\text{to-bl } w)$
 ⟨proof⟩

lemma *word-rev-rev* [simp] : $\text{word-reverse } (\text{word-reverse } w) = w$
 ⟨proof⟩

lemma *word-rev-gal*: $\text{word-reverse } w = u ==> \text{word-reverse } u = w$
 ⟨proof⟩

lemmas *word-rev-gal'* = sym [THEN *word-rev-gal*, symmetric, standard]

lemmas *length-bl-gt-0* [iff] = xtr1 [OF *word-bl.Rep'* len-gt-0, standard]

lemmas *bl-not-Nil* [iff] =
 $\text{length-bl-gt-0 } [THEN \text{length-greater-0-conv } [THEN \text{iffD1}], \text{standard}]$

lemmas *length-bl-neq-0* [iff] = *length-bl-gt-0* [THEN *gr-implies-not0*]

lemma *hd-bl-sign-sint*: $\text{hd } (\text{to-bl } w) = (\text{bin-sign } (\text{sint } w) = \text{Numeral.Min})$
 ⟨proof⟩

lemma *of-bl-drop'*:
 $\text{lend} = \text{length } \text{bl} - \text{len-of TYPE } ('a :: \text{len0}) ==>$
 $\text{of-bl } (\text{drop } \text{lend } \text{bl}) = (\text{of-bl } \text{bl} :: 'a \text{ word})$
 ⟨proof⟩

lemmas *of-bl-no* = *of-bl-def* [folded *word-number-of-def*]

lemma *test-bit-of-bl*:
 $(\text{of-bl } \text{bl} :: 'a :: \text{len0} \text{ word}) !! n = (\text{rev } \text{bl} ! n \wedge n < \text{len-of TYPE } ('a) \wedge n < \text{length } \text{bl})$
 ⟨proof⟩

lemma *no-of-bl*:
 $(\text{number-of bin} :: 'a :: \text{len0} \text{ word}) = \text{of-bl } (\text{bin-to-bl } (\text{len-of TYPE } ('a)) \text{ bin})$
 ⟨proof⟩

lemma *uint-bl*: $\text{to-bl } w == \text{bin-to-bl } (\text{size } w) (\text{uint } w)$
 ⟨proof⟩

lemma *to-bl-bin*: $\text{bl-to-bin } (\text{to-bl } w) = \text{uint } w$
 ⟨proof⟩

lemma *to-bl-of-bin*:

to-bl (*word-of-int* *bin*::'*a*::*len0* *word*) = *bin-to-bl* (*len-of TYPE*('a)) *bin*
 ⟨*proof*⟩

lemmas *to-bl-no-bin* [*simp*] = *to-bl-of-bin* [*folded word-number-of-def*]

lemma *to-bl-to-bin* [*simp*] : *bl-to-bin* (*to-bl* *w*) = *uint* *w*
 ⟨*proof*⟩

lemmas *uint-bl-bin* [*simp*] = *trans* [*OF bin-bl-bin word-ubin.norm-Rep, standard*]

lemmas *num-AB-u* [*simp*] = *word-uint.Rep-inverse*
 [*unfolded o-def word-number-of-def [symmetric], standard*]

lemmas *num-AB-s* [*simp*] = *word-sint.Rep-inverse*
 [*unfolded o-def word-number-of-def [symmetric], standard*]

lemma *uints-unats*: *uints* *n* = *int* ‘ *unats* *n*
 ⟨*proof*⟩

lemma *unats-uints*: *unats* *n* = *nat* ‘ *uints* *n*
 ⟨*proof*⟩

lemmas *bintr-num* = *word-ubin.norm-eq-iff*
 [*symmetric, folded word-number-of-def, standard*]

lemmas *sbintr-num* = *word-sbin.norm-eq-iff*
 [*symmetric, folded word-number-of-def, standard*]

lemmas *num-of-bintr* = *word-ubin.Abs-norm* [*folded word-number-of-def, standard*]

lemmas *num-of-sbintr* = *word-sbin.Abs-norm* [*folded word-number-of-def, standard*]

lemma *num-of-bintr'*:
bintrunc (*len-of TYPE*('a :: *len0*)) *a* = *b* ==>
number-of *a* = (*number-of* *b* :: 'a *word*)
 ⟨*proof*⟩

lemma *num-of-sbintr'*:
sbintrunc (*len-of TYPE*('a :: *len*) - 1) *a* = *b* ==>
number-of *a* = (*number-of* *b* :: 'a *word*)
 ⟨*proof*⟩

lemmas *num-abs-bintr* = *sym* [*THEN trans*,
OF num-of-bintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]

lemmas *num-abs-sbintr* = *sym* [*THEN trans*,
OF num-of-sbintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]

lemma *ucast-id*: *ucast w = w*
 ⟨*proof*⟩

lemma *scast-id*: *scast w = w*
 ⟨*proof*⟩

lemma *ucast-bl*: *ucast w == of-bl (to-bl w)*
 ⟨*proof*⟩

lemma *nth-ucast*:
 (*ucast w :: 'a :: len0 word*) !! *n* = (*w* !! *n* & *n* < *len-of TYPE('a)*)
 ⟨*proof*⟩

lemma *ucast-bintr* [*simp*]:
ucast (number-of w :: 'a :: len0 word) =
number-of (bintrunc (len-of TYPE('a)) w)
 ⟨*proof*⟩

lemma *scast-sbintr* [*simp*]:
scast (number-of w :: 'a :: len word) =
number-of (sbintrunc (len-of TYPE('a) - Suc 0) w)
 ⟨*proof*⟩

lemmas *source-size = source-size-def* [*unfolded Let-def word-size*]
lemmas *target-size = target-size-def* [*unfolded Let-def word-size*]
lemmas *is-down = is-down-def* [*unfolded source-size target-size*]
lemmas *is-up = is-up-def* [*unfolded source-size target-size*]

lemmas *is-up-down =*
trans [OF is-up [THEN meta-eq-to-obj-eq]
is-down [THEN meta-eq-to-obj-eq, symmetric],
standard]

lemma *down-cast-same'*: *uc = ucast ==> is-down uc ==> uc = scast*
 ⟨*proof*⟩

lemma *word-rev-tf'*:
r = to-bl (of-bl bl) ==> r = rev (takefill False (length r) (rev bl))
 ⟨*proof*⟩

lemmas *word-rev-tf = refl [THEN word-rev-tf', unfolded word-bl.Rep', standard]*

lemmas *word-rep-drop = word-rev-tf* [*simplified takefill-alt,*
simplified, simplified rev-take, simplified]

lemma *to-bl-ucast*:

to-bl (*ucast* (*w*::'b::len0 *word*) ::'a::len0 *word*) =
replicate (*len-of TYPE*('a') – *len-of TYPE*('b')) *False* @
drop (*len-of TYPE*('b') – *len-of TYPE*('a')) (*to-bl w*)
 ⟨*proof*⟩

lemma *ucast-up-app'*:

uc = *ucast* ==> *source-size uc* + *n* = *target-size uc* ==>
to-bl (*uc w*) = *replicate n False* @ (*to-bl w*)
 ⟨*proof*⟩

lemma *ucast-down-drop'*:

uc = *ucast* ==> *source-size uc* = *target-size uc* + *n* ==>
to-bl (*uc w*) = *drop n* (*to-bl w*)
 ⟨*proof*⟩

lemma *scast-down-drop'*:

sc = *scast* ==> *source-size sc* = *target-size sc* + *n* ==>
to-bl (*sc w*) = *drop n* (*to-bl w*)
 ⟨*proof*⟩

lemma *sint-up-scast'*:

sc = *scast* ==> *is-up sc* ==> *sint* (*sc w*) = *sint w*
 ⟨*proof*⟩

lemma *uint-up-ucast'*:

uc = *ucast* ==> *is-up uc* ==> *uint* (*uc w*) = *uint w*
 ⟨*proof*⟩

lemmas *down-cast-same* = *refl* [*THEN down-cast-same*']

lemmas *ucast-up-app* = *refl* [*THEN ucast-up-app*']

lemmas *ucast-down-drop* = *refl* [*THEN ucast-down-drop*']

lemmas *scast-down-drop* = *refl* [*THEN scast-down-drop*']

lemmas *uint-up-ucast* = *refl* [*THEN uint-up-ucast*']

lemmas *sint-up-scast* = *refl* [*THEN sint-up-scast*']

lemma *ucast-up-ucast'*: *uc* = *ucast* ==> *is-up uc* ==> *ucast* (*uc w*) = *ucast w*
 ⟨*proof*⟩

lemma *scast-up-scast'*: *sc* = *scast* ==> *is-up sc* ==> *scast* (*sc w*) = *scast w*
 ⟨*proof*⟩

lemma *ucast-of-bl-up'*:

w = *of-bl bl* ==> *size bl* <= *size w* ==> *ucast w* = *of-bl bl*
 ⟨*proof*⟩

lemmas *ucast-up-ucast* = *refl* [*THEN ucast-up-ucast*']

lemmas *scast-up-scast* = *refl* [*THEN scast-up-scast*']

lemmas *ucast-of-bl-up* = *refl* [*THEN ucast-of-bl-up*']

lemmas *ucast-up-ucast-id* = *trans* [*OF ucast-up-ucast ucast-id*]

lemmas *scast-up-scast-id* = *trans* [*OF scast-up-scast scast-id*]

lemmas *isduu* = *is-up-down* [**where** *c* = *ucast*, *THEN iffD2*]

lemmas *isdus* = *is-up-down* [**where** *c* = *scast*, *THEN iffD2*]

lemmas *ucast-down-ucast-id* = *isduu* [*THEN ucast-up-ucast-id*]

lemmas *scast-down-scast-id* = *isdus* [*THEN ucast-up-ucast-id*]

lemma *up-ucast-surj*:

is-up (*ucast* :: 'b::len0 word => 'a::len0 word) ==>

surj (*ucast* :: 'a word => 'b word)

<proof>

lemma *up-scast-surj*:

is-up (*scast* :: 'b::len word => 'a::len word) ==>

surj (*scast* :: 'a word => 'b word)

<proof>

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word => 'a::len word) ==>

inj-on (*ucast* :: 'a word => 'b word) *A*

<proof>

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word => 'a::len0 word) ==>

inj-on (*ucast* :: 'a word => 'b word) *A*

<proof>

lemma *of-bl-append-same*: *of-bl* (*X @ to-bl w*) = *w*

<proof>

lemma *ucast-down-no'*:

uc = *ucast* ==> *is-down uc* ==> *uc (number-of bin) = number-of bin*

<proof>

lemmas *ucast-down-no* = *ucast-down-no'* [*OF refl*]

lemma *ucast-down-bl'*: *uc* = *ucast* ==> *is-down uc* ==> *uc (of-bl bl) = of-bl bl*

<proof>

lemmas *ucast-down-bl* = *ucast-down-bl'* [*OF refl*]

lemmas *slice-def'* = *slice-def* [*unfolded word-size*]

lemmas *test-bit-def'* = *word-test-bit-def* [*THEN meta-eq-to-obj-eq, THEN fun-cong*]

lemmas *word-log-defs* = *word-and-def word-or-def word-xor-def word-not-def*

lemmas *word-log-bin-defs* = *word-log-defs*

end

11 WordArith: Word Arithmetic

theory *WordArith* **imports** *WordDefinition* **begin**

lemma *word-less-alt*: $(a < b) = (\text{uint } a < \text{uint } b)$
 $\langle \text{proof} \rangle$

lemma *signed-linorder*: *linorder* *word-sle* *word-sless*
 $\langle \text{proof} \rangle$

interpretation *signed*: *linorder* [*word-sle* *word-sless*]
 $\langle \text{proof} \rangle$

lemmas *word-arith-wis* [*THEN meta-eq-to-obj-eq*] =
word-add-def *word-mult-def* *word-minus-def*
word-succ-def *word-pred-def* *word-0-wi* *word-1-wi*

lemma *udvdI*:
 $0 \leq n \implies \text{uint } b = n * \text{uint } a \implies a \text{ udvd } b$
 $\langle \text{proof} \rangle$

lemmas *word-div-no* [*simp*] =
word-div-def [*of number-of a number-of b, standard*]

lemmas *word-mod-no* [*simp*] =
word-mod-def [*of number-of a number-of b, standard*]

lemmas *word-less-no* [*simp*] =
word-less-def [*of number-of a number-of b, standard*]

lemmas *word-le-no* [*simp*] =
word-le-def [*of number-of a number-of b, standard*]

lemmas *word-sless-no* [*simp*] =
word-sless-def [*of number-of a number-of b, standard*]

lemmas *word-sle-no* [*simp*] =
word-sle-def [*of number-of a number-of b, standard*]

lemmas *word-0-wi-Pls* = *word-0-wi* [*folded Pls-def*]

lemmas *word-0-no* = *word-0-wi-Pls* [*folded word-no-wi*]

lemma *int-one-bin*: $(1 :: \text{int}) == (\text{Numeral.Pls } \text{BIT } \text{bit.B1})$
 $\langle \text{proof} \rangle$

lemma *word-1-no*:

$(1 :: 'a :: \text{len0 word}) == \text{number-of } (\text{Numeral.Pls BIT bit.B1})$
 $\langle \text{proof} \rangle$

lemma *word-m1-wi*: $-1 == \text{word-of-int } -1$

$\langle \text{proof} \rangle$

lemma *word-m1-wi-Min*: $-1 = \text{word-of-int Numeral.Min}$

$\langle \text{proof} \rangle$

lemma *word-0-bl*: $\text{of-bl } [] = 0$

$\langle \text{proof} \rangle$

lemma *word-1-bl*: $\text{of-bl } [\text{True}] = 1$

$\langle \text{proof} \rangle$

lemma *uint-0 [simp]*: $(\text{uint } 0 = 0)$

$\langle \text{proof} \rangle$

lemma *of-bl-0 [simp]*: $\text{of-bl } (\text{replicate } n \text{ False}) = 0$

$\langle \text{proof} \rangle$

lemma *to-bl-0*:

