

Machine Words in Isabelle/HOL

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Abstract

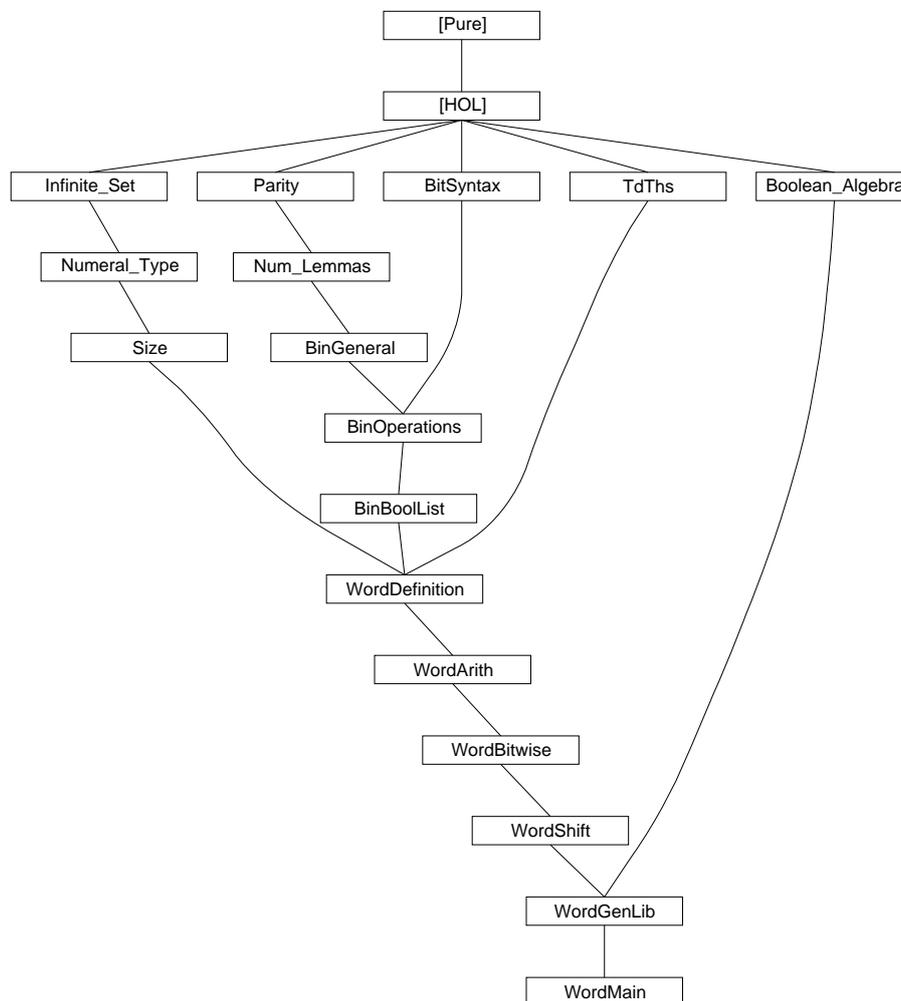
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
  imports Infinite-Set
begin
```

1.1 Preliminary lemmas

```
lemma inj-Inl [simp]: inj-on Inl A
  <proof>
```

```
lemma inj-Inr [simp]: inj-on Inr A
  <proof>
```

```
lemma inj-Some [simp]: inj-on Some A
  <proof>
```

```
lemma card-Plus:
  [| finite A; finite B |] ==> card (A <+> B) = card A + card B
  <proof>
```

```
lemma (in type-definition) univ:
  UNIV = Abs ' A
  <proof>
```

```
lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  <proof>
```

1.2 Cardinalities of types

```
syntax -type-card :: type => nat ((1CARD/(1'(-))))
```

```
translations CARD(t) => card (UNIV::t set)
```

```
<ML>
```

```
lemma card-unit: CARD(unit) = 1
  <proof>
```

```
lemma card-bool: CARD(bool) = 2
  <proof>
```

```
lemma card-prod: CARD('a::finite × 'b::finite) = CARD('a) * CARD('b)
  <proof>
```

```
lemma card-sum: CARD('a::finite + 'b::finite) = CARD('a) + CARD('b)
  <proof>
```

```
lemma card-option: CARD('a::finite option) = Suc CARD('a)
  <proof>
```

lemma *card-set*: $CARD('a::finite\ set) = 2 \wedge CARD('a)$
 ⟨*proof*⟩

1.3 Numeral Types

typedef (**open**) *num0* = *UNIV* :: *nat set* ⟨*proof*⟩
typedef (**open**) *num1* = *UNIV* :: *unit set* ⟨*proof*⟩
typedef (**open**) *'a bit0* = *UNIV* :: (*bool * 'a*) *set* ⟨*proof*⟩
typedef (**open**) *'a bit1* = *UNIV* :: (*bool * 'a*) *option set* ⟨*proof*⟩

instance *num1* :: *finite*
 ⟨*proof*⟩

instance *bit0* :: (*finite*) *finite*
 ⟨*proof*⟩

instance *bit1* :: (*finite*) *finite*
 ⟨*proof*⟩

lemma *card-num1*: $CARD(num1) = 1$
 ⟨*proof*⟩

lemma *card-bit0*: $CARD('a::finite\ bit0) = 2 * CARD('a)$
 ⟨*proof*⟩

lemma *card-bit1*: $CARD('a::finite\ bit1) = Suc\ (2 * CARD('a))$
 ⟨*proof*⟩

lemma *card-num0*: $CARD\ (num0) = 0$
 ⟨*proof*⟩

lemmas *card-univ-simps* [*simp*] =
card-unit
card-bool
card-prod
card-sum
card-option
card-set
card-num1
card-bit0
card-bit1
card-num0

1.4 Syntax

syntax
 -*NumeralType* :: *num-const* => *type* (-)
 -*NumeralType0* :: *type* (0)
 -*NumeralType1* :: *type* (1)

translations

```
-NumeralType1 == (type) num1
-NumeralType0 == (type) num0
```

⟨ML⟩

1.5 Classes with at least 1 and 2

Class `finite` already captures ”at least 1”

```
lemma zero-less-card-finite [simp]:
  0 < CARD('a::finite)
⟨proof⟩
```

```
lemma one-le-card-finite [simp]:
  Suc 0 <= CARD('a::finite)
⟨proof⟩
```

Class for cardinality ”at least 2”

```
class card2 = finite +
  assumes two-le-card: 2 <= CARD('a)
```

```
lemma one-less-card: Suc 0 < CARD('a::card2)
⟨proof⟩
```

```
instance bit0 :: (finite) card2
⟨proof⟩
```

```
instance bit1 :: (finite) card2
⟨proof⟩
```

1.6 Examples

```
term TYPE(10)
```

```
lemma CARD(0) = 0 ⟨proof⟩
lemma CARD(17) = 17 ⟨proof⟩
```

end

2 Size: The size class

```
theory Size
imports Numeral-Type
begin
```

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in `Numeral_Type`.

```
axclass len0 < type
```

```
consts
```

```
  len-of :: ('a :: len0 itself) => nat
```

Some theorems are only true on words with length greater 0.