$\text{to-bl } (0 :: 'a :: \text{len0 word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ False}$
 $\langle \text{proof} \rangle$

lemma *uint-0-iff*: $(\text{uint } x = 0) = (x = 0)$

$\langle \text{proof} \rangle$

lemma *unat-0-iff*: $(\text{unat } x = 0) = (x = 0)$

$\langle \text{proof} \rangle$

lemma *unat-0 [simp]*: $\text{unat } 0 = 0$

$\langle \text{proof} \rangle$

lemma *size-0-same'*: $\text{size } w = 0 ==> w = (v :: 'a :: \text{len0 word})$

$\langle \text{proof} \rangle$

lemmas *size-0-same* = *size-0-same'* [folded word-size]

lemmas *unat-eq-0* = *unat-0-iff*

lemmas *unat-eq-zero* = *unat-0-iff*

lemma *unat-gt-0*: $(0 < \text{unat } x) = (x \sim= 0)$

$\langle \text{proof} \rangle$

lemma *ucast-0 [simp]*: $\text{ucast } 0 = 0$

$\langle \text{proof} \rangle$

lemma *sint-0* [*simp*] : *sint* 0 = 0
 ⟨*proof*⟩

lemma *scast-0* [*simp*] : *scast* 0 = 0
 ⟨*proof*⟩

lemma *sint-n1* [*simp*] : *sint* -1 = -1
 ⟨*proof*⟩

lemma *scast-n1* [*simp*] : *scast* -1 = -1
 ⟨*proof*⟩

lemma *uint-1* [*simp*] : *uint* (1 :: 'a :: len word) = 1
 ⟨*proof*⟩

lemma *unat-1* [*simp*] : *unat* (1 :: 'a :: len word) = 1
 ⟨*proof*⟩

lemma *ucast-1* [*simp*] : *ucast* (1 :: 'a :: len word) = 1
 ⟨*proof*⟩

lemmas *ariths* =
bintr-ariths [THEN *word-ubin.norm-eq-iff* [THEN *iffD1*],
folded word-ubin.eq-norm, standard]

lemma *wi-homs*:
shows
wi-hom-add: *word-of-int* a + *word-of-int* b = *word-of-int* (a + b) **and**
wi-hom-mult: *word-of-int* a * *word-of-int* b = *word-of-int* (a * b) **and**
wi-hom-neg: - *word-of-int* a = *word-of-int* (- a) **and**
wi-hom-succ: *word-succ* (*word-of-int* a) = *word-of-int* (*Numeral.succ* a) **and**
wi-hom-pred: *word-pred* (*word-of-int* a) = *word-of-int* (*Numeral.pred* a)
 ⟨*proof*⟩

lemmas *wi-hom-syms* = *wi-homs* [*symmetric*]

lemma *word-sub-def*: a - b == a + - (b :: 'a :: len0 word)
 ⟨*proof*⟩

lemmas *word-diff-minus* = *word-sub-def* [THEN *meta-eq-to-obj-eq, standard*]

lemma *word-of-int-sub-hom*:
 (*word-of-int* a) - *word-of-int* b = *word-of-int* (a - b)
 ⟨*proof*⟩

lemmas *new-word-of-int-homs* =

word-of-int-sub-hom wi-homs word-0-wi word-1-wi

lemmas *new-word-of-int-hom-syms = new-word-of-int-homs [symmetric, standard]*

lemmas *word-of-int-hom-syms =*
new-word-of-int-hom-syms [unfolded succ-def pred-def]

lemmas *word-of-int-homs =*
new-word-of-int-homs [unfolded succ-def pred-def]

lemmas *word-of-int-add-hom = word-of-int-homs (2)*
lemmas *word-of-int-mult-hom = word-of-int-homs (3)*
lemmas *word-of-int-minus-hom = word-of-int-homs (4)*
lemmas *word-of-int-succ-hom = word-of-int-homs (5)*
lemmas *word-of-int-pred-hom = word-of-int-homs (6)*
lemmas *word-of-int-0-hom = word-of-int-homs (7)*
lemmas *word-of-int-1-hom = word-of-int-homs (8)*

lemmas *word-arith-alt =*
word-sub-wi [unfolded succ-def pred-def, THEN meta-eq-to-obj-eq, standard]
word-arith-wis [unfolded succ-def pred-def, standard]

lemmas *word-sub-alt = word-arith-alt (1)*
lemmas *word-add-alt = word-arith-alt (2)*
lemmas *word-mult-alt = word-arith-alt (3)*
lemmas *word-minus-alt = word-arith-alt (4)*
lemmas *word-succ-alt = word-arith-alt (5)*
lemmas *word-pred-alt = word-arith-alt (6)*
lemmas *word-0-alt = word-arith-alt (7)*
lemmas *word-1-alt = word-arith-alt (8)*

11.1 Transferring goals from words to ints

lemma *word-ths:*

shows

word-succ-p1: word-succ a = a + 1 and
word-pred-m1: word-pred a = a - 1 and
word-pred-succ: word-pred (word-succ a) = a and
word-succ-pred: word-succ (word-pred a) = a and
*word-mult-succ: word-succ a * b = b + a * b*
<proof>

lemmas *uint-cong = arg-cong [where f = uint]*

lemmas *uint-word-ariths =*
word-arith-alt [THEN trans [OF uint-cong int-word-uint], standard]

lemmas *uint-word-arith-bintrs* = *uint-word-ariths* [*folded bintrunc-mod2p*]

lemmas *sint-word-ariths* = *uint-word-arith-bintrs*
 [*THEN uint-sint* [*symmetric*, *THEN trans*],
unfolded uint-sint bintr-arith1s bintr-ariths
len-gt-0 [*THEN bin-sbin-eq-iff* ¹] *word-sbin.norm-Rep*, *standard*]

lemmas *uint-div-alt* = *word-div-def*
 [*THEN meta-eq-to-obj-eq* [*THEN trans* [*OF uint-cong int-word-uint*]], *standard*]

lemmas *uint-mod-alt* = *word-mod-def*
 [*THEN meta-eq-to-obj-eq* [*THEN trans* [*OF uint-cong int-word-uint*]], *standard*]

lemma *word-pred-0-n1*: *word-pred 0 = word-of-int -1*
 ⟨*proof*⟩

lemma *word-pred-0-Min*: *word-pred 0 = word-of-int Numeral.Min*
 ⟨*proof*⟩

lemma *word-m1-Min*: *- 1 = word-of-int Numeral.Min*
 ⟨*proof*⟩

lemma *succ-pred-no* [*simp*]:
word-succ (number-of bin) = number-of (Numeral.succ bin) &
word-pred (number-of bin) = number-of (Numeral.pred bin)
 ⟨*proof*⟩

lemma *word-sp-01* [*simp*] :
word-succ -1 = 0 & word-succ 0 = 1 & word-pred 0 = -1 & word-pred 1 = 0
 ⟨*proof*⟩

lemma *word-of-int-Ex*:
 $\exists y. x = \text{word-of-int } y$
 ⟨*proof*⟩

lemma *word-arith-egs*:
fixes *a* :: 'a::len0 word
fixes *b* :: 'a::len0 word
shows
word-add-0: $0 + a = a$ **and**
word-add-0-right: $a + 0 = a$ **and**
word-mult-1: $1 * a = a$ **and**
word-mult-1-right: $a * 1 = a$ **and**
word-add-commute: $a + b = b + a$ **and**
word-add-assoc: $a + b + c = a + (b + c)$ **and**
word-add-left-commute: $a + (b + c) = b + (a + c)$ **and**
word-mult-commute: $a * b = b * a$ **and**
word-mult-assoc: $a * b * c = a * (b * c)$ **and**

word-mult-left-commute: $a * (b * c) = b * (a * c)$ **and**
word-left-distrib: $(a + b) * c = a * c + b * c$ **and**
word-right-distrib: $a * (b + c) = a * b + a * c$ **and**
word-left-minus: $- a + a = 0$ **and**
word-diff-0-right: $a - 0 = a$ **and**
word-diff-self: $a - a = 0$
 ⟨*proof*⟩

lemmas *word-add-ac* = *word-add-commute word-add-assoc word-add-left-commute*
lemmas *word-mult-ac* = *word-mult-commute word-mult-assoc word-mult-left-commute*

lemmas *word-plus-ac0* = *word-add-0 word-add-0-right word-add-ac*
lemmas *word-times-ac1* = *word-mult-1 word-mult-1-right word-mult-ac*

11.2 Order on fixed-length words

lemma *word-order-trans*: $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-order-refl*: $z \leq (z :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-order-antisym*: $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-order-linear*:
 $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-zero-le* [*simp*] :
 $0 \leq (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

instance *word* :: (*len0*) *semigroup-add*
 ⟨*proof*⟩

instance *word* :: (*len0*) *linorder*
 ⟨*proof*⟩

instance *word* :: (*len0*) *ring*
 ⟨*proof*⟩

lemma *word-m1-ge* [*simp*] : *word-pred* 0 >= *y*
 ⟨*proof*⟩

lemmas *word-n1-ge* [*simp*] = *word-m1-ge* [*simplified word-sp-01*]

lemmas *word-not-simps* [*simp*] =
word-zero-le [*THEN leD*] *word-m1-ge* [*THEN leD*] *word-n1-ge* [*THEN leD*]

lemma *word-gt-0*: $0 < y = (0 \sim= (y :: 'a :: \text{len0 word}))$
 $\langle \text{proof} \rangle$

lemmas *word-gt-0-no* [*simp*] = *word-gt-0* [*of number-of y, standard*]

lemma *word-sless-alt*: $(a <_s b) == (\text{sint } a < \text{sint } b)$
 $\langle \text{proof} \rangle$

lemma *word-le-nat-alt*: $(a <= b) = (\text{unat } a <= \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *word-less-nat-alt*: $(a < b) = (\text{unat } a < \text{unat } b)$
 $\langle \text{proof} \rangle$

lemma *wi-less*:
 $(\text{word-of-int } n < (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} < m \bmod 2 \wedge \text{len-of TYPE('a)})$
 $\langle \text{proof} \rangle$

lemma *wi-le*:
 $(\text{word-of-int } n <= (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} <= m \bmod 2 \wedge \text{len-of TYPE('a)})$
 $\langle \text{proof} \rangle$

lemma *udvd-nat-alt*: $a \text{ udvd } b = (\exists n \geq 0. \text{unat } b = n * \text{unat } a)$
 $\langle \text{proof} \rangle$

lemma *udvd-iff-dvd*: $x \text{ udvd } y \iff \text{unat } x \text{ dvd } \text{unat } y$
 $\langle \text{proof} \rangle$

lemmas *unat-mono* = *word-less-nat-alt* [*THEN iffD1, standard*]

lemma *word-zero-neq-one*: $0 < \text{len-of TYPE('a :: len0)} ==> (0 :: 'a \text{ word}) \sim= 1$
 $\langle \text{proof} \rangle$

lemmas *lenw1-zero-neq-one* = *len-gt-0* [*THEN word-zero-neq-one*]

lemma *no-no* [*simp*] : $\text{number-of } (\text{number-of } b) = \text{number-of } b$
 $\langle \text{proof} \rangle$

lemma *unat-minus-one*: $x \sim= 0 ==> \text{unat } (x - 1) = \text{unat } x - 1$
 $\langle \text{proof} \rangle$

lemma *measure-unat*: $p \sim= 0 ==> \text{unat } (p - 1) < \text{unat } p$
 $\langle \text{proof} \rangle$

lemmas *uint-add-ge0* [*simp*] =

add-nonneg-nonneg [*OF uint-ge-0 uint-ge-0, standard*]
lemmas *uint-mult-ge0* [*simp*] =
mult-nonneg-nonneg [*OF uint-ge-0 uint-ge-0, standard*]

lemma *uint-sub-lt2p* [*simp*]:
 $\text{uint } (x :: 'a :: \text{len0 word}) - \text{uint } (y :: 'b :: \text{len0 word}) <$
 $2 \wedge \text{len-of TYPE('a)}$
 ⟨*proof*⟩

11.3 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:
 $(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x + y :: 'a :: \text{len0 word}) = \text{uint } x + \text{uint } y)$
 ⟨*proof*⟩

lemma *uint-mult-lem*:
 $(\text{uint } x * \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$
 $(\text{uint } (x * y :: 'a :: \text{len0 word}) = \text{uint } x * \text{uint } y)$
 ⟨*proof*⟩

lemma *uint-sub-lem*:
 $(\text{uint } x \geq \text{uint } y) = (\text{uint } (x - y) = \text{uint } x - \text{uint } y)$
 ⟨*proof*⟩

lemma *uint-add-le*: $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$
 ⟨*proof*⟩

lemma *uint-sub-ge*: $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$
 ⟨*proof*⟩

lemmas *uint-sub-if'* =
trans [*OF uint-word-ariths(1) mod-sub-if-z, simplified, standard*]
lemmas *uint-plus-if'* =
trans [*OF uint-word-ariths(2) mod-add-if-z, simplified, standard*]

11.4 Definition of uint_arith

lemma *word-of-int-inverse*:
 $\text{word-of-int } r = a \implies 0 \leq r \implies r < 2 \wedge \text{len-of TYPE('a)} \implies$
 $\text{uint } (a :: 'a :: \text{len0 word}) = r$
 ⟨*proof*⟩

lemma *uint-split*:
fixes $x :: 'a :: \text{len0 word}$
shows $P (\text{uint } x) =$
 $(\text{ALL } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2 \wedge \text{len-of TYPE('a)} \implies P \ i)$
 ⟨*proof*⟩

lemma *uint-split-asm*:
fixes $x :: 'a :: \text{len0 word}$
shows $P (\text{uint } x) =$
 $(\sim (EX i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P i))$
 $\langle \text{proof} \rangle$

lemmas *uint-splits* = *uint-split uint-split-asm*

lemmas *uint-arith-simps* =
word-le-def word-less-alt
word-uint.Rep-inject [symmetric]
uint-sub-if' uint-plus-if'

lemma *power-False-cong*: $\text{False} ==> a \wedge b = c \wedge d$
 $\langle \text{proof} \rangle$

$\langle ML \rangle$

11.5 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*:
 $((x :: 'a :: \text{len0 word}) \leq x + y) = (\text{uint } x + \text{uint } y < 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemmas *no-olen-add* = *no-plus-overflow-uint-size [unfolded word-size]*

lemma *no-ulen-sub*: $((x :: 'a :: \text{len0 word}) >= x - y) = (\text{uint } y \leq \text{uint } x)$
 $\langle \text{proof} \rangle$

lemma *no-olen-add'*:
fixes $x :: 'a :: \text{len0 word}$
shows $(x \leq y + x) = (\text{uint } y + \text{uint } x < 2^{\text{len-of TYPE('a)}})$
 $\langle \text{proof} \rangle$

lemmas *olen-add-equiv* = *trans [OF no-olen-add no-olen-add' [symmetric], standard]*

lemmas *uint-plus-simple-iff* = *trans [OF no-olen-add uint-add-lem, standard]*

lemmas *uint-plus-simple* = *uint-plus-simple-iff [THEN iffD1, standard]*

lemmas *uint-minus-simple-iff* = *trans [OF no-ulen-sub uint-sub-lem, standard]*

lemmas *uint-minus-simple-alt* = *uint-sub-lem [folded word-le-def]*

lemmas *word-sub-le-iff* = *no-ulen-sub [folded word-le-def]*

lemmas *word-sub-le* = *word-sub-le-iff [THEN iffD2, standard]*

lemma *word-less-sub1*:
 $(x :: 'a :: \text{len word}) \sim 0 ==> (1 < x) = (0 < x - 1)$

$\langle proof \rangle$

lemma *word-le-sub1*:

$$(x :: 'a :: len0\ word) \sim = 0 ==> (1 <= x) = (0 <= x - 1)$$

$\langle proof \rangle$

lemma *sub-wrap-lt*:

$$((x :: 'a :: len0\ word) < x - z) = (x < z)$$

$\langle proof \rangle$

lemma *sub-wrap*:

$$((x :: 'a :: len0\ word) <= x - z) = (z = 0 \mid x < z)$$

$\langle proof \rangle$

lemma *plus-minus-not-NULL-ab*:

$$(x :: 'a :: len0\ word) <= ab - c ==> c <= ab ==> c \sim = 0 ==> x + c \sim = 0$$

$\langle proof \rangle$

lemma *plus-minus-no-overflow-ab*:

$$(x :: 'a :: len0\ word) <= ab - c ==> c <= ab ==> x <= x + c$$

$\langle proof \rangle$

lemma *le-minus'*:

$$(a :: 'a :: len0\ word) + c <= b ==> a <= a + c ==> c <= b - a$$

$\langle proof \rangle$

lemma *le-plus'*:

$$(a :: 'a :: len0\ word) <= b ==> c <= b - a ==> a + c <= b$$

$\langle proof \rangle$

lemmas *le-plus = le-plus'* [rotated]

lemmas *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*, standard]

lemma *word-plus-mono-right*:

$$(y :: 'a :: len0\ word) <= z ==> x <= x + z ==> x + y <= x + z$$

$\langle proof \rangle$

lemma *word-less-minus-cancel*:

$$y - x < z - x ==> x <= z ==> (y :: 'a :: len0\ word) < z$$

$\langle proof \rangle$

lemma *word-less-minus-mono-left*:

$$(y :: 'a :: len0\ word) < z ==> x <= y ==> y - x < z - x$$

$\langle proof \rangle$

lemma *word-less-minus-mono*:

$$a < c ==> d < b ==> a - b < a ==> c - d < c$$

$$==> a - b < c - (d :: 'a :: len0\ word)$$

$\langle proof \rangle$

lemma *word-le-minus-cancel*:

$y - x \leq z - x \implies x \leq z \implies (y :: 'a :: len0\ word) \leq z$
 $\langle proof \rangle$

lemma *word-le-minus-mono-left*:

$(y :: 'a :: len0\ word) \leq z \implies x \leq y \implies y - x \leq z - x$
 $\langle proof \rangle$

lemma *word-le-minus-mono*:

$a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c$
 $\implies a - b \leq c - (d :: 'a :: len\ word)$
 $\langle proof \rangle$

lemma *plus-le-left-cancel-wrap*:

$(x :: 'a :: len0\ word) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$
 $\langle proof \rangle$

lemma *plus-le-left-cancel-nowrap*:

$(x :: 'a :: len0\ word) \leq x + y' \implies x \leq x + y \implies$
 $(x + y' < x + y) = (y' < y)$
 $\langle proof \rangle$

lemma *word-plus-mono-right2*:

$(a :: 'a :: len0\ word) \leq a + b \implies c \leq b \implies a \leq a + c$
 $\langle proof \rangle$

lemma *word-less-add-right*:

$(x :: 'a :: len0\ word) < y - z \implies z \leq y \implies x + z < y$
 $\langle proof \rangle$

lemma *word-less-sub-right*:

$(x :: 'a :: len0\ word) < y + z \implies y \leq x \implies x - y < z$
 $\langle proof \rangle$

lemma *word-le-plus-either*:

$(x :: 'a :: len0\ word) \leq y \mid x \leq z \implies y \leq y + z \implies x \leq y + z$
 $\langle proof \rangle$

lemma *word-less-nowrapI*:

$(x :: 'a :: len0\ word) < z - k \implies k \leq z \implies 0 < k \implies x < x + k$
 $\langle proof \rangle$

lemma *inc-le*: $(i :: 'a :: len\ word) < m \implies i + 1 \leq m$

$\langle proof \rangle$

lemma *inc-i*:

$(1 :: 'a :: \text{len word}) \leq i \implies i < m \implies 1 \leq (i + 1) \ \& \ i + 1 \leq m$
 <proof>

lemma *udvd-incr-lem*:

$up < uq \implies up = ua + n * \text{uint } K \implies$
 $uq = ua + n' * \text{uint } K \implies up + \text{uint } K \leq uq$
 <proof>

lemma *udvd-incr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p + K \leq q$
 <proof>

lemma *udvd-decr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p \leq q - K$
 <proof>

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [where *ua=0, simplified*]

lemmas *udvd-incr0* = *udvd-incr'* [where *ua=0, simplified*]

lemmas *udvd-decr0* = *udvd-decr'* [where *ua=0, simplified*]

lemma *udvd-minus-le'*:

$xy < k \implies z \text{ udvd } xy \implies z \text{ udvd } k \implies xy \leq k - z$
 <proof>

lemma *udvd-incr2-K*:

$p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq$
 $p \implies$
 $0 < K \implies p \leq p + K \ \& \ p + K \leq a + s$
 <proof>

lemma *word-succ-rbl*:

$\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-succ } w) = (\text{rev } (\text{rbl-succ } (\text{rev } \text{bl})))$
 <proof>

lemma *word-pred-rbl*:

$\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-pred } w) = (\text{rev } (\text{rbl-pred } (\text{rev } \text{bl})))$
 <proof>

lemma *word-add-rbl*:

$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$
 $\text{to-bl } (v + w) = (\text{rev } (\text{rbl-add } (\text{rev } \text{vbl}) (\text{rev } \text{wbl})))$
 <proof>

lemma *word-mult-rbl*:

$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$
 $\text{to-bl } (v * w) = (\text{rev } (\text{rbl-mult } (\text{rev } \text{vbl}) (\text{rev } \text{wbl})))$

$\langle proof \rangle$

lemma *rtb-rbl-ariths*:

$rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-succ\ w)) = rbl-succ\ ys$

$rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-pred\ w)) = rbl-pred\ ys$

$[| rev\ (to-bl\ v) = ys; rev\ (to-bl\ w) = xs |]$
 $\implies rev\ (to-bl\ (v * w)) = rbl-mult\ ys\ xs$

$[| rev\ (to-bl\ v) = ys; rev\ (to-bl\ w) = xs |]$
 $\implies rev\ (to-bl\ (v + w)) = rbl-add\ ys\ xs$
 $\langle proof \rangle$

11.6 Arithmetic type class instantiations

instance *word* :: (len0) comm-monoid-add $\langle proof \rangle$

instance *word* :: (len0) comm-monoid-mult
 $\langle proof \rangle$

instance *word* :: (len0) comm-semiring
 $\langle proof \rangle$

instance *word* :: (len0) ab-group-add $\langle proof \rangle$

instance *word* :: (len0) comm-ring $\langle proof \rangle$

instance *word* :: (len) comm-semiring-1
 $\langle proof \rangle$

instance *word* :: (len) comm-ring-1 $\langle proof \rangle$

instance *word* :: (len0) comm-semiring-0 $\langle proof \rangle$

instance *word* :: (len0) order $\langle proof \rangle$

instance *word* :: (len) recpower
 $\langle proof \rangle$

lemma *zero-bintrunc*:

$iszero\ (number-of\ x :: 'a :: len\ word) =$
 $(bintrunc\ (len-of\ TYPE('a))\ x = Numeral.Pls)$
 $\langle proof \rangle$

lemmas *word-le-0-iff* [simp] =
 $word-zero-le\ [THEN\ leD,\ THEN\ linorder-antisym-conv1]$

lemma *word-of-nat*: $of_nat\ n = word_of_int\ (int\ n)$
 $\langle proof \rangle$

lemma *word-of-int*: $of_int = word_of_int$
 $\langle proof \rangle$

lemma *word-of-int-nat*:
 $0 \leq x \implies word_of_int\ x = of_nat\ (nat\ x)$
 $\langle proof \rangle$

lemma *word-number-of-eq*:
 $number_of\ w = (of_int\ w :: 'a :: len\ word)$
 $\langle proof \rangle$

instance *word* :: $(len)\ number_ring$
 $\langle proof \rangle$

lemma *iszero-word-no* [simp] :
 $iszero\ (number_of\ bin :: 'a :: len\ word) =$
 $iszero\ (number_of\ (bintrunc\ (len_of\ TYPE('a))\ bin) :: int)$
 $\langle proof \rangle$

11.7 Word and nat

lemma *td-ext-unat'*:
 $n = len_of\ TYPE\ ('a :: len) \implies$
 $td_ext\ (unat :: 'a\ word \Rightarrow nat)\ of_nat$
 $(unats\ n)\ (\%i.\ i\ mod\ 2\ ^\ n)$
 $\langle proof \rangle$

lemmas *td-ext-unat* = refl [THEN *td-ext-unat'*]

lemmas *unat-of-nat* = *td-ext-unat* [THEN *td-ext.eq-norm*, *standard*]

interpretation *word-unat*:
 $td_ext\ [unat :: 'a :: len\ word \Rightarrow nat$
 of_nat
 $unats\ (len_of\ TYPE('a :: len))$
 $\%i.\ i\ mod\ 2\ ^\ len_of\ TYPE('a :: len)]$
 $\langle proof \rangle$

lemmas *td-unat* = *word-unat.td-thm*

lemmas *unat-lt2p* [iff] = *word-unat.Rep* [unfolded *unats-def* *mem-Collect-eq*]

lemma *unat-le*: $y \leq unat\ (z :: 'a :: len\ word) \implies y : unats\ (len_of\ TYPE('a))$
 $\langle proof \rangle$

lemma *word-nchotomy*:

ALL w . *EX* n . $(w :: 'a :: \text{len word}) = \text{of-nat } n \ \& \ n < 2^{\text{len-of TYPE } ('a)}$
 $\langle \text{proof} \rangle$

lemma *of-nat-eq*:

fixes $w :: 'a :: \text{len word}$

shows $(\text{of-nat } n = w) = (\exists q. n = \text{unat } w + q * 2^{\text{len-of TYPE } ('a)})$

$\langle \text{proof} \rangle$

lemma *of-nat-eq-size*:

$(\text{of-nat } n = w) = (\text{EX } q. n = \text{unat } w + q * 2^{\text{size } w})$

$\langle \text{proof} \rangle$

lemma *of-nat-0*:

$(\text{of-nat } m = (0 :: 'a :: \text{len word})) = (\exists q. m = q * 2^{\text{len-of TYPE } ('a)})$

$\langle \text{proof} \rangle$

lemmas *of-nat-2p = mult-1 [symmetric, THEN iffD2 [OF of-nat-0 exI]]*

lemma *of-nat-gt-0*: $\text{of-nat } k \sim 0 ==> 0 < k$

$\langle \text{proof} \rangle$

lemma *of-nat-neq-0*:

$0 < k ==> k < 2^{\text{len-of TYPE } ('a :: \text{len})} ==> \text{of-nat } k \sim (0 :: 'a \text{ word})$

$\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-add*:

$\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$

$\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-mult*:

$\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$

$\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-Suc*:

$\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$

$\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-0*: $(0 :: 'a :: \text{len word}) = \text{of-nat } 0$

$\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-1*: $(1 :: 'a :: \text{len word}) = \text{of-nat } (\text{Suc } 0)$

$\langle \text{proof} \rangle$

lemmas *Abs-fnat-homs =*

Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc

Abs-fnat-hom-0 Abs-fnat-hom-1

lemma *word-arith-nat-add*:

$a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$

$\langle \text{proof} \rangle$

lemma *word-arith-nat-mult*:

$a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$

$\langle \text{proof} \rangle$

lemma *word-arith-nat-Suc*:

$\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$

$\langle \text{proof} \rangle$

lemma *word-arith-nat-div*:

$a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div } \text{unat } b)$

$\langle \text{proof} \rangle$

lemma *word-arith-nat-mod*:

$a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod } \text{unat } b)$

$\langle \text{proof} \rangle$

lemmas *word-arith-nat-defs* =

word-arith-nat-add word-arith-nat-mult

word-arith-nat-Suc Abs-fnat-hom-0

Abs-fnat-hom-1 word-arith-nat-div

word-arith-nat-mod

lemmas *unat-cong* = *arg-cong* [where $f = \text{unat}$]

lemmas *unat-word-ariths* = *word-arith-nat-defs*

[*THEN trans* [*OF unat-cong unat-of-nat*], *standard*]

lemmas *word-sub-less-iff* = *word-sub-le-iff*

[*simplified linorder-not-less* [*symmetric*], *simplified*]

lemma *unat-add-lem*:

$(\text{unat } x + \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$

$(\text{unat } (x + y :: 'a :: \text{len word}) = \text{unat } x + \text{unat } y)$

$\langle \text{proof} \rangle$

lemma *unat-mult-lem*:

$(\text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$

$(\text{unat } (x * y :: 'a :: \text{len word}) = \text{unat } x * \text{unat } y)$

$\langle \text{proof} \rangle$

lemmas *unat-plus-if'* =

trans [*OF unat-word-ariths(1) mod-nat-add, simplified, standard*]

lemma *le-no-overflow*:

$x \leq b \implies a \leq a + b \implies x \leq a + (b :: 'a :: \text{len0 word})$

$\langle \text{proof} \rangle$

lemmas *un-ui-le = trans*
[OF word-le-nat-alt [symmetric]
word-le-def [THEN meta-eq-to-obj-eq],
standard]

lemma *unat-sub-if-size:*
unat (x - y) = (if unat y <= unat x
then unat x - unat y
else unat x + 2 ^ size x - unat y)
<proof>

lemmas *unat-sub-if' = unat-sub-if-size [unfolded word-size]*

lemma *unat-div: unat ((x :: 'a :: len word) div y) = unat x div unat y*
<proof>

lemma *unat-mod: unat ((x :: 'a :: len word) mod y) = unat x mod unat y*
<proof>

lemma *uint-div: uint ((x :: 'a :: len word) div y) = uint x div uint y*
<proof>

lemma *uint-mod: uint ((x :: 'a :: len word) mod y) = uint x mod uint y*
<proof>

11.8 Definition of unat_arith tactic

lemma *unat-split:*
fixes x::'a::len word
shows P (unat x) =
(ALL n. of-nat n = x & n < 2^len-of TYPE('a) --> P n)
<proof>

lemma *unat-split-asm:*
fixes x::'a::len word
shows P (unat x) =
(~(EX n. of-nat n = x & n < 2^len-of TYPE('a) & ~ P n))
<proof>

lemmas *of-nat-inverse =*
word-unat.Abs-inverse' [rotated, unfolded unats-def, simplified]

lemmas *unat-splits = unat-split unat-split-asm*

lemmas *unat-arith-simps =*
word-le-nat-alt word-less-nat-alt
word-unat.Rep-inject [symmetric]
unat-sub-if' unat-plus-if' unat-div unat-mod

$\langle ML \rangle$

lemma *no-plus-overflow-unat-size*:

$((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemma *unat-sub*: $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemmas *no-olen-add-nat* = *no-plus-overflow-unat-size* [unfolded word-size]

lemmas *unat-plus-simple* = trans [OF no-olen-add-nat unat-add-lem, standard]

lemma *word-div-mult*:

$(0 :: 'a :: \text{len word}) < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies$
 $x * y \text{ div } y = x$
 $\langle \text{proof} \rangle$

lemma *div-lt'*: $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies$
 $\text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemmas *div-lt''* = *order-less-imp-le* [THEN div-lt']

lemma *div-lt-mult*: $(i :: 'a :: \text{len word}) < k \text{ div } x \implies 0 < x \implies i * x < k$
 $\langle \text{proof} \rangle$

lemma *div-le-mult*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$
 $\langle \text{proof} \rangle$

lemma *div-lt-uint'*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemmas *div-lt-uint''* = *order-less-imp-le* [THEN div-lt-uint']

lemma *word-le-exists'*:

$(x :: 'a :: \text{len0 word}) \leq y \implies$
 $(\exists z. y = x + z \ \& \ \text{uint } x + \text{uint } z < 2^{\text{len-of TYPE('a)}})$
 $\langle \text{proof} \rangle$

lemmas *plus-minus-not-NULL* = *order-less-imp-le* [THEN plus-minus-not-NULL-ab]

lemmas *plus-minus-no-overflow* =
order-less-imp-le [THEN plus-minus-no-overflow-ab]

lemmas *mcs* = *word-less-minus-cancel* *word-less-minus-mono-left*

word-le-minus-cancel word-le-minus-mono-left

lemmas *word-l-diffs* = *mcs* [**where** $y = w + x$, *unfolded add-diff-cancel, standard*]
lemmas *word-diff-ls* = *mcs* [**where** $z = w + x$, *unfolded add-diff-cancel, standard*]
lemmas *word-plus-mcs* = *word-diff-ls*
 [**where** $y = v + x$, *unfolded add-diff-cancel, standard*]

lemmas *le-unat-uo* = *unat-le* [*THEN word-unat.Abs-inverse*]

lemmas *thd* = *refl* [*THEN* [2] *split-div-lemma* [*THEN iffD2*], *THEN conjunct1*]

lemma *thd1*:
 $a \text{ div } b * b \leq (a :: \text{nat})$
<proof>

lemmas *uno-simps* [*THEN le-unat-uo*, *standard*] =
mod-le-divisor div-le-dividend thd1

lemma *word-mod-div-equality*:
 $(n \text{ div } b) * b + (n \text{ mod } b) = (n :: 'a :: \text{len word})$
<proof>

lemma *word-div-mult-le*: $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$
<proof>

lemma *word-mod-less-divisor*: $0 < n \implies m \text{ mod } n < (n :: 'a :: \text{len word})$
<proof>

lemma *word-of-int-power-hom*:
 $\text{word-of-int } a ^ n = (\text{word-of-int } (a ^ n) :: 'a :: \text{len word})$
<proof>

lemma *word-arith-power-alt*:
 $a ^ n = (\text{word-of-int } (\text{uint } a ^ n) :: 'a :: \text{len word})$
<proof>

lemma *of-bl-length-less*:
 $\text{length } x = k \implies k < \text{len-of TYPE('a)} \implies (\text{of-bl } x :: 'a :: \text{len word}) < 2 ^ k$
<proof>

11.9 Cardinality, finiteness of set of words

lemmas *card-lessThan'* = *card-lessThan* [*unfolded lessThan-def*]

lemmas *card-eq* = *word-unat.Abs-inj-on* [*THEN card-image*,
unfolded word-unat.image, unfolded unats-def, standard]

lemmas *card-word* = *trans* [*OF card-eq card-lessThan'*, *standard*]

lemma *finite-word-UNIV*: *finite* (*UNIV* :: 'a :: len word set)
 ⟨*proof*⟩

lemma *card-word-size*:
 $\text{card } (\text{UNIV} :: 'a :: \text{len word set}) = (2 \wedge \text{size } (x :: 'a \text{ word}))$
 ⟨*proof*⟩

end

12 WordBitwise: Bitwise Operations on Words

theory *WordBitwise* **imports** *WordArith* **begin**

lemmas *bin-log-bintrs* = *bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or*

lemmas *wils1* = *bin-log-bintrs* [*THEN word-ubin.norm-eq-iff* [*THEN iffD1*],
folded word-ubin.eq-norm, THEN eq-reflection, standard]

lemmas *word-log-binary-defs* =
word-and-def word-or-def word-xor-def

lemmas *word-no-log-defs* [*simp*] =
word-not-def [**where** *a=number-of a*,
unfolded word-no-wi wils1, folded word-no-wi, standard]
word-log-binary-defs [**where** *a=number-of a and b=number-of b*,
unfolded word-no-wi wils1, folded word-no-wi, standard]

lemmas *word-wi-log-defs* = *word-no-log-defs* [*unfolded word-no-wi*]

lemma *uint-or*: *uint* (*x OR y*) = (*uint x*) *OR* (*uint y*)
 ⟨*proof*⟩

lemma *uint-and*: *uint* (*x AND y*) = (*uint x*) *AND* (*uint y*)
 ⟨*proof*⟩

lemma *word-ops-nth-size*:
 $n < \text{size } (x :: 'a :: \text{len0 word}) \implies$
 $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \ \&$
 $(x \text{ AND } y) !! n = (x !! n \ \& y !! n) \ \&$
 $(x \text{ XOR } y) !! n = (x !! n \ \sim y !! n) \ \&$
 $(\text{NOT } x) !! n = (\sim x !! n)$
 ⟨*proof*⟩

lemma *word-ao-nth*:
fixes $x :: 'a::len0 \text{ word}$
shows $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \ \&$
 $(x \text{ AND } y) !! n = (x !! n \ \& \ y !! n)$
 $\langle \text{proof} \rangle$

lemmas *bwsimps* =
word-of-int-homs(2)
word-0-wi-Pls
word-m1-wi-Min
word-wi-log-defs

lemma *word-bw-assocs*:
fixes $x :: 'a::len0 \text{ word}$
shows
 $(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 $\langle \text{proof} \rangle$

lemma *word-bw-comms*:
fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
 $\langle \text{proof} \rangle$

lemma *word-bw-lcs*:
fixes $x :: 'a::len0 \text{ word}$
shows
 $y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 $\langle \text{proof} \rangle$

lemma *word-log-esimps* [*simp*]:
fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$

$0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$
 $\langle \text{proof} \rangle$

lemma *word-not-dist*:
fixes $x :: 'a::\text{len0 word}$
shows
 $\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$
 $\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$
 $\langle \text{proof} \rangle$

lemma *word-bw-same*:
fixes $x :: 'a::\text{len0 word}$
shows
 $x \text{ AND } x = x$
 $x \text{ OR } x = x$
 $x \text{ XOR } x = 0$
 $\langle \text{proof} \rangle$

lemma *word-ao-absorbs* [simp]:
fixes $x :: 'a::\text{len0 word}$
shows
 $x \text{ AND } (y \text{ OR } x) = x$
 $x \text{ OR } y \text{ AND } x = x$
 $x \text{ AND } (x \text{ OR } y) = x$
 $y \text{ AND } x \text{ OR } x = x$
 $(y \text{ OR } x) \text{ AND } x = x$
 $x \text{ OR } x \text{ AND } y = x$
 $(x \text{ OR } y) \text{ AND } x = x$
 $x \text{ AND } y \text{ OR } x = x$
 $\langle \text{proof} \rangle$

lemma *word-not-not* [simp]:
 $\text{NOT } \text{NOT } (x :: 'a::\text{len0 word}) = x$
 $\langle \text{proof} \rangle$

lemma *word-ao-dist*:
fixes $x :: 'a::\text{len0 word}$
shows $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$
 $\langle \text{proof} \rangle$

lemma *word-oa-dist*:
fixes $x :: 'a::\text{len0 word}$
shows $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
 $\langle \text{proof} \rangle$

lemma *word-add-not* [simp]:

fixes $x :: 'a::len0 \text{ word}$
shows $x + NOT\ x = -1$
 $\langle proof \rangle$

lemma *word-plus-and-or* [simp]:
fixes $x :: 'a::len0 \text{ word}$
shows $(x\ AND\ y) + (x\ OR\ y) = x + y$
 $\langle proof \rangle$

lemma *leoa*:
fixes $x :: 'a::len0 \text{ word}$
shows $(w = (x\ OR\ y)) ==> (y = (w\ AND\ y))\ \langle proof \rangle$

lemma *leao*:
fixes $x' :: 'a::len0 \text{ word}$
shows $(w' = (x'\ AND\ y')) ==> (x' = (x'\ OR\ w'))\ \langle proof \rangle$

lemmas *word-ao-equiv* = *leao* [COMP *leoa* [COMP *iffI*]]

lemma *le-word-or2*: $x \leq x\ OR\ (y :: 'a::len0 \text{ word})$
 $\langle proof \rangle$

lemmas *le-word-or1* = *xtr3* [OF *word-bw-comms* (2) *le-word-or2*, *standard*]

lemmas *word-and-le1* =
xtr3 [OF *word-ao-absorbs* (4) [symmetric] *le-word-or2*, *standard*]

lemmas *word-and-le2* =
xtr3 [OF *word-ao-absorbs* (8) [symmetric] *le-word-or2*, *standard*]

lemma *bl-word-not*: $to_bl\ (NOT\ w) = map\ Not\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *bl-word-xor*: $to_bl\ (v\ XOR\ w) = app2\ op\ \sim = (to_bl\ v)\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *bl-word-or*: $to_bl\ (v\ OR\ w) = app2\ op\ | \ (to_bl\ v)\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *bl-word-and*: $to_bl\ (v\ AND\ w) = app2\ op\ \& \ (to_bl\ v)\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *word-lsb-alt*: $lsb\ (w :: 'a::len0 \text{ word}) = test_bit\ w\ 0$
 $\langle proof \rangle$

lemma *word-lsb-1-0*: $lsb\ (1 :: 'a::len \text{ word}) \& \sim lsb\ (0 :: 'b::len0 \text{ word})$
 $\langle proof \rangle$

lemma *word-lsb-last*: $lsb\ (w :: 'a::len \text{ word}) = last\ (to_bl\ w)$
 $\langle proof \rangle$

lemma *word-lsb-int*: $lsb\ w = (uint\ w\ mod\ 2 = 1)$

<proof>

lemma *word-msb-sint*: $msb\ w = (sint\ w < 0)$
<proof>

lemma *word-msb-no'*:
 $w = number-of\ bin \implies msb\ (w::'a::len\ word) = bin_nth\ bin\ (size\ w - 1)$
<proof>

lemmas *word-msb-no* = *refl* [*THEN* *word-msb-no'*, *unfolded word-size*]

lemma *word-msb-nth'*: $msb\ (w::'a::len\ word) = bin_nth\ (uint\ w)\ (size\ w - 1)$
<proof>

lemmas *word-msb-nth* = *word-msb-nth'* [*unfolded word-size*]

lemma *word-msb-alt*: $msb\ (w::'a::len\ word) = hd\ (to_bl\ w)$
<proof>

lemma *word-set-nth*:
 $set_bit\ w\ n\ (test_bit\ w\ n) = (w::'a::len0\ word)$
<proof>

lemma *bin-nth-uint'*:
 $bin_nth\ (uint\ w)\ n = (rev\ (bin_to_bl\ (size\ w)\ (uint\ w))\ !\ n \ \&\ n < size\ w)$
<proof>

lemmas *bin-nth-uint* = *bin-nth-uint'* [*unfolded word-size*]

lemma *test-bit-bl*: $w\ !!\ n = (rev\ (to_bl\ w)\ !\ n \ \&\ n < size\ w)$
<proof>

lemma *to-bl-nth*: $n < size\ w \implies to_bl\ w\ !\ n = w\ !!\ (size\ w - Suc\ n)$
<proof>

lemma *test-bit-set*:
fixes $w :: 'a::len0\ word$
shows $(set_bit\ w\ n\ x)\ !!\ n = (n < size\ w \ \&\ x)$
<proof>

lemma *test-bit-set-gen*:
fixes $w :: 'a::len0\ word$
shows $test_bit\ (set_bit\ w\ n\ x)\ m =$
 $(if\ m = n\ then\ n < size\ w \ \&\ x\ else\ test_bit\ w\ m)$
<proof>

lemma *of-bl-rep-False*: $of_bl\ (replicate\ n\ False\ @\ bs) = of_bl\ bs$
<proof>

lemma *msb-nth'*:

fixes $w :: 'a::len\ word$

shows $msb\ w = w\ !!\ (size\ w - 1)$

<proof>

lemmas $msb-nth = msb-nth'\ [unfolded\ word-size]$

lemmas $msb0 = len-gt-0\ [THEN\ diff-Suc-less,\ THEN\ word-ops-nth-size\ [unfolded\ word-size],\ standard]$

lemmas $msb1 = msb0\ [where\ i = 0]$

lemmas $word-ops-msb = msb1\ [unfolded\ msb-nth\ [symmetric,\ unfolded\ One-nat-def]]$

lemmas $lsb0 = len-gt-0\ [THEN\ word-ops-nth-size\ [unfolded\ word-size],\ standard]$

lemmas $word-ops-lsb = lsb0\ [unfolded\ word-lsb-alt]$

lemma *td-ext-nth'*:

$n = size\ (w::'a::len0\ word) ==> ofn = set-bits ==> [w,\ ofn\ g] = l ==>$

$td-ext\ test-bit\ ofn\ \{f.\ ALL\ i.\ f\ i \longrightarrow i < n\}\ (\%h\ i.\ h\ i \ \&\ i < n)$

<proof>

lemmas $td-ext-nth = td-ext-nth'\ [OF\ refl\ refl\ refl,\ unfolded\ word-size]$

interpretation *test-bit*:

$td-ext\ [op\ !! :: 'a::len0\ word ==> nat ==> bool$

$set-bits$

$\{f.\ \forall i.\ f\ i \longrightarrow i < len-of\ TYPE('a::len0)\}$

$(\lambda h\ i.\ h\ i \wedge i < len-of\ TYPE('a::len0))]$

<proof>

declare *test-bit.Rep'* [simp del]

declare *test-bit.Rep'* [rule del]

lemmas $td-nth = test-bit.td-thm$

lemma *word-set-set-same*:

fixes $w :: 'a::len0\ word$

shows $set-bit\ (set-bit\ w\ n\ x)\ n\ y = set-bit\ w\ n\ y$

<proof>

lemma *word-set-set-diff*:

fixes $w :: 'a::len0\ word$

assumes $m \sim n$

shows $set-bit\ (set-bit\ w\ m\ x)\ n\ y = set-bit\ (set-bit\ w\ n\ y)\ m\ x$

<proof>

lemma *test-bit-no'*:

fixes $w :: 'a::len0\ word$

shows $w = number-of\ bin ==> test-bit\ w\ n = (n < size\ w \ \&\ bin-nth\ bin\ n)$

<proof>

lemmas *test-bit-no* =
refl [THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard]

lemma *nth-0*: $\sim (0::'a::len0 \text{ word}) !! n$
<proof>

lemma *nth-sint*:
fixes *w* :: *'a*::*len* *word*
defines *l* \equiv *len-of TYPE ('a)*
shows *bin-nth (sint w) n* = (*if* *n* < *l* - 1 *then* *w* !! *n* *else* *w* !! (*l* - 1))
<proof>

lemma *word-lsb-no*:
lsb (number-of bin :: 'a :: len word) = (*bin-last bin* = *bit.B1*)
<proof>

lemma *word-set-no*:
set-bit (number-of bin::'a::len0 word) n b =
number-of (bin-sc n (if b then bit.B1 else bit.B0) bin)
<proof>

lemmas *setBit-no* = *setBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*
simplified if-simps, THEN eq-reflection, standard]

lemmas *clearBit-no* = *clearBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*
simplified if-simps, THEN eq-reflection, standard]

lemma *to-bl-n1*:
to-bl (-1::'a::len0 word) = *replicate (len-of TYPE ('a)) True*
<proof>

lemma *word-msb-n1*: *msb (-1::'a::len word)*
<proof>

declare *word-set-set-same* [*simp*] *word-set-nth* [*simp*]
test-bit-no [*simp*] *word-set-no* [*simp*] *nth-0* [*simp*]
setBit-no [*simp*] *clearBit-no* [*simp*]
word-lsb-no [*simp*] *word-msb-no* [*simp*] *word-msb-n1* [*simp*] *word-lsb-1-0* [*simp*]

lemma *word-set-nth-iff*:
(*set-bit w n b* = *w*) = (*w* !! *n* = *b* | *n* >= *size (w::'a::len0 word)*)
<proof>

lemma *test-bit-2p'*:
w = *word-of-int (2 ^ n)* ==>
w !! *m* = (*m* = *n* & *m* < *size (w :: 'a :: len word)*)
<proof>

lemmas *test-bit-2p* = *refl [THEN test-bit-2p', unfolded word-size]*

lemmas *nth-w2p* = *test-bit-2p* [*unfolded of-int-number-of-eq*
word-of-int [*symmetric*] *of-int-power*]

lemma *uint-2p*:
 $(0 :: 'a :: \text{len } \text{word}) < 2^{\wedge} n \implies \text{uint } (2^{\wedge} n :: 'a :: \text{len } \text{word}) = 2^{\wedge} n$
 ⟨*proof*⟩

lemma *word-of-int-2p*: $(\text{word-of-int } (2^{\wedge} n) :: 'a :: \text{len } \text{word}) = 2^{\wedge} n$
 ⟨*proof*⟩

lemma *bang-is-le*: $x \ll m \implies 2^{\wedge} m \leq (x :: 'a :: \text{len } \text{word})$
 ⟨*proof*⟩

lemma *word-clr-le*:
fixes $w :: 'a :: \text{len } 0 \text{ word}$
shows $w \geq \text{set-bit } w \ n \ \text{False}$
 ⟨*proof*⟩

lemma *word-set-ge*:
fixes $w :: 'a :: \text{len } \text{word}$
shows $w \leq \text{set-bit } w \ n \ \text{True}$
 ⟨*proof*⟩

end

13 WordShift: Shifting, Rotating, and Splitting Words

theory *WordShift* **imports** *WordBitwise* **begin**

13.1 Bit shifting

lemma *shiftrl1-number* [*simp*] :
 $\text{shiftrl1 } (\text{number-of } w) = \text{number-of } (w \ \text{BIT } \text{bit.B0})$
 ⟨*proof*⟩

lemma *shiftrl1-0* [*simp*] : $\text{shiftrl1 } 0 = 0$
 ⟨*proof*⟩

lemmas *shiftrl1-def-u* = *shiftrl1-def* [*folded word-number-of-def*]

lemma *shiftrl1-def-s*: $\text{shiftrl1 } w = \text{number-of } (\text{sint } w \ \text{BIT } \text{bit.B0})$
 ⟨*proof*⟩

lemma *shiftr1-0* [*simp*] : $\text{shiftr1 } 0 = 0$

$\langle proof \rangle$

lemma *sshiftr1-0* [simp] : *sshiftr1 0 = 0*
 $\langle proof \rangle$

lemma *sshiftr1-n1* [simp] : *sshiftr1 -1 = -1*
 $\langle proof \rangle$

lemma *shiftr1-0* [simp] : $(0::'a::len0 \text{ word}) << n = 0$
 $\langle proof \rangle$

lemma *shiftr1-0* [simp] : $(0::'a::len0 \text{ word}) >> n = 0$
 $\langle proof \rangle$

lemma *sshiftr1-0* [simp] : $0 >>> n = 0$
 $\langle proof \rangle$

lemma *sshiftr1-n1* [simp] : $-1 >>> n = -1$
 $\langle proof \rangle$

lemma *nth-shiftr1*: *shiftr1 w !! n = (n < size w & n > 0 & w !! (n - 1))*
 $\langle proof \rangle$

lemma *nth-shiftr1'* [rule-format]:
 $ALL n. ((w::'a::len0 \text{ word}) << m) !! n = (n < size w \& n \geq m \& w !! (n - m))$
 $\langle proof \rangle$

lemmas *nth-shiftr1 = nth-shiftr1'* [unfolded word-size]

lemma *nth-shiftr1*: *shiftr1 w !! n = w !! Suc n*
 $\langle proof \rangle$

lemma *nth-shiftr*:
 $\bigwedge n. ((w::'a::len0 \text{ word}) >> m) !! n = w !! (n + m)$
 $\langle proof \rangle$

lemma *uint-shiftr1*: *uint (shiftr1 w) = bin-rest (uint w)*
 $\langle proof \rangle$

lemma *nth-sshiftr1*:
 $shiftr1 w !! n = (\text{if } n = \text{size } w - 1 \text{ then } w !! n \text{ else } w !! \text{Suc } n)$
 $\langle proof \rangle$

lemma *nth-sshiftr* [rule-format] :
 $ALL n. sshiftr w m !! n = (n < \text{size } w \& (\text{if } n + m \geq \text{size } w \text{ then } w !! (\text{size } w - 1) \text{ else } w !! (n + m)))$