```
axclass len < len0
```

```
  len-gt-0 [iff]: 0 < len-of TYPE ('a :: len0)
```

```
instance num0 :: len0 <proof>
```

```
instance num1 :: len0 <proof>
```

```
instance bit0 :: (len0) len0 <proof>
```

```
instance bit1 :: (len0) len0 <proof>
```

```
defs (overloaded)
```

```
  len-num0: len-of (x::num0 itself) == 0
```

```
  len-num1: len-of (x::num1 itself) == 1
```

```
  len-bit0: len-of (x::'a::len0 bit0 itself) == 2 * len-of TYPE ('a)
```

```
  len-bit1: len-of (x::'a::len0 bit1 itself) == 2 * len-of TYPE ('a) + 1
```

```
lemmas len-of-numeral-defs [simp] = len-num0 len-num1 len-bit0 len-bit1
```

```
instance num1 :: len <proof>
```

```
instance bit0 :: (len) len <proof>
```

```
instance bit1 :: (len0) len <proof>
```

```
lemma len-of TYPE(17) = 17 <proof>
```

```
lemma len-of TYPE(0) = 0 <proof>
```

```
lemma len-of TYPE('a::len0) = x  
  <proof>
```

```
end
```

3 Num-Lemmas: Useful Numerical Lemmas

```
theory Num-Lemmas imports Parity begin
```

```
lemma contentsI: y = {x} ==> contents y = x  
  <proof>
```

```
lemma prod-case-split: prod-case = split  
  <proof>
```

```
lemmas split-split = prod.split [unfolded prod-case-split]
```

lemmas *split-split-asm* = *prod.split-asm* [*unfolded prod-case-split*]
lemmas *split.splits* = *split-split split-split-asm*

lemmas *funpow-0* = *funpow.simps(1)*
lemmas *funpow-Suc* = *funpow.simps(2)*

lemma *nonemptyE*: $S \sim = \{\} \implies (!x. x : S \implies R) \implies R$
 ⟨*proof*⟩

lemma *gt-or-eq-0*: $0 < y \vee 0 = (y::nat)$ ⟨*proof*⟩

constdefs

mod-alt :: 'a => 'a => 'a :: *Divides.div*
mod-alt n m == n mod m

— alternative way of defining *bin-last*, *bin-rest*

bin-rl :: int => int * bit
bin-rl w == SOME (r, l). w = r BIT l

declare *iszero-0* [*iff*]

lemmas *xtr1* = *xtrans(1)*
lemmas *xtr2* = *xtrans(2)*
lemmas *xtr3* = *xtrans(3)*
lemmas *xtr4* = *xtrans(4)*
lemmas *xtr5* = *xtrans(5)*
lemmas *xtr6* = *xtrans(6)*
lemmas *xtr7* = *xtrans(7)*
lemmas *xtr8* = *xtrans(8)*

lemma *Min-ne-Pls* [*iff*]:
Numeral.Min $\sim =$ *Numeral.Pl*
 ⟨*proof*⟩

lemmas *Pls-ne-Min* [*iff*] = *Min-ne-Pls* [*symmetric*]

lemmas *PlsMin-defs* [*intro!*] =
Pls-def Min-def Pls-def [*symmetric*] *Min-def* [*symmetric*]

lemmas *PlsMin-simps* [*simp*] = *PlsMin-defs* [*THEN Eq-TrueI*]

lemma *number-of-False-cong*:
False \implies *number-of* x = *number-of* y
 ⟨*proof*⟩

lemmas *nat-simps* = *diff-add-inverse2 diff-add-inverse*
lemmas *nat-iffs* = *le-add1 le-add2*

lemma *sum-imp-diff*: $j = k + i \implies j - i = (k :: nat)$

<proof>

lemma nobm1:

$0 < (\text{number-of } w :: \text{nat}) ==>$
 $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Numeral.pred } w)$
<proof>

lemma of-int-power:

$\text{of-int } (a \wedge n) = (\text{of-int } a \wedge n :: 'a :: \{\text{recpower, comm-ring-1}\})$
<proof>

lemma zless2: $0 < (2 :: \text{int})$

<proof>

lemmas zless2p [simp] = zless2 [THEN zero-less-power]

lemmas zle2p [simp] = zless2p [THEN order-less-imp-le]

lemmas pos-mod-sign2 = zless2 [THEN pos-mod-sign [where b = 2::int]]

lemmas pos-mod-bound2 = zless2 [THEN pos-mod-bound [where b = 2::int]]

— the inverse(s) of *number-of*

lemma nmod2: $n \bmod (2 :: \text{int}) = 0 \mid n \bmod 2 = 1$

<proof>

lemma emep1:

$\text{even } n ==> \text{even } d ==> 0 \leq d ==> (n + 1) \bmod (d :: \text{int}) = (n \bmod d) + 1$

<proof>

lemmas eme1p = emep1 [simplified add-commute]

lemma le-diff-eq': $(a \leq c - b) = (b + a \leq (c :: \text{int}))$

<proof>

lemma less-diff-eq': $(a < c - b) = (b + a < (c :: \text{int}))$

<proof>

lemma diff-le-eq': $(a - b \leq c) = (a \leq b + (c :: \text{int}))$

<proof>

lemma diff-less-eq': $(a - b < c) = (a < b + (c :: \text{int}))$

<proof>

lemmas m1mod2k = zless2p [THEN zmod-minus1]

lemmas m1mod2k = mult-pos-pos [OF zless2 zless2p, THEN zmod-minus1]

lemmas p1mod2k' = zless2p [THEN order-less-imp-le, THEN pos-zmod-mult-2]

lemmas z1pmod2' = zero-le-one [THEN pos-zmod-mult-2, simplified]

lemmas z1pdiv2' = zero-le-one [THEN pos-zdiv-mult-2, simplified]

lemma *p1mod22k*:

$$(2 * b + 1) \text{ mod } (2 * 2 ^ n) = 2 * (b \text{ mod } 2 ^ n) + (1::\text{int})$$

<proof>

lemma *z1pmod2*:

$$(2 * b + 1) \text{ mod } 2 = (1::\text{int})$$

<proof>

lemma *z1pdiv2*:

$$(2 * b + 1) \text{ div } 2 = (b::\text{int})$$

<proof>

lemmas *zdiv-le-dividend = xtr3 [OF zdiv-1 [symmetric] zdiv-mono2, simplified int-one-le-iff-zero-less, simplified, standard]*

lemma *BIT-eq*: $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$
<proof>

lemmas *BIT-eqE [elim!] = BIT-eq [THEN conjE, standard]*

lemma *BIT-eq-iff [simp]*:

$$(u \text{ BIT } b = v \text{ BIT } c) = (u = v \ \wedge \ b = c)$$

<proof>

lemmas *BIT-eqI [intro!] = conjI [THEN BIT-eq-iff [THEN iffD2]]*

lemma *less-Bits*:

$$(v \text{ BIT } b < w \text{ BIT } c) = (v < w \ | \ v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$$

<proof>

lemma *le-Bits*:

$$(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \ | \ v \leq w \ \& \ (b \ \sim = \text{bit.B1} \ | \ c \ \sim = \text{bit.B0}))$$

<proof>

lemma *neB1E [elim!]*:

assumes *ne*: $y \neq \text{bit.B1}$
assumes *y*: $y = \text{bit.B0} \implies P$
shows *P*
<proof>

lemma *bin-ex-rl*: $EX \ w \ b. \ w \text{ BIT } b = \text{bin}$
<proof>

lemma *bin-exhaust*:

assumes *Q*: $\bigwedge x \ b. \ \text{bin} = x \text{ BIT } b \implies Q$
shows *Q*

<proof>

lemma *bin-rl-char*: $(\text{bin-rl } w = (r, l)) = (r \text{ BIT } l = w)$
<proof>

lemmas *bin-rl-simps* [*THEN bin-rl-char* [*THEN iffD2*], *standard*, *simp*] =
Pls-0-eq Min-1-eq refl

lemma *bin-abs-lem*:
 $\text{bin} = (w \text{ BIT } b) \implies \sim \text{bin} = \text{Numeral.Min} \dashrightarrow \sim \text{bin} = \text{Numeral.Pls} \dashrightarrow$
 $\text{nat } (\text{abs } w) < \text{nat } (\text{abs } \text{bin})$
<proof>

lemma *bin-induct*:
assumes *PPls*: $P \text{ Numeral.Pls}$
and *PMin*: $P \text{ Numeral.Min}$
and *PBit*: $\forall \text{bin bit. } P \text{ bin} \implies P (\text{bin BIT bit})$
shows $P \text{ bin}$
<proof>

lemma *no-no* [*simp*]: $\text{number-of } (\text{number-of } i) = i$
<proof>

lemma *Bit-B0*:
 $k \text{ BIT bit.B0} = k + k$
<proof>

lemma *Bit-B1*:
 $k \text{ BIT bit.B1} = k + k + 1$
<proof>

lemma *Bit-B0-2t*: $k \text{ BIT bit.B0} = 2 * k$
<proof>

lemma *Bit-B1-2t*: $k \text{ BIT bit.B1} = 2 * k + 1$
<proof>

lemma *B-mod-2'*:
 $X = 2 \implies (w \text{ BIT bit.B1}) \text{ mod } X = 1 \ \& \ (w \text{ BIT bit.B0}) \text{ mod } X = 0$
<proof>

lemmas *B1-mod-2* [*simp*] = *B-mod-2'* [*OF refl*, *THEN conjunct1*, *standard*]

lemmas *B0-mod-2* [*simp*] = *B-mod-2'* [*OF refl*, *THEN conjunct2*, *standard*]

lemma *axbbyy*:
 $a + m + m = b + n + n \implies (a = 0 \mid a = 1) \implies (b = 0 \mid b = 1) \implies$
 $a = b \ \& \ m = (n :: \text{int})$
<proof>

lemma *axxmod2*:

$$(1 + x + x) \text{ mod } 2 = (1 :: \text{int}) \ \& \ (0 + x + x) \text{ mod } 2 = (0 :: \text{int})$$

<proof>

lemma *axxdiv2*:

$$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int})$$

<proof>

lemmas *iszero-minus* = *trans* [THEN *trans*,

OF iszero-def neg-equal-0-iff-equal iszero-def [*symmetric*], *standard*]

lemmas *zadd-diff-inverse* = *trans* [*OF diff-add-cancel* [*symmetric*] *add-commute*,
standard]

lemmas *add-diff-cancel2* = *add-commute* [THEN *diff-eq-eq* [THEN *iffD2*], *standard*]

lemma *zmod-uminus*: $-(a :: \text{int}) \text{ mod } b \text{ mod } b = -a \text{ mod } b$
<proof>

lemma *zmod-zsub-distrib*: $((a :: \text{int}) - b) \text{ mod } c = (a \text{ mod } c - b \text{ mod } c) \text{ mod } c$
<proof>

lemma *zmod-zsub-right-eq*: $((a :: \text{int}) - b) \text{ mod } c = (a - b \text{ mod } c) \text{ mod } c$
<proof>

lemma *zmod-zsub-left-eq*: $((a :: \text{int}) - b) \text{ mod } c = (a \text{ mod } c - b) \text{ mod } c$
<proof>

lemma *zmod-zsub-self* [*simp*]:

$$((b :: \text{int}) - a) \text{ mod } a = b \text{ mod } a$$

<proof>

lemma *zmod-zmult1-eq-rev*:

$$b * a \text{ mod } c = b \text{ mod } c * a \text{ mod } (c :: \text{int})$$

<proof>

lemmas *rdmods* [*symmetric*] = *zmod-uminus* [*symmetric*]

zmod-zsub-left-eq zmod-zsub-right-eq zmod-zadd-left-eq
zmod-zadd-right-eq zmod-zmult1-eq zmod-zmult1-eq-rev

lemma *mod-plus-right*:

$$((a + x) \text{ mod } m = (b + x) \text{ mod } m) = (a \text{ mod } m = b \text{ mod } (m :: \text{nat}))$$

<proof>

lemma *nat-minus-mod*: $(n - n \text{ mod } m) \text{ mod } m = (0 :: \text{nat})$
<proof>

lemmas *nat-minus-mod-plus-right* = *trans* [*OF nat-minus-mod mod-0* [*symmetric*],

THEN mod-plus-right [*THEN iffD2*], *standard*, *simplified*]

lemmas *push-mods'* = *zmod-zadd1-eq* [*standard*]
zmod-zmult-distrib [*standard*] *zmod-zsub-distrib* [*standard*]
zmod-uminus [*symmetric*, *standard*]

lemmas *push-mods* = *push-mods'* [*THEN eq-reflection*, *standard*]

lemmas *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN eq-reflection*, *standard*]

lemmas *mod-simps* =
zmod-zmult-self1 [*THEN eq-reflection*] *zmod-zmult-self2* [*THEN eq-reflection*]
mod-mod-trivial [*THEN eq-reflection*]

lemma *nat-mod-eq*:

!!*b. b < n ==> a mod n = b mod n ==> a mod n = (b :: nat)*
<proof>

lemmas *nat-mod-eq'* = *refl* [*THEN* [2] *nat-mod-eq*]

lemma *nat-mod-lem*:

(0 :: nat) < n ==> b < n = (b mod n = b)
<proof>

lemma *mod-nat-add*:

(x :: nat) < z ==> y < z ==>
(x + y) mod z = (if x + y < z then x + y else x + y - z)
<proof>

lemma *mod-nat-sub*:

(x :: nat) < z ==> (x - y) mod z = x - y
<proof>

lemma *int-mod-lem*:

(0 :: int) < n ==> (0 <= b & b < n) = (b mod n = b)
<proof>

lemma *int-mod-eq*:

(0 :: int) <= b ==> b < n ==> a mod n = b mod n ==> a mod n = b
<proof>

lemmas *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

lemma *int-mod-le*: *0 <= a ==> 0 < (n :: int) ==> a mod n <= a*
<proof>

lemma *int-mod-le'*: *0 <= b - n ==> 0 < (n :: int) ==> b mod n <= b - n*
<proof>

lemma *int-mod-ge*: *a < n ==> 0 < (n :: int) ==> a <= a mod n*

<proof>

lemma *int-mod-ge'*: $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$
<proof>

lemma *mod-add-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
<proof>

lemma *mod-sub-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
<proof>

lemmas *zmde = zmod-zdiv-equality* [THEN *diff-eq-eq* [THEN *iffD2*], *symmetric*]

lemmas *mcl = mult-cancel-left* [THEN *iffD1*, THEN *make-pos-rule*]

lemma *zdiv-mult-self*: $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$
<proof>

lemma *eqne*: $\text{equiv } A \ r \implies X : A // r \implies X \sim = \{\}$
<proof>

lemmas *Rep-Integ-ne = Integ.Rep-Integ*

[THEN *equiv-intrel* [THEN *eqne*, *simplified Integ-def* [symmetric]], *standard*]

lemmas *riq = Integ.Rep-Integ* [*simplified Integ-def*]

lemmas *intrel-refl = refl* [THEN *equiv-intrel-iff* [THEN *iffD1*], *standard*]

lemmas *Rep-Integ-equiv = quotient-eq-iff*

[OF *equiv-intrel riq riq*, *simplified Integ.Rep-Integ-inject*, *standard*]

lemmas *Rep-Integ-same =*

Rep-Integ-equiv [THEN *intrel-refl* [THEN *rev-iffD2*], *standard*]

lemma *RI-int*: $(a, 0) : \text{Rep-Integ } (int \ a)$

<proof>

lemmas *RI-intrel [simp] = UNIV-I* [THEN *quotientI*,

THEN *Integ.Abs-Integ-inverse* [*simplified Integ-def*], *standard*]

lemma *RI-minus*: $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (- \ x)$

<proof>

lemma *RI-add*:

$(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$

$(a + c, b + d) : \text{Rep-Integ } (x + y)$

<proof>

lemma *mem-same*: $a : S \implies a = b \implies b : S$
 ⟨proof⟩

lemma *RI-eq-diff'*: $(a, b) : \text{Rep-Integ } (\text{int } a - \text{int } b)$
 ⟨proof⟩

lemma *RI-eq-diff*: $((a, b) : \text{Rep-Integ } x) = (\text{int } a - \text{int } b = x)$
 ⟨proof⟩

lemma *mod-power-lem*:
 $a > 1 \implies a^n \bmod a^m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int})^n)$
 ⟨proof⟩

lemma *min-pm* [*simp*]: $\text{min } a \ b + (a - b) = (a :: \text{nat})$
 ⟨proof⟩

lemmas *min-pm1* [*simp*] = *trans* [*OF add-commute min-pm*]

lemma *rev-min-pm* [*simp*]: $\text{min } b \ a + (a - b) = (a :: \text{nat})$
 ⟨proof⟩

lemmas *rev-min-pm1* [*simp*] = *trans* [*OF add-commute rev-min-pm*]

lemma *pl-pl-rels*:
 $a + b = c + d \implies$
 $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$
 ⟨proof⟩

lemmas *pl-pl-rels'* = *add-commute* [*THEN* [2] *trans*, *THEN pl-pl-rels*]

lemma *minus-eq*: $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$
 ⟨proof⟩

lemma *pl-pl-mm*: $(a :: \text{nat}) + b = c + d \implies a - c = d - b$
 ⟨proof⟩

lemmas *pl-pl-mm'* = *add-commute* [*THEN* [2] *trans*, *THEN pl-pl-mm*]

lemma *min-minus* [*simp*]: $\text{min } m \ (m - k) = (m - k :: \text{nat})$
 ⟨proof⟩

lemmas *min-minus'* [*simp*] = *trans* [*OF min-max.inf-commute min-minus*]

lemma *nat-no-eq-iff*:
 $(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$
 $(\text{number-of } b = (\text{number-of } c :: \text{nat})) = (b = c)$
 ⟨proof⟩

lemmas *dme* = *box-equals* [*OF* *div-mod-equality* *add-0-right* *add-0-right*]
lemmas *dtle* = *xtr3* [*OF* *dme* [*symmetric*] *le-add1*]
lemmas *th2* = *order-trans* [*OF* *order-refl* [*THEN* [2] *mult-le-mono*] *dtle*]

lemma *td-gal*:

$0 < c \implies (a \geq b * c) = (a \text{ div } c \geq (b :: \text{nat}))$
 ⟨*proof*⟩

lemmas *td-gal-lt* = *td-gal* [*simplified le-def*, *simplified*]

lemma *div-mult-le*: $(a :: \text{nat}) \text{ div } b * b \leq a$
 ⟨*proof*⟩

lemmas *sdl* = *split-div-lemma* [*THEN* *iffD1*, *symmetric*]

lemma *given-quot*: $f > (0 :: \text{nat}) \implies (f * l + (f - 1)) \text{ div } f = l$
 ⟨*proof*⟩

lemma *given-quot-alt*: $f > (0 :: \text{nat}) \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$
 ⟨*proof*⟩

lemma *diff-mod-le*: $(a :: \text{nat}) < d \implies b \text{ dvd } d \implies a - a \text{ mod } b \leq d - b$
 ⟨*proof*⟩

lemma *less-le-mult'*:

$w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * (c :: \text{int})$
 ⟨*proof*⟩

lemmas *less-le-mult* = *less-le-mult'* [*simplified left-distrib*, *simplified*]

lemmas *less-le-mult-minus* = *iffD2* [*OF* *le-diff-eq* *less-le-mult*,
simplified left-diff-distrib, *standard*]

lemma *lrlem'*:

assumes *d*: $(i :: \text{nat}) \leq j \vee m < j'$
assumes *R1*: $i * k \leq j * k \implies R$
assumes *R2*: $\text{Suc } m * k' \leq j' * k' \implies R$
shows *R* ⟨*proof*⟩

lemma *lrlem*: $(0 :: \text{nat}) < sc \implies$

$(sc - n + (n + lb * n) \leq m * n) = (sc + lb * n \leq m * n)$
 ⟨*proof*⟩

lemma *gen-minus*: $0 < n \implies f n = f (\text{Suc } (n - 1))$
 ⟨*proof*⟩

lemma *mpl-lem*: $j \leq (i :: \text{nat}) \implies k < j \implies i - j + k < i$
 ⟨*proof*⟩

lemma *nonneg-mod-div*:

$0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b :: \text{int}) \ \& \ 0 \leq a \text{ div } b$
 ⟨proof⟩

end

4 BinGeneral: Basic Definitions for Binary Integers

theory *BinGeneral* **imports** *Num-Lemmas*

begin

4.1 Recursion combinator for binary integers

lemma *brlem*: $(\text{bin} = \text{Numeral.Min}) = (- \text{bin} + \text{Numeral.pred } 0 = 0)$
 ⟨proof⟩

function

$\text{bin-rec}' :: \text{int} * 'a * 'a * (\text{int} \implies \text{bit} \implies 'a \implies 'a) \implies 'a$

where

$\text{bin-rec}' (\text{bin}, f1, f2, f3) = (\text{if } \text{bin} = \text{Numeral.Plus} \text{ then } f1$

$\text{else if } \text{bin} = \text{Numeral.Min} \text{ then } f2$

$\text{else case bin-rl bin of } (w, b) \implies f3 \ w \ b \ (\text{bin-rec}' (w, f1, f2, f3)))$

⟨proof⟩

termination

⟨proof⟩

constdefs

$\text{bin-rec} :: 'a \implies 'a \implies (\text{int} \implies \text{bit} \implies 'a \implies 'a) \implies \text{int} \implies 'a$

$\text{bin-rec } f1 \ f2 \ f3 \ \text{bin} == \text{bin-rec}' (\text{bin}, f1, f2, f3)$

lemma *bin-rec-PM*:

$f = \text{bin-rec } f1 \ f2 \ f3 \implies f \ \text{Numeral.Plus} = f1 \ \& \ f \ \text{Numeral.Min} = f2$

⟨proof⟩

lemmas *bin-rec-Plus* = *refl* [THEN *bin-rec-PM*, THEN *conjunct1*, *standard*]

lemmas *bin-rec-Min* = *refl* [THEN *bin-rec-PM*, THEN *conjunct2*, *standard*]

lemma *bin-rec-Bit*:

$f = \text{bin-rec } f1 \ f2 \ f3 \implies f3 \ \text{Numeral.Plus} \ \text{bit.B0} \ f1 = f1 \implies$

$f3 \ \text{Numeral.Min} \ \text{bit.B1} \ f2 = f2 \implies f (w \ \text{BIT} \ b) = f3 \ w \ b (f \ w)$

⟨proof⟩

lemmas *bin-rec-simps* = *refl* [THEN *bin-rec-Bit*] *bin-rec-Plus bin-rec-Min*

4.2 Destructors for binary integers

consts

— corresponding operations analysing bins
 $bin_last :: int \Rightarrow bit$
 $bin_rest :: int \Rightarrow int$
 $bin_sign :: int \Rightarrow int$
 $bin_nth :: int \Rightarrow nat \Rightarrow bool$

primrec

$Z : bin_nth\ w\ 0 = (bin_last\ w = bit.B1)$
 $Suc : bin_nth\ w\ (Suc\ n) = bin_nth\ (bin_rest\ w)\ n$

defs

$bin_rest_def : bin_rest\ w == fst\ (bin_rl\ w)$
 $bin_last_def : bin_last\ w == snd\ (bin_rl\ w)$
 $bin_sign_def : bin_sign == bin_rec\ Numeral.Pls\ Numeral.Min\ (\%w\ b\ s.\ s)$

lemma bin_rl : $bin_rl\ w = (bin_rest\ w, bin_last\ w)$
 $\langle proof \rangle$

lemmas bin_rl_simp $[simp] = iffD1$ $[OF\ bin_rl_char\ bin_rl]$

lemma bin_rest_simps $[simp]$:

$bin_rest\ Numeral.Pls = Numeral.Pls$
 $bin_rest\ Numeral.Min = Numeral.Min$
 $bin_rest\ (w\ BIT\ b) = w$
 $\langle proof \rangle$

lemma bin_last_simps $[simp]$:

$bin_last\ Numeral.Pls = bit.B0$
 $bin_last\ Numeral.Min = bit.B1$
 $bin_last\ (w\ BIT\ b) = b$
 $\langle proof \rangle$

lemma bin_sign_simps $[simp]$:

$bin_sign\ Numeral.Pls = Numeral.Pls$
 $bin_sign\ Numeral.Min = Numeral.Min$
 $bin_sign\ (w\ BIT\ b) = bin_sign\ w$
 $\langle proof \rangle$

lemma $bin_r_l_extras$ $[simp]$:

$bin_last\ 0 = bit.B0$
 $bin_last\ (-1) = bit.B1$
 $bin_last\ -1 = bit.B1$
 $bin_last\ 1 = bit.B1$
 $bin_rest\ 1 = 0$
 $bin_rest\ 0 = 0$
 $bin_rest\ (-1) = -1$
 $bin_rest\ -1 = -1$

<proof>

lemma *bin-last-mod*:

bin-last $w = (\text{if } w \text{ mod } 2 = 0 \text{ then bit.B0 else bit.B1})$

<proof>

lemma *bin-rest-div*:

bin-rest $w = w \text{ div } 2$

<proof>

lemma *Bit-div2* [*simp*]: $(w \text{ BIT } b) \text{ div } 2 = w$

<proof>

lemma *bin-nth-lem* [*rule-format*]:

ALL y . *bin-nth* $x = \text{bin-nth } y \dashrightarrow x = y$

<proof>

lemma *bin-nth-eq-iff*: $(\text{bin-nth } x = \text{bin-nth } y) = (x = y)$

<proof>

lemmas *bin-eqI* = *ext* [*THEN bin-nth-eq-iff* [*THEN iffD1*], *standard*]

lemma *bin-nth-Pls* [*simp*]: $\sim \text{bin-nth Numeral.Pls } n$

<proof>

lemma *bin-nth-Min* [*simp*]: *bin-nth Numeral.Min* n

<proof>

lemma *bin-nth-0-BIT*: *bin-nth* $(w \text{ BIT } b) 0 = (b = \text{bit.B1})$

<proof>

lemma *bin-nth-Suc-BIT*: *bin-nth* $(w \text{ BIT } b) (\text{Suc } n) = \text{bin-nth } w \ n$

<proof>

lemma *bin-nth-minus* [*simp*]: $0 < n \implies \text{bin-nth } (w \text{ BIT } b) \ n = \text{bin-nth } w \ (n - 1)$

<proof>

lemmas *bin-nth-0* = *bin-nth.simps(1)*

lemmas *bin-nth-Suc* = *bin-nth.simps(2)*

lemmas *bin-nth-simps* =

bin-nth-0 bin-nth-Suc bin-nth-Pls bin-nth-Min bin-nth-minus

lemma *bin-sign-rest* [*simp*]:

bin-sign $(\text{bin-rest } w) = (\text{bin-sign } w)$

<proof>

4.3 Truncating binary integers

consts

bintrunc :: nat => int => int

primrec

Z : *bintrunc* 0 *bin* = *Numeral.Pls*

Suc : *bintrunc* (*Suc* *n*) *bin* = *bintrunc* *n* (*bin-rest* *bin*) *BIT* (*bin-last* *bin*)

consts

sbintrunc :: nat => int => int

primrec

Z : *sbintrunc* 0 *bin* =

(*case bin-last bin of bit.B1 => Numeral.Min* | *bit.B0 => Numeral.Pls*)

Suc : *sbintrunc* (*Suc* *n*) *bin* = *sbintrunc* *n* (*bin-rest* *bin*) *BIT* (*bin-last* *bin*)

lemma *sign-bintr*:

!!*w*. *bin-sign* (*bintrunc* *n* *w*) = *Numeral.Pls*

<*proof*>

lemma *bintrunc-mod2p*:

!!*w*. *bintrunc* *n* *w* = (*w mod* 2 ^ *n* :: int)

<*proof*>

lemma *sbintrunc-mod2p*:

!!*w*. *sbintrunc* *n* *w* = ((*w* + 2 ^ *n*) *mod* 2 ^ (*Suc* *n*) - 2 ^ *n* :: int)

<*proof*>

4.4 Simplifications for (s)bintrunc

lemma *bit-bool*:

(*b* = (*b'* = *bit.B1*)) = (*b'* = (*if* *b* *then bit.B1* *else bit.B0*))

<*proof*>

lemmas *bit-bool1* [*simp*] = *refl* [*THEN bit-bool* [*THEN iffD1*], *symmetric*]

lemma *bin-sign-lem*:

!!*bin*. (*bin-sign* (*sbintrunc* *n* *bin*) = *Numeral.Min*) = *bin-nth* *bin* *n*

<*proof*>

lemma *nth-bintr*:

!!*w m*. *bin-nth* (*bintrunc* *m* *w*) *n* = (*n* < *m* & *bin-nth* *w* *n*)

<*proof*>

lemma *nth-sbintr*:

!!*w m*. *bin-nth* (*sbintrunc* *m* *w*) *n* =

(*if* *n* < *m* *then bin-nth* *w* *n* *else bin-nth* *w* *m*)

<*proof*>

lemma *bin-nth-Bit*:

bin-nth (*w BIT b*) *n* = (*n* = 0 & *b* = *bit.B1* | (*EX m*. *n* = *Suc* *m* & *bin-nth* *w*

m)
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-l*:
 $n \leq m \implies (\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } n w)$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-sbintrunc-l*:
 $n \leq m \implies (\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } n w)$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-ge*:
 $n \leq m \implies (\text{bintrunc } n (\text{bintrunc } m w) = \text{bintrunc } n w)$
 $\langle \text{proof} \rangle$

lemma *bintrunc-bintrunc-min* [simp]:
 $\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } (\min m n) w$
 $\langle \text{proof} \rangle$

lemma *sbintrunc-sbintrunc-min* [simp]:
 $\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } (\min m n) w$
 $\langle \text{proof} \rangle$

lemmas *bintrunc-Pls* =
 bintrunc.Suc [where $\text{bin} = \text{Numeral.Pls}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *bintrunc-Min* [simp] =
 bintrunc.Suc [where $\text{bin} = \text{Numeral.Min}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *bintrunc-BIT* [simp] =
 bintrunc.Suc [where $\text{bin} = w \text{ BIT } b$, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

lemmas *sbintrunc-Suc-Pls* =
 sbintrunc.Suc [where $\text{bin} = \text{Numeral.Pls}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *sbintrunc-Suc-Min* =
 sbintrunc.Suc [where $\text{bin} = \text{Numeral.Min}$, *simplified bin-last-simps bin-rest-simps*,
standard]

lemmas *sbintrunc-Suc-BIT* [simp] =
 sbintrunc.Suc [where $\text{bin} = w \text{ BIT } b$, *simplified bin-last-simps bin-rest-simps*, *standard*]

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls* *sbintrunc-Suc-Min* *sbintrunc-Suc-BIT*

lemmas *sbintrunc-Pls* =
sbintrunc.Z [**where** *bin=Numeral.Pls*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-Min* =
sbintrunc.Z [**where** *bin=Numeral.Min*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-BIT-B0* [*simp*] =
sbintrunc.Z [**where** *bin=w BIT bit.B0*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-BIT-B1* [*simp*] =
sbintrunc.Z [**where** *bin=w BIT bit.B1*,
simplified bin-last-simps bin-rest-simps bit.simps, standard]

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls *sbintrunc-Min* *sbintrunc-0-BIT-B0* *sbintrunc-0-BIT-B1*

lemmas *bintrunc-simps* = *bintrunc.Z* *bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps* *sbintrunc-Sucs*

lemma *bintrunc-minus*:
 $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) w = \text{bintrunc } n w$
<proof>

lemma *sbintrunc-minus*:
 $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) w = \text{sbintrunc } n w$
<proof>

lemmas *bintrunc-minus-simps* =
bintrunc-Sucs [*THEN* [2] *bintrunc-minus* [*symmetric, THEN trans*], *standard*]

lemmas *sbintrunc-minus-simps* =
sbintrunc-Sucs [*THEN* [2] *sbintrunc-minus* [*symmetric, THEN trans*], *standard*]

lemma *bintrunc-n-Pls* [*simp*]:
 $\text{bintrunc } n \text{ Numeral.Pls} = \text{Numeral.Pls}$
<proof>

lemma *sbintrunc-n-PM* [*simp*]:
 $\text{sbintrunc } n \text{ Numeral.Pls} = \text{Numeral.Pls}$
 $\text{sbintrunc } n \text{ Numeral.Min} = \text{Numeral.Min}$
<proof>

lemmas *thobini1* = *arg-cong* [**where** $f = \%w$. $w \text{ BIT } b$, *standard*]

lemmas *bintrunc-BIT-I* = *trans* [*OF* *bintrunc-BIT* *thobini1*]

lemmas *bintrunc-Min-I* = *trans* [*OF bintrunc-Min thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps* [*THEN thobini1* [*THEN* [2] *trans*], *standard*]

lemmas *bintrunc-Pls-minus-I* = *bmsts*(1)

lemmas *bintrunc-Min-minus-I* = *bmsts*(2)

lemmas *bintrunc-BIT-minus-I* = *bmsts*(3)

lemma *bintrunc-0-Min*: *bintrunc 0 Numeral.Min* = *Numeral.Pls*
 ⟨*proof*⟩

lemma *bintrunc-0-BIT*: *bintrunc 0 (w BIT b)* = *Numeral.Pls*
 ⟨*proof*⟩

lemma *bintrunc-Suc-lem*:

bintrunc (Suc n) x = y ==> m = Suc n ==> bintrunc m x = y
 ⟨*proof*⟩

lemmas *bintrunc-Suc-Ialts* =

bintrunc-Min-I bintrunc-BIT-I [*THEN bintrunc-Suc-lem, standard*]

lemmas *sbintrunc-BIT-I* = *trans* [*OF sbintrunc-Suc-BIT thobini1*]

lemmas *sbintrunc-Suc-Is* =

sbintrunc-Sucs [*THEN thobini1* [*THEN* [2] *trans*], *standard*]

lemmas *sbintrunc-Suc-minus-Is* =

sbintrunc-minus-simps [*THEN thobini1* [*THEN* [2] *trans*], *standard*]

lemma *sbintrunc-Suc-lem*:

sbintrunc (Suc n) x = y ==> m = Suc n ==> sbintrunc m x = y
 ⟨*proof*⟩

lemmas *sbintrunc-Suc-Ialts* =

sbintrunc-Suc-Is [*THEN sbintrunc-Suc-lem, standard*]

lemma *sbintrunc-bintrunc-lt*:

m > n ==> sbintrunc n (bintrunc m w) = sbintrunc n w
 ⟨*proof*⟩

lemma *bintrunc-sbintrunc-le*:

m <= Suc n ==> bintrunc m (sbintrunc n w) = bintrunc m w
 ⟨*proof*⟩

lemmas *bintrunc-sbintrunc* [*simp*] = *order-refl* [*THEN bintrunc-sbintrunc-le*]

lemmas *sbintrunc-bintrunc* [*simp*] = *lessI* [*THEN sbintrunc-bintrunc-lt*]

lemmas *bintrunc-bintrunc* [*simp*] = *order-refl* [*THEN bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc* [*simp*] = *order-refl* [*THEN sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc'* [*simp*]:

$0 < n \implies \text{bintrunc } n \ (\text{sbintrunc } (n - 1) \ w) = \text{bintrunc } n \ w$
 ⟨proof⟩

lemma *sbintrunc-bintrunc'* [simp]:

$0 < n \implies \text{sbintrunc } (n - 1) \ (\text{bintrunc } n \ w) = \text{sbintrunc } (n - 1) \ w$
 ⟨proof⟩

lemma *bin-sbin-eq-iff*:

$\text{bintrunc } (\text{Suc } n) \ x = \text{bintrunc } (\text{Suc } n) \ y \iff$
 $\text{sbintrunc } n \ x = \text{sbintrunc } n \ y$
 ⟨proof⟩

lemma *bin-sbin-eq-iff'*:

$0 < n \implies \text{bintrunc } n \ x = \text{bintrunc } n \ y \iff$
 $\text{sbintrunc } (n - 1) \ x = \text{sbintrunc } (n - 1) \ y$
 ⟨proof⟩

lemmas *bintrunc-sbintruncS0* [simp] = *bintrunc-sbintrunc'* [unfolded One-nat-def]

lemmas *sbintrunc-bintruncS0* [simp] = *sbintrunc-bintrunc'* [unfolded One-nat-def]

lemmas *bintrunc-bintrunc-l'* = *le-add1* [THEN *bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc-l'* = *le-add1* [THEN *sbintrunc-sbintrunc-l*]

lemmas *nat-non0-gr* =

trans [OF *iszero-def* [THEN *Not-eq-iff* [THEN *iffD2*]] *refl, standard*]

lemmas *bintrunc-pred-simps* [simp] =

bintrunc-minus-simps [of *number-of bin, simplified nobm1, standard*]

lemmas *sbintrunc-pred-simps* [simp] =

sbintrunc-minus-simps [of *number-of bin, simplified nobm1, standard*]

lemma *no-bintr-alt*:

$\text{number-of } (\text{bintrunc } n \ w) = w \ \text{mod } 2 \ ^n$
 ⟨proof⟩

lemma *no-bintr-alt1*: $\text{bintrunc } n = (\%w. \ w \ \text{mod } 2 \ ^n \ :: \ \text{int})$

⟨proof⟩

lemma *range-bintrunc*: $\text{range } (\text{bintrunc } n) = \{i. \ 0 \leq i \ \& \ i < 2 \ ^n\}$

⟨proof⟩

lemma *no-bintr*:

$\text{number-of } (\text{bintrunc } n \ w) = (\text{number-of } w \ \text{mod } 2 \ ^n \ :: \ \text{int})$
 ⟨proof⟩

lemma *no-sbintr-alt2*:

$sbintrunc\ n = (\%w. (w + 2^{\wedge}n) \bmod 2^{\wedge}Suc\ n - 2^{\wedge}n :: int)$
 ⟨proof⟩

lemma *no-sbintr*:

$number-of\ (sbintrunc\ n\ w) =$
 $((number-of\ w + 2^{\wedge}n) \bmod 2^{\wedge}Suc\ n - 2^{\wedge}n :: int)$
 ⟨proof⟩

lemma *range-sbintrunc*:

$range\ (sbintrunc\ n) = \{i. - (2^{\wedge}n) <= i \ \& \ i < 2^{\wedge}n\}$
 ⟨proof⟩

lemma *sb-inc-lem*:

$(a::int) + 2^{\wedge}k < 0 \implies a + 2^{\wedge}k + 2^{\wedge}(Suc\ k) <= (a + 2^{\wedge}k) \bmod 2^{\wedge}(Suc\ k)$
 ⟨proof⟩

lemma *sb-inc-lem'*:

$(a::int) < - (2^{\wedge}k) \implies a + 2^{\wedge}k + 2^{\wedge}(Suc\ k) <= (a + 2^{\wedge}k) \bmod 2^{\wedge}(Suc\ k)$
 ⟨proof⟩

lemma *sbintrunc-inc*:

$x < - (2^{\wedge}n) \implies x + 2^{\wedge}(Suc\ n) <= sbintrunc\ n\ x$
 ⟨proof⟩

lemma *sb-dec-lem*:

$(0::int) <= - (2^{\wedge}k) + a \implies (a + 2^{\wedge}k) \bmod (2 * 2^{\wedge}k) <= - (2^{\wedge}k) + a$
 ⟨proof⟩

lemma *sb-dec-lem'*:

$(2::int)^{\wedge}k <= a \implies (a + 2^{\wedge}k) \bmod (2 * 2^{\wedge}k) <= - (2^{\wedge}k) + a$
 ⟨proof⟩

lemma *sbintrunc-dec*:

$x >= (2^{\wedge}n) \implies x - 2^{\wedge}(Suc\ n) >= sbintrunc\ n\ x$
 ⟨proof⟩

lemmas *zmod-uminus'* = *zmod-uminus* [where $b=c$, standard]

lemmas *zpower-zmod'* = *zpower-zmod* [where $m=c$ and $y=k$, standard]

lemmas *brdmod1s'* [symmetric] =

zmod-zadd-left-eq *zmod-zadd-right-eq*
zmod-zsub-left-eq *zmod-zsub-right-eq*
zmod-zmult1-eq *zmod-zmult1-eq-rev*

lemmas *brdmods'* [symmetric] =

zpower-zmod' [symmetric]
trans [OF *zmod-zadd-left-eq* *zmod-zadd-right-eq*]
trans [OF *zmod-zsub-left-eq* *zmod-zsub-right-eq*]
trans [OF *zmod-zmult1-eq* *zmod-zmult1-eq-rev*]

zmod-uminus' [symmetric]
zmod-zadd-left-eq [where $b = 1$]
zmod-zsub-left-eq [where $b = 1$]

lemmas *bintr-arith1s* =
brdmod1s' [where $c=2^n$, folded pred-def succ-def bintrunc-mod2p, standard]

lemmas *bintr-ariths* =
brdmods' [where $c=2^n$, folded pred-def succ-def bintrunc-mod2p, standard]

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [OF *zless2p*, standard]

lemma *bintr-ge0*: $(0 :: int) \leq \text{number-of } (\text{bintrunc } n \ w)$
 ⟨proof⟩

lemma *bintr-lt2p*: $\text{number-of } (\text{bintrunc } n \ w) < (2^n :: int)$
 ⟨proof⟩

lemma *bintr-Min*:
 $\text{number-of } (\text{bintrunc } n \ \text{Numeral.Min}) = (2^n :: int) - 1$
 ⟨proof⟩

lemma *sbintr-ge*: $(-(2^n :: int) \leq \text{number-of } (\text{sbintrunc } n \ w))$
 ⟨proof⟩

lemma *sbintr-lt*: $\text{number-of } (\text{sbintrunc } n \ w) < (2^n :: int)$
 ⟨proof⟩

lemma *bintrunc-Suc*:
 $\text{bintrunc } (\text{Suc } n) \ \text{bin} = \text{bintrunc } n \ (\text{bin-rest } \text{bin}) \ \text{BIT } \text{bin-last } \text{bin}$
 ⟨proof⟩

lemma *sign-Pls-ge-0*:
 $(\text{bin-sign } \text{bin} = \text{Numeral.Plus}) = (\text{number-of } \text{bin} \geq (0 :: int))$
 ⟨proof⟩

lemma *sign-Min-lt-0*:
 $(\text{bin-sign } \text{bin} = \text{Numeral.Min}) = (\text{number-of } \text{bin} < (0 :: int))$
 ⟨proof⟩

lemmas *sign-Min-neg* = *trans* [OF *sign-Min-lt-0 neg-def* [symmetric]]

lemma *bin-rest-trunc*:
 $!!\text{bin}. (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bintrunc } (n - 1) \ (\text{bin-rest } \text{bin})$
 ⟨proof⟩

lemma *bin-rest-power-trunc* [rule-format] :
 $(\text{bin-rest } ^k) \ (\text{bintrunc } n \ \text{bin}) =$
 $\text{bintrunc } (n - k) \ ((\text{bin-rest } ^k) \ \text{bin})$
 ⟨proof⟩

lemma *bin-rest-trunc-i*:

$$\text{bintrunc } n \text{ (bin-rest bin)} = \text{bin-rest (bintrunc (Suc n) bin)}$$

<proof>

lemma *bin-rest-strunc*:

$$\text{!!bin. bin-rest (sbintrunc (Suc n) bin)} = \text{sbintrunc } n \text{ (bin-rest bin)}$$

<proof>

lemma *bintrunc-rest [simp]*:

$$\text{!!bin. bintrunc } n \text{ (bin-rest (bintrunc n bin))} = \text{bin-rest (bintrunc n bin)}$$

<proof>

lemma *sbintrunc-rest [simp]*:

$$\text{!!bin. sbintrunc } n \text{ (bin-rest (sbintrunc n bin))} = \text{bin-rest (sbintrunc n bin)}$$

<proof>

lemma *bintrunc-rest'*:

$$\text{bintrunc } n \text{ o bin-rest o bintrunc } n = \text{bin-rest o bintrunc } n$$

<proof>

lemma *sbintrunc-rest'*:

$$\text{sbintrunc } n \text{ o bin-rest o sbintrunc } n = \text{bin-rest o sbintrunc } n$$

<proof>

lemma *rco-lem*:

$$f \circ g \circ f = g \circ f \implies f \circ (g \circ f) \wedge n = g \wedge n \circ f$$

<proof>

lemma *rco-alt*: $(f \circ g) \wedge n \circ f = f \circ (g \circ f) \wedge n$

<proof>

lemmas *rco-bintr = bintrunc-rest'*

[*THEN rco-lem [THEN fun-cong], unfolded o-def*]

lemmas *rco-sbintr = sbintrunc-rest'*

[*THEN rco-lem [THEN fun-cong], unfolded o-def*]

4.5 Splitting and concatenation

consts

bin-split :: $\text{nat} \Rightarrow \text{int} \Rightarrow \text{int} * \text{int}$

primrec

Z : *bin-split* 0 *w* = (*w*, *Numeral.Pls*)

Suc : *bin-split* (*Suc n*) *w* = (*let* (*w1*, *w2*) = *bin-split* *n* (*bin-rest* *w*)
in (*w1*, *w2 BIT bin-last* *w*))

consts

bin-cat :: $\text{int} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int}$

primrec

$Z : \text{bin-cat } w \ 0 \ v = w$
 $\text{Suc} : \text{bin-cat } w \ (\text{Suc } n) \ v = \text{bin-cat } w \ n \ (\text{bin-rest } v) \ \text{BIT } \text{bin-last } v$

4.6 Miscellaneous lemmas

lemmas *funpow-minus-simp* =
trans [OF gen-minus [where f = power f] funpow-Suc, standard]

lemmas *funpow-pred-simp [simp]* =
funpow-minus-simp [of number-of bin, simplified nobm1, standard]

lemmas *replicate-minus-simp* =
trans [OF gen-minus [where f = %n. replicate n x] replicate.replicate-Suc, standard]

lemmas *replicate-pred-simp [simp]* =
replicate-minus-simp [of number-of bin, simplified nobm1, standard]

lemmas *power-Suc-no [simp]* = *power-Suc [of number-of a, standard]*

lemmas *power-minus-simp* =
trans [OF gen-minus [where f = power f] power-Suc, standard]

lemmas *power-pred-simp* =
power-minus-simp [of number-of bin, simplified nobm1, standard]

lemmas *power-pred-simp-no [simp]* = *power-pred-simp [where f = number-of f, standard]*

lemma *list-exhaust-size-gt0*:
assumes $y : \bigwedge a \ \text{list}. y = a \ \# \ \text{list} \implies P$
shows $0 < \text{length } y \implies P$
<proof>

lemma *list-exhaust-size-eq0*:
assumes $y : y = [] \implies P$
shows $\text{length } y = 0 \implies P$
<proof>

lemma *size-Cons-lem-eq*:
 $y = xa \ \# \ \text{list} \implies \text{size } y = \text{Suc } k \implies \text{size } \text{list} = k$
<proof>

lemma *size-Cons-lem-eq-bin*:
 $y = xa \ \# \ \text{list} \implies \text{size } y = \text{number-of } (\text{Numeral.succ } k) \implies$
 $\text{size } \text{list} = \text{number-of } k$
<proof>

lemmas *ls-splits* =
prod.split split-split prod.split-asm split-split-asm split-if-asm

lemma *not-B1-is-B0*: $y \neq \text{bit}.B1 \implies y = \text{bit}.B0$
 ⟨*proof*⟩

lemma *B1-ass-B0*:
assumes $y: y = \text{bit}.B0 \implies y = \text{bit}.B1$
shows $y = \text{bit}.B1$
 ⟨*proof*⟩

lemmas *n2s-ths* [*THEN eq-reflection*] = *add-2-eq-Suc add-2-eq-Suc'*

lemmas *s2n-ths* = *n2s-ths* [*symmetric*]

end

5 BitSyntax: Syntactic class for bitwise operations

theory *BitSyntax*
imports *Main*
begin

class *bit* = *type* +
fixes *bitNOT* :: 'a ⇒ 'a
and *bitAND* :: 'a ⇒ 'a ⇒ 'a
and *bitOR* :: 'a ⇒ 'a ⇒ 'a
and *bitXOR* :: 'a ⇒ 'a ⇒ 'a

We want the bitwise operations to bind slightly weaker than + and −, but ~ to bind slightly stronger than *.

notation
bitNOT (*NOT* - [70] 71) **and**
bitAND (**infixr** *AND* 64) **and**
bitOR (**infixr** *OR* 59) **and**
bitXOR (**infixr** *XOR* 59)

Testing and shifting operations.

consts
test-bit :: 'a::bit ⇒ nat ⇒ bool (**infixl** !! 100)
lsb :: 'a::bit ⇒ bool
msb :: 'a::bit ⇒ bool
set-bit :: 'a::bit ⇒ nat ⇒ bool ⇒ 'a
set-bits :: (nat ⇒ bool) ⇒ 'a::bit (**binder** *BITS* 10)
shiffl :: 'a::bit ⇒ nat ⇒ 'a (**infixl** << 55)
shiftr :: 'a::bit ⇒ nat ⇒ 'a (**infixl** >> 55)

5.1 Bitwise operations on type *bit*

instance *bit* :: *bit*

NOT-bit-def: $NOT\ x \equiv \text{case } x \text{ of } bit.B0 \Rightarrow bit.B1 \mid bit.B1 \Rightarrow bit.B0$

AND-bit-def: $x\ AND\ y \equiv \text{case } x \text{ of } bit.B0 \Rightarrow bit.B0 \mid bit.B1 \Rightarrow y$

OR-bit-def: $x\ OR\ y :: bit \equiv NOT\ (NOT\ x\ AND\ NOT\ y)$

XOR-bit-def: $x\ XOR\ y :: bit \equiv (x\ AND\ NOT\ y)\ OR\ (NOT\ x\ AND\ y)$

<proof>

lemma *bit-simps* [*simp*]:

$NOT\ bit.B0 = bit.B1$

$NOT\ bit.B1 = bit.B0$

$bit.B0\ AND\ y = bit.B0$

$bit.B1\ AND\ y = y$

$bit.B0\ OR\ y = y$

$bit.B1\ OR\ y = bit.B1$

$bit.B0\ XOR\ y = y$

$bit.B1\ XOR\ y = NOT\ y$

<proof>

lemma *bit-extra-simps* [*simp*]:

$x\ AND\ bit.B0 = bit.B0$

$x\ AND\ bit.B1 = x$

$x\ OR\ bit.B1 = bit.B1$

$x\ OR\ bit.B0 = x$

$x\ XOR\ bit.B1 = NOT\ x$

$x\ XOR\ bit.B0 = x$

<proof>

lemma *bit-ops-comm*:

$(x::bit)\ AND\ y = y\ AND\ x$

$(x::bit)\ OR\ y = y\ OR\ x$

$(x::bit)\ XOR\ y = y\ XOR\ x$

<proof>

lemma *bit-ops-same* [*simp*]:

$(x::bit)\ AND\ x = x$

$(x::bit)\ OR\ x = x$

$(x::bit)\ XOR\ x = bit.B0$

<proof>

lemma *bit-not-not* [*simp*]: $NOT\ (NOT\ (x::bit)) = x$

<proof>

end

6 BinOperations: Bitwise Operations on Binary Integers

theory *BinOperations* **imports** *BinGeneral BitSyntax*

begin

6.1 Logical operations

bit-wise logical operations on the int type

instance *int* :: *bit*

int-not-def: $\text{bitNOT} \equiv \text{bin-rec Numeral.Min Numeral.Pls}$
 $(\lambda w b s. s \text{ BIT } (\text{NOT } b))$
int-and-def: $\text{bitAND} \equiv \text{bin-rec } (\lambda x. \text{Numeral.Pls}) (\lambda y. y)$
 $(\lambda w b s y. s (\text{bin-rest } y) \text{ BIT } (b \text{ AND } \text{bin-last } y))$
int-or-def: $\text{bitOR} \equiv \text{bin-rec } (\lambda x. x) (\lambda y. \text{Numeral.Min})$
 $(\lambda w b s y. s (\text{bin-rest } y) \text{ BIT } (b \text{ OR } \text{bin-last } y))$
int-xor-def: $\text{bitXOR} \equiv \text{bin-rec } (\lambda x. x) \text{ bitNOT}$
 $(\lambda w b s y. s (\text{bin-rest } y) \text{ BIT } (b \text{ XOR } \text{bin-last } y))$
 $\langle \text{proof} \rangle$

lemma *int-not-simps* [*simp*]:

$\text{NOT Numeral.Pls} = \text{Numeral.Min}$
 $\text{NOT Numeral.Min} = \text{Numeral.Pls}$
 $\text{NOT } (w \text{ BIT } b) = (\text{NOT } w) \text{ BIT } (\text{NOT } b)$
 $\langle \text{proof} \rangle$

lemma *int-xor-Pls* [*simp*]:

$\text{Numeral.Pls XOR } x = x$
 $\langle \text{proof} \rangle$

lemma *int-xor-Min* [*simp*]:

$\text{Numeral.Min XOR } x = \text{NOT } x$
 $\langle \text{proof} \rangle$

lemma *int-xor-Bits* [*simp*]:

$(x \text{ BIT } b) \text{ XOR } (y \text{ BIT } c) = (x \text{ XOR } y) \text{ BIT } (b \text{ XOR } c)$
 $\langle \text{proof} \rangle$

lemma *int-xor-x-simps'*:

$w \text{ XOR } (\text{Numeral.Pls BIT bit.B0}) = w$
 $w \text{ XOR } (\text{Numeral.Min BIT bit.B1}) = \text{NOT } w$
 $\langle \text{proof} \rangle$

lemmas *int-xor-extra-simps* [*simp*] = *int-xor-x-simps'* [*simplified arith-simps*]

lemma *int-or-Pls* [*simp*]:

$\text{Numeral.Pls OR } x = x$

<proof>

lemma *int-or-Min* [*simp*]:
Numeral.Min OR x = Numeral.Min
<proof>

lemma *int-or-Bits* [*simp*]:
 $(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \text{ OR } c)$
<proof>

lemma *int-or-x-simps'*:
 $w \text{ OR } (\text{Numeral.Plus BIT bit.B0}) = w$
 $w \text{ OR } (\text{Numeral.Min BIT bit.B1}) = \text{Numeral.Min}$
<proof>

lemmas *int-or-extra-simps* [*simp*] = *int-or-x-simps'* [*simplified arith-simps*]

lemma *int-and-Pls* [*simp*]:
Numeral.Plus AND x = Numeral.Plus
<proof>

lemma *int-and-Min* [*simp*]:
Numeral.Min AND x = x
<proof>

lemma *int-and-Bits* [*simp*]:
 $(x \text{ BIT } b) \text{ AND } (y \text{ BIT } c) = (x \text{ AND } y) \text{ BIT } (b \text{ AND } c)$
<proof>

lemma *int-and-x-simps'*:
 $w \text{ AND } (\text{Numeral.Plus BIT bit.B0}) = \text{Numeral.Plus}$
 $w \text{ AND } (\text{Numeral.Min BIT bit.B1}) = w$
<proof>

lemmas *int-and-extra-simps* [*simp*] = *int-and-x-simps'* [*simplified arith-simps*]

lemma *bin-ops-comm*:
shows
int-and-comm: !!y::int. x AND y = y AND x and
int-or-comm: !!y::int. x OR y = y OR x and
int-xor-comm: !!y::int. x XOR y = y XOR x
<proof>

lemma *bin-ops-same* [*simp*]:
 $(x::\text{int}) \text{ AND } x = x$
 $(x::\text{int}) \text{ OR } x = x$
 $(x::\text{int}) \text{ XOR } x = \text{Numeral.Plus}$

<proof>

lemma *int-not-not* [*simp*]: $NOT (NOT (x::int)) = x$
<proof>

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min

lemma *bbw-ao-absorb*:
 $!!y::int. x AND (y OR x) = x \& x OR (y AND x) = x$
<proof>

lemma *bbw-ao-absorbs-other*:
 $x AND (x OR y) = x \wedge (y AND x) OR x = (x::int)$
 $(y OR x) AND x = x \wedge x OR (x AND y) = (x::int)$
 $(x OR y) AND x = x \wedge (x AND y) OR x = (x::int)$
<proof>

lemmas *bbw-ao-absorbs* [*simp*] = *bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:
 $!!y::int. (NOT x) XOR y = NOT (x XOR y) \&$
 $x XOR (NOT y) = NOT (x XOR y)$
<proof>

lemma *bbw-assocs'*:
 $!!y z::int. (x AND y) AND z = x AND (y AND z) \&$
 $(x OR y) OR z = x OR (y OR z) \&$
 $(x XOR y) XOR z = x XOR (y XOR z)$
<proof>

lemma *int-and-assoc*:
 $(x AND y) AND (z::int) = x AND (y AND z)$
<proof>

lemma *int-or-assoc*:
 $(x OR y) OR (z::int) = x OR (y OR z)$
<proof>

lemma *int-xor-assoc*:
 $(x XOR y) XOR (z::int) = x XOR (y XOR z)$
<proof>

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [*simp*]:

$(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$
 $(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$
 $(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$
 ⟨proof⟩

lemma *bbw-not-dist*:

$!!y::int. \text{ NOT } (x \text{ OR } y) = (\text{ NOT } x) \text{ AND } (\text{ NOT } y)$
 $!!y::int. \text{ NOT } (x \text{ AND } y) = (\text{ NOT } x) \text{ OR } (\text{ NOT } y)$
 ⟨proof⟩

lemma *bbw-oa-dist*:

$!!y \ z::int. (x \text{ AND } y) \text{ OR } z =$
 $(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
 ⟨proof⟩

lemma *bbw-ao-dist*:

$!!y \ z::int. (x \text{ OR } y) \text{ AND } z =$
 $(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$
 ⟨proof⟩

lemma *plus-and-or* [rule-format]:

$ALL \ y::int. (x \text{ AND } y) + (x \text{ OR } y) = x + y$
 ⟨proof⟩

lemma *le-int-or*:

$!!x. \text{ bin-sign } y = \text{ Numeral.Pl } ==> x \leq x \text{ OR } y$
 ⟨proof⟩

lemmas *int-and-le* =

xtr3 [OF *bbw-ao-absorbs* (2) [THEN *conjunct2*, *symmetric*] *le-int-or*]

lemma *bin-nth-ops*:

$!!x \ y. \text{ bin-nth } (x \text{ AND } y) \ n = (\text{ bin-nth } x \ n \ \& \ \text{ bin-nth } y \ n)$
 $!!x \ y. \text{ bin-nth } (x \text{ OR } y) \ n = (\text{ bin-nth } x \ n \ | \ \text{ bin-nth } y \ n)$
 $!!x \ y. \text{ bin-nth } (x \text{ XOR } y) \ n = (\text{ bin-nth } x \ n \ \sim = \ \text{ bin-nth } y \ n)$
 $!!x. \text{ bin-nth } (\text{ NOT } x) \ n = (\sim \ \text{ bin-nth } x \ n)$
 ⟨proof⟩

lemma *bin-add-not*: $x + \text{ NOT } x = \text{ Numeral.Min}$

⟨proof⟩

lemma *bin-trunc-ao*:

$!!x \ y. (\text{ bintrunc } n \ x) \ \text{ AND } (\text{ bintrunc } n \ y) = \text{ bintrunc } n \ (x \ \text{ AND } \ y)$
 $!!x \ y. (\text{ bintrunc } n \ x) \ \text{ OR } (\text{ bintrunc } n \ y) = \text{ bintrunc } n \ (x \ \text{ OR } \ y)$

$\langle proof \rangle$

lemma *bin-trunc-xor*:

$!!x\ y. \text{bintrunc } n (\text{bintrunc } n\ x\ \text{XOR}\ \text{bintrunc } n\ y) =$
 $\text{bintrunc } n (x\ \text{XOR}\ y)$

$\langle proof \rangle$

lemma *bin-trunc-not*:

$!!x. \text{bintrunc } n (\text{NOT } (\text{bintrunc } n\ x)) = \text{bintrunc } n (\text{NOT } x)$

$\langle proof \rangle$

lemma *bintr-bintr-i*:

$x = \text{bintrunc } n\ y \implies \text{bintrunc } n\ x = \text{bintrunc } n\ y$

$\langle proof \rangle$

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

6.2 Setting and clearing bits

consts

bin-sc :: *nat* => *bit* => *int* => *int*

primrec

Z : *bin-sc* 0 *b w* = *bin-rest w BIT b*

Suc :

bin-sc (*Suc n*) *b w* = *bin-sc n b (bin-rest w) BIT bin-last w*

lemma *bin-nth-sc* [*simp*]:

$!!w. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ n = (b = \text{bit.B1})$

$\langle proof \rangle$

lemma *bin-sc-sc-same* [*simp*]:

$!!w. \text{bin-sc } n\ c (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ c\ w$

$\langle proof \rangle$

lemma *bin-sc-sc-diff*:

$!!w\ m. m \sim n \implies$

$\text{bin-sc } m\ c (\text{bin-sc } n\ b\ w) = \text{bin-sc } n\ b (\text{bin-sc } m\ c\ w)$

$\langle proof \rangle$

lemma *bin-nth-sc-gen*:

$!!w\ m. \text{bin-nth } (\text{bin-sc } n\ b\ w)\ m = (\text{if } m = n \text{ then } b = \text{bit.B1} \text{ else } \text{bin-nth } w\ m)$

$\langle proof \rangle$

lemma *bin-sc-nth* [*simp*]:

!!w. (bin-sc n (If (bin-nth w n) bit.B1 bit.B0) w) = w
 ⟨proof⟩

lemma bin-sign-sc [simp]:
 !!w. bin-sign (bin-sc n b w) = bin-sign w
 ⟨proof⟩

lemma bin-sc-bintr [simp]:
 !!w m. bintrunc m (bin-sc n x (bintrunc m (w))) = bintrunc m (bin-sc n x w)
 ⟨proof⟩

lemma bin-clr-le:
 !!w. bin-sc n bit.B0 w <= w
 ⟨proof⟩

lemma bin-set-ge:
 !!w. bin-sc n bit.B1 w >= w
 ⟨proof⟩

lemma bintr-bin-clr-le:
 !!w m. bintrunc n (bin-sc m bit.B0 w) <= bintrunc n w
 ⟨proof⟩

lemma bintr-bin-set-ge:
 !!w m. bintrunc n (bin-sc m bit.B1 w) >= bintrunc n w
 ⟨proof⟩

lemma bin-sc-FP [simp]: bin-sc n bit.B0 Numeral.Pls = Numeral.Pls
 ⟨proof⟩

lemma bin-sc-TM [simp]: bin-sc n bit.B1 Numeral.Min = Numeral.Min
 ⟨proof⟩

lemmas bin-sc-simps = bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP

lemma bin-sc-minus:
 0 < n ==> bin-sc (Suc (n - 1)) b w = bin-sc n b w
 ⟨proof⟩

lemmas bin-sc-Suc-minus =
 trans [OF bin-sc-minus [symmetric] bin-sc.Suc, standard]

lemmas bin-sc-Suc-pred [simp] =
 bin-sc-Suc-minus [of number-of bin, simplified nobm1, standard]

6.3 Operations on lists of booleans

consts

bin-to-bl :: nat => int => bool list

$bin\text{-}to\text{-}bl\text{-}aux :: nat \Rightarrow int \Rightarrow bool\ list \Rightarrow bool\ list$
 $bl\text{-}to\text{-}bin :: bool\ list \Rightarrow int$
 $bl\text{-}to\text{-}bin\text{-}aux :: int \Rightarrow bool\ list \Rightarrow int$

$bl\text{-}of\text{-}nth :: nat \Rightarrow (nat \Rightarrow bool) \Rightarrow bool\ list$

primrec

$Nil : bl\text{-}to\text{-}bin\text{-}aux\ w\ [] = w$
 $Cons : bl\text{-}to\text{-}bin\text{-}aux\ w\ (b \# bs) =$
 $bl\text{-}to\text{-}bin\text{-}aux\ (w\ BIT\ (if\ b\ then\ bit.B1\ else\ bit.B0))\ bs$

primrec

$Z : bin\text{-}to\text{-}bl\text{-}aux\ 0\ w\ bl = bl$
 $Suc : bin\text{-}to\text{-}bl\text{-}aux\ (Suc\ n)\ w\ bl =$
 $bin\text{-}to\text{-}bl\text{-}aux\ n\ (bin\text{-}rest\ w)\ ((bin\text{-}last\ w = bit.B1) \# bl)$

defs

$bin\text{-}to\text{-}bl\text{-}def : bin\text{-}to\text{-}bl\ n\ w == bin\text{-}to\text{-}bl\text{-}aux\ n\ w\ []$
 $bl\text{-}to\text{-}bin\text{-}def : bl\text{-}to\text{-}bin\ bs == bl\text{-}to\text{-}bin\text{-}aux\ Numeral.Pls\ bs$

primrec

$Suc : bl\text{-}of\text{-}nth\ (Suc\ n)\ f = f\ n \# bl\text{-}of\text{-}nth\ n\ f$
 $Z : bl\text{-}of\text{-}nth\ 0\ f = []$

consts

$takefill :: 'a \Rightarrow nat \Rightarrow 'a\ list \Rightarrow 'a\ list$
 $app2 :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a\ list \Rightarrow 'b\ list \Rightarrow 'c\ list$

— takefill - like take but if argument list too short,
— extends result to get requested length

primrec

$Z : takefill\ fill\ 0\ xs = []$
 $Suc : takefill\ fill\ (Suc\ n)\ xs = ($
 $case\ xs\ of\ [] \Rightarrow fill \# takefill\ fill\ n\ xs$
 $| y \# ys \Rightarrow y \# takefill\ fill\ n\ ys)$

defs

$app2\text{-}def : app2\ f\ as\ bs == map\ (split\ f)\ (zip\ as\ bs)$

6.4 Splitting and concatenation

— rcat and rsplit

consts

$bin\text{-}rcat :: nat \Rightarrow int\ list \Rightarrow int$
 $bin\text{-}rsplit\text{-}aux :: nat * int\ list * nat * int \Rightarrow int\ list$
 $bin\text{-}rsplit :: nat \Rightarrow (nat * int) \Rightarrow int\ list$
 $bin\text{-}rsplitl\text{-}aux :: nat * int\ list * nat * int \Rightarrow int\ list$
 $bin\text{-}rsplitl :: nat \Rightarrow (nat * int) \Rightarrow int\ list$

```

recdef bin-rsplit-aux measure (fst o snd o snd)
  bin-rsplit-aux (n, bs, (m, c)) =
    (if m = 0 | n = 0 then bs else
     let (a, b) = bin-split n c
     in bin-rsplit-aux (n, b # bs, (m - n, a)))

recdef bin-rsplittl-aux measure (fst o snd o snd)
  bin-rsplittl-aux (n, bs, (m, c)) =
    (if m = 0 | n = 0 then bs else
     let (a, b) = bin-split (min m n) c
     in bin-rsplittl-aux (n, b # bs, (m - n, a)))

defs
  bin-rcat-def : bin-rcat n bs == foldl (%u v. bin-cat u n v) Numeral.Pls bs
  bin-rsplit-def : bin-rsplit n w == bin-rsplit-aux (n, [], w)
  bin-rsplittl-def : bin-rsplittl n w == bin-rsplittl-aux (n, [], w)

declare bin-rsplit-aux.simps [simp del]
declare bin-rsplittl-aux.simps [simp del]

lemma bin-sign-cat:
  !!y. bin-sign (bin-cat x n y) = bin-sign x
  <proof>

lemma bin-cat-Suc-Bit:
  bin-cat w (Suc n) (v BIT b) = bin-cat w n v BIT b
  <proof>

lemma bin-nth-cat:
  !!n y. bin-nth (bin-cat x k y) n =
    (if n < k then bin-nth y n else bin-nth x (n - k))
  <proof>

lemma bin-nth-split:
  !!b c. bin-split n c = (a, b) ==>
    (ALL k. bin-nth a k = bin-nth c (n + k)) &
    (ALL k. bin-nth b k = (k < n & bin-nth c k))
  <proof>

lemma bin-cat-assoc:
  !!z. bin-cat (bin-cat x m y) n z = bin-cat x (m + n) (bin-cat y n z)
  <proof>