$\langle proof \rangle$

lemma *shiftr1-div-2*: $uint\ (shiftr1\ w) = uint\ w\ div\ 2$
 $\langle proof \rangle$

lemma *sshiftr1-div-2*: $sint\ (sshiftr1\ w) = sint\ w\ div\ 2$
 $\langle proof \rangle$

lemma *shiftr-div-2n*: $uint\ (shiftr\ w\ n) = uint\ w\ div\ 2\ ^\ n$
 $\langle proof \rangle$

lemma *sshiftr-div-2n*: $sint\ (sshiftr\ w\ n) = sint\ w\ div\ 2\ ^\ n$
 $\langle proof \rangle$

13.1.1 shift functions in terms of lists of bools

lemmas *bshiftr1-no-bin* [simp] =
bshiftr1-def [where *w=number-of w, unfolded to-bl-no-bin, standard*]

lemma *bshiftr1-bl*: $to-bl\ (bshiftr1\ b\ w) = b\ \# \ butlast\ (to-bl\ w)$
 $\langle proof \rangle$

lemma *shiftrl1-of-bl*: $shiftrl1\ (of-bl\ bl) = of-bl\ (bl\ @\ [False])$
 $\langle proof \rangle$

lemma *shiftrl1-bl*: $shiftrl1\ (w :: 'a :: len0\ word) = of-bl\ (to-bl\ w\ @\ [False])$
 $\langle proof \rangle$

lemma *bl-shiftrl1*:
 $to-bl\ (shiftrl1\ (w :: 'a :: len\ word)) = tl\ (to-bl\ w)\ @\ [False]$
 $\langle proof \rangle$

lemma *shiftr1-bl*: $shiftr1\ w = of-bl\ (butlast\ (to-bl\ w))$
 $\langle proof \rangle$

lemma *bl-shiftr1*:
 $to-bl\ (shiftr1\ (w :: 'a :: len\ word)) = False\ \# \ butlast\ (to-bl\ w)$
 $\langle proof \rangle$

lemma *shiftrl1-rev*:
 $shiftrl1\ (w :: 'a :: len\ word) = word-reverse\ (shiftr1\ (word-reverse\ w))$
 $\langle proof \rangle$

lemma *shiftrl-rev*:
 $shiftrl\ (w :: 'a :: len\ word)\ n = word-reverse\ (shiftr\ (word-reverse\ w)\ n)$
 $\langle proof \rangle$

lemmas *rev-shiffl* =

shiffl-rev [where $w = \text{word-reverse } w$, *simplified*, *standard*]

lemmas *shiftr-rev* = *rev-shiffl* [THEN *word-rev-gal'*, *standard*]

lemmas *rev-shiftr* = *shiffl-rev* [THEN *word-rev-gal'*, *standard*]

lemma *bl-sshiftr1*:

$\text{to-bl } (\text{sshiftr1 } (w :: 'a :: \text{len word})) = \text{hd } (\text{to-bl } w) \# \text{butlast } (\text{to-bl } w)$
 ⟨proof⟩

lemma *drop-shiftr*:

$\text{drop } n \ (\text{to-bl } ((w :: 'a :: \text{len word}) >> n)) = \text{take } (\text{size } w - n) \ (\text{to-bl } w)$
 ⟨proof⟩

lemma *drop-sshiftr*:

$\text{drop } n \ (\text{to-bl } ((w :: 'a :: \text{len word}) >>> n)) = \text{take } (\text{size } w - n) \ (\text{to-bl } w)$
 ⟨proof⟩

lemma *take-shiftr* [rule-format] :

$n \leq \text{size } (w :: 'a :: \text{len word}) \longrightarrow \text{take } n \ (\text{to-bl } (w >> n)) =$
 $\text{replicate } n \ \text{False}$
 ⟨proof⟩

lemma *take-sshiftr'* [rule-format] :

$n \leq \text{size } (w :: 'a :: \text{len word}) \longrightarrow \text{hd } (\text{to-bl } (w >>> n)) = \text{hd } (\text{to-bl } w) \ \&$
 $\text{take } n \ (\text{to-bl } (w >>> n)) = \text{replicate } n \ (\text{hd } (\text{to-bl } w))$
 ⟨proof⟩

lemmas *hd-sshiftr* = *take-sshiftr'* [THEN *conjunct1*, *standard*]

lemmas *take-sshiftr* = *take-sshiftr'* [THEN *conjunct2*, *standard*]

lemma *atd-lem*: $\text{take } n \ xs = t \implies \text{drop } n \ xs = d \implies xs = t @ d$

⟨proof⟩

lemmas *bl-shiftr* = *atd-lem* [OF *take-shiftr* *drop-shiftr*]

lemmas *bl-sshiftr* = *atd-lem* [OF *take-sshiftr* *drop-sshiftr*]

lemma *shiffl-of-bl*: $\text{of-bl } bl \ll n = \text{of-bl } (bl @ \text{replicate } n \ \text{False})$

⟨proof⟩

lemma *shiffl-bl*:

$(w :: 'a :: \text{len0 word}) \ll (n :: \text{nat}) = \text{of-bl } (\text{to-bl } w @ \text{replicate } n \ \text{False})$
 ⟨proof⟩

lemmas *shiffl-number* [simp] = *shiffl-def* [where $w = \text{number-of } w$, *standard*]

lemma *bl-shiffl*:

$\text{to-bl } (w \ll n) = \text{drop } n \ (\text{to-bl } w) @ \text{replicate } (\min (\text{size } w) \ n) \ \text{False}$
 ⟨proof⟩

lemma *shiffl-zero-size*:

fixes $x :: 'a::len0 \text{ word}$

shows $\text{size } x \leq n \implies x \ll n = 0$

<proof>

lemma *shiffl1-2t*: $\text{shiffl1 } (w :: 'a :: len \text{ word}) = 2 * w$

<proof>

lemma *shiffl1-p*: $\text{shiffl1 } (w :: 'a :: len \text{ word}) = w + w$

<proof>

lemma *shiffl-t2n*: $\text{shiffl } (w :: 'a :: len \text{ word}) \ n = 2 ^ n * w$

<proof>

lemma *shiftr1-bintr* [simp]:

$(\text{shiftr1 } (\text{number-of } w) :: 'a :: len0 \text{ word}) =$

$\text{number-of } (\text{bin-rest } (\text{bintrunc } (\text{len-of TYPE } ('a)) \ w))$

<proof>

lemma *sshiftr1-sbintr* [simp] :

$(\text{sshiftr1 } (\text{number-of } w) :: 'a :: len \text{ word}) =$

$\text{number-of } (\text{bin-rest } (\text{sbintrunc } (\text{len-of TYPE } ('a) - 1) \ w))$

<proof>

lemma *shiftr-no'*:

$w = \text{number-of bin} \implies$

$(w :: 'a :: len0 \text{ word}) \gg n = \text{number-of } ((\text{bin-rest } ^ n) (\text{bintrunc } (\text{size } w) \ \text{bin}))$

<proof>

lemma *sshiftr-no'*:

$w = \text{number-of bin} \implies w \ggg n = \text{number-of } ((\text{bin-rest } ^ n)$

$(\text{sbintrunc } (\text{size } w - 1) \ \text{bin}))$

<proof>

lemmas *sshiftr-no* [simp] =

sshiftr-no' [where $w = \text{number-of } w$, OF *refl*, *unfolded word-size*, *standard*]

lemmas *shiftr-no* [simp] =

shiftr-no' [where $w = \text{number-of } w$, OF *refl*, *unfolded word-size*, *standard*]

lemma *shiftr1-bl-of'*:

$us = \text{shiftr1 } (\text{of-bl } bl) \implies \text{length } bl \leq \text{size } us \implies$

$us = \text{of-bl } (\text{butlast } bl)$

<proof>

lemmas *shiftr1-bl-of* = *refl* [THEN *shiftr1-bl-of'*, *unfolded word-size*]

lemma *shiftr-bl-of'* [rule-format]:

us = of-bl bl >> n ==> length bl <= size us -->
us = of-bl (take (length bl - n) bl)
 ⟨proof⟩

lemmas *shiftr-bl-of* = *refl* [THEN *shiftr-bl-of'*, *unfolded word-size*]

lemmas *shiftr-bl* = *word-bl.Rep'* [THEN *eq-imp-le*, THEN *shiftr-bl-of*,
simplified word-size, *simplified*, THEN *eq-reflection*, *standard*]

lemma *msb-shift'*: *msb (w::'a::len word) <-> (w >> (size w - 1)) ~ = 0*
 ⟨proof⟩

lemmas *msb-shift* = *msb-shift'* [*unfolded word-size*]

lemma *align-lem-or* [rule-format] :

ALL x m. length x = n + m --> length y = n + m -->
drop m x = replicate n False --> take m y = replicate m False -->
app2 op | x y = take m x @ drop m y
 ⟨proof⟩

lemma *align-lem-and* [rule-format] :

ALL x m. length x = n + m --> length y = n + m -->
drop m x = replicate n False --> take m y = replicate m False -->
app2 op & x y = replicate (n + m) False
 ⟨proof⟩

lemma *aligned-bl-add-size'*:

size x - n = m ==> n <= size x ==> drop m (to-bl x) = replicate n False
==>
take m (to-bl y) = replicate m False ==>
to-bl (x + y) = take m (to-bl x) @ drop m (to-bl y)
 ⟨proof⟩

lemmas *aligned-bl-add-size* = *refl* [THEN *aligned-bl-add-size'*]

13.1.2 Mask

lemma *nth-mask'*: *m = mask n ==> test-bit m i = (i < n & i < size m)*
 ⟨proof⟩

lemmas *nth-mask* [*simp*] = *refl* [THEN *nth-mask'*]

lemma *mask-bl*: *mask n = of-bl (replicate n True)*
 ⟨proof⟩

lemma *mask-bin*: *mask n = number-of (bintrunc n Numeral.Min)*
 ⟨proof⟩

lemma *and-mask-bintr*: $w \text{ AND mask } n = \text{number-of } (\text{bintrunc } n \text{ (uint } w))$
 ⟨proof⟩

lemma *and-mask-no*: $\text{number-of } i \text{ AND mask } n = \text{number-of } (\text{bintrunc } n \text{ } i)$
 ⟨proof⟩

lemmas *and-mask-wi* = *and-mask-no* [unfolded word-number-of-def]

lemma *bl-and-mask*:
 $\text{to-bl } (w \text{ AND mask } n :: 'a :: \text{len word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$
 $\text{drop } (\text{len-of TYPE('a)} - n) (\text{to-bl } w)$
 ⟨proof⟩

lemmas *and-mask-mod-2p* =
and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]

lemma *and-mask-lt-2p*: $\text{uint } (w \text{ AND mask } n) < 2 \wedge n$
 ⟨proof⟩

lemmas *eq-mod-iff* = *trans* [symmetric, OF int-mod-lem eq-sym-conv]

lemma *mask-eq-iff*: $(w \text{ AND mask } n) = w \iff \text{uint } w < 2 \wedge n$
 ⟨proof⟩

lemma *and-mask-dvd*: $2 \wedge n \text{ dvd uint } w = (w \text{ AND mask } n = 0)$
 ⟨proof⟩

lemma *and-mask-dvd-nat*: $2 \wedge n \text{ dvd unat } w = (w \text{ AND mask } n = 0)$
 ⟨proof⟩

lemma *word-2p-lem*:
 $n < \text{size } w \implies w < 2 \wedge n = (\text{uint } (w :: 'a :: \text{len word}) < 2 \wedge n)$
 ⟨proof⟩

lemma *less-mask-eq*: $x < 2 \wedge n \implies x \text{ AND mask } n = (x :: 'a :: \text{len word})$
 ⟨proof⟩

lemmas *mask-eq-iff-w2p* =
trans [OF mask-eq-iff word-2p-lem [symmetric], standard]

lemmas *and-mask-less'* =
iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]

lemma *and-mask-less-size*: $n < \text{size } x \implies x \text{ AND mask } n < 2 \wedge n$
 ⟨proof⟩

lemma *word-mod-2p-is-mask'*:

$c = 2 \wedge n \implies c > 0 \implies x \bmod c = (x :: 'a :: \text{len word}) \text{ AND mask } n$
 <proof>

lemmas *word-mod-2p-is-mask* = refl [THEN *word-mod-2p-is-mask'*]

lemma *mask-egs*:

$(a \text{ AND mask } n) + b \text{ AND mask } n = a + b \text{ AND mask } n$
 $a + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - b \text{ AND mask } n = a - b \text{ AND mask } n$
 $a - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $a * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $(b \text{ AND mask } n) * a \text{ AND mask } n = b * a \text{ AND mask } n$
 $(a \text{ AND mask } n) + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $(a \text{ AND mask } n) * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $-(a \text{ AND mask } n) \text{ AND mask } n = -a \text{ AND mask } n$
 $\text{word-succ } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-succ } a \text{ AND mask } n$
 $\text{word-pred } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-pred } a \text{ AND mask } n$
 <proof>

lemma *mask-power-eq*:

$(x \text{ AND mask } n) \wedge^k \text{ AND mask } n = x \wedge^k \text{ AND mask } n$
 <proof>

13.1.3 Recast

lemmas *revcast-def'* = *revcast-def* [simplified]

lemmas *revcast-def''* = *revcast-def'* [simplified word-size]

lemmas *revcast-no-def* [simp] =
revcast-def' [where $w = \text{number-of } w, \text{ unfolded word-size, standard}$]

lemma *to-bl-revcast*:

$\text{to-bl } (\text{revcast } w :: 'a :: \text{len0 word}) =$
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$
 <proof>

lemma *revcast-rev-ucast'*:

$cs = [rc, uc] \implies rc = \text{revcast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$
 $rc = \text{word-reverse } uc$
 <proof>

lemmas *revcast-rev-ucast* = *revcast-rev-ucast'* [OF refl refl refl]

lemmas *revcast-ucast* = *revcast-rev-ucast*

[where $w = \text{word-reverse } w, \text{ simplified word-rev-rev, standard}$]

lemmas *ucast-revcast* = *revcast-rev-ucast* [THEN *word-rev-gal'*, standard]

lemmas *ucast-rev-revcast* = *revcast-ucast* [THEN *word-rev-gal'*, standard]