lemma bin-cat-assoc-sym: !!z m.
  bin-cat x m (bin-cat y n z) = bin-cat (bin-cat x (m - n) y) (min m n) z
  <proof>

```

lemma *bin-cat-Pls* [simp]:

!!w. *bin-cat Numeral.Pls* n w = *bintrunc* n w

⟨proof⟩

lemma *bintr-cat1*:

!!b. *bintrunc* (k + n) (*bin-cat* a n b) = *bin-cat* (*bintrunc* k a) n b

⟨proof⟩

lemma *bintr-cat*: *bintrunc* m (*bin-cat* a n b) =

bin-cat (*bintrunc* (m - n) a) n (*bintrunc* (min m n) b)

⟨proof⟩

lemma *bintr-cat-same* [simp]:

bintrunc n (*bin-cat* a n b) = *bintrunc* n b

⟨proof⟩

lemma *cat-bintr* [simp]:

!!b. *bin-cat* a n (*bintrunc* n b) = *bin-cat* a n b

⟨proof⟩

lemma *split-bintrunc*:

!!b c. *bin-split* n c = (a, b) ==> b = *bintrunc* n c

⟨proof⟩

lemma *bin-cat-split*:

!!v w. *bin-split* n w = (u, v) ==> w = *bin-cat* u n v

⟨proof⟩

lemma *bin-split-cat*:

!!w. *bin-split* n (*bin-cat* v n w) = (v, *bintrunc* n w)

⟨proof⟩

lemma *bin-split-Pls* [simp]:

bin-split n *Numeral.Pls* = (*Numeral.Pls*, *Numeral.Pls*)

⟨proof⟩

lemma *bin-split-Min* [simp]:

bin-split n *Numeral.Min* = (*Numeral.Min*, *bintrunc* n *Numeral.Min*)

⟨proof⟩

lemma *bin-split-trunc*:

!!m b c. *bin-split* (min m n) c = (a, b) ==>

bin-split n (*bintrunc* m c) = (*bintrunc* (m - n) a, b)

⟨proof⟩

lemma *bin-split-trunc1*:

!!m b c. *bin-split* n c = (a, b) ==>

bin-split n (*bintrunc* m c) = (*bintrunc* (m - n) a, *bintrunc* m b)

⟨proof⟩

lemma *bin-cat-num*:

!!*b*. *bin-cat* *a n b* = *a* * $2^{\wedge} n$ + *bintrunc* *n b*
 ⟨*proof*⟩

lemma *bin-split-num*:

!!*b*. *bin-split* *n b* = (*b div* $2^{\wedge} n$, *b mod* $2^{\wedge} n$)
 ⟨*proof*⟩

6.5 Miscellaneous lemmas

lemma *nth-2p-bin*:

!!*m*. *bin-nth* ($2^{\wedge} n$) *m* = (*m = n*)
 ⟨*proof*⟩

lemma *ex-eq-or*:

(*EX m*. *n = Suc m* & (*m = k* | *P m*)) = (*n = Suc k* | (*EX m*. *n = Suc m* & *P m*))
 ⟨*proof*⟩

end

7 BinBoolList: Bool lists and integers

theory *BinBoolList* **imports** *BinOperations* **begin**

7.1 Arithmetic in terms of bool lists

consts

rbl-succ :: *bool list* => *bool list*
rbl-pred :: *bool list* => *bool list*
rbl-add :: *bool list* => *bool list* => *bool list*
rbl-mult :: *bool list* => *bool list* => *bool list*

primrec

Nil: *rbl-succ Nil* = *Nil*
Cons: *rbl-succ* (*x # xs*) = (*if x then False # rbl-succ xs else True # xs*)

primrec

Nil : *rbl-pred Nil* = *Nil*
Cons : *rbl-pred* (*x # xs*) = (*if x then False # xs else True # rbl-pred xs*)

primrec

Nil : *rbl-add Nil x* = *Nil*
Cons : *rbl-add* (*y # ys*) *x* = (*let ws = rbl-add ys (tl x) in*

$(y \sim = hd\ x) \# (if\ hd\ x \ \&\ y\ then\ rbl\ succ\ ws\ else\ ws))$

primrec

$Nil : rbl\ mult\ Nil\ x = Nil$

$Cons : rbl\ mult\ (y \# ys)\ x = (let\ ws = False \# rbl\ mult\ ys\ x\ in\ if\ y\ then\ rbl\ add\ ws\ x\ else\ ws)$

lemma *tl-take*: $tl\ (take\ n\ l) = take\ (n - 1)\ (tl\ l)$
 $\langle proof \rangle$

lemma *take-butlast* [rule-format] :
 $ALL\ n.\ n < length\ l \ \dashrightarrow\ take\ n\ (butlast\ l) = take\ n\ l$
 $\langle proof \rangle$

lemma *butlast-take* [rule-format] :
 $ALL\ n.\ n \leq length\ l \ \dashrightarrow\ butlast\ (take\ n\ l) = take\ (n - 1)\ l$
 $\langle proof \rangle$

lemma *butlast-drop* [rule-format] :
 $ALL\ n.\ butlast\ (drop\ n\ l) = drop\ n\ (butlast\ l)$
 $\langle proof \rangle$

lemma *butlast-power*:
 $(butlast\ ^\ n)\ bl = take\ (length\ bl - n)\ bl$
 $\langle proof \rangle$

lemma *bin-to-bl-aux-Pls-minus-simp*:
 $0 < n \ \Longrightarrow\ bin\ to\ bl\ aux\ n\ Numeral.Pls\ bl =$
 $bin\ to\ bl\ aux\ (n - 1)\ Numeral.Pls\ (False \# bl)$
 $\langle proof \rangle$

lemma *bin-to-bl-aux-Min-minus-simp*:
 $0 < n \ \Longrightarrow\ bin\ to\ bl\ aux\ n\ Numeral.Min\ bl =$
 $bin\ to\ bl\ aux\ (n - 1)\ Numeral.Min\ (True \# bl)$
 $\langle proof \rangle$

lemma *bin-to-bl-aux-Bit-minus-simp*:
 $0 < n \ \Longrightarrow\ bin\ to\ bl\ aux\ n\ (w\ BIT\ b)\ bl =$
 $bin\ to\ bl\ aux\ (n - 1)\ w\ ((b = bit.B1) \# bl)$
 $\langle proof \rangle$

declare *bin-to-bl-aux-Pls-minus-simp* [simp]
bin-to-bl-aux-Min-minus-simp [simp]
bin-to-bl-aux-Bit-minus-simp [simp]

lemma *bl-to-bin-aux-append* [rule-format] :
 $ALL\ w.\ bl\ to\ bin\ aux\ w\ (bs\ @\ cs) = bl\ to\ bin\ aux\ (bl\ to\ bin\ aux\ w\ bs)\ cs$

<proof>

lemma *bin-to-bl-aux-append* [rule-format] :

ALL w bs. bin-to-bl-aux n w bs @ cs = bin-to-bl-aux n w (bs @ cs)

<proof>

lemma *bl-to-bin-append*:

bl-to-bin (bs @ cs) = bl-to-bin-aux (bl-to-bin bs) cs

<proof>

lemma *bin-to-bl-aux-alt*:

bin-to-bl-aux n w bs = bin-to-bl n w @ bs

<proof>

lemma *bin-to-bl-0*: *bin-to-bl 0 bs = []*

<proof>

lemma *size-bin-to-bl-aux* [rule-format] :

ALL w bs. size (bin-to-bl-aux n w bs) = n + length bs

<proof>

lemma *size-bin-to-bl*: *size (bin-to-bl n w) = n*

<proof>

lemma *bin-bl-bin'* [rule-format] :

ALL w bs. bl-to-bin (bin-to-bl-aux n w bs) =

bl-to-bin-aux (bintrunc n w) bs

<proof>

lemma *bin-bl-bin*: *bl-to-bin (bin-to-bl n w) = bintrunc n w*

<proof>

lemma *bl-bin-bl'* [rule-format] :

ALL w n. bin-to-bl (n + length bs) (bl-to-bin-aux w bs) =

bin-to-bl-aux n w bs

<proof>

lemma *bl-bin-bl*: *bin-to-bl (length bs) (bl-to-bin bs) = bs*

<proof>

declare

bin-to-bl-0 [simp]

size-bin-to-bl [simp]

bin-bl-bin [simp]

bl-bin-bl [simp]

lemma *bl-to-bin-inj*:

bl-to-bin bs = bl-to-bin cs ==> length bs = length cs ==> bs = cs

<proof>

lemma *bl-to-bin-False*: $bl\text{-to}\text{-bin} (False \# bl) = bl\text{-to}\text{-bin} bl$
 ⟨proof⟩

lemma *bl-to-bin-Nil*: $bl\text{-to}\text{-bin} [] = Numeral.Pls$
 ⟨proof⟩

lemma *bin-to-bl-Pls-aux* [rule-format] :
 ALL *bl*. $bin\text{-to}\text{-bl}\text{-aux} n Numeral.Pls bl = replicate\ n\ False @ bl$
 ⟨proof⟩

lemma *bin-to-bl-Pls*: $bin\text{-to}\text{-bl} n Numeral.Pls = replicate\ n\ False$
 ⟨proof⟩

lemma *bin-to-bl-Min-aux* [rule-format] :
 ALL *bl*. $bin\text{-to}\text{-bl}\text{-aux} n Numeral.Min bl = replicate\ n\ True @ bl$
 ⟨proof⟩

lemma *bin-to-bl-Min*: $bin\text{-to}\text{-bl} n Numeral.Min = replicate\ n\ True$
 ⟨proof⟩

lemma *bl-to-bin-rep-F*:
 $bl\text{-to}\text{-bin} (replicate\ n\ False @ bl) = bl\text{-to}\text{-bin} bl$
 ⟨proof⟩

lemma *bin-to-bl-trunc*:
 $n \leq m \implies bin\text{-to}\text{-bl} n (bintrunc\ m\ w) = bin\text{-to}\text{-bl} n\ w$
 ⟨proof⟩

declare
bin-to-bl-trunc [simp]
bl-to-bin-False [simp]
bl-to-bin-Nil [simp]

lemma *bin-to-bl-aux-bintr* [rule-format] :
 ALL *m bin bl*. $bin\text{-to}\text{-bl}\text{-aux} n (bintrunc\ m\ bin) bl =$
 $replicate\ (n - m)\ False @ bin\text{-to}\text{-bl}\text{-aux} (min\ n\ m) bin bl$
 ⟨proof⟩

lemmas *bin-to-bl-bintr* =
bin-to-bl-aux-bintr [where *bl* = [], folded *bin-to-bl-def*]

lemma *bl-to-bin-rep-False*: $bl\text{-to}\text{-bin} (replicate\ n\ False) = Numeral.Pls$
 ⟨proof⟩

lemma *len-bin-to-bl-aux* [rule-format] :
 ALL *w bs*. $length (bin\text{-to}\text{-bl}\text{-aux} n\ w\ bs) = n + length\ bs$
 ⟨proof⟩

lemma *len-bin-to-bl* [*simp*]: $\text{length } (\text{bin-to-bl } n \ w) = n$
 ⟨*proof*⟩

lemma *sign-bl-bin'* [*rule-format*] :
 ALL w . $\text{bin-sign } (\text{bl-to-bin-aux } w \ bs) = \text{bin-sign } w$
 ⟨*proof*⟩

lemma *sign-bl-bin*: $\text{bin-sign } (\text{bl-to-bin } bs) = \text{Numeral.Pl}s$
 ⟨*proof*⟩

lemma *bl-sbin-sign-aux* [*rule-format*] :
 ALL $w \ bs$. $\text{hd } (\text{bin-to-bl-aux } (\text{Suc } n) \ w \ bs) =$
 $(\text{bin-sign } (\text{sbintrunc } n \ w) = \text{Numeral.Min})$
 ⟨*proof*⟩

lemma *bl-sbin-sign*:
 $\text{hd } (\text{bin-to-bl } (\text{Suc } n) \ w) = (\text{bin-sign } (\text{sbintrunc } n \ w) = \text{Numeral.Min})$
 ⟨*proof*⟩

lemma *bin-nth-of-bl-aux* [*rule-format*] :
 ALL w . $\text{bin-nth } (\text{bl-to-bin-aux } w \ bl) \ n =$
 $(n < \text{size } bl \ \& \ \text{rev } bl \ ! \ n \ | \ n >= \text{length } bl \ \& \ \text{bin-nth } w \ (n - \text{size } bl))$
 ⟨*proof*⟩

lemma *bin-nth-of-bl*: $\text{bin-nth } (\text{bl-to-bin } bl) \ n = (n < \text{length } bl \ \& \ \text{rev } bl \ ! \ n)$
 ⟨*proof*⟩

lemma *bin-nth-bl* [*rule-format*] : ALL $m \ w$. $n < m \ \dashrightarrow$
 $\text{bin-nth } w \ n = \text{nth } (\text{rev } (\text{bin-to-bl } m \ w)) \ n$
 ⟨*proof*⟩

lemma *nth-rev* [*rule-format*] :
 $n < \text{length } xs \ \dashrightarrow \ \text{rev } xs \ ! \ n = xs \ ! \ (\text{length } xs - 1 - n)$
 ⟨*proof*⟩

lemmas *nth-rev-alt* = *nth-rev* [**where** $xs = \text{rev } ys$, *simplified*, *standard*]

lemma *nth-bin-to-bl-aux* [*rule-format*] :
 ALL $w \ n \ bl$. $n < m + \text{length } bl \ \dashrightarrow \ (\text{bin-to-bl-aux } m \ w \ bl) \ ! \ n =$
 $(\text{if } n < m \ \text{then } \text{bin-nth } w \ (m - 1 - n) \ \text{else } bl \ ! \ (n - m))$
 ⟨*proof*⟩

lemma *nth-bin-to-bl*: $n < m \ \Longrightarrow \ (\text{bin-to-bl } m \ w) \ ! \ n = \text{bin-nth } w \ (m - \text{Suc } n)$
 ⟨*proof*⟩

lemma *bl-to-bin-lt2p-aux* [*rule-format*] :
 ALL w . $\text{bl-to-bin-aux } w \ bs < (w + 1) * (2 \ ^ \ \text{length } bs)$
 ⟨*proof*⟩

lemma *bl-to-bin-lt2p*: $bl\text{-to-bin } bs < (2 \wedge \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge2p-aux* [rule-format] :
 ALL w . $bl\text{-to-bin-aux } w \text{ } bs \geq w * (2 \wedge \text{length } bs)$
 ⟨proof⟩

lemma *bl-to-bin-ge0*: $bl\text{-to-bin } bs \geq 0$
 ⟨proof⟩

lemma *butlast-rest-bin*:
 $butlast (bin\text{-to-bl } n \ w) = bin\text{-to-bl } (n - 1) (bin\text{-rest } w)$
 ⟨proof⟩

lemmas *butlast-bin-rest = butlast-rest-bin*
 [where $w = bl\text{-to-bin } bl$ and $n = \text{length } bl$, simplified, standard]

lemma *butlast-rest-bl2bin-aux* [rule-format] :
 ALL w . $bl \sim = [] \text{ } \dashrightarrow$
 $bl\text{-to-bin-aux } w (butlast \ bl) = bin\text{-rest } (bl\text{-to-bin-aux } w \ bl)$
 ⟨proof⟩

lemma *butlast-rest-bl2bin*:
 $bl\text{-to-bin } (butlast \ bl) = bin\text{-rest } (bl\text{-to-bin } bl)$
 ⟨proof⟩

lemma *trunc-bl2bin-aux* [rule-format] :
 ALL w . $bintrunc \ m (bl\text{-to-bin-aux } w \ bl) =$
 $bl\text{-to-bin-aux } (bintrunc \ (m - \text{length } bl) \ w) (drop \ (\text{length } bl - m) \ bl)$
 ⟨proof⟩

lemma *trunc-bl2bin*:
 $bintrunc \ m (bl\text{-to-bin } bl) = bl\text{-to-bin } (drop \ (\text{length } bl - m) \ bl)$
 ⟨proof⟩

lemmas *trunc-bl2bin-len* [simp] =
 $trunc\text{-bl2bin} [of \ \text{length } bl \ bl, \ \text{simplified}, \ \text{standard}]$

lemma *bl2bin-drop*:
 $bl\text{-to-bin } (drop \ k \ bl) = bintrunc \ (\text{length } bl - k) (bl\text{-to-bin } bl)$
 ⟨proof⟩

lemma *nth-rest-power-bin* [rule-format] :
 ALL n . $bin\text{-nth } ((bin\text{-rest } \wedge \ k) \ w) \ n = bin\text{-nth } w \ (n + k)$
 ⟨proof⟩

lemma *take-rest-power-bin*:
 $m \leq n \implies take \ m (bin\text{-to-bl } n \ w) = bin\text{-to-bl } m ((bin\text{-rest } \wedge \ (n - m)) \ w)$
 ⟨proof⟩

lemma *hd-butlast*: $\text{size } xs > 1 \implies \text{hd } (\text{butlast } xs) = \text{hd } xs$
 ⟨proof⟩

lemma *last-bin-last'* [rule-format] :
 ALL w . $\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last } (\text{bl-to-bin-aux } w \ xs) = \text{bit.B1})$
 ⟨proof⟩

lemma *last-bin-last*:
 $\text{size } xs > 0 \implies \text{last } xs = (\text{bin-last } (\text{bl-to-bin } xs) = \text{bit.B1})$
 ⟨proof⟩

lemma *bin-last-last*:
 $\text{bin-last } w = (\text{if } \text{last } (\text{bin-to-bl } (\text{Suc } n) \ w) \ \text{then } \text{bit.B1} \ \text{else } \text{bit.B0})$
 ⟨proof⟩

lemma *app2-Nil* [simp]: $\text{app2 } f \ [] \ ys = []$
 ⟨proof⟩

lemma *app2-Cons* [simp]:
 $\text{app2 } f \ (x \# \ xs) \ (y \# \ ys) = f \ x \ y \# \ \text{app2 } f \ xs \ ys$
 ⟨proof⟩

lemma *bl-xor-aux-bin* [rule-format] : ALL $v \ w \ bs \ cs$.
 $\text{app2 } (\%x \ y. \ x \ \sim = \ y) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ \text{XOR } w) \ (\text{app2 } (\%x \ y. \ x \ \sim = \ y) \ bs \ cs)$
 ⟨proof⟩

lemma *bl-or-aux-bin* [rule-format] : ALL $v \ w \ bs \ cs$.
 $\text{app2 } (\text{op } |) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ \text{OR } w) \ (\text{app2 } (\text{op } |) \ bs \ cs)$
 ⟨proof⟩

lemma *bl-and-aux-bin* [rule-format] : ALL $v \ w \ bs \ cs$.
 $\text{app2 } (\text{op } \&) \ (\text{bin-to-bl-aux } n \ v \ bs) \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (v \ \text{AND } w) \ (\text{app2 } (\text{op } \&) \ bs \ cs)$
 ⟨proof⟩

lemma *bl-not-aux-bin* [rule-format] :
 ALL $w \ cs$. $\text{map } \text{Not} \ (\text{bin-to-bl-aux } n \ w \ cs) =$
 $\text{bin-to-bl-aux } n \ (\text{NOT } w) \ (\text{map } \text{Not} \ cs)$
 ⟨proof⟩

lemmas *bl-not-bin = bl-not-aux-bin*
 [where $cs = []$, *unfolded bin-to-bl-def* [symmetric] *map.simps*]

lemmas *bl-and-bin = bl-and-aux-bin* [where $bs=[]$ and $cs=[]$,

unfolded app2-Nil, folded bin-to-bl-def]

lemmas *bl-or-bin* = *bl-or-aux-bin* [**where** *bs*=[] **and** *cs*=[],
unfolded app2-Nil, folded bin-to-bl-def]

lemmas *bl-xor-bin* = *bl-xor-aux-bin* [**where** *bs*=[] **and** *cs*=[],
unfolded app2-Nil, folded bin-to-bl-def]

lemma *drop-bin2bl-aux* [*rule-format*] :
ALL m bin bs. drop m (bin-to-bl-aux n bin bs) =
bin-to-bl-aux (n - m) bin (drop (m - n) bs)
<proof>

lemma *drop-bin2bl*: *drop m (bin-to-bl n bin) = bin-to-bl (n - m) bin*
<proof>

lemma *take-bin2bl-lem1* [*rule-format*] :
ALL w bs. take m (bin-to-bl-aux m w bs) = bin-to-bl m w
<proof>

lemma *take-bin2bl-lem* [*rule-format*] :
ALL w bs. take m (bin-to-bl-aux (m + n) w bs) =
take m (bin-to-bl (m + n) w)
<proof>

lemma *bin-split-take* [*rule-format*] :
ALL b c. bin-split n c = (a, b) -->
bin-to-bl m a = take m (bin-to-bl (m + n) c)
<proof>

lemma *bin-split-take1*:
k = m + n ==> bin-split n c = (a, b) ==>
bin-to-bl m a = take m (bin-to-bl k c)
<proof>

lemma *nth-takefill* [*rule-format*] : *ALL m l. m < n -->*
takefill fill n l ! m = (if m < length l then l ! m else fill)
<proof>

lemma *takefill-alt* [*rule-format*] :
ALL l. takefill fill n l = take n l @ replicate (n - length l) fill
<proof>

lemma *takefill-replicate* [*simp*]:
takefill fill n (replicate m fill) = replicate n fill
<proof>

lemma *takefill-le'* [*rule-format*] :
ALL l n. n = m + k --> takefill x m (takefill x n l) = takefill x m l

<proof>

lemma *length-takefill* [*simp*]: $\text{length } (\text{takefill } \text{fill } n \ l) = n$
<proof>

lemma *take-takefill'*:
 $!!w \ n. \ n = k + m \implies \text{take } k \ (\text{takefill } \text{fill } n \ w) = \text{takefill } \text{fill } k \ w$
<proof>

lemma *drop-takefill*:
 $!!w. \ \text{drop } k \ (\text{takefill } \text{fill } (m + k) \ w) = \text{takefill } \text{fill } m \ (\text{drop } k \ w)$
<proof>

lemma *takefill-le* [*simp*]:
 $m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
<proof>

lemma *take-takefill* [*simp*]:
 $m \leq n \implies \text{take } m \ (\text{takefill } \text{fill } n \ w) = \text{takefill } \text{fill } m \ w$
<proof>

lemma *takefill-append*:
 $\text{takefill } \text{fill } (m + \text{length } xs) \ (xs \ @ \ w) = xs \ @ \ (\text{takefill } \text{fill } m \ w)$
<proof>

lemma *takefill-same'*:
 $l = \text{length } xs \implies \text{takefill } \text{fill } l \ xs = xs$
<proof>

lemmas *takefill-same* [*simp*] = *takefill-same'* [*OF refl*]

lemma *takefill-bintrunc*:
 $\text{takefill } \text{False } n \ bl = \text{rev } (\text{bin-to-bl } n \ (\text{bl-to-bin } (\text{rev } bl)))$
<proof>

lemma *bl-bin-bl-rtf*:
 $\text{bin-to-bl } n \ (\text{bl-to-bin } bl) = \text{rev } (\text{takefill } \text{False } n \ (\text{rev } bl))$
<proof>

lemmas *bl-bin-bl-rep-drop* =
bl-bin-bl-rtf [*simplified takefill-alt*,
simplified, simplified rev-take, simplified]

lemma *tf-rev*:
 $n + k = m + \text{length } bl \implies \text{takefill } x \ m \ (\text{rev } (\text{takefill } y \ n \ bl)) =$
 $\text{rev } (\text{takefill } y \ m \ (\text{rev } (\text{takefill } x \ k \ (\text{rev } bl))))$
<proof>

lemma *takefill-minus*:

$0 < n \implies \text{takefill fill (Suc (n - 1)) w} = \text{takefill fill n w}$
 ⟨proof⟩

lemmas *takefill-Suc-cases* =
list.cases [THEN takefill.Suc [THEN trans], standard]

lemmas *takefill-Suc-Nil* = *takefill-Suc-cases (1)*
lemmas *takefill-Suc-Cons* = *takefill-Suc-cases (2)*

lemmas *takefill-minus-simps* = *takefill-Suc-cases [THEN [2]*
takefill-minus [symmetric, THEN trans], standard]

lemmas *takefill-pred-simps* [*simp*] =
takefill-minus-simps [where n=number-of bin, simplified nobm1, standard]

lemma *bl-to-bin-aux-cat*:
 !!*nv v. bl-to-bin-aux (bin-cat w nv v) bs* =
bin-cat w (nv + length bs) (bl-to-bin-aux v bs)
 ⟨proof⟩

lemma *bin-to-bl-aux-cat*:
 !!*w bs. bin-to-bl-aux (nv + nw) (bin-cat v nw w) bs* =
bin-to-bl-aux nv v (bin-to-bl-aux nw w bs)
 ⟨proof⟩

lemmas *bl-to-bin-aux-alt* =
bl-to-bin-aux-cat [where nv = 0 and v = Numeral.Pls,
simplified bl-to-bin-def [symmetric], simplified]

lemmas *bin-to-bl-cat* =
bin-to-bl-aux-cat [where bs = [], folded bin-to-bl-def]

lemmas *bl-to-bin-aux-app-cat* =
trans [OF bl-to-bin-aux-append bl-to-bin-aux-alt]

lemmas *bin-to-bl-aux-cat-app* =
trans [OF bin-to-bl-aux-cat bin-to-bl-aux-alt]

lemmas *bl-to-bin-app-cat* = *bl-to-bin-aux-app-cat*
 [*where w = Numeral.Pls, folded bl-to-bin-def*]

lemmas *bin-to-bl-cat-app* = *bin-to-bl-aux-cat-app*
 [*where bs = [], folded bin-to-bl-def*]

lemma *bl-to-bin-app-cat-alt*:
bin-cat (bl-to-bin cs) n w = bl-to-bin (cs @ bin-to-bl n w)

<proof>

lemma *mask-lem*: $(bl\text{-to-bin } (True \# replicate\ n\ False)) =$
 $Numeral.succ\ (bl\text{-to-bin } (replicate\ n\ True))$
<proof>

lemma *length-bl-of-nth* [*simp*]: $length\ (bl\text{-of-nth } n\ f) = n$
<proof>

lemma *nth-bl-of-nth* [*simp*]:
 $m < n \implies rev\ (bl\text{-of-nth } n\ f) ! m = f\ m$
<proof>

lemma *bl-of-nth-inj*:
 $(!!k. k < n \implies f\ k = g\ k) \implies bl\text{-of-nth } n\ f = bl\text{-of-nth } n\ g$
<proof>

lemma *bl-of-nth-nth-le* [*rule-format*]: *ALL xs.*
 $length\ xs \geq n \implies bl\text{-of-nth } n\ (nth\ (rev\ xs)) = drop\ (length\ xs - n)\ xs$
<proof>

lemmas *bl-of-nth-nth* [*simp*] = *order-refl* [*THEN bl-of-nth-nth-le, simplified*]

lemma *size-rbl-pred*: $length\ (rbl\text{-pred } bl) = length\ bl$
<proof>

lemma *size-rbl-succ*: $length\ (rbl\text{-succ } bl) = length\ bl$
<proof>

lemma *size-rbl-add*:
 $!!cl. length\ (rbl\text{-add } bl\ cl) = length\ bl$
<proof>

lemma *size-rbl-mult*:
 $!!cl. length\ (rbl\text{-mult } bl\ cl) = length\ bl$
<proof>

lemmas *rbl-sizes* [*simp*] =
 $size\text{-rbl-pred } size\text{-rbl-succ } size\text{-rbl-add } size\text{-rbl-mult}$

lemmas *rbl-Nils* =
 $rbl\text{-pred}.Nil\ rbl\text{-succ}.Nil\ rbl\text{-add}.Nil\ rbl\text{-mult}.Nil$

lemma *rbl-pred*:
 $!!bin. rbl\text{-pred } (rev\ (bin\text{-to-bl } n\ bin)) = rev\ (bin\text{-to-bl } n\ (Numeral.pred\ bin))$
<proof>

lemma *rbl-succ*:

!!bin. rbl-succ (rev (bin-to-bl n bin)) = rev (bin-to-bl n (Numeral.succ bin))
 ⟨proof⟩

lemma rbl-add:

!!bina binb. rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =
 rev (bin-to-bl n (bina + binb))
 ⟨proof⟩

lemma rbl-add-app2:

!!blb. length blb >= length bla ==>
 rbl-add bla (blb @ blc) = rbl-add bla blb
 ⟨proof⟩

lemma rbl-add-take2:

!!blb. length blb >= length bla ==>
 rbl-add bla (take (length bla) blb) = rbl-add bla blb
 ⟨proof⟩

lemma rbl-add-long:

m >= n ==> rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
 rev (bin-to-bl n (bina + binb))
 ⟨proof⟩

lemma rbl-mult-app2:

!!blb. length blb >= length bla ==>
 rbl-mult bla (blb @ blc) = rbl-mult bla blb
 ⟨proof⟩

lemma rbl-mult-take2:

length blb >= length bla ==>
 rbl-mult bla (take (length bla) blb) = rbl-mult bla blb
 ⟨proof⟩

lemma rbl-mult-gt1:

m >= length bl ==> rbl-mult bl (rev (bin-to-bl m binb)) =
 rbl-mult bl (rev (bin-to-bl (length bl) binb))
 ⟨proof⟩

lemma rbl-mult-gt:

m > n ==> rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
 rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb))
 ⟨proof⟩

lemmas rbl-mult-Suc = lessI [THEN rbl-mult-gt]

lemma rdbl-Cons:

b # rev (bin-to-bl n x) = rev (bin-to-bl (Suc n) (x BIT If b bit.B1 bit.B0))
 ⟨proof⟩

lemma *rbl-mult*: !!*bina binb*.

$$\begin{aligned} & \text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \text{ bina})) (\text{rev } (\text{bin-to-bl } n \text{ binb})) = \\ & \text{rev } (\text{bin-to-bl } n \text{ (bina * binb)}) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *rbl-add-split*:

$$\begin{aligned} & P (\text{rbl-add } (y \# ys) (x \# xs)) = \\ & (\text{ALL } ws. \text{length } ws = \text{length } ys \text{ ---> } ws = \text{rbl-add } ys \text{ xs --->} \\ & (y \text{ ---> } ((x \text{ ---> } P (\text{False} \# \text{rbl-succ } ws)) \ \& \ (\sim x \text{ ---> } P (\text{True} \# ws)))) \\ & \& \\ & (\sim y \text{ ---> } P (x \# ws)) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *rbl-mult-split*:

$$\begin{aligned} & P (\text{rbl-mult } (y \# ys) \text{ xs}) = \\ & (\text{ALL } ws. \text{length } ws = \text{Suc } (\text{length } ys) \text{ ---> } ws = \text{False} \# \text{rbl-mult } ys \text{ xs --->} \\ & (y \text{ ---> } P (\text{rbl-add } ws \text{ xs})) \ \& \ (\sim y \text{ ---> } P \text{ ws})) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *and-len*: $xs = ys \implies xs = ys \ \& \ \text{length } xs = \text{length } ys$

$\langle \text{proof} \rangle$

lemma *size-if*: $\text{size } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{size } xs \text{ else } \text{size } ys)$

$\langle \text{proof} \rangle$

lemma *tl-if*: $\text{tl } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{tl } xs \text{ else } \text{tl } ys)$

$\langle \text{proof} \rangle$

lemma *hd-if*: $\text{hd } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{hd } xs \text{ else } \text{hd } ys)$

$\langle \text{proof} \rangle$

lemma *if-Not-x*: $(\text{if } p \text{ then } \sim x \text{ else } x) = (p = (\sim x))$

$\langle \text{proof} \rangle$

lemma *if-x-Not*: $(\text{if } p \text{ then } x \text{ else } \sim x) = (p = x)$

$\langle \text{proof} \rangle$

lemma *if-same-and*: $(\text{If } p \text{ } x \text{ } y \ \& \ \text{If } p \text{ } u \text{ } v) = (\text{if } p \text{ then } x \ \& \ u \text{ else } y \ \& \ v)$

$\langle \text{proof} \rangle$

lemma *if-same-eq*: $(\text{If } p \text{ } x \text{ } y = (\text{If } p \text{ } u \text{ } v)) = (\text{if } p \text{ then } x = (u) \text{ else } y = (v))$

$\langle \text{proof} \rangle$

lemma *if-same-eq-not*:

$$(\text{If } p \text{ } x \text{ } y = (\sim \text{If } p \text{ } u \text{ } v)) = (\text{if } p \text{ then } x = (\sim u) \text{ else } y = (\sim v))$$

$\langle \text{proof} \rangle$

lemma *if-Cons*: $(\text{if } p \text{ then } x \# xs \text{ else } y \# ys) = \text{If } p \ x \ y \# \text{If } p \ xs \ ys$
 ⟨proof⟩

lemma *if-single*:
 $(\text{if } xc \text{ then } [xab] \text{ else } [an]) = [\text{if } xc \text{ then } xab \text{ else } an]$
 ⟨proof⟩

lemma *if-bool-simps*:
 $\text{If } p \ \text{True} \ y = (p \mid y) \ \& \ \text{If } p \ \text{False} \ y = (\sim p \ \& \ y) \ \&$
 $\text{If } p \ y \ \text{True} = (p \ \dashrightarrow y) \ \& \ \text{If } p \ y \ \text{False} = (p \ \& \ y)$
 ⟨proof⟩

lemmas *if-simps = if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

lemmas *segr = eq-reflection* [where $x = \text{size } w$, *standard*]

lemmas *tl-Nil = tl.simps (1)*
lemmas *tl-Cons = tl.simps (2)*

7.2 Repeated splitting or concatenation

lemma *sclcm*:
 $\text{size } (\text{concat } (\text{map } (\text{bin-to-bl } n) \ xs)) = \text{length } xs * n$
 ⟨proof⟩

lemma *bin-cat-foldl-lem* [rule-format] :
 $\text{ALL } x. \text{foldl } (\%u. \text{bin-cat } u \ n) \ x \ xs =$
 $\text{bin-cat } x \ (\text{size } xs * n) \ (\text{foldl } (\%u. \text{bin-cat } u \ n) \ y \ xs)$
 ⟨proof⟩

lemma *bin-rcat-bl*:
 $(\text{bin-rcat } n \ wl) = \text{bl-to-bin } (\text{concat } (\text{map } (\text{bin-to-bl } n) \ wl))$
 ⟨proof⟩

lemmas *bin-rsplit-aux-simps = bin-rsplit-aux.simps bin-rsplitl-aux.simps*
lemmas *rsplit-aux-simps = bin-rsplit-aux-simps*

lemmas *th-if-simp1 = split-if* [where $P = op = l$,
 THEN *iffD1*, THEN *conjunct1*, THEN *mp*, *standard*]
lemmas *th-if-simp2 = split-if* [where $P = op = l$,
 THEN *iffD1*, THEN *conjunct2*, THEN *mp*, *standard*]

lemmas *rsplit-aux-simp1s = rsplit-aux-simps* [THEN *th-if-simp1*]

lemmas *rsplit-aux-simp2ls = rsplit-aux-simps* [THEN *th-if-simp2*]

lemmas *bin-rsplit-aux-simp2s* [simp] = *rsplit-aux-simp2ls* [unfolded *Let-def*]
lemmas *rbscl = bin-rsplit-aux-simp2s (2)*

lemmas *rsplit-aux-0-simps* [*simp*] =
rsplit-aux-simp1s [*OF disjI1*] *rsplit-aux-simp1s* [*OF disjI2*]

lemma *bin-rsplit-aux-append*:
bin-rsplit-aux (*n*, *bs* @ *cs*, *m*, *c*) = *bin-rsplit-aux* (*n*, *bs*, *m*, *c*) @ *cs*
 ⟨*proof*⟩

lemma *bin-rsplitl-aux-append*:
bin-rsplitl-aux (*n*, *bs* @ *cs*, *m*, *c*) = *bin-rsplitl-aux* (*n*, *bs*, *m*, *c*) @ *cs*
 ⟨*proof*⟩

lemmas *rsplit-aux-apps* [where *bs* = []] =
bin-rsplit-aux-append *bin-rsplitl-aux-append*

lemmas *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplitl-def*

lemmas *rsplit-aux-alt*s = *rsplit-aux-apps*
 [*unfolded append-Nil* *rsplit-def-auxs* [*symmetric*]]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) w = \text{bin-split } n w$
 ⟨*proof*⟩

lemmas *bin-split-minus-simp* =
bin-split.Suc [*THEN* [2] *bin-split-minus* [*symmetric*, *THEN trans*], *standard*]

lemma *bin-split-pred-simp* [*simp*]:
 $(0 :: \text{nat}) < \text{number-of bin} \implies$
 $\text{bin-split } (\text{number-of bin}) w =$
 $(\text{let } (w1, w2) = \text{bin-split } (\text{number-of } (\text{Numeral.pred bin})) (\text{bin-rest } w)$
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$
 ⟨*proof*⟩

declare *bin-split-pred-simp* [*simp*]

lemma *bin-rsplit-aux-simp-alt*:
bin-rsplit-aux (*n*, *bs*, *m*, *c*) =
 (if $m = 0 \vee n = 0$
 then *bs*
 else let (*a*, *b*) = *bin-split* *n* *c* in *bin-rsplit* *n* (*m* - *n*, *a*) @ *b* # *bs*)
 ⟨*proof*⟩

lemmas *bin-rsplit-simp-alt* =
trans [*OF bin-rsplit-def* [*THEN meta-eq-to-obj-eq*]
bin-rsplit-aux-simp-alt, *standard*]

lemmas *bthrs* = *bin-rsplit-simp-alt* [*THEN* [2] *trans*]

lemma *bin-rsplit-size-sign'* [*rule-format*] :
 $n > 0 \implies (\text{ALL } nw \ w. \text{rev } sw = \text{bin-rsplit } n \ (nw, w) \dashrightarrow$

(*ALL* v : *set* sw . $bintrunc\ n\ v = v$)
 ⟨*proof*⟩

lemmas $bin-rsplit-size-sign = bin-rsplit-size-sign'$ [*OF* *asm-rl*
rev-rev-ident [*THEN* *trans*] *set-rev* [*THEN* *equalityD2* [*THEN* *subsetD*]],
standard]

lemma $bin-nth-rsplit$ [*rule-format*] :
 $n > 0 \implies m < n \implies (ALL\ w\ k\ nw.\ rev\ sw = bin-rsplit\ n\ (nw,\ w) \dashrightarrow$
 $k < size\ sw \dashrightarrow bin-nth\ (sw\ !\ k)\ m = bin-nth\ w\ (k * n + m))$
 ⟨*proof*⟩

lemma $bin-rsplit-all$:
 $0 < nw \implies nw \leq n \implies bin-rsplit\ n\ (nw,\ w) = [bintrunc\ n\ w]$
 ⟨*proof*⟩

lemma $bin-rsplit-l$ [*rule-format*] :
 $ALL\ bin.\ bin-rsplitl\ n\ (m,\ bin) = bin-rsplit\ n\ (m,\ bintrunc\ m\ bin)$
 ⟨*proof*⟩

lemma $bin-rsplit-rcat$ [*rule-format*] :
 $n > 0 \dashrightarrow bin-rsplit\ n\ (n * size\ ws,\ bin-rcat\ n\ ws) = map\ (bintrunc\ n)\ ws$
 ⟨*proof*⟩

lemma $bin-rsplit-aux-len-le$ [*rule-format*] :
 $ALL\ ws\ m.\ n \neq 0 \dashrightarrow ws = bin-rsplit-aux\ (n,\ bs,\ nw,\ w) \dashrightarrow$
 $(length\ ws \leq m) = (nw + length\ bs * n \leq m * n)$
 ⟨*proof*⟩

lemma $bin-rsplit-len-le$:
 $n \neq 0 \dashrightarrow ws = bin-rsplit\ n\ (nw,\ w) \dashrightarrow (length\ ws \leq m) = (nw \leq m * n)$
 ⟨*proof*⟩

lemma $bin-rsplit-aux-len$ [*rule-format*] :
 $n \neq 0 \dashrightarrow length\ (bin-rsplit-aux\ (n,\ cs,\ nw,\ w)) =$
 $(nw + n - 1) \div n + length\ cs$
 ⟨*proof*⟩

lemma $bin-rsplit-len$:
 $n \neq 0 \implies length\ (bin-rsplit\ n\ (nw,\ w)) = (nw + n - 1) \div n$
 ⟨*proof*⟩

lemma $bin-rsplit-aux-len-indep$ [*rule-format*] :
 $n \neq 0 \implies (ALL\ v\ bs.\ length\ bs = length\ cs \dashrightarrow$
 $length\ (bin-rsplit-aux\ (n,\ bs,\ nw,\ v)) =$
 $length\ (bin-rsplit-aux\ (n,\ cs,\ nw,\ w)))$
 ⟨*proof*⟩

lemma *bin-rsplit-len-indep*:

$n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, v)) = \text{length } (\text{bin-rsplit } n \ (nw, w))$
 ⟨proof⟩

end

8 TdThs: Type Definition Theorems

theory *TdThs* **imports** *Main* **begin**

9 More lemmas about normal type definitions

lemma

tdD1: type-definition *Rep Abs A* $\implies \forall x. \text{Rep } x \in A$ **and**
tdD2: type-definition *Rep Abs A* $\implies \forall x. \text{Abs } (\text{Rep } x) = x$ **and**
tdD3: type-definition *Rep Abs A* $\implies \forall y. y \in A \longrightarrow \text{Rep } (\text{Abs } y) = y$
 ⟨proof⟩

lemma *td-nat-int*:

type-definition *int nat* (*Collect* (*op* ≤ 0))
 ⟨proof⟩

context *type-definition*

begin

lemmas *Rep'* [*iff*] = *Rep* [*simplified*]

declare *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

lemma *Abs-eqD*: *Abs* $x = \text{Abs } y \implies x \in A \implies y \in A \implies x = y$
 ⟨proof⟩

lemma *Abs-inverse'*:

$r : A \implies \text{Abs } r = a \implies \text{Rep } a = r$
 ⟨proof⟩

lemma *Rep-comp-inverse*:

Rep $o f = g \implies \text{Abs } o g = f$
 ⟨proof⟩

lemma *Rep-eqD* [*elim!*]: *Rep* $x = \text{Rep } y \implies x = y$

⟨proof⟩

lemma *Rep-inverse'*: *Rep* $a = r \implies \text{Abs } r = a$

⟨proof⟩

lemma *comp-Abs-inverse*:

$f \circ \text{Abs} = g \implies g \circ \text{Rep} = f$
 ⟨proof⟩

lemma *set-Rep*:

$A = \text{range Rep}$
 ⟨proof⟩

lemma *set-Rep-Abs*: $A = \text{range (Rep o Abs)}$

⟨proof⟩

lemma *Abs-inj-on*: *inj-on Abs A*

⟨proof⟩

lemma *image*: $\text{Abs } \cdot A = \text{UNIV}$

⟨proof⟩

lemmas *td-thm* = *type-definition-axioms*

lemma *fns1*:

$\text{Rep o fa} = \text{fr o Rep} \mid \text{fa o Abs} = \text{Abs o fr} \implies \text{Abs o fr o Rep} = \text{fa}$
 ⟨proof⟩

lemmas *fns1a* = *disjI1 [THEN fns1]*

lemmas *fns1b* = *disjI2 [THEN fns1]*

lemma *fns4*:

$\text{Rep o fa o Abs} = \text{fr} \implies$
 $\text{Rep o fa} = \text{fr o Rep} \ \& \ \text{fa o Abs} = \text{Abs o fr}$
 ⟨proof⟩

end

interpretation *nat-int*: *type-definition [int nat Collect (op <= 0)]*

⟨proof⟩

declare *Nat.induct* [*case-names 0 Suc, induct type*]

declare *Nat.exhaust* [*case-names 0 Suc, cases type*]

9.1 Extended form of type definition predicate

lemma *td-conds*:

$\text{norm o norm} = \text{norm} \implies (\text{fr o norm} = \text{norm o fr}) =$
 $(\text{norm o fr o norm} = \text{fr o norm} \ \& \ \text{norm o fr o norm} = \text{norm o fr})$
 ⟨proof⟩

lemma *fn-comm-power*:

$\text{fa o tr} = \text{tr o fr} \implies \text{fa } ^n \text{ o tr} = \text{tr o fr } ^n$
 ⟨proof⟩

lemmas *fn-comm-power'* =
ext [THEN *fn-comm-power*, THEN *fun-cong*, *unfolded o-def*, *standard*]

locale *td-ext* = *type-definition* +
fixes *norm*
assumes *eq-norm*: $\bigwedge x. \text{Rep } (\text{Abs } x) = \text{norm } x$
begin

lemma *Abs-norm* [*simp*]:
 $\text{Abs } (\text{norm } x) = \text{Abs } x$
<proof>

lemma *td-th*:
 $g \circ \text{Abs} = f \implies f (\text{Rep } x) = g x$
<proof>

lemma *eq-norm'*: $\text{Rep } o \text{Abs} = \text{norm}$
<proof>

lemma *norm-Rep* [*simp*]: $\text{norm } (\text{Rep } x) = \text{Rep } x$
<proof>

lemmas *td* = *td-thm*

lemma *set-iff-norm*: $w : A \longleftrightarrow w = \text{norm } w$
<proof>

lemma *inverse-norm*:
 $(\text{Abs } n = w) = (\text{Rep } w = \text{norm } n)$
<proof>

lemma *norm-eq-iff*:
 $(\text{norm } x = \text{norm } y) = (\text{Abs } x = \text{Abs } y)$
<proof>

lemma *norm-comps*:
 $\text{Abs } o \text{norm} = \text{Abs}$
 $\text{norm } o \text{Rep} = \text{Rep}$
 $\text{norm } o \text{norm} = \text{norm}$
<proof>

lemmas *norm-norm* [*simp*] = *norm-comps*

lemma *fns5*:
 $\text{Rep } o \text{fa } o \text{Abs} = \text{fr} \implies$
 $\text{fr } o \text{norm} = \text{fr} \ \& \ \text{norm } o \text{fr} = \text{fr}$
<proof>

```

lemma fns2:
  Abs o fr o Rep = fa ==>
  (norm o fr o norm = fr o norm) = (Rep o fa = fr o Rep)
  <proof>

lemma fns3:
  Abs o fr o Rep = fa ==>
  (norm o fr o norm = norm o fr) = (fa o Abs = Abs o fr)
  <proof>

lemma fns:
  fr o norm = norm o fr ==>
  (fa o Abs = Abs o fr) = (Rep o fa = fr o Rep)
  <proof>

lemma range-norm:
  range (Rep o Abs) = A
  <proof>

end

lemmas td-ext-def' =
  td-ext-def [unfolded type-definition-def td-ext-axioms-def]

```

```

end

```

10 WordDefinition: Definition of Word Type

```

theory WordDefinition imports Size BinBoolList TdThs begin

```

```

typedef (open word) 'a word
  = {(0::int) ..< 2len-of TYPE('a::len0)}} <proof>

```

```

instance word :: (len0) number <proof>
instance word :: (type) minus <proof>
instance word :: (type) plus <proof>
instance word :: (type) one <proof>
instance word :: (type) zero <proof>
instance word :: (type) times <proof>
instance word :: (type) Divides.div <proof>
instance word :: (type) power <proof>
instance word :: (type) ord <proof>
instance word :: (type) size <proof>
instance word :: (type) inverse <proof>
instance word :: (type) bit <proof>

```

10.1 Type conversions and casting

constdefs

— representation of words using unsigned or signed bins, only difference in these is the type class

```
word-of-int :: int => 'a :: len0 word
word-of-int w == Abs-word (bintrunc (len-of TYPE ('a)) w)
```

— uint and sint cast a word to an integer, uint treats the word as unsigned, sint treats the most-significant-bit as a sign bit

```
uint :: 'a :: len0 word => int
uint w == Rep-word w
sint :: 'a :: len word => int
sint-uint: sint w == sbintrunc (len-of TYPE ('a) - 1) (uint w)
unat :: 'a :: len0 word => nat
unat w == nat (uint w)
```

— the sets of integers representing the words

```
uints :: nat => int set
uints n == range (bintrunc n)
sints :: nat => int set
sints n == range (sbintrunc (n - 1))
unats :: nat => nat set
unats n == {i. i < 2 ^ n}
norm-sint :: nat => int => int
norm-sint n w == (w + 2 ^ (n - 1)) mod 2 ^ n - 2 ^ (n - 1)
```

— cast a word to a different length

```
scast :: 'a :: len word => 'b :: len word
scast w == word-of-int (sint w)
ucast :: 'a :: len0 word => 'b :: len0 word
ucast w == word-of-int (uint w)
```

— whether a cast (or other) function is to a longer or shorter length

```
source-size :: ('a :: len0 word => 'b) => nat
source-size c == let arb = arbitrary ; x = c arb in size arb
target-size :: ('a => 'b :: len0 word) => nat
target-size c == size (c arbitrary)
is-up :: ('a :: len0 word => 'b :: len0 word) => bool
is-up c == source-size c <= target-size c
is-down :: ('a :: len0 word => 'b :: len0 word) => bool
is-down c == target-size c <= source-size c
```

constdefs