— linking revcast and cast via shift

lemmas *wsst-TYs* = *source-size target-size word-size*

lemma *revcast-down-uu'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w >> n)$
 $\langle \text{proof} \rangle$

lemma *revcast-down-us'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w >>> n)$
 $\langle \text{proof} \rangle$

lemma *revcast-down-su'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{scast } (w >> n)$
 $\langle \text{proof} \rangle$

lemma *revcast-down-ss'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{scast } (w >>> n)$
 $\langle \text{proof} \rangle$

lemmas *revcast-down-uu* = *refl [THEN revcast-down-uu']*

lemmas *revcast-down-us* = *refl [THEN revcast-down-us']*

lemmas *revcast-down-su* = *refl [THEN revcast-down-su']*

lemmas *revcast-down-ss* = *refl [THEN revcast-down-ss']*

lemma *cast-down-rev*:

$uc = \text{ucast} \implies \text{source-size } uc = \text{target-size } uc + n \implies$
 $uc w = \text{revcast } ((w :: 'a :: \text{len word}) << n)$
 $\langle \text{proof} \rangle$

lemma *revcast-up'*:

$rc = \text{revcast} \implies \text{source-size } rc + n = \text{target-size } rc \implies$
 $rc w = (\text{ucast } w :: 'a :: \text{len word}) << n$
 $\langle \text{proof} \rangle$

lemmas *revcast-up* = *refl [THEN revcast-up']*

lemmas *rc1* = *revcast-up [THEN*

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *rc2* = *revcast-down-uu [THEN*

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *ucast-up* =

rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]

lemmas *ucast-down* =
rc2 [*simplified rev-shiftr revcast-ucast* [*symmetric*]]

13.1.4 Slices

lemmas *slice1-no-bin* [*simp*] =
slice1-def [**where** *w=number-of w, unfolded to-bl-no-bin, standard*]

lemmas *slice-no-bin* [*simp*] =
trans [*OF slice-def* [*THEN meta-eq-to-obj-eq*]
slice1-no-bin [*THEN meta-eq-to-obj-eq*],
unfolded word-size, standard]

lemma *slice1-0* [*simp*] : *slice1 n 0 = 0*
<proof>

lemma *slice-0* [*simp*] : *slice n 0 = 0*
<proof>

lemma *slice-take'*: *slice n w = of-bl (take (size w - n) (to-bl w))*
<proof>

lemmas *slice-take = slice-take'* [*unfolded word-size*]

— shiftr to a word of the same size is just slice, slice is just shiftr then ucast

lemmas *shiftr-slice = trans*
[OF shiftr-bl [*THEN meta-eq-to-obj-eq*] *slice-take* [*symmetric*], *standard*]

lemma *slice-shiftr*: *slice n w = ucast (w >> n)*
<proof>

lemma *nth-slice*:
(slice n w :: 'a :: len0 word) !! m =
(w !! (m + n) & m < len-of TYPE ('a))
<proof>

lemma *slice1-down-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>
to-bl sl = takefill False fs (drop k (to-bl w))
<proof>

lemma *slice1-up-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>
to-bl sl = takefill False fs (replicate k False @ (to-bl w))
<proof>

lemmas *sd1 = slice1-down-alt'* [*OF refl refl, unfolded word-size*]

lemmas *su1 = slice1-up-alt'* [*OF refl refl, unfolded word-size*]

lemmas *slice1-down-alt = le-add-diff-inverse* [*THEN sd1*]

lemmas *slice1-up-alt*s =

le-add-diff-inverse [*symmetric*, *THEN su1*]

le-add-diff-inverse2 [*symmetric*, *THEN su1*]

lemma *ucast-slice1*: *ucast w = slice1 (size w) w*
 ⟨*proof*⟩

lemma *ucast-slice*: *ucast w = slice 0 w*
 ⟨*proof*⟩

lemmas *slice-id* = *trans* [*OF ucast-slice* [*symmetric*] *ucast-id*]

lemma *revcast-slice1'*:
rc = revcast w ==> slice1 (size rc) w = rc
 ⟨*proof*⟩

lemmas *revcast-slice1* = *refl* [*THEN revcast-slice1'*]

lemma *slice1-tf-tf'*:
to-bl (slice1 n w :: 'a :: len0 word) =
rev (takefill False (len-of TYPE('a)) (rev (takefill False n (to-bl w))))
 ⟨*proof*⟩

lemmas *slice1-tf-tf' = slice1-tf-tf'*
 [*THEN word-bl.Rep-inverse'*, *symmetric*, *standard*]

lemma *rev-slice1*:
n + k = len-of TYPE('a) + len-of TYPE('b) ==>
slice1 n (word-reverse w :: 'b :: len0 word) =
word-reverse (slice1 k w :: 'a :: len0 word)
 ⟨*proof*⟩

lemma *rev-slice'*:
res = slice n (word-reverse w) ==> n + k + size res = size w ==>
res = word-reverse (slice k w)
 ⟨*proof*⟩

lemmas *rev-slice* = *refl* [*THEN rev-slice'*, *unfolded word-size*]

lemmas *sym-notr* =
not-iff [*THEN iffD2*, *THEN not-sym*, *THEN not-iff* [*THEN iffD1*]]

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:
(sint (x :: 'a :: len word) + sint y = sint (x + y)) =
((((x+y) XOR x) AND ((x+y) XOR y)) >> (size x - 1) = 0)
 ⟨*proof*⟩

13.2 Split and cat

lemmas *word-split-bin'* = *word-split-def* [*THEN meta-eq-to-obj-eq, standard*]

lemmas *word-cat-bin'* = *word-cat-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-rsplit-no*:

(*word-rsplit* (*number-of* *bin* :: 'b :: len0 word) :: 'a word list) =
 map *number-of* (*bin-rsplit* (*len-of TYPE*('a :: len))
 (*len-of TYPE*('b), *bintrunc* (*len-of TYPE*('b)) *bin*))
 ⟨*proof*⟩

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*
 [*unfolded bin-rsplittl-def bin-rsplit-l [symmetric]*]

lemma *test-bit-cat*:

wc = *word-cat* *a b* ==> *wc* !! *n* = (*n* < *size wc* &
 (*if n* < *size b* then *b* !! *n* else *a* !! (*n* - *size b*)))
 ⟨*proof*⟩

lemma *word-cat-bl*: *word-cat* *a b* = *of-bl* (*to-bl* *a* @ *to-bl* *b*)
 ⟨*proof*⟩

lemma *of-bl-append*:

(*of-bl* (*xs* @ *ys*) :: 'a :: len word) = *of-bl* *xs* * 2^(*length ys*) + *of-bl* *ys*
 ⟨*proof*⟩

lemma *of-bl-False* [*simp*]:

of-bl (*False*#*xs*) = *of-bl* *xs*
 ⟨*proof*⟩

lemma *of-bl-True*:

(*of-bl* (*True*#*xs*)::'a::len word) = 2^{*length xs*} + *of-bl* *xs*
 ⟨*proof*⟩

lemma *of-bl-Cons*:

of-bl (*x*#*xs*) = *of-bool* *x* * 2^{*length xs*} + *of-bl* *xs*
 ⟨*proof*⟩

lemma *split-uint-lem*: *bin-split* *n* (*uint* (*w* :: 'a :: len0 word)) = (*a*, *b*) ==>
a = *bintrunc* (*len-of TYPE*('a) - *n*) *a* & *b* = *bintrunc* (*len-of TYPE*('a)) *b*
 ⟨*proof*⟩

lemma *word-split-bl'*:

std = *size c* - *size b* ==> (*word-split* *c* = (*a*, *b*)) ==>
 (*a* = *of-bl* (*take std* (*to-bl* *c*)) & *b* = *of-bl* (*drop std* (*to-bl* *c*)))
 ⟨*proof*⟩

lemma *word-split-bl*: *std* = *size c* - *size b* ==>

(*a* = *of-bl* (*take std* (*to-bl* *c*)) & *b* = *of-bl* (*drop std* (*to-bl* *c*))) <->
word-split *c* = (*a*, *b*)

<proof>

lemma *word-split-bl-eq*:

(*word-split* (*c*::'a::len *word*) :: ('c :: len0 *word* * 'd :: len0 *word*)) =
 (of-bl (take (len-of TYPE('a::len) - len-of TYPE('d::len0)) (to-bl c)),
 of-bl (drop (len-of TYPE('a) - len-of TYPE('d)) (to-bl c)))
<proof>

lemma *test-bit-split'*:

word-split *c* = (*a*, *b*) $\dashv\dashv$ (*ALL* *n m*. *b* !! *n* = (*n* < size *b* & *c* !! *n*) &
a !! *m* = (*m* < size *a* & *c* !! (*m* + size *b*)))
<proof>

lemmas *test-bit-split* =

test-bit-split' [*THEN mp, simplified all-simps, standard*]

lemma *test-bit-split-eq*: *word-split* *c* = (*a*, *b*) $\dashv\dashv$

((*ALL* *n*::nat. *b* !! *n* = (*n* < size *b* & *c* !! *n*)) &
 (*ALL* *m*::nat. *a* !! *m* = (*m* < size *a* & *c* !! (*m* + size *b*))))
<proof>

lemma *word-cat-id*: *word-cat* *a b* = *b*

<proof>

lemma *word-cat-hom*:

len-of TYPE('a::len0) <= len-of TYPE('b::len0) + len-of TYPE('c::len0)
 \implies
 (*word-cat* (*word-of-int* *w* :: 'b *word*) (*b* :: 'c *word*) :: 'a *word*) =
word-of-int (*bin-cat* *w* (size *b*) (*uint* *b*))
<proof>

lemma *word-cat-split-alt*:

size *w* <= size *u* + size *v* \implies *word-split* *w* = (*u*, *v*) \implies *word-cat* *u v* = *w*
<proof>

lemmas *word-cat-split-size* =

sym [*THEN* [2] *word-cat-split-alt* [*symmetric*], *standard*]

13.2.1 Split and slice

lemma *split-slices*:

word-split *w* = (*u*, *v*) \implies *u* = *slice* (size *v*) *w* & *v* = *slice* 0 *w*
<proof>

lemma *slice-cat1'*:

wc = *word-cat* *a b* \implies size *wc* >= size *a* + size *b* \implies *slice* (size *b*) *wc* = *a*
<proof>

lemmas *slice-cat1* = *refl* [*THEN slice-cat1'*]

lemmas *slice-cat2* = *trans* [*OF slice-id word-cat-id*]

lemma *cat-slices*:

$a = \text{slice } n \ c ==> b = \text{slice } 0 \ c ==> n = \text{size } b ==>$
 $\text{size } a + \text{size } b >= \text{size } c ==> \text{word-cat } a \ b = c$
 ⟨proof⟩

lemma *word-split-cat-alt*:

$w = \text{word-cat } u \ v ==> \text{size } u + \text{size } v <= \text{size } w ==> \text{word-split } w = (u, v)$
 ⟨proof⟩

lemmas *word-cat-bl-no-bin* [simp] =

word-cat-bl [where $a = \text{number-of } a$
 and $b = \text{number-of } b$,
 unfolded to-bl-no-bin, standard]

lemmas *word-split-bl-no-bin* [simp] =

word-split-bl-eq [where $c = \text{number-of } c$, unfolded to-bl-no-bin, standard]

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

lemma *word-rsplit-same*: $\text{word-rsplit } w = [w]$

⟨proof⟩

lemma *word-rsplit-empty-iff-size*:

$(\text{word-rsplit } w = []) = (\text{size } w = 0)$
 ⟨proof⟩

lemma *test-bit-rsplit*:

$sw = \text{word-rsplit } w ==> m < \text{size } (\text{hd } sw :: 'a :: \text{len word}) ==>$
 $k < \text{length } sw ==> (\text{rev } sw ! k) !! m = (w !! (k * \text{size } (\text{hd } sw) + m))$
 ⟨proof⟩

lemma *word-rcat-bl*: $\text{word-rcat } wl == \text{of-bl } (\text{concat } (\text{map to-bl } wl))$

⟨proof⟩

lemma *size-rcat-lem'*:

$\text{size } (\text{concat } (\text{map to-bl } wl)) = \text{length } wl * \text{size } (\text{hd } wl)$
 ⟨proof⟩

lemmas *size-rcat-lem* = *size-rcat-lem'* [unfolded word-size]

lemmas *td-gal-lt-len* = *len-gt-0* [THEN *td-gal-lt*, standard]

lemma *nth-rcat-lem'* [rule-format] :

$sw = \text{size } (\text{hd } wl :: 'a :: \text{len word}) ==> (\text{ALL } n. n < \text{size } wl * sw \longrightarrow$
 $\text{rev } (\text{concat } (\text{map to-bl } wl)) ! n =$
 $\text{rev } (\text{to-bl } (\text{rev } wl ! (n \text{ div } sw))) ! (n \text{ mod } sw))$
 ⟨proof⟩

lemmas *nth-rcat-lem* = refl [THEN *nth-rcat-lem'*, unfolded *word-size*]

lemma *test-bit-rcat*:

sw = size (hd *wl* :: 'a :: len *word*) ==> *rc* = word-rcat *wl* ==> *rc* !! *n* =
 (*n* < size *rc* & *n* div *sw* < size *wl* & (rev *wl*) ! (*n* div *sw*) !! (*n* mod *sw*))
 <proof>

lemma *foldl-eq-foldr* [rule-format] :

ALL *x*. foldl *op* + *x* *xs* = foldr *op* + (*x* # *xs*) (0 :: 'a :: comm-monoid-add)
 <proof>

lemmas *test-bit-cong* = arg-cong [where *f* = *test-bit*, THEN *fun-cong*]

lemmas *test-bit-rsplit-alt* =

trans [OF *nth-rev-alt* [THEN *test-bit-cong*]
test-bit-rsplit [OF refl *asm-rl* *diff-Suc-less*]]

— lazy way of expressing that *u* and *v*, and *su* and *sv*, have same types

lemma *word-rsplit-len-indep'*:

[*u,v*] = *p* ==> [*su,sv*] = *q* ==> word-rsplit *u* = *su* ==>
 word-rsplit *v* = *sv* ==> length *su* = length *sv*
 <proof>

lemmas *word-rsplit-len-indep* = *word-rsplit-len-indep'* [OF refl refl refl refl]

lemma *length-word-rsplit-size*:

n = len-of TYPE ('a :: len) ==>
 (length (word-rsplit *w* :: 'a word list) <= *m*) = (size *w* <= *m* * *n*)
 <proof>

lemmas *length-word-rsplit-lt-size* =

length-word-rsplit-size [unfolded *Not-eq-iff* *linorder-not-less* [symmetric]]

lemma *length-word-rsplit-exp-size*:

n = len-of TYPE ('a :: len) ==>
 length (word-rsplit *w* :: 'a word list) = (size *w* + *n* - 1) div *n*
 <proof>

lemma *length-word-rsplit-even-size*:

n = len-of TYPE ('a :: len) ==> size *w* = *m* * *n* ==>
 length (word-rsplit *w* :: 'a word list) = *m*
 <proof>

lemmas *length-word-rsplit-exp-size'* = refl [THEN *length-word-rsplit-exp-size*]

lemmas *tdle* = iffD2 [OF *split-div-lemma* refl, THEN *conjunct1*]

lemmas *dtle* = xtr4 [OF *tdle* *mult-commute*]

lemma *word-rcat-rsplit*: *word-rcat (word-rsplit w) = w*
 ⟨*proof*⟩

lemma *size-word-rsplit-rcat-size'*:
word-rcat (ws :: 'a :: len word list) = frcw ==>
*size frcw = length ws * len-of TYPE ('a) ==>*
size (hd [word-rsplit frcw, ws]) = size ws
 ⟨*proof*⟩

lemmas *size-word-rsplit-rcat-size =*
size-word-rsplit-rcat-size' [simplified]

lemma *msrevs*:
fixes *n::nat*
shows $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$
and $(k * n + m) \text{ mod } n = m \text{ mod } n$
 ⟨*proof*⟩

lemma *word-rsplit-rcat-size'*:
word-rcat (ws :: 'a :: len word list) = frcw ==>
*size frcw = length ws * len-of TYPE ('a) ==> word-rsplit frcw = ws*
 ⟨*proof*⟩

lemmas *word-rsplit-rcat-size = refl [THEN word-rsplit-rcat-size']*

13.3 Rotation

lemmas *rotater-0'* [*simp*] = *rotater-def [where n = 0, simplified]*

lemmas *word-rot-defs = word-roti-def word-rotr-def word-rotl-def*

lemma *rotate-eq-mod*:
m mod length xs = n mod length xs ==> rotate m xs = rotate n xs
 ⟨*proof*⟩

lemmas *rotate-egs [standard] =*
trans [OF rotate0 [THEN fun-cong] id-apply]
rotate-rotate [symmetric]
rotate-id
rotate-conv-mod
rotate-eq-mod