```
of-bl :: bool list => 'a :: len0 word
of-bl bl == word-of-int (bl-to-bin bl)
to-bl :: 'a :: len0 word => bool list
to-bl w ==
bin-to-bl (len-of TYPE ('a)) (uint w)
```

word-reverse :: 'a :: len0 word => 'a word
word-reverse w == of-bl (rev (to-bl w))

defs (overloaded)

word-size: size (w :: 'a :: len0 word) == len-of TYPE('a)
word-number-of-def: number-of w == word-of-int w

constdefs

word-int-case :: (int => 'b) => ('a :: len0 word) => 'b
word-int-case f w == f (uint w)

syntax

of-int :: int => 'a

translations

case x of of-int y => b == word-int-case (%y. b) x

10.2 Arithmetic operations**defs (overloaded)**

word-1-wi: (1 :: ('a :: len0) word) == word-of-int 1
word-0-wi: (0 :: ('a :: len0) word) == word-of-int 0

word-le-def: a <= b == uint a <= uint b
word-less-def: x < y == x <= y & x ~ = (y :: 'a :: len0 word)

constdefs

word-succ :: 'a :: len0 word => 'a word
word-succ a == word-of-int (Numeral.succ (uint a))

word-pred :: 'a :: len0 word => 'a word
word-pred a == word-of-int (Numeral.pred (uint a))

udvd :: 'a::len word => 'a::len word => bool (**infixl** udvd 50)
a udvd b == EX n>=0. uint b = n * uint a

word-sle :: 'a :: len word => 'a word => bool ((-/ <=s -) [50, 51] 50)
a <=s b == sint a <= sint b

word-sless :: 'a :: len word => 'a word => bool ((-/ <s -) [50, 51] 50)
(x <s y) == (x <=s y & x ~ = y)

consts

word-power :: 'a :: len0 word => nat => 'a word

primrec

word-power a 0 = 1
word-power a (Suc n) = a * *word-power* a n

defs (overloaded)

word-pow: power == word-power

word-add-def: $a + b == \text{word-of-int } (\text{uint } a + \text{uint } b)$
word-sub-wi: $a - b == \text{word-of-int } (\text{uint } a - \text{uint } b)$
word-minus-def: $- a == \text{word-of-int } (- \text{uint } a)$
word-mult-def: $a * b == \text{word-of-int } (\text{uint } a * \text{uint } b)$
word-div-def: $a \text{ div } b == \text{word-of-int } (\text{uint } a \text{ div } \text{uint } b)$
word-mod-def: $a \text{ mod } b == \text{word-of-int } (\text{uint } a \text{ mod } \text{uint } b)$

10.3 Bit-wise operations

defs (overloaded)

word-and-def:
 $(a::'a::\text{len0 word}) \text{ AND } b == \text{word-of-int } (\text{uint } a \text{ AND } \text{uint } b)$

word-or-def:
 $(a::'a::\text{len0 word}) \text{ OR } b == \text{word-of-int } (\text{uint } a \text{ OR } \text{uint } b)$

word-xor-def:
 $(a::'a::\text{len0 word}) \text{ XOR } b == \text{word-of-int } (\text{uint } a \text{ XOR } \text{uint } b)$

word-not-def:
 $\text{NOT } (a::'a::\text{len0 word}) == \text{word-of-int } (\text{NOT } (\text{uint } a))$

word-test-bit-def:
 $\text{test-bit } (a::'a::\text{len0 word}) == \text{bin-nth } (\text{uint } a)$

word-set-bit-def:
 $\text{set-bit } (a::'a::\text{len0 word}) \text{ } n \text{ } x ==$
 $\text{word-of-int } (\text{bin-sc } n \text{ (If } x \text{ bit.B1 bit.B0)} (\text{uint } a))$

word-set-bits-def:
 $(\text{BITS } n. f n)::'a::\text{len0 word} == \text{of-bl } (\text{bl-of-nth } (\text{len-of TYPE } ('a)) f)$

word-lsb-def:
 $\text{lsb } (a::'a::\text{len0 word}) == \text{bin-last } (\text{uint } a) = \text{bit.B1}$

word-msb-def:
 $\text{msb } (a::'a::\text{len word}) == \text{bin-sign } (\text{sint } a) = \text{Numeral.Min}$

constdefs

$\text{setBit} :: 'a :: \text{len0 word} ==> \text{nat} ==> 'a \text{ word}$
 $\text{setBit } w \text{ } n == \text{set-bit } w \text{ } n \text{ True}$

$\text{clearBit} :: 'a :: \text{len0 word} ==> \text{nat} ==> 'a \text{ word}$
 $\text{clearBit } w \text{ } n == \text{set-bit } w \text{ } n \text{ False}$

10.4 Shift operations

constdefs

$\text{shifl1} :: 'a :: \text{len0 word} ==> 'a \text{ word}$

shiftr1 w == word-of-int (uint w BIT bit.B0)

— shift right as unsigned or as signed, ie logical or arithmetic

shiftr1 :: 'a :: len0 word => 'a word

shiftr1 w == word-of-int (bin-rest (uint w))

sshiftr1 :: 'a :: len word => 'a word

sshiftr1 w == word-of-int (bin-rest (sint w))

bshiftr1 :: bool => 'a :: len word => 'a word

bshiftr1 b w == of-bl (b # butlast (to-bl w))

sshiftr :: 'a :: len word => nat => 'a word (infixl >>> 55)

w >>> n == (sshiftr1 ^ n) w

mask :: nat => 'a::len word

mask n == (1 << n) - 1

revcast :: 'a :: len0 word => 'b :: len0 word

revcast w == of-bl (takefill False (len-of TYPE('b)) (to-bl w))

slice1 :: nat => 'a :: len0 word => 'b :: len0 word

slice1 n w == of-bl (takefill False n (to-bl w))

slice :: nat => 'a :: len0 word => 'b :: len0 word

slice n w == slice1 (size w - n) w

defs (overloaded)

shiftr-def: (w::'a::len0 word) << n == (shiftr1 ^ n) w

shiftr-def: (w::'a::len0 word) >> n == (shiftr1 ^ n) w

10.5 Rotation

constdefs

rotater1 :: 'a list => 'a list

rotater1 ys ==

case ys of [] => [] | x # xs => last ys # butlast ys

rotater :: nat => 'a list => 'a list

rotater n == rotater1 ^ n

word-rotr :: nat => 'a :: len0 word => 'a :: len0 word

word-rotr n w == of-bl (rotater n (to-bl w))

word-rotl :: nat => 'a :: len0 word => 'a :: len0 word

word-rotl n w == of-bl (rotate n (to-bl w))

word-roti :: int => 'a :: len0 word => 'a :: len0 word

word-roti i w == if $i \geq 0$ then *word-rotr* (nat i) w
 else *word-rotl* (nat ($- i$)) w

10.6 Split and cat operations

constdefs

word-cat :: 'a :: len0 word => 'b :: len0 word => 'c :: len0 word
word-cat a b == *word-of-int* (bin-cat (uint a) (len-of TYPE ('b)) (uint b))

word-split :: 'a :: len0 word => ('b :: len0 word) * ('c :: len0 word)
word-split a ==
 case bin-split (len-of TYPE ('c)) (uint a) of
 (u , v) => (*word-of-int* u , *word-of-int* v)

word-rcat :: 'a :: len0 word list => 'b :: len0 word
word-rcat ws ==
word-of-int (bin-rcat (len-of TYPE ('a)) (map uint ws))

word-rsplit :: 'a :: len0 word => 'b :: len word list
word-rsplit w ==
 map *word-of-int* (bin-rsplit (len-of TYPE ('b)) (len-of TYPE ('a), uint w))

constdefs

— Largest representable machine integer.
max-word :: 'a::len word
max-word \equiv *word-of-int* ($2^{\text{len-of TYPE('a)}} - 1$)

consts

of-bool :: bool \Rightarrow 'a::len word

primrec

of-bool False = 0
of-bool True = 1

lemmas *of-nth-def* = *word-set-bits-def*

lemmas *word-size-gt-0* [iff] =

xtr1 [OF *word-size* [THEN *meta-eq-to-obj-eq*] *len-gt-0*, *standard*]

lemmas *lens-gt-0* = *word-size-gt-0* *len-gt-0*

lemmas *lens-not-0* [iff] = *lens-gt-0* [THEN *gr-implies-not0*, *standard*]

lemma *uints-num*: $\text{uints } n = \{i. 0 \leq i \wedge i < 2^n\}$
 <proof>

lemma *sints-num*: $\text{sints } n = \{i. -(2^{n-1}) \leq i \wedge i < 2^{n-1}\}$
 <proof>

lemmas *atLeastLessThan-alt* = *atLeastLessThan-def* [unfolded]

atLeast-def lessThan-def Collect-conj-eq [symmetric]

lemma *mod-in-reps*: $m > 0 \implies y \bmod m : \{0::\text{int} \ ..< m\}$
 ⟨proof⟩

lemma
Rep-word-0:0 \leq *Rep-word x* **and**
Rep-word-lt: *Rep-word* ($x::'a::\text{len0}$ word) $< 2 \wedge \text{len-of TYPE}('a)$
 ⟨proof⟩

lemma *Rep-word-mod-same*:
Rep-word $x \bmod 2 \wedge \text{len-of TYPE}('a) = \text{Rep-word}$ ($x::'a::\text{len0}$ word)
 ⟨proof⟩

lemma *td-ext-uint*:
td-ext ($\text{uint} :: 'a$ word $\implies \text{int}$) *word-of-int* (uints ($\text{len-of TYPE}('a::\text{len0})$))
 ($\%w::\text{int}.$ $w \bmod 2 \wedge \text{len-of TYPE}('a)$)
 ⟨proof⟩

lemmas *int-word-uint* = *td-ext-uint* [THEN *td-ext.eq-norm*, *standard*]

interpretation *word-uint*:
td-ext [$\text{uint}::'a::\text{len0}$ word $\implies \text{int}$
word-of-int
uints ($\text{len-of TYPE}('a::\text{len0})$)
 $\lambda w. w \bmod 2 \wedge \text{len-of TYPE}('a::\text{len0})$]
 ⟨proof⟩

lemmas *td-uint* = *word-uint.td-thm*

lemmas *td-ext-ubin* = *td-ext-uint*
 [*simplified len-gt-0 no-bintr-alt1* [symmetric]]

interpretation *word-ubin*:
td-ext [$\text{uint}::'a::\text{len0}$ word $\implies \text{int}$
word-of-int
uints ($\text{len-of TYPE}('a::\text{len0})$)
bintrunc ($\text{len-of TYPE}('a::\text{len0})$)]
 ⟨proof⟩

lemma *sint-sbintrunc'*:
sint (*word-of-int* bin $:: 'a$ word) =
 (*sbintrunc* ($\text{len-of TYPE}('a :: \text{len}) - 1$) bin)
 ⟨proof⟩

lemma *uint-sint*:
uint $w = \text{bintrunc}$ ($\text{len-of TYPE}('a)$) (*sint* ($w :: 'a :: \text{len}$ word))
 ⟨proof⟩

lemma *bintr-uint'*:

$n \geq \text{size } w \implies \text{bintrunc } n (\text{uint } w) = \text{uint } w$
 ⟨proof⟩

lemma *wi-bintr'*:

$wb = \text{word-of-int } bin \implies n \geq \text{size } wb \implies$
 $\text{word-of-int } (\text{bintrunc } n \text{ bin}) = wb$
 ⟨proof⟩

lemmas *bintr-uint = bintr-uint'* [unfolded word-size]

lemmas *wi-bintr = wi-bintr'* [unfolded word-size]

lemma *td-ext-sbin*:

$\text{td-ext } (\text{sint} :: 'a \text{ word} \implies \text{int}) \text{ word-of-int } (\text{sints } (\text{len-of TYPE}('a::\text{len})))$
 $(\text{sbintrunc } (\text{len-of TYPE}('a) - 1))$
 ⟨proof⟩

lemmas *td-ext-sint = td-ext-sbin*

[simplified len-gt-0 no-sbintr-alt2 Suc-pred' [symmetric]]

interpretation *word-sint*:

$\text{td-ext } [\text{sint} :: 'a::\text{len} \text{ word} \implies \text{int}]$
 word-of-int
 $\text{sints } (\text{len-of TYPE}('a::\text{len}))$
 $\%w. (w + 2^{(\text{len-of TYPE}('a::\text{len}) - 1)}) \bmod 2^{(\text{len-of TYPE}('a::\text{len}) - 1)}$
 ⟨proof⟩

interpretation *word-sbin*:

$\text{td-ext } [\text{sint} :: 'a::\text{len} \text{ word} \implies \text{int}]$
 word-of-int
 $\text{sints } (\text{len-of TYPE}('a::\text{len}))$
 $\text{sbintrunc } (\text{len-of TYPE}('a::\text{len}) - 1]$
 ⟨proof⟩

lemmas *int-word-sint = td-ext-sint* [THEN *td-ext.eq-norm*, *standard*]

lemmas *td-sint = word-sint.td*

lemma *word-number-of-alt*: $\text{number-of } b == \text{word-of-int } (\text{number-of } b)$

⟨proof⟩

lemma *word-no-wi*: $\text{number-of} = \text{word-of-int}$

⟨proof⟩

lemma *to-bl-def'*:

$(\text{to-bl} :: 'a :: \text{len0} \text{ word} \implies \text{bool list}) =$
 $\text{bin-to-bl } (\text{len-of TYPE}('a)) \circ \text{uint}$

<proof>

lemmas *word-reverse-no-def* [*simp*] = *word-reverse-def* [*of number-of w, standard*]

lemmas *uints-mod* = *uints-def* [*unfolded no-bintr-alt1*]

lemma *uint-bintrunc*: *uint* (*number-of bin* :: 'a word) =
number-of (*bintrunc* (*len-of TYPE* ('a :: len0)) *bin*)
<proof>

lemma *sint-sbintrunc*: *sint* (*number-of bin* :: 'a word) =
number-of (*sbintrunc* (*len-of TYPE* ('a :: len) - 1) *bin*)
<proof>

lemma *unat-bintrunc*:
unat (*number-of bin* :: 'a :: len0 word) =
number-of (*bintrunc* (*len-of TYPE*('a)) *bin*)
<proof>

declare

uint-bintrunc [*simp*]
sint-sbintrunc [*simp*]
unat-bintrunc [*simp*]

lemma *size-0-eq*: *size* (*w* :: 'a :: len0 word) = 0 ==> *v* = *w*
<proof>

lemmas *uint-lem* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]

lemmas *sint-lem* = *word-sint.Rep* [*unfolded sints-num mem-Collect-eq*]

lemmas *uint-ge-0* [*iff*] = *uint-lem* [*THEN conjunct1, standard*]

lemmas *uint-lt2p* [*iff*] = *uint-lem* [*THEN conjunct2, standard*]

lemmas *sint-ge* = *sint-lem* [*THEN conjunct1, standard*]

lemmas *sint-lt* = *sint-lem* [*THEN conjunct2, standard*]

lemma *sign-uint-Pls* [*simp*]:
bin-sign (*uint x*) = *Numeral.Pl*
<proof>

lemmas *uint-m2p-neg* = *iffD2* [*OF diff-less-0-iff-less uint-lt2p, standard*]

lemmas *uint-m2p-not-non-neg* =
iffD2 [*OF linorder-not-le uint-m2p-neg, standard*]

lemma *lt2p-lem*:
len-of TYPE('a) <= *n* ==> *uint* (*w* :: 'a :: len0 word) < 2 ^ *n*
<proof>

lemmas *uint-le-0-iff* [*simp*] =
uint-ge-0 [*THEN leD, THEN linorder-antisym-conv1, standard*]

lemma *uint-nat*: $uint\ w == int\ (unat\ w)$
 ⟨proof⟩

lemma *uint-number-of*:
 $uint\ (number-of\ b :: 'a :: len0\ word) = number-of\ b\ mod\ 2^{\wedge}\ len-of\ TYPE('a)$
 ⟨proof⟩

lemma *unat-number-of*:
 $bin-sign\ b = Numeral.Pls ==>$
 $unat\ (number-of\ b :: 'a :: len0\ word) = number-of\ b\ mod\ 2^{\wedge}\ len-of\ TYPE\ ('a)$
 ⟨proof⟩

lemma *sint-number-of*: $sint\ (number-of\ b :: 'a :: len\ word) = (number-of\ b +$
 $2^{\wedge}\ (len-of\ TYPE('a) - 1))\ mod\ 2^{\wedge}\ len-of\ TYPE('a) -$
 $2^{\wedge}\ (len-of\ TYPE('a) - 1)$
 ⟨proof⟩

lemma *word-of-int-bin* [*simp*] :
 $(word-of-int\ (number-of\ bin) :: 'a :: len0\ word) = (number-of\ bin)$
 ⟨proof⟩

lemma *word-int-case-wi*:
 $word-int-case\ f\ (word-of-int\ i :: 'b\ word) =$
 $f\ (i\ mod\ 2^{\wedge}\ len-of\ TYPE('b::len0))$
 ⟨proof⟩

lemma *word-int-split*:
 $P\ (word-int-case\ f\ x) =$
 $(ALL\ i.\ x = (word-of-int\ i :: 'b :: len0\ word) \&$
 $0 <= i \& i < 2^{\wedge}\ len-of\ TYPE('b) \longrightarrow P\ (f\ i))$
 ⟨proof⟩

lemma *word-int-split-asm*:
 $P\ (word-int-case\ f\ x) =$
 $(\sim\ (EX\ n.\ x = (word-of-int\ n :: 'b::len0\ word) \&$
 $0 <= n \& n < 2^{\wedge}\ len-of\ TYPE('b::len0) \& \sim\ P\ (f\ n)))$
 ⟨proof⟩

lemmas *uint-range'* =
 $word-uint.Rep\ [unfolded\ uints-num\ mem-Collect-eq,\ standard]$

lemmas *sint-range'* = $word-sint.Rep\ [unfolded\ One-nat-def$
 $sints-num\ mem-Collect-eq,\ standard]$

lemma *uint-range-size*: $0 <= uint\ w \& uint\ w < 2^{\wedge}\ size\ w$
 ⟨proof⟩

lemma *sint-range-size*:
 $(2^{\wedge}\ (size\ w - Suc\ 0)) <= sint\ w \& sint\ w < 2^{\wedge}\ (size\ w - Suc\ 0)$

<proof>

lemmas *sint-above-size = sint-range-size*
 [THEN *conjunct2*, THEN [2] *xtr8*, folded *One-nat-def*, *standard*]

lemmas *sint-below-size = sint-range-size*
 [THEN *conjunct1*, THEN [2] *order-trans*, folded *One-nat-def*, *standard*]

lemma *test-bit-eq-iff*: $(\text{test-bit } (u::'a::\text{len0 } \text{word}) = \text{test-bit } v) = (u = v)$
<proof>

lemma *test-bit-size* [rule-format] : $(w::'a::\text{len0 } \text{word}) !! n \dashrightarrow n < \text{size } w$
<proof>

lemma *word-eqI* [rule-format] :
fixes $u :: 'a::\text{len0 } \text{word}$
shows $(\text{ALL } n. n < \text{size } u \dashrightarrow u !! n = v !! n) \implies u = v$
<proof>

lemmas *word-eqD = test-bit-eq-iff* [THEN *iffD2*, THEN *fun-cong*, *standard*]

lemma *test-bit-bin'*: $w !! n = (n < \text{size } w \ \& \ \text{bin-nth } (\text{uint } w) \ n)$
<proof>

lemmas *test-bit-bin = test-bit-bin'* [unfolded *word-size*]

lemma *bin-nth-uint-imp'*: $\text{bin-nth } (\text{uint } w) \ n \dashrightarrow n < \text{size } w$
<proof>

lemma *bin-nth-sint'*:
 $n \geq \text{size } w \dashrightarrow \text{bin-nth } (\text{sint } w) \ n = \text{bin-nth } (\text{sint } w) \ (\text{size } w - 1)$
<proof>

lemmas *bin-nth-uint-imp = bin-nth-uint-imp'* [rule-format, unfolded *word-size*]

lemmas *bin-nth-sint = bin-nth-sint'* [rule-format, unfolded *word-size*]

lemma *td-bl*:
type-definition $(\text{to-bl} :: 'a::\text{len0 } \text{word} \Rightarrow \text{bool list})$
of-bl
 $\{\text{bl. length } \text{bl} = \text{len-of } \text{TYPE}('a)\}$
<proof>

interpretation *word-bl*:
type-definition $[\text{to-bl} :: 'a::\text{len0 } \text{word} \Rightarrow \text{bool list}]$
of-bl
 $\{\text{bl. length } \text{bl} = \text{len-of } \text{TYPE}('a::\text{len0})\}$
<proof>

lemma *word-size-bl*: $size\ w == size\ (to-bl\ w)$
 ⟨proof⟩

lemma *to-bl-use-of-bl*:
 $(to-bl\ w = bl) = (w = of-bl\ bl \wedge length\ bl = length\ (to-bl\ w))$
 ⟨proof⟩

lemma *to-bl-word-rev*: $to-bl\ (word-reverse\ w) = rev\ (to-bl\ w)$
 ⟨proof⟩

lemma *word-rev-rev* [simp] : $word-reverse\ (word-reverse\ w) = w$
 ⟨proof⟩

lemma *word-rev-gal*: $word-reverse\ w = u ==> word-reverse\ u = w$
 ⟨proof⟩

lemmas *word-rev-gal'* = sym [THEN *word-rev-gal*, symmetric, standard]

lemmas *length-bl-gt-0* [iff] = xtr1 [OF *word-bl.Rep'* len-gt-0, standard]

lemmas *bl-not-Nil* [iff] =

length-bl-gt-0 [THEN *length-greater-0-conv* [THEN *iffD1*], standard]

lemmas *length-bl-neq-0* [iff] = *length-bl-gt-0* [THEN *gr-implies-not0*]

lemma *hd-bl-sign-sint*: $hd\ (to-bl\ w) = (bin-sign\ (sint\ w) = Numeral.Min)$
 ⟨proof⟩

lemma *of-bl-drop'*:
 $lend = length\ bl - len-of\ TYPE\ ('a :: len0) ==>$
 $of-bl\ (drop\ lend\ bl) = (of-bl\ bl :: 'a\ word)$
 ⟨proof⟩

lemmas *of-bl-no* = *of-bl-def* [folded *word-number-of-def*]

lemma *test-bit-of-bl*:
 $(of-bl\ bl :: 'a :: len0\ word) !! n = (rev\ bl ! n \wedge n < len-of\ TYPE\ ('a) \wedge n < length\ bl)$
 ⟨proof⟩

lemma *no-of-bl*:
 $(number-of\ bin :: 'a :: len0\ word) = of-bl\ (bin-to-bl\ (len-of\ TYPE\ ('a))\ bin)$
 ⟨proof⟩

lemma *uint-bl*: $to-bl\ w == bin-to-bl\ (size\ w)\ (uint\ w)$
 ⟨proof⟩

lemma *to-bl-bin*: $bl-to-bin\ (to-bl\ w) = uint\ w$
 ⟨proof⟩

lemma *to-bl-of-bin*:

to-bl (*word-of-int* *bin::'a::len0* *word*) = *bin-to-bl* (*len-of TYPE('a)*) *bin*
 ⟨*proof*⟩

lemmas *to-bl-no-bin* [*simp*] = *to-bl-of-bin* [*folded word-number-of-def*]

lemma *to-bl-to-bin* [*simp*] : *bl-to-bin* (*to-bl* *w*) = *uint* *w*
 ⟨*proof*⟩

lemmas *uint-bl-bin* [*simp*] = *trans* [*OF bin-bl-bin word-ubin.norm-Rep, standard*]

lemmas *num-AB-u* [*simp*] = *word-uint.Rep-inverse*
 [*unfolded o-def word-number-of-def [symmetric], standard*]

lemmas *num-AB-s* [*simp*] = *word-sint.Rep-inverse*
 [*unfolded o-def word-number-of-def [symmetric], standard*]

lemma *uints-unats*: *uints* *n* = *int* ‘ *unats* *n*
 ⟨*proof*⟩

lemma *unats-uints*: *unats* *n* = *nat* ‘ *uints* *n*
 ⟨*proof*⟩

lemmas *bintr-num* = *word-ubin.norm-eq-iff*
 [*symmetric, folded word-number-of-def, standard*]

lemmas *sbintr-num* = *word-sbin.norm-eq-iff*
 [*symmetric, folded word-number-of-def, standard*]

lemmas *num-of-bintr* = *word-ubin.Abs-norm* [*folded word-number-of-def, standard*]

lemmas *num-of-sbintr* = *word-sbin.Abs-norm* [*folded word-number-of-def, standard*]

lemma *num-of-bintr'*:
bintrunc (*len-of TYPE('a :: len0)*) *a* = *b* ==>
number-of *a* = (*number-of* *b* :: '*a* *word*)
 ⟨*proof*⟩

lemma *num-of-sbintr'*:
sbintrunc (*len-of TYPE('a :: len) - 1*) *a* = *b* ==>
number-of *a* = (*number-of* *b* :: '*a* *word*)
 ⟨*proof*⟩

lemmas *num-abs-bintr* = *sym* [*THEN trans,*
OF num-of-bintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]

lemmas *num-abs-sbintr* = *sym* [*THEN trans,*
OF num-of-sbintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]

lemma *ucast-id*: $ucast\ w = w$
 ⟨*proof*⟩

lemma *scast-id*: $scast\ w = w$
 ⟨*proof*⟩

lemma *ucast-bl*: $ucast\ w ==\ of-bl\ (to-bl\ w)$
 ⟨*proof*⟩

lemma *nth-ucast*:
 $(ucast\ w :: 'a :: len0\ word)\ !!\ n = (w\ !!\ n \ \&\ n < len-of\ TYPE('a))$
 ⟨*proof*⟩

lemma *ucast-bintr* [*simp*]:
 $ucast\ (number-of\ w :: 'a :: len0\ word) =$
 $number-of\ (bintrunc\ (len-of\ TYPE('a))\ w)$
 ⟨*proof*⟩

lemma *scast-sbintr* [*simp*]:
 $scast\ (number-of\ w :: 'a :: len\ word) =$
 $number-of\ (sbintrunc\ (len-of\ TYPE('a) - Suc\ 0)\ w)$
 ⟨*proof*⟩

lemmas *source-size = source-size-def* [*unfolded Let-def word-size*]

lemmas *target-size = target-size-def* [*unfolded Let-def word-size*]

lemmas *is-down = is-down-def* [*unfolded source-size target-size*]

lemmas *is-up = is-up-def* [*unfolded source-size target-size*]

lemmas *is-up-down =*
 $trans\ [OF\ is-up\ [THEN\ meta-eq-to-obj-eq]$
 $is-down\ [THEN\ meta-eq-to-obj-eq,\ symmetric],$
 $standard]$

lemma *down-cast-same'*: $uc = ucast ==> is-down\ uc ==> uc = scast$
 ⟨*proof*⟩

lemma *word-rev-tf'*:
 $r = to-bl\ (of-bl\ bl) ==> r = rev\ (takefill\ False\ (length\ r)\ (rev\ bl))$
 ⟨*proof*⟩

lemmas *word-rev-tf = refl* [*THEN word-rev-tf', unfolded word-bl.Rep', standard*]

lemmas *word-rep-drop = word-rev-tf* [*simplified takefill-alt,*
simplified, simplified rev-take, simplified]

lemma *to-bl-ucast*:

$to-bl (ucast (w::'b::len0\ word) ::'a::len0\ word) =$
 $replicate (len-of\ TYPE('a) - len-of\ TYPE('b))\ False\ @$
 $drop (len-of\ TYPE('b) - len-of\ TYPE('a)) (to-bl\ w)$
 $\langle proof \rangle$

lemma *ucast-up-app'*:

$uc = ucast ==> source-size\ uc + n = target-size\ uc ==>$
 $to-bl (uc\ w) = replicate\ n\ False\ @ (to-bl\ w)$
 $\langle proof \rangle$

lemma *ucast-down-drop'*:

$uc = ucast ==> source-size\ uc = target-size\ uc + n ==>$
 $to-bl (uc\ w) = drop\ n (to-bl\ w)$
 $\langle proof \rangle$

lemma *scast-down-drop'*:

$sc = scast ==> source-size\ sc = target-size\ sc + n ==>$
 $to-bl (sc\ w) = drop\ n (to-bl\ w)$
 $\langle proof \rangle$

lemma *sint-up-scast'*:

$sc = scast ==> is-up\ sc ==> sint (sc\ w) = sint\ w$
 $\langle proof \rangle$

lemma *uint-up-ucast'*:

$uc = ucast ==> is-up\ uc ==> uint (uc\ w) = uint\ w$
 $\langle proof \rangle$

lemmas *down-cast-same* = refl [THEN down-cast-same']

lemmas *ucast-up-app* = refl [THEN ucast-up-app']

lemmas *ucast-down-drop* = refl [THEN ucast-down-drop']

lemmas *scast-down-drop* = refl [THEN scast-down-drop']

lemmas *uint-up-ucast* = refl [THEN uint-up-ucast']

lemmas *sint-up-scast* = refl [THEN sint-up-scast']

lemma *ucast-up-ucast'*: $uc = ucast ==> is-up\ uc ==> ucast (uc\ w) = ucast\ w$
 $\langle proof \rangle$

lemma *scast-up-scast'*: $sc = scast ==> is-up\ sc ==> scast (sc\ w) = scast\ w$
 $\langle proof \rangle$

lemma *ucast-of-bl-up'*:

$w = of-bl\ bl ==> size\ bl <= size\ w ==> ucast\ w = of-bl\ bl$
 $\langle proof \rangle$

lemmas *ucast-up-ucast* = refl [THEN ucast-up-ucast']

lemmas *scast-up-scast* = refl [THEN scast-up-scast']

lemmas *ucast-of-bl-up* = refl [THEN ucast-of-bl-up']

lemmas *ucast-up-ucast-id* = *trans* [*OF ucast-up-ucast ucast-id*]

lemmas *scast-up-scast-id* = *trans* [*OF scast-up-scast scast-id*]

lemmas *isduu* = *is-up-down* [**where** *c* = *ucast*, *THEN iffD2*]

lemmas *isdus* = *is-up-down* [**where** *c* = *scast*, *THEN iffD2*]

lemmas *ucast-down-ucast-id* = *isduu* [*THEN ucast-up-ucast-id*]

lemmas *scast-down-scast-id* = *isdus* [*THEN ucast-up-ucast-id*]

lemma *up-ucast-surj*:

is-up (*ucast* :: 'b::len0 word => 'a::len0 word) ==>

surj (*ucast* :: 'a word => 'b word)

<proof>

lemma *up-scast-surj*:

is-up (*scast* :: 'b::len word => 'a::len word) ==>

surj (*scast* :: 'a word => 'b word)

<proof>

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word => 'a::len word) ==>

inj-on (*ucast* :: 'a word => 'b word) *A*

<proof>

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word => 'a::len0 word) ==>

inj-on (*ucast* :: 'a word => 'b word) *A*

<proof>

lemma *of-bl-append-same*: *of-bl* (*X @ to-bl w*) = *w*

<proof>

lemma *ucast-down-no'*:

uc = *ucast* ==> *is-down uc* ==> *uc* (*number-of bin*) = *number-of bin*

<proof>

lemmas *ucast-down-no* = *ucast-down-no'* [*OF refl*]

lemma *ucast-down-bl'*: *uc* = *ucast* ==> *is-down uc* ==> *uc* (*of-bl bl*) = *of-bl bl*

<proof>

lemmas *ucast-down-bl* = *ucast-down-bl'* [*OF refl*]

lemmas *slice-def'* = *slice-def* [*unfolded word-size*]

lemmas *test-bit-def'* = *word-test-bit-def* [*THEN meta-eq-to-obj-eq, THEN fun-cong*]

lemmas *word-log-defs* = *word-and-def word-or-def word-xor-def word-not-def*

lemmas *word-log-bin-defs* = *word-log-defs*

end

11 WordArith: Word Arithmetic

theory *WordArith* imports *WordDefinition* begin

lemma *word-less-alt*: $(a < b) = (\text{uint } a < \text{uint } b)$
 ⟨*proof*⟩

lemma *signed-linorder*: *linorder* *word-sle* *word-sless*
 ⟨*proof*⟩

interpretation *signed*: *linorder* [*word-sle* *word-sless*]
 ⟨*proof*⟩

lemmas *word-arith-wis* [*THEN meta-eq-to-obj-eq*] =
word-add-def *word-mult-def* *word-minus-def*
word-succ-def *word-pred-def* *word-0-wi* *word-1-wi*

lemma *udvdI*:
 $0 \leq n \implies \text{uint } b = n * \text{uint } a \implies a \text{ udvd } b$
 ⟨*proof*⟩

lemmas *word-div-no* [*simp*] =
word-div-def [*of number-of a number-of b, standard*]

lemmas *word-mod-no* [*simp*] =
word-mod-def [*of number-of a number-of b, standard*]

lemmas *word-less-no* [*simp*] =
word-less-def [*of number-of a number-of b, standard*]

lemmas *word-le-no* [*simp*] =
word-le-def [*of number-of a number-of b, standard*]

lemmas *word-sless-no* [*simp*] =
word-sless-def [*of number-of a number-of b, standard*]

lemmas *word-sle-no* [*simp*] =
word-sle-def [*of number-of a number-of b, standard*]

lemmas *word-0-wi-Pls* = *word-0-wi* [*folded Pls-def*]

lemmas *word-0-no* = *word-0-wi-Pls* [*folded word-no-wi*]

lemma *int-one-bin*: $(1 :: \text{int}) == (\text{Numeral.Pls BIT bit.B1})$
 ⟨*proof*⟩

lemma *word-1-no*:

$(1 :: 'a :: \text{len0 word}) == \text{number-of (Numeral.Pls BIT bit.B1)}$
 $\langle \text{proof} \rangle$

lemma *word-m1-wi*: $-1 == \text{word-of-int } -1$

$\langle \text{proof} \rangle$

lemma *word-m1-wi-Min*: $-1 = \text{word-of-int Numeral.Min}$

$\langle \text{proof} \rangle$

lemma *word-0-bl*: $\text{of-bl } [] = 0$

$\langle \text{proof} \rangle$

lemma *word-1-bl*: $\text{of-bl } [\text{True}] = 1$

$\langle \text{proof} \rangle$

lemma *uint-0 [simp]*: $(\text{uint } 0 = 0)$

$\langle \text{proof} \rangle$

lemma *of-bl-0 [simp]*: $\text{of-bl (replicate } n \text{ False)} = 0$

$\langle \text{proof} \rangle$

lemma *to-bl-0*:

$\text{to-bl } (0 :: 'a :: \text{len0 word}) = \text{replicate (len-of TYPE('a)) False}$
 $\langle \text{proof} \rangle$

lemma *uint-0-iff*: $(\text{uint } x = 0) = (x = 0)$

$\langle \text{proof} \rangle$

lemma *unat-0-iff*: $(\text{unat } x = 0) = (x = 0)$

$\langle \text{proof} \rangle$

lemma *unat-0 [simp]*: $\text{unat } 0 = 0$

$\langle \text{proof} \rangle$

lemma *size-0-same'*: $\text{size } w = 0 ==> w = (v :: 'a :: \text{len0 word})$

$\langle \text{proof} \rangle$

lemmas *size-0-same* = *size-0-same'* [folded word-size]