13.3.1 Rotation of list to right

lemma *rotate1-rl'*: *rotater1 (l @ [a]) = a # l*
 ⟨*proof*⟩

lemma *rotate1-rl [simp]* : *rotater1 (rotate1 l) = l*
 ⟨*proof*⟩

lemma *rotate1-lr* [*simp*] : *rotate1* (*rotater1* *l*) = *l*
 ⟨*proof*⟩

lemma *rotater1-rev'*: *rotater1* (*rev xs*) = *rev* (*rotate1 xs*)
 ⟨*proof*⟩

lemma *rotater-rev'*: *rotater* *n* (*rev xs*) = *rev* (*rotate* *n xs*)
 ⟨*proof*⟩

lemmas *rotater-rev* = *rotater-rev'* [where *xs* = *rev ys*, *simplified*, *standard*]

lemma *rotater-drop-take*:
rotater *n xs* =
 drop (*length xs* − *n mod length xs*) *xs* @
 take (*length xs* − *n mod length xs*) *xs*
 ⟨*proof*⟩

lemma *rotater-Suc* [*simp*] :
rotater (*Suc n*) *xs* = *rotater1* (*rotater* *n xs*)
 ⟨*proof*⟩

lemma *rotate-inv-plus* [*rule-format*] :
 ALL *k*. *k* = *m* + *n* \longrightarrow *rotater* *k* (*rotate* *n xs*) = *rotater* *m xs* &
 rotate *k* (*rotater* *n xs*) = *rotate* *m xs* &
 rotater *n* (*rotate* *k xs*) = *rotate* *m xs* &
 rotate *n* (*rotater* *k xs*) = *rotater* *m xs*
 ⟨*proof*⟩

lemmas *rotate-inv-rel* = *le-add-diff-inverse2* [*symmetric*, *THEN rotate-inv-plus*]

lemmas *rotate-inv-eq* = *order-refl* [*THEN rotate-inv-rel*, *simplified*]

lemmas *rotate-lr* [*simp*] = *rotate-inv-eq* [*THEN conjunct1*, *standard*]

lemmas *rotate-rl* [*simp*] =
 rotate-inv-eq [*THEN conjunct2*, *THEN conjunct1*, *standard*]

lemma *rotate-gal*: (*rotater* *n xs* = *ys*) = (*rotate* *n ys* = *xs*)
 ⟨*proof*⟩

lemma *rotate-gal'*: (*ys* = *rotater* *n xs*) = (*xs* = *rotate* *n ys*)
 ⟨*proof*⟩

lemma *length-rotater* [*simp*]:
length (*rotater* *n xs*) = *length xs*
 ⟨*proof*⟩

lemmas *rrs0* = *rotate-eqs* [*THEN restrict-to-left*,
 simplified rotate-gal [*symmetric*] *rotate-gal'* [*symmetric*], *standard*]

lemmas $rrs1 = rrs0$ [THEN refl [THEN rev-iffD1]]
lemmas $rotater-eqs = rrs1$ [simplified length-rotater, standard]
lemmas $rotater-0 = rotater-eqs$ (1)
lemmas $rotater-add = rotater-eqs$ (2)

13.3.2 map, app2, commuting with rotate(r)

lemma $last-map$: $xs \sim [] ==> last (map f xs) = f (last xs)$
 <proof>

lemma $butlast-map$:
 $xs \sim [] ==> butlast (map f xs) = map f (butlast xs)$
 <proof>

lemma $rotater1-map$: $rotater1 (map f xs) = map f (rotater1 xs)$
 <proof>

lemma $rotater-map$:
 $rotater n (map f xs) = map f (rotater n xs)$
 <proof>

lemma $but-last-zip$ [rule-format] :
 ALL ys . $length xs = length ys --> xs \sim [] -->$
 $last (zip xs ys) = (last xs, last ys) \&$
 $butlast (zip xs ys) = zip (butlast xs) (butlast ys)$
 <proof>

lemma $but-last-app2$ [rule-format] :
 ALL ys . $length xs = length ys --> xs \sim [] -->$
 $last (app2 f xs ys) = f (last xs) (last ys) \&$
 $butlast (app2 f xs ys) = app2 f (butlast xs) (butlast ys)$
 <proof>

lemma $rotater1-zip$:
 $length xs = length ys ==>$
 $rotater1 (zip xs ys) = zip (rotater1 xs) (rotater1 ys)$
 <proof>

lemma $rotater1-app2$:
 $length xs = length ys ==>$
 $rotater1 (app2 f xs ys) = app2 f (rotater1 xs) (rotater1 ys)$
 <proof>

lemmas $lrth =$
 $box-equals$ [OF $asm-rl$ length-rotater [symmetric]
 $length-rotater$ [symmetric],
 THEN $rotater1-app2$]

lemma $rotater-app2$:

$length\ xs = length\ ys ==>$
 $rotater\ n\ (app2\ f\ xs\ ys) = app2\ f\ (rotater\ n\ xs)\ (rotater\ n\ ys)$
 $\langle proof \rangle$

lemma *rotate1-app2*:
 $length\ xs = length\ ys ==>$
 $rotate1\ (app2\ f\ xs\ ys) = app2\ f\ (rotate1\ xs)\ (rotate1\ ys)$
 $\langle proof \rangle$

lemmas *lth = box-equals* [*OF asm-rl length-rotate* [*symmetric*]
 $length-rotate$ [*symmetric*], *THEN rotate1-app2*]

lemma *rotate-app2*:
 $length\ xs = length\ ys ==>$
 $rotate\ n\ (app2\ f\ xs\ ys) = app2\ f\ (rotate\ n\ xs)\ (rotate\ n\ ys)$
 $\langle proof \rangle$

lemma *to-bl-rotl*:
 $to-bl\ (word-rotl\ n\ w) = rotate\ n\ (to-bl\ w)$
 $\langle proof \rangle$

lemmas *blrs0 = rotate-egs* [*THEN to-bl-rotl* [*THEN trans*]]

lemmas *word-rotl-egs =*
 $blrs0$ [*simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotl* [*symmetric*]]

lemma *to-bl-rotr*:
 $to-bl\ (word-rotr\ n\ w) = rotater\ n\ (to-bl\ w)$
 $\langle proof \rangle$

lemmas *brrs0 = rotater-egs* [*THEN to-bl-rotr* [*THEN trans*]]

lemmas *word-rotr-egs =*
 $brrs0$ [*simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotr* [*symmetric*]]

declare *word-rotr-egs* (1) [*simp*]

declare *word-rotl-egs* (1) [*simp*]

lemma
 $word-rot-rl$ [*simp*]:
 $word-rotl\ k\ (word-rotr\ k\ v) = v$ **and**
 $word-rot-lr$ [*simp*]:
 $word-rotr\ k\ (word-rotl\ k\ v) = v$
 $\langle proof \rangle$

lemma
 $word-rot-gal$:
 $(word-rotr\ n\ v = w) = (word-rotl\ n\ w = v)$ **and**
 $word-rot-gal'$:

$(w = \text{word-rottr } n \ v) = (v = \text{word-rotl } n \ w)$
 $\langle \text{proof} \rangle$

lemma *word-rottr-rev*:

$\text{word-rottr } n \ w = \text{word-reverse } (\text{word-rotl } n \ (\text{word-reverse } w))$
 $\langle \text{proof} \rangle$

lemma *word-roti-0* [simp]: $\text{word-roti } 0 \ w = w$
 $\langle \text{proof} \rangle$

lemmas *abl-cong* = *arg-cong* [where $f = \text{of-bl}$]

lemma *word-roti-add*:

$\text{word-roti } (m + n) \ w = \text{word-roti } m \ (\text{word-roti } n \ w)$
 $\langle \text{proof} \rangle$

lemma *word-roti-conv-mod'*: $\text{word-roti } n \ w = \text{word-roti } (n \bmod \text{int } (\text{size } w)) \ w$
 $\langle \text{proof} \rangle$

lemmas *word-roti-conv-mod* = *word-roti-conv-mod'* [unfolded *word-size*]

13.3.3 Word rotation commutes with bit-wise operations

locale *word-rotate*

context *word-rotate*

begin

lemmas *word-rot-defs'* = *to-bl-rotl to-bl-rottr*

lemmas *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

lemmas *lbl-lbl* = *trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

lemmas *ths-app2* [OF *lbl-lbl*] = *rotate-app2 rotater-app2*

lemmas *ths-map* [where $xs = \text{to-bl } v$] = *rotate-map rotater-map*

lemmas *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-app2 ths-map*

lemma *word-rot-logs*:

$\text{word-rotl } n \ (\text{NOT } v) = \text{NOT } \text{word-rotl } n \ v$
 $\text{word-rottr } n \ (\text{NOT } v) = \text{NOT } \text{word-rottr } n \ v$
 $\text{word-rotl } n \ (x \ \text{AND } y) = \text{word-rotl } n \ x \ \text{AND } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{AND } y) = \text{word-rottr } n \ x \ \text{AND } \text{word-rottr } n \ y$
 $\text{word-rotl } n \ (x \ \text{OR } y) = \text{word-rotl } n \ x \ \text{OR } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{OR } y) = \text{word-rottr } n \ x \ \text{OR } \text{word-rottr } n \ y$
 $\text{word-rotl } n \ (x \ \text{XOR } y) = \text{word-rotl } n \ x \ \text{XOR } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{XOR } y) = \text{word-rottr } n \ x \ \text{XOR } \text{word-rottr } n \ y$


```

    <proof>
end

lemmas word-rot-logs = word-rotate.word-rot-logs

lemmas bl-word-rotl-dt = trans [OF to-bl-rotl rotate-drop-take,
    simplified word-bl.Rep', standard]

lemmas bl-word-rotr-dt = trans [OF to-bl-rotr rotater-drop-take,
    simplified word-bl.Rep', standard]

lemma bl-word-roti-dt':
  n = nat ((- i) mod int (size (w :: 'a :: len word))) ==>
    to-bl (word-roti i w) = drop n (to-bl w) @ take n (to-bl w)
  <proof>

lemmas bl-word-roti-dt = bl-word-roti-dt' [unfolded word-size]

lemmas word-rotl-dt = bl-word-rotl-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-rotr-dt = bl-word-rotr-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-roti-dt = bl-word-roti-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]

lemma word-rotx-0 [simp] : word-rotr i 0 = 0 & word-rotl i 0 = 0
  <proof>

lemma word-roti-0' [simp] : word-roti n 0 = 0
  <proof>

lemmas word-rotr-dt-no-bin' [simp] =
  word-rotr-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

lemmas word-rotl-dt-no-bin' [simp] =
  word-rotl-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

declare word-roti-def [simp]

end

```

14 Boolean-Algebra: Boolean Algebras

```

theory Boolean-Algebra
imports Main
begin

```

```

locale boolean =
  fixes conj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\sqcap$  70)
  fixes disj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\sqcup$  65)
  fixes compl :: 'a  $\Rightarrow$  'a ( $\sim$  - [81] 80)
  fixes zero :: 'a (0)
  fixes one  :: 'a (1)
  assumes conj-assoc:  $(x \sqcap y) \sqcap z = x \sqcap (y \sqcap z)$ 
  assumes disj-assoc:  $(x \sqcup y) \sqcup z = x \sqcup (y \sqcup z)$ 
  assumes conj-commute:  $x \sqcap y = y \sqcap x$ 
  assumes disj-commute:  $x \sqcup y = y \sqcup x$ 
  assumes conj-disj-distrib:  $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$ 
  assumes disj-conj-distrib:  $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$ 
  assumes conj-one-right [simp]:  $x \sqcap \mathbf{1} = x$ 
  assumes disj-zero-right [simp]:  $x \sqcup \mathbf{0} = x$ 
  assumes conj-cancel-right [simp]:  $x \sqcap \sim x = \mathbf{0}$ 
  assumes disj-cancel-right [simp]:  $x \sqcup \sim x = \mathbf{1}$ 
begin

lemmas disj-ac =
  disj-assoc disj-commute
  mk-left-commute [where 'a = 'a, of disj, OF disj-assoc disj-commute]

lemmas conj-ac =
  conj-assoc conj-commute
  mk-left-commute [where 'a = 'a, of conj, OF conj-assoc conj-commute]

lemma dual: boolean disj conj compl one zero
  <proof>

14.1 Complement

lemma complement-unique:
  assumes 1:  $a \sqcap x = \mathbf{0}$ 
  assumes 2:  $a \sqcup x = \mathbf{1}$ 
  assumes 3:  $a \sqcap y = \mathbf{0}$ 
  assumes 4:  $a \sqcup y = \mathbf{1}$ 
  shows  $x = y$ 
  <proof>

lemma compl-unique:  $\llbracket x \sqcap y = \mathbf{0}; x \sqcup y = \mathbf{1} \rrbracket \Longrightarrow \sim x = y$ 
  <proof>

lemma double-compl [simp]:  $\sim (\sim x) = x$ 
  <proof>

lemma compl-eq-compl-iff [simp]:  $(\sim x = \sim y) = (x = y)$ 
  <proof>

```

14.2 Conjunction

lemma *conj-absorb* [simp]: $x \sqcap x = x$
 $\langle proof \rangle$

lemma *conj-zero-right* [simp]: $x \sqcap \mathbf{0} = \mathbf{0}$
 $\langle proof \rangle$

lemma *compl-one* [simp]: $\sim \mathbf{1} = \mathbf{0}$
 $\langle proof \rangle$

lemma *conj-zero-left* [simp]: $\mathbf{0} \sqcap x = \mathbf{0}$
 $\langle proof \rangle$

lemma *conj-one-left* [simp]: $\mathbf{1} \sqcap x = x$
 $\langle proof \rangle$

lemma *conj-cancel-left* [simp]: $\sim x \sqcap x = \mathbf{0}$
 $\langle proof \rangle$

lemma *conj-left-absorb* [simp]: $x \sqcap (x \sqcap y) = x \sqcap y$
 $\langle proof \rangle$

lemma *conj-disj-distrib2*:
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$
 $\langle proof \rangle$

lemmas *conj-disj-distrib* =
conj-disj-distrib conj-disj-distrib2

14.3 Disjunction

lemma *disj-absorb* [simp]: $x \sqcup x = x$
 $\langle proof \rangle$

lemma *disj-one-right* [simp]: $x \sqcup \mathbf{1} = \mathbf{1}$
 $\langle proof \rangle$

lemma *compl-zero* [simp]: $\sim \mathbf{0} = \mathbf{1}$
 $\langle proof \rangle$

lemma *disj-zero-left* [simp]: $\mathbf{0} \sqcup x = x$
 $\langle proof \rangle$

lemma *disj-one-left* [simp]: $\mathbf{1} \sqcup x = \mathbf{1}$
 $\langle proof \rangle$

lemma *disj-cancel-left* [simp]: $\sim x \sqcup x = \mathbf{1}$
 $\langle proof \rangle$

lemma *disj-left-absorb* [simp]: $x \sqcup (x \sqcup y) = x \sqcup y$
 $\langle proof \rangle$

lemma *disj-conj-distrib2*:
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$
 $\langle proof \rangle$

lemmas *disj-conj-distrib* =
disj-conj-distrib disj-conj-distrib2

14.4 De Morgan’s Laws

lemma *de-Morgan-conj* [simp]: $\sim (x \sqcap y) = \sim x \sqcup \sim y$
 $\langle proof \rangle$

lemma *de-Morgan-disj* [simp]: $\sim (x \sqcup y) = \sim x \sqcap \sim y$
 $\langle proof \rangle$

end

14.5 Symmetric Difference

locale *boolean-xor* = *boolean* +
fixes *xor* :: $'a \Rightarrow 'a \Rightarrow 'a$ (**infixr** \oplus 65)
assumes *xor-def*: $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$
begin

lemma *xor-def2*:
 $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$
 $\langle proof \rangle$

lemma *xor-commute*: $x \oplus y = y \oplus x$
 $\langle proof \rangle$

lemma *xor-assoc*: $(x \oplus y) \oplus z = x \oplus (y \oplus z)$
 $\langle proof \rangle$

lemmas *xor-ac* =
xor-assoc xor-commute
mk-left-commute [**where** $'a = 'a$, *of xor*, *OF xor-assoc xor-commute*]

lemma *xor-zero-right* [simp]: $x \oplus \mathbf{0} = x$
 $\langle proof \rangle$

lemma *xor-zero-left* [simp]: $\mathbf{0} \oplus x = x$
 $\langle proof \rangle$

lemma *xor-one-right* [simp]: $x \oplus \mathbf{1} = \sim x$
 $\langle proof \rangle$

lemma *xor-one-left* [*simp*]: $\mathbf{1} \oplus x = \sim x$
 $\langle proof \rangle$

lemma *xor-self* [*simp*]: $x \oplus x = \mathbf{0}$
 $\langle proof \rangle$

lemma *xor-left-self* [*simp*]: $x \oplus (x \oplus y) = y$
 $\langle proof \rangle$

lemma *xor-compl-left*: $\sim x \oplus y = \sim (x \oplus y)$
 $\langle proof \rangle$

lemma *xor-compl-right*: $x \oplus \sim y = \sim (x \oplus y)$
 $\langle proof \rangle$

lemma *xor-cancel-right* [*simp*]: $x \oplus \sim x = \mathbf{1}$
 $\langle proof \rangle$

lemma *xor-cancel-left* [*simp*]: $\sim x \oplus x = \mathbf{1}$
 $\langle proof \rangle$

lemma *conj-xor-distrib*: $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$
 $\langle proof \rangle$

lemma *conj-xor-distrib2*:
 $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$
 $\langle proof \rangle$

lemmas *conj-xor-distrib* =
conj-xor-distrib conj-xor-distrib2

end

end

15 WordGenLib: Miscellaneous Library for Words

theory *WordGenLib* **imports** *WordShift Boolean-Algebra*
begin

declare *of-nat-2p* [*simp*]

lemma *word-int-cases*:
 $\llbracket \bigwedge n. \llbracket (x :: 'a :: \text{len } 0 \text{ word}) = \text{word-of-int } n; 0 \leq n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$
 $\implies P$
 $\langle proof \rangle$

lemma *word-nat-cases* [*cases type: word*]:

$$\llbracket \bigwedge n. \llbracket (x :: 'a :: \text{len } \text{word}) = \text{of-nat } n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$$

$$\implies P$$

$$\langle \text{proof} \rangle$$

lemma *max-word-eq*:

$$(\text{max-word} :: 'a :: \text{len } \text{word}) = 2^{\text{len-of TYPE('a)}} - 1$$

$$\langle \text{proof} \rangle$$

lemma *max-word-max* [*simp, intro!*]:

$$n \leq \text{max-word}$$

$$\langle \text{proof} \rangle$$

lemma *word-of-int-2p-len*:

$$\text{word-of-int } (2^{\text{len-of TYPE('a)}}) = (0 :: 'a :: \text{len } 0 \text{ word})$$

$$\langle \text{proof} \rangle$$

lemma *word-pow-0*:

$$(2 :: 'a :: \text{len } \text{word})^{\text{len-of TYPE('a)}} = 0$$

$$\langle \text{proof} \rangle$$

lemma *max-word-wrap*: $x + 1 = 0 \implies x = \text{max-word}$

$$\langle \text{proof} \rangle$$

lemma *max-word-minus*:

$$\text{max-word} = (-1 :: 'a :: \text{len } \text{word})$$

$$\langle \text{proof} \rangle$$

lemma *max-word-bl* [*simp*]:

$$\text{to-bl } (\text{max-word} :: 'a :: \text{len } \text{word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ True}$$

$$\langle \text{proof} \rangle$$

lemma *max-test-bit* [*simp*]:

$$(\text{max-word} :: 'a :: \text{len } \text{word}) !! n = (n < \text{len-of TYPE('a)})$$

$$\langle \text{proof} \rangle$$

lemma *word-and-max* [*simp*]:

$$x \text{ AND } \text{max-word} = x$$

$$\langle \text{proof} \rangle$$

lemma *word-or-max* [*simp*]:

$$x \text{ OR } \text{max-word} = \text{max-word}$$

$$\langle \text{proof} \rangle$$

lemma *word-ao-dist2*:

$$x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } (z :: 'a :: \text{len } 0 \text{ word})$$

$$\langle \text{proof} \rangle$$

lemma *word-oa-dist2*:

$x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z::'a::len0 \text{ word}))$
 $\langle \text{proof} \rangle$

lemma *word-and-not* [simp]:
 $x \text{ AND NOT } x = (0::'a::len0 \text{ word})$
 $\langle \text{proof} \rangle$

lemma *word-or-not* [simp]:
 $x \text{ OR NOT } x = \text{max-word}$
 $\langle \text{proof} \rangle$

lemma *word-boolean*:
 $\text{boolean } (op \text{ AND}) (op \text{ OR}) \text{ bitNOT } 0 \text{ max-word}$
 $\langle \text{proof} \rangle$

interpretation *word-bool-alg*:
 $\text{boolean } [op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word}]$
 $\langle \text{proof} \rangle$

lemma *word-xor-and-or*:
 $x \text{ XOR } y = x \text{ AND NOT } y \text{ OR NOT } x \text{ AND } (y::'a::len0 \text{ word})$
 $\langle \text{proof} \rangle$

interpretation *word-bool-alg*:
 $\text{boolean-xor } [op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word } op \text{ XOR}]$
 $\langle \text{proof} \rangle$

lemma *shiftr-0* [iff]:
 $(x::'a::len0 \text{ word}) >> 0 = x$
 $\langle \text{proof} \rangle$

lemma *shiftr-0* [simp]:
 $(x :: 'a :: len \text{ word}) << 0 = x$
 $\langle \text{proof} \rangle$

lemma *shiftr-1* [simp]:
 $(1::'a::len \text{ word}) << n = 2^n$
 $\langle \text{proof} \rangle$

lemma *uint-lt-0* [simp]:
 $\text{uint } x < 0 = \text{False}$
 $\langle \text{proof} \rangle$

lemma *shiftr1-1* [simp]:
 $\text{shiftr1 } (1::'a::len \text{ word}) = 0$
 $\langle \text{proof} \rangle$

lemma *shiftr-1* [simp]:
 $(1::'a::len \text{ word}) >> n = (\text{if } n = 0 \text{ then } 1 \text{ else } 0)$

$\langle \text{proof} \rangle$

lemma *word-less-1* [simp]:
 $((x :: 'a :: \text{len word}) < 1) = (x = 0)$
 $\langle \text{proof} \rangle$

lemma *to-bl-mask*:
 $\text{to-bl } (\text{mask } n :: 'a :: \text{len word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$
 $\text{replicate } (\text{min } (\text{len-of TYPE('a)}) n) \text{ True}$
 $\langle \text{proof} \rangle$

lemma *map-replicate-True*:
 $n = \text{length } xs ==>$
 $\text{map } (\lambda(x,y). x \ \& \ y) (\text{zip } xs (\text{replicate } n \text{ True})) = xs$
 $\langle \text{proof} \rangle$

lemma *map-replicate-False*:
 $n = \text{length } xs ==> \text{map } (\lambda(x,y). x \ \& \ y)$
 $(\text{zip } xs (\text{replicate } n \text{ False})) = \text{replicate } n \text{ False}$
 $\langle \text{proof} \rangle$

lemma *bl-and-mask*:
fixes $w :: 'a :: \text{len word}$
fixes n
defines $n' \equiv \text{len-of TYPE('a)} - n$
shows $\text{to-bl } (w \text{ AND mask } n) = \text{replicate } n' \text{ False } @ \text{drop } n' (\text{to-bl } w)$
 $\langle \text{proof} \rangle$

lemma *drop-rev-takefill*:
 $\text{length } xs \leq n ==>$
 $\text{drop } (n - \text{length } xs) (\text{rev } (\text{takefill False } n (\text{rev } xs))) = xs$
 $\langle \text{proof} \rangle$

lemma *map-nth-0* [simp]:
 $\text{map } (\text{op } !! (0 :: 'a :: \text{len0 word})) xs = \text{replicate } (\text{length } xs) \text{ False}$
 $\langle \text{proof} \rangle$

lemma *uint-plus-if-size*:
 $\text{uint } (x + y) =$
 $(\text{if } \text{uint } x + \text{uint } y < 2^{\text{size } x} \text{ then}$
 $\text{uint } x + \text{uint } y$
 else
 $\text{uint } x + \text{uint } y - 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemma *unat-plus-if-size*:
 $\text{unat } (x + (y :: 'a :: \text{len word})) =$
 $(\text{if } \text{unat } x + \text{unat } y < 2^{\text{size } x} \text{ then}$


```

      unat x + unat y
    else
      unat x + unat y - 2^size x)
  ⟨proof⟩

```

lemma *word-neq-0-conv* [simp]:
fixes $w :: 'a :: \text{len word}$
shows $(w \neq 0) = (0 < w)$
 ⟨proof⟩

lemma *max-lt*:
 $\text{unat } (\max a b \text{ div } c) = \text{unat } (\max a b) \text{ div } \text{unat } (c :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *uint-sub-if-size*:
 $\text{uint } (x - y) =$
 (if $\text{uint } y \leq \text{uint } x$ then
 $\text{uint } x - \text{uint } y$
 else
 $\text{uint } x - \text{uint } y + 2^{\text{size } x}$)
 ⟨proof⟩

lemma *unat-sub-simple*:
 $x \leq y \implies \text{unat } (y - x) = \text{unat } y - \text{unat } x$
 ⟨proof⟩

lemmas *unat-sub = unat-sub-simple*

lemma *word-less-sub1*:
fixes $x :: 'a :: \text{len word}$
shows $x \neq 0 \implies 1 < x = (0 < x - 1)$
 ⟨proof⟩

lemma *word-le-sub1*:
fixes $x :: 'a :: \text{len word}$
shows $x \neq 0 \implies 1 \leq x = (0 \leq x - 1)$
 ⟨proof⟩

lemmas *word-less-sub1-numberof* [simp] =
 $\text{word-less-sub1 } [\text{of number-of } w, \text{ standard}]$
lemmas *word-le-sub1-numberof* [simp] =
 $\text{word-le-sub1 } [\text{of number-of } w, \text{ standard}]$

lemma *word-of-int-minus*:
 $\text{word-of-int } (2^{\text{len-of TYPE('a)}} - i) = (\text{word-of-int } (-i) :: 'a :: \text{len word})$
 ⟨proof⟩

lemmas *word-of-int-inj* =
 $\text{word-uint.Abs-inject } [\text{unfolded uints-num, simplified}]$

lemma *word-le-less-eq*:

$(x :: 'a :: \text{len word}) \leq y = (x = y \vee x < y)$
 $\langle \text{proof} \rangle$

lemma *mod-plus-cong*:

assumes $1: (b :: \text{int}) = b'$
and $2: x \bmod b' = x' \bmod b'$
and $3: y \bmod b' = y' \bmod b'$
and $4: x' + y' = z'$
shows $(x + y) \bmod b = z' \bmod b'$
 $\langle \text{proof} \rangle$

lemma *mod-minus-cong*:

assumes $1: (b :: \text{int}) = b'$
and $2: x \bmod b' = x' \bmod b'$
and $3: y \bmod b' = y' \bmod b'$
and $4: x' - y' = z'$
shows $(x - y) \bmod b = z' \bmod b'$
 $\langle \text{proof} \rangle$

lemma *word-induct-less*:

$\llbracket P (0 :: 'a :: \text{len word}); \bigwedge n. \llbracket n < m; P n \rrbracket \implies P (1 + n) \rrbracket \implies P m$
 $\langle \text{proof} \rangle$

lemma *word-induct*:

$\llbracket P (0 :: 'a :: \text{len word}); \bigwedge n. P n \implies P (1 + n) \rrbracket \implies P m$
 $\langle \text{proof} \rangle$

lemma *word-induct2* [*induct type*]:

$\llbracket P 0; \bigwedge n. \llbracket 1 + n \neq 0; P n \rrbracket \implies P (1 + n) \rrbracket \implies P (n :: 'b :: \text{len word})$
 $\langle \text{proof} \rangle$

constdefs

$\text{word-rec} :: 'a \Rightarrow ('b :: \text{len word} \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b \text{ word} \Rightarrow 'a$
 $\text{word-rec forZero forSuc } n \equiv \text{nat-rec forZero (forSuc } \circ \text{ of-nat)} (\text{unat } n)$

lemma *word-rec-0*: $\text{word-rec } z \text{ } s \text{ } 0 = z$

$\langle \text{proof} \rangle$

lemma *word-rec-Suc*:

$1 + n \neq (0 :: 'a :: \text{len word}) \implies \text{word-rec } z \text{ } s \text{ } (1 + n) = s \text{ } n (\text{word-rec } z \text{ } s \text{ } n)$
 $\langle \text{proof} \rangle$

lemma *word-rec-Pred*:

$n \neq 0 \implies \text{word-rec } z \text{ } s \text{ } n = s \text{ } (n - 1) (\text{word-rec } z \text{ } s \text{ } (n - 1))$
 $\langle \text{proof} \rangle$

lemma *word-rec-in*:

$f \text{ (word-rec } z \text{ (}\lambda\text{-. } f) \text{ } n) = \text{word-rec (} f \text{ } z) \text{ (}\lambda\text{-. } f) \text{ } n$
 $\langle \text{proof} \rangle$

lemma *word-rec-in2*:

$f \text{ } n \text{ (word-rec } z \text{ } f \text{ } n) = \text{word-rec (} f \text{ } 0 \text{ } z) \text{ (} f \circ \text{op} + 1) \text{ } n$
 $\langle \text{proof} \rangle$

lemma *word-rec-twice*:

$m \leq n \implies \text{word-rec } z \text{ } f \text{ } n = \text{word-rec (word-rec } z \text{ } f \text{ (} n - m)) \text{ (} f \circ \text{op} + (n - m)) \text{ } m$
 $\langle \text{proof} \rangle$

lemma *word-rec-id*: $\text{word-rec } z \text{ (}\lambda\text{-. id) } n = z$
 $\langle \text{proof} \rangle$

lemma *word-rec-id-eq*: $\forall m < n. f \text{ } m = \text{id} \implies \text{word-rec } z \text{ } f \text{ } n = z$
 $\langle \text{proof} \rangle$

lemma *word-rec-max*:

$\forall m \geq n. m \neq -1 \longrightarrow f \text{ } m = \text{id} \implies \text{word-rec } z \text{ } f \text{ } -1 = \text{word-rec } z \text{ } f \text{ } n$
 $\langle \text{proof} \rangle$

lemma *unatSuc*:

$1 + n \neq (0::'a::\text{len word}) \implies \text{unat (} 1 + n) = \text{Suc (unat } n)$
 $\langle \text{proof} \rangle$

end

16 WordMain: Main Word Library

theory *WordMain* **imports** *WordGenLib*
begin

lemmas *word-no-1* $[\text{simp}] = \text{word-1-no} [\text{symmetric}]$

lemmas *word-no-0* $[\text{simp}] = \text{word-0-no} [\text{symmetric}]$

declare *word-0-bl* $[\text{simp}]$

declare *bin-to-bl-def* $[\text{simp}]$

declare *to-bl-0* $[\text{simp}]$

declare *of-bl-True* $[\text{simp}]$

Examples

types *word32* = 32 *word*

types *word8* = 8 *word*

types *byte* = *word8*

for more see WordExampes.thy

end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.