lemmas *unat-eq-0* = *unat-0-iff*

lemmas *unat-eq-zero* = *unat-0-iff*

lemma *unat-gt-0*: $(0 < \text{unat } x) = (x \sim = 0)$

$\langle \text{proof} \rangle$

lemma *ucast-0 [simp]*: $\text{ucast } 0 = 0$

$\langle \text{proof} \rangle$

lemma *sint-0* [*simp*] : *sint 0 = 0*
 ⟨*proof*⟩

lemma *scast-0* [*simp*] : *scast 0 = 0*
 ⟨*proof*⟩

lemma *sint-n1* [*simp*] : *sint -1 = -1*
 ⟨*proof*⟩

lemma *scast-n1* [*simp*] : *scast -1 = -1*
 ⟨*proof*⟩

lemma *wint-1* [*simp*] : *wint (1 :: 'a :: len word) = 1*
 ⟨*proof*⟩

lemma *unat-1* [*simp*] : *unat (1 :: 'a :: len word) = 1*
 ⟨*proof*⟩

lemma *ucast-1* [*simp*] : *ucast (1 :: 'a :: len word) = 1*
 ⟨*proof*⟩

lemmas *ariths* =
bintr-ariths [*THEN word-ubin.norm-eq-iff* [*THEN iffD1*],
folded word-ubin.eq-norm, standard]

lemma *wi-homs*:

shows

wi-hom-add: *word-of-int a + word-of-int b = word-of-int (a + b)* **and**

wi-hom-mult: *word-of-int a * word-of-int b = word-of-int (a * b)* **and**

wi-hom-neg: *- word-of-int a = word-of-int (- a)* **and**

wi-hom-succ: *word-succ (word-of-int a) = word-of-int (Numeral.succ a)* **and**

wi-hom-pred: *word-pred (word-of-int a) = word-of-int (Numeral.pred a)*

⟨*proof*⟩

lemmas *wi-hom-syms* = *wi-homs* [*symmetric*]

lemma *word-sub-def*: *a - b == a + - (b :: 'a :: len0 word)*
 ⟨*proof*⟩

lemmas *word-diff-minus* = *word-sub-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-of-int-sub-hom*:

(word-of-int a) - word-of-int b = word-of-int (a - b)

⟨*proof*⟩

lemmas *new-word-of-int-homs* =

word-of-int-sub-hom wi-homs word-0-wi word-1-wi

lemmas *new-word-of-int-hom-syms = new-word-of-int-homs [symmetric, standard]*

lemmas *word-of-int-hom-syms =
new-word-of-int-hom-syms [unfolded succ-def pred-def]*

lemmas *word-of-int-homs =
new-word-of-int-homs [unfolded succ-def pred-def]*

lemmas *word-of-int-add-hom = word-of-int-homs (2)*
lemmas *word-of-int-mult-hom = word-of-int-homs (3)*
lemmas *word-of-int-minus-hom = word-of-int-homs (4)*
lemmas *word-of-int-succ-hom = word-of-int-homs (5)*
lemmas *word-of-int-pred-hom = word-of-int-homs (6)*
lemmas *word-of-int-0-hom = word-of-int-homs (7)*
lemmas *word-of-int-1-hom = word-of-int-homs (8)*

lemmas *word-arith-alt =
word-sub-wi [unfolded succ-def pred-def, THEN meta-eq-to-obj-eq, standard]
word-arith-wis [unfolded succ-def pred-def, standard]*

lemmas *word-sub-alt = word-arith-alt (1)*
lemmas *word-add-alt = word-arith-alt (2)*
lemmas *word-mult-alt = word-arith-alt (3)*
lemmas *word-minus-alt = word-arith-alt (4)*
lemmas *word-succ-alt = word-arith-alt (5)*
lemmas *word-pred-alt = word-arith-alt (6)*
lemmas *word-0-alt = word-arith-alt (7)*
lemmas *word-1-alt = word-arith-alt (8)*

11.1 Transferring goals from words to ints

lemma *word-ths:*

shows

word-succ-p1: word-succ a = a + 1 and
word-pred-m1: word-pred a = a - 1 and
word-pred-succ: word-pred (word-succ a) = a and
word-succ-pred: word-succ (word-pred a) = a and
*word-mult-succ: word-succ a * b = b + a * b*
<proof>

lemmas *uint-cong = arg-cong [where f = uint]*

lemmas *uint-word-ariths =
word-arith-alt [THEN trans [OF uint-cong int-word-uint], standard]*

lemmas *uint-word-arith-bintrs* = *uint-word-ariths* [folded bintrunc-mod2p]

lemmas *sint-word-ariths* = *uint-word-arith-bintrs*
 [THEN *uint-sint* [symmetric, THEN trans],
 unfolded *uint-sint* *bintr-arith1s* *bintr-ariths*
len-gt-0 [THEN *bin-sbin-eq-iff*] *word-sbin.norm-Rep*, standard]

lemmas *uint-div-alt* = *word-div-def*
 [THEN *meta-eq-to-obj-eq* [THEN trans [OF *uint-cong* *int-word-uint*]], standard]

lemmas *uint-mod-alt* = *word-mod-def*
 [THEN *meta-eq-to-obj-eq* [THEN trans [OF *uint-cong* *int-word-uint*]], standard]

lemma *word-pred-0-n1*: *word-pred 0* = *word-of-int -1*
 ⟨proof⟩

lemma *word-pred-0-Min*: *word-pred 0* = *word-of-int Numeral.Min*
 ⟨proof⟩

lemma *word-m1-Min*: *- 1* = *word-of-int Numeral.Min*
 ⟨proof⟩

lemma *succ-pred-no* [simp]:
word-succ (number-of bin) = *number-of (Numeral.succ bin)* &
word-pred (number-of bin) = *number-of (Numeral.pred bin)*
 ⟨proof⟩

lemma *word-sp-01* [simp] :
word-succ -1 = *0* & *word-succ 0* = *1* & *word-pred 0* = *-1* & *word-pred 1* = *0*
 ⟨proof⟩

lemma *word-of-int-Ex*:
 $\exists y. x = \text{word-of-int } y$
 ⟨proof⟩

lemma *word-arith-eqs*:
fixes *a* :: 'a::len0 word
fixes *b* :: 'a::len0 word
shows
word-add-0: $0 + a = a$ **and**
word-add-0-right: $a + 0 = a$ **and**
word-mult-1: $1 * a = a$ **and**
word-mult-1-right: $a * 1 = a$ **and**
word-add-commute: $a + b = b + a$ **and**
word-add-assoc: $a + b + c = a + (b + c)$ **and**
word-add-left-commute: $a + (b + c) = b + (a + c)$ **and**
word-mult-commute: $a * b = b * a$ **and**
word-mult-assoc: $a * b * c = a * (b * c)$ **and**

word-mult-left-commute: $a * (b * c) = b * (a * c)$ **and**
word-left-distrib: $(a + b) * c = a * c + b * c$ **and**
word-right-distrib: $a * (b + c) = a * b + a * c$ **and**
word-left-minus: $- a + a = 0$ **and**
word-diff-0-right: $a - 0 = a$ **and**
word-diff-self: $a - a = 0$
 ⟨*proof*⟩

lemmas *word-add-ac* = *word-add-commute word-add-assoc word-add-left-commute*
lemmas *word-mult-ac* = *word-mult-commute word-mult-assoc word-mult-left-commute*

lemmas *word-plus-ac0* = *word-add-0 word-add-0-right word-add-ac*
lemmas *word-times-ac1* = *word-mult-1 word-mult-1-right word-mult-ac*

11.2 Order on fixed-length words

lemma *word-order-trans*: $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-order-refl*: $z \leq (z :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-order-antisym*: $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-order-linear*:
 $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-zero-le* [*simp*] :
 $0 \leq (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

instance *word* :: (*len0*) *semigroup-add*
 ⟨*proof*⟩

instance *word* :: (*len0*) *linorder*
 ⟨*proof*⟩

instance *word* :: (*len0*) *ring*
 ⟨*proof*⟩

lemma *word-m1-ge* [*simp*] : *word-pred* 0 $\geq y$
 ⟨*proof*⟩

lemmas *word-n1-ge* [*simp*] = *word-m1-ge* [*simplified word-sp-01*]

lemmas *word-not-simps* [*simp*] =
word-zero-le [*THEN leD*] *word-m1-ge* [*THEN leD*] *word-n1-ge* [*THEN leD*]

lemma *word-gt-0*: $0 < y = (0 \sim = (y :: 'a :: \text{len0 word}))$
 ⟨*proof*⟩

lemmas *word-gt-0-no* [*simp*] = *word-gt-0* [*of number-of y, standard*]

lemma *word-sless-alt*: $(a <_s b) == (sint a < sint b)$
 ⟨*proof*⟩

lemma *word-le-nat-alt*: $(a <= b) = (unat a <= unat b)$
 ⟨*proof*⟩

lemma *word-less-nat-alt*: $(a < b) = (unat a < unat b)$
 ⟨*proof*⟩

lemma *wi-less*:
 $(\text{word-of-int } n < (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} < m \bmod 2 \wedge \text{len-of TYPE('a)})$
 ⟨*proof*⟩

lemma *wi-le*:
 $(\text{word-of-int } n <= (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} <= m \bmod 2 \wedge \text{len-of TYPE('a)})$
 ⟨*proof*⟩

lemma *udvd-nat-alt*: $a \text{ udvd } b = (EX n \geq 0. unat b = n * unat a)$
 ⟨*proof*⟩

lemma *udvd-iff-dvd*: $x \text{ udvd } y \leftrightarrow unat x \text{ dvd } unat y$
 ⟨*proof*⟩

lemmas *unat-mono* = *word-less-nat-alt* [*THEN iffD1, standard*]

lemma *word-zero-neq-one*: $0 < \text{len-of TYPE('a :: len0)} ==> (0 :: 'a word) \sim = 1$
 ⟨*proof*⟩

lemmas *lenw1-zero-neq-one* = *len-gt-0* [*THEN word-zero-neq-one*]

lemma *no-no* [*simp*] : $\text{number-of } (\text{number-of } b) = \text{number-of } b$
 ⟨*proof*⟩

lemma *unat-minus-one*: $x \sim = 0 ==> unat (x - 1) = unat x - 1$
 ⟨*proof*⟩

lemma *measure-unat*: $p \sim = 0 ==> unat (p - 1) < unat p$
 ⟨*proof*⟩

lemmas *uint-add-ge0* [*simp*] =

add-nonneg-nonneg [*OF uint-ge-0 uint-ge-0, standard*]
lemmas *uint-mult-ge0* [*simp*] =
mult-nonneg-nonneg [*OF uint-ge-0 uint-ge-0, standard*]

lemma *uint-sub-lt2p* [*simp*]:
 $uint\ (x :: 'a :: len0\ word) - uint\ (y :: 'b :: len0\ word) <$
 $2 \wedge len-of\ TYPE('a)$
 ⟨*proof*⟩

11.3 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:
 $(uint\ x + uint\ y < 2 \wedge len-of\ TYPE('a)) =$
 $(uint\ (x + y :: 'a :: len0\ word) = uint\ x + uint\ y)$
 ⟨*proof*⟩

lemma *uint-mult-lem*:
 $(uint\ x * uint\ y < 2 \wedge len-of\ TYPE('a)) =$
 $(uint\ (x * y :: 'a :: len0\ word) = uint\ x * uint\ y)$
 ⟨*proof*⟩

lemma *uint-sub-lem*:
 $(uint\ x \geq uint\ y) = (uint\ (x - y) = uint\ x - uint\ y)$
 ⟨*proof*⟩

lemma *uint-add-le*: $uint\ (x + y) \leq uint\ x + uint\ y$
 ⟨*proof*⟩

lemma *uint-sub-ge*: $uint\ (x - y) \geq uint\ x - uint\ y$
 ⟨*proof*⟩

lemmas *uint-sub-if'* =
trans [*OF uint-word-ariths(1) mod-sub-if-z, simplified, standard*]

lemmas *uint-plus-if'* =
trans [*OF uint-word-ariths(2) mod-add-if-z, simplified, standard*]

11.4 Definition of uint_arith

lemma *word-of-int-inverse*:
 $word-of-int\ r = a \implies 0 \leq r \implies r < 2 \wedge len-of\ TYPE('a) \implies$
 $uint\ (a :: 'a :: len0\ word) = r$
 ⟨*proof*⟩

lemma *uint-split*:
fixes $x :: 'a :: len0\ word$
shows $P\ (uint\ x) =$
 $(ALL\ i.\ word-of-int\ i = x \ \&\ 0 \leq i \ \&\ i < 2 \wedge len-of\ TYPE('a) \implies P\ i)$
 ⟨*proof*⟩

lemma *uint-split-asm*:

fixes $x :: 'a :: \text{len0 word}$

shows $P (\text{uint } x) =$

$(\sim (EX i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ i))$

$\langle \text{proof} \rangle$

lemmas *uint-splits* = *uint-split uint-split-asm*

lemmas *uint-arith-simps* =

word-le-def word-less-alt

word-uint.Rep-inject [*symmetric*]

uint-sub-if' uint-plus-if'

lemma *power-False-cong*: $\text{False} \implies a \hat{=} b = c \hat{=} d$

$\langle \text{proof} \rangle$

$\langle \text{ML} \rangle$

11.5 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*:

$((x :: 'a :: \text{len0 word}) \leq x + y) = (\text{uint } x + \text{uint } y < 2^{\text{size } x})$

$\langle \text{proof} \rangle$

lemmas *no-olen-add* = *no-plus-overflow-uint-size* [*unfolded word-size*]

lemma *no-ulen-sub*: $((x :: 'a :: \text{len0 word}) \geq x - y) = (\text{uint } y \leq \text{uint } x)$

$\langle \text{proof} \rangle$

lemma *no-olen-add'*:

fixes $x :: 'a :: \text{len0 word}$

shows $(x \leq y + x) = (\text{uint } y + \text{uint } x < 2^{\text{len-of TYPE('a)}})$

$\langle \text{proof} \rangle$

lemmas *olen-add-quiv* = *trans* [*OF no-olen-add no-olen-add'* [*symmetric*], *standard*]

lemmas *uint-plus-simple-iff* = *trans* [*OF no-olen-add uint-add-lem, standard*]

lemmas *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1, standard*]

lemmas *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem, standard*]

lemmas *uint-minus-simple-alt* = *uint-sub-lem* [*folded word-le-def*]

lemmas *word-sub-le-iff* = *no-ulen-sub* [*folded word-le-def*]

lemmas *word-sub-le* = *word-sub-le-iff* [*THEN iffD2, standard*]

lemma *word-less-sub1*:

$(x :: 'a :: \text{len word}) \sim = 0 \implies (1 < x) = (0 < x - 1)$

<proof>

lemma *word-le-sub1*:

$(x :: 'a :: \text{len0 word}) \sim = 0 \implies (1 \leq x) = (0 \leq x - 1)$
<proof>

lemma *sub-wrap-lt*:

$((x :: 'a :: \text{len0 word}) < x - z) = (x < z)$
<proof>

lemma *sub-wrap*:

$((x :: 'a :: \text{len0 word}) \leq x - z) = (z = 0 \mid x < z)$
<proof>

lemma *plus-minus-not-NULL-ab*:

$(x :: 'a :: \text{len0 word}) \leq ab - c \implies c \leq ab \implies c \sim = 0 \implies x + c \sim = 0$
<proof>

lemma *plus-minus-no-overflow-ab*:

$(x :: 'a :: \text{len0 word}) \leq ab - c \implies c \leq ab \implies x \leq x + c$
<proof>

lemma *le-minus'*:

$(a :: 'a :: \text{len0 word}) + c \leq b \implies a \leq a + c \implies c \leq b - a$
<proof>

lemma *le-plus'*:

$(a :: 'a :: \text{len0 word}) \leq b \implies c \leq b - a \implies a + c \leq b$
<proof>

lemmas *le-plus = le-plus'* [rotated]

lemmas *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*, standard]

lemma *word-plus-mono-right*:

$(y :: 'a :: \text{len0 word}) \leq z \implies x \leq x + z \implies x + y \leq x + z$
<proof>

lemma *word-less-minus-cancel*:

$y - x < z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) < z$
<proof>

lemma *word-less-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) < z \implies x \leq y \implies y - x < z - x$
<proof>

lemma *word-less-minus-mono*:

$a < c \implies d < b \implies a - b < a \implies c - d < c$
 $\implies a - b < c - (d :: 'a :: \text{len word})$

<proof>

lemma *word-le-minus-cancel:*

$y - x \leq z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) \leq z$
<proof>

lemma *word-le-minus-mono-left:*

$(y :: 'a :: \text{len0 word}) \leq z \implies x \leq y \implies y - x \leq z - x$
<proof>

lemma *word-le-minus-mono:*

$a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c$
 $\implies a - b \leq c - (d :: 'a :: \text{len word})$
<proof>

lemma *plus-le-left-cancel-wrap:*

$(x :: 'a :: \text{len0 word}) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$
<proof>

lemma *plus-le-left-cancel-nowrap:*

$(x :: 'a :: \text{len0 word}) \leq x + y' \implies x \leq x + y \implies (x + y' < x + y) = (y' < y)$
<proof>

lemma *word-plus-mono-right2:*

$(a :: 'a :: \text{len0 word}) \leq a + b \implies c \leq b \implies a \leq a + c$
<proof>

lemma *word-less-add-right:*

$(x :: 'a :: \text{len0 word}) < y - z \implies z \leq y \implies x + z < y$
<proof>

lemma *word-less-sub-right:*

$(x :: 'a :: \text{len0 word}) < y + z \implies y \leq x \implies x - y < z$
<proof>

lemma *word-le-plus-either:*

$(x :: 'a :: \text{len0 word}) \leq y \mid x \leq z \implies y \leq y + z \implies x \leq y + z$
<proof>

lemma *word-less-nowrapI:*

$(x :: 'a :: \text{len0 word}) < z - k \implies k \leq z \implies 0 < k \implies x < x + k$
<proof>

lemma *inc-le:* $(i :: 'a :: \text{len word}) < m \implies i + 1 \leq m$

<proof>

lemma *inc-i:*

$(1 :: 'a :: \text{len word}) \leq i \implies i < m \implies 1 \leq (i + 1) \ \& \ i + 1 \leq m$
 ⟨proof⟩

lemma *udvd-incr-lem*:

$up < uq \implies up = ua + n * \text{uint } K \implies$
 $uq = ua + n' * \text{uint } K \implies up + \text{uint } K \leq uq$
 ⟨proof⟩

lemma *udvd-incr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p + K \leq q$
 ⟨proof⟩

lemma *udvd-decr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p \leq q - K$
 ⟨proof⟩

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [where *ua=0, simplified*]

lemmas *udvd-incr0* = *udvd-incr'* [where *ua=0, simplified*]

lemmas *udvd-decr0* = *udvd-decr'* [where *ua=0, simplified*]

lemma *udvd-minus-le'*:

$xy < k \implies z \text{ udvd } xy \implies z \text{ udvd } k \implies xy \leq k - z$
 ⟨proof⟩

lemma *udvd-incr2-K*:

$p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq$
 $p \implies$
 $0 < K \implies p \leq p + K \ \& \ p + K \leq a + s$
 ⟨proof⟩

lemma *word-succ-rbl*:

$\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-succ } w) = (\text{rev } (\text{rbl-succ } (\text{rev } \text{bl})))$
 ⟨proof⟩

lemma *word-pred-rbl*:

$\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-pred } w) = (\text{rev } (\text{rbl-pred } (\text{rev } \text{bl})))$
 ⟨proof⟩

lemma *word-add-rbl*:

$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$
 $\text{to-bl } (v + w) = (\text{rev } (\text{rbl-add } (\text{rev } \text{vbl}) (\text{rev } \text{wbl})))$
 ⟨proof⟩

lemma *word-mult-rbl*:

$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$
 $\text{to-bl } (v * w) = (\text{rev } (\text{rbl-mult } (\text{rev } \text{vbl}) (\text{rev } \text{wbl})))$

<proof>

lemma *rtb-rbl-ariths*:

$rev (to-bl w) = ys \implies rev (to-bl (word-succ w)) = rbl-succ ys$

$rev (to-bl w) = ys \implies rev (to-bl (word-pred w)) = rbl-pred ys$

$[| rev (to-bl v) = ys; rev (to-bl w) = xs |]$
 $\implies rev (to-bl (v * w)) = rbl-mult ys xs$

$[| rev (to-bl v) = ys; rev (to-bl w) = xs |]$
 $\implies rev (to-bl (v + w)) = rbl-add ys xs$
<proof>

11.6 Arithmetic type class instantiations

instance *word* :: (*len0*) *comm-monoid-add* *<proof>*

instance *word* :: (*len0*) *comm-monoid-mult*
<proof>

instance *word* :: (*len0*) *comm-semiring*
<proof>

instance *word* :: (*len0*) *ab-group-add* *<proof>*

instance *word* :: (*len0*) *comm-ring* *<proof>*

instance *word* :: (*len*) *comm-semiring-1*
<proof>

instance *word* :: (*len*) *comm-ring-1* *<proof>*

instance *word* :: (*len0*) *comm-semiring-0* *<proof>*

instance *word* :: (*len0*) *order* *<proof>*

instance *word* :: (*len*) *recpower*
<proof>

lemma *zero-bintrunc*:

$iszero (number-of x :: 'a :: len word) =$
 $(bintrunc (len-of TYPE('a)) x = Numeral.Pls)$
<proof>

lemmas *word-le-0-iff* [*simp*] =
 $word-zero-le [THEN leD, THEN linorder-antisym-conv1]$

lemma *word-of-nat*: $of\text{-}nat\ n = word\text{-}of\text{-}int\ (int\ n)$
 ⟨*proof*⟩

lemma *word-of-int*: $of\text{-}int = word\text{-}of\text{-}int$
 ⟨*proof*⟩

lemma *word-of-int-nat*:
 $0 \leq x \implies word\text{-}of\text{-}int\ x = of\text{-}nat\ (nat\ x)$
 ⟨*proof*⟩

lemma *word-number-of-eq*:
 $number\text{-}of\ w = (of\text{-}int\ w :: 'a :: len\ word)$
 ⟨*proof*⟩

instance *word* :: $(len)\ number\text{-}ring$
 ⟨*proof*⟩

lemma *iszero-word-no* [*simp*] :
 $iszero\ (number\text{-}of\ bin :: 'a :: len\ word) =$
 $iszero\ (number\text{-}of\ (bintrunc\ (len\text{-}of\ TYPE('a))\ bin) :: int)$
 ⟨*proof*⟩

11.7 Word and nat

lemma *td-ext-unat'*:
 $n = len\text{-}of\ TYPE\ ('a :: len) \implies$
 $td\text{-}ext\ (unat :: 'a\ word \Rightarrow nat)\ of\text{-}nat$
 $(unats\ n)\ (\%i.\ i\ mod\ 2\ ^\ n)$
 ⟨*proof*⟩

lemmas *td-ext-unat* = *refl* [*THEN* *td-ext-unat'*]

lemmas *unat-of-nat* = *td-ext-unat* [*THEN* *td-ext.eq-norm*, *standard*]

interpretation *word-unat*:
 $td\text{-}ext\ [unat :: 'a :: len\ word \Rightarrow nat$
 $of\text{-}nat$
 $unats\ (len\text{-}of\ TYPE('a :: len))$
 $\%i.\ i\ mod\ 2\ ^\ len\text{-}of\ TYPE('a :: len)]$
 ⟨*proof*⟩

lemmas *td-unat* = *word-unat.td-thm*

lemmas *unat-lt2p* [*iff*] = *word-unat.Rep* [*unfolded* *unats-def* *mem-Collect-eq*]

lemma *unat-le*: $y \leq unat\ (z :: 'a :: len\ word) \implies y : unats\ (len\text{-}of\ TYPE('a))$
 ⟨*proof*⟩

lemma *word-nchotomy*:

ALL w . *EX* n . ($w :: 'a :: \text{len word}$) = *of-nat* n & $n < 2 \wedge \text{len-of TYPE } ('a)$
 ⟨*proof*⟩

lemma *of-nat-eq*:

fixes $w :: 'a :: \text{len word}$

shows (*of-nat* $n = w$) = ($\exists q$. $n = \text{unat } w + q * 2 \wedge \text{len-of TYPE } ('a)$)
 ⟨*proof*⟩

lemma *of-nat-eq-size*:

(*of-nat* $n = w$) = (*EX* q . $n = \text{unat } w + q * 2 \wedge \text{size } w$)
 ⟨*proof*⟩

lemma *of-nat-0*:

(*of-nat* $m = (0 :: 'a :: \text{len word})$) = ($\exists q$. $m = q * 2 \wedge \text{len-of TYPE } ('a)$)
 ⟨*proof*⟩

lemmas *of-nat-2p = mult-1 [symmetric, THEN iffD2 [OF of-nat-0 exI]]*

lemma *of-nat-gt-0*: *of-nat* $k \sim = 0 \implies 0 < k$

⟨*proof*⟩

lemma *of-nat-neq-0*:

$0 < k \implies k < 2 \wedge \text{len-of TYPE } ('a :: \text{len}) \implies \text{of-nat } k \sim = (0 :: 'a \text{ word})$
 ⟨*proof*⟩

lemma *Abs-fnat-hom-add*:

of-nat $a + \text{of-nat } b = \text{of-nat } (a + b)$
 ⟨*proof*⟩

lemma *Abs-fnat-hom-mult*:

of-nat $a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$
 ⟨*proof*⟩

lemma *Abs-fnat-hom-Suc*:

word-succ (*of-nat* a) = *of-nat* (*Suc* a)
 ⟨*proof*⟩

lemma *Abs-fnat-hom-0*: ($0 :: 'a :: \text{len word}$) = *of-nat* 0

⟨*proof*⟩

lemma *Abs-fnat-hom-1*: ($1 :: 'a :: \text{len word}$) = *of-nat* (*Suc* 0)

⟨*proof*⟩

lemmas *Abs-fnat-homs =*

Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc
Abs-fnat-hom-0 Abs-fnat-hom-1

lemma *word-arith-nat-add*:

$a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$

<proof>

lemma *word-arith-nat-mult:*

$a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$

<proof>

lemma *word-arith-nat-Suc:*

$\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$

<proof>

lemma *word-arith-nat-div:*

$a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div } \text{unat } b)$

<proof>

lemma *word-arith-nat-mod:*

$a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod } \text{unat } b)$

<proof>

lemmas *word-arith-nat-defs =*

word-arith-nat-add word-arith-nat-mult

word-arith-nat-Suc Abs-fnat-hom-0

Abs-fnat-hom-1 word-arith-nat-div

word-arith-nat-mod

lemmas *unat-cong = arg-cong [where f = unat]*

lemmas *unat-word-ariths = word-arith-nat-defs*

[THEN trans [OF unat-cong unat-of-nat], standard]

lemmas *word-sub-less-iff = word-sub-le-iff*

[simplified linorder-not-less [symmetric], simplified]

lemma *unat-add-lem:*

$(\text{unat } x + \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$

$(\text{unat } (x + y :: 'a :: \text{len word}) = \text{unat } x + \text{unat } y)$

<proof>

lemma *unat-mult-lem:*

$(\text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$

$(\text{unat } (x * y :: 'a :: \text{len word}) = \text{unat } x * \text{unat } y)$

<proof>

lemmas *unat-plus-if' =*

trans [OF unat-word-ariths(1) mod-nat-add, simplified, standard]

lemma *le-no-overflow:*

$x \leq b \iff a \leq a + b \iff x \leq a + (b :: 'a :: \text{len0 word})$

<proof>

lemmas *un-ui-le = trans*
 [OF *word-le-nat-alt* [symmetric]
word-le-def [THEN *meta-eq-to-obj-eq*],
standard]

lemma *unat-sub-if-size*:
 $unat (x - y) = (if\ unat\ y \leq\ unat\ x$
 $then\ unat\ x - unat\ y$
 $else\ unat\ x + 2^{\text{size } x} - unat\ y)$
 ⟨proof⟩

lemmas *unat-sub-if' = unat-sub-if-size* [unfolded word-size]

lemma *unat-div*: $unat ((x :: 'a :: len\ word)\ div\ y) = unat\ x\ div\ unat\ y$
 ⟨proof⟩

lemma *unat-mod*: $unat ((x :: 'a :: len\ word)\ mod\ y) = unat\ x\ mod\ unat\ y$
 ⟨proof⟩

lemma *uint-div*: $uint ((x :: 'a :: len\ word)\ div\ y) = uint\ x\ div\ uint\ y$
 ⟨proof⟩

lemma *uint-mod*: $uint ((x :: 'a :: len\ word)\ mod\ y) = uint\ x\ mod\ uint\ y$
 ⟨proof⟩

11.8 Definition of unat_arith tactic

lemma *unat-split*:
 fixes $x :: 'a :: len\ word$
 shows $P (unat\ x) =$
 $(ALL\ n.\ of_nat\ n = x \ \&\ n < 2^{\text{len-of } TYPE('a)} \ \longrightarrow\ P\ n)$
 ⟨proof⟩

lemma *unat-split-asm*:
 fixes $x :: 'a :: len\ word$
 shows $P (unat\ x) =$
 $(\sim (EX\ n.\ of_nat\ n = x \ \&\ n < 2^{\text{len-of } TYPE('a)} \ \&\ \sim\ P\ n))$
 ⟨proof⟩

lemmas *of-nat-inverse =*
word-unat.Abs-inverse' [rotated, unfolded unats-def, simplified]

lemmas *unat-splits = unat-split unat-split-asm*

lemmas *unat-arith-simps =*
word-le-nat-alt word-less-nat-alt
word-unat.Rep-inject [symmetric]
unat-sub-if' unat-plus-if' unat-div unat-mod

⟨ML⟩

lemma *no-plus-overflow-unat-size*:

$((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2 \wedge \text{size } x)$
 ⟨proof⟩

lemma *unat-sub*: $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$
 ⟨proof⟩

lemmas *no-olen-add-nat* = *no-plus-overflow-unat-size* [unfolded word-size]

lemmas *unat-plus-simple* = *trans* [OF *no-olen-add-nat* *unat-add-lem*, *standard*]

lemma *word-div-mult*:

$(0 :: 'a :: \text{len word}) < y \implies \text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE('a)} \implies$
 $x * y \text{ div } y = x$
 ⟨proof⟩

lemma *div-lt'*: $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies$
 $\text{unat } i * \text{unat } x < 2 \wedge \text{len-of TYPE('a)}$
 ⟨proof⟩

lemmas *div-lt''* = *order-less-imp-le* [THEN *div-lt'*]

lemma *div-lt-mult*: $(i :: 'a :: \text{len word}) < k \text{ div } x \implies 0 < x \implies i * x < k$
 ⟨proof⟩

lemma *div-le-mult*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$
 ⟨proof⟩

lemma *div-lt-uint'*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2 \wedge \text{len-of TYPE('a)}$
 ⟨proof⟩

lemmas *div-lt-uint''* = *order-less-imp-le* [THEN *div-lt-uint'*]

lemma *word-le-exists'*:

$(x :: 'a :: \text{len0 word}) \leq y \implies$
 $(\exists z. y = x + z \ \& \ \text{uint } x + \text{uint } z < 2 \wedge \text{len-of TYPE('a)})$
 ⟨proof⟩

lemmas *plus-minus-not-NULL* = *order-less-imp-le* [THEN *plus-minus-not-NULL-ab*]

lemmas *plus-minus-no-overflow* =
order-less-imp-le [THEN *plus-minus-no-overflow-ab*]

lemmas *mcs* = *word-less-minus-cancel* *word-less-minus-mono-left*

word-le-minus-cancel word-le-minus-mono-left

lemmas *word-l-diffs* = *mcs* [**where** $y = w + x$, *unfolded add-diff-cancel, standard*]
lemmas *word-diff-ls* = *mcs* [**where** $z = w + x$, *unfolded add-diff-cancel, standard*]
lemmas *word-plus-mcs* = *word-diff-ls*
 [**where** $y = v + x$, *unfolded add-diff-cancel, standard*]

lemmas *le-unat-voi* = *unat-le* [*THEN word-unat.Abs-inverse*]

lemmas *thd* = *refl* [*THEN* [2] *split-div-lemma* [*THEN iffD2*], *THEN conjunct1*]

lemma *thd1*:
 $a \text{ div } b * b \leq (a :: \text{nat})$
<proof>

lemmas *uno-simps* [*THEN le-unat-voi, standard*] =
mod-le-divisor div-le-dividend thd1

lemma *word-mod-div-equality*:
 $(n \text{ div } b) * b + (n \text{ mod } b) = (n :: 'a :: \text{len word})$
<proof>

lemma *word-div-mult-le*: $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$
<proof>

lemma *word-mod-less-divisor*: $0 < n \implies m \text{ mod } n < (n :: 'a :: \text{len word})$
<proof>

lemma *word-of-int-power-hom*:
 $\text{word-of-int } a ^ n = (\text{word-of-int } (a ^ n) :: 'a :: \text{len word})$
<proof>

lemma *word-arith-power-alt*:
 $a ^ n = (\text{word-of-int } (\text{uint } a ^ n) :: 'a :: \text{len word})$
<proof>

lemma *of-bl-length-less*:
 $\text{length } x = k \implies k < \text{len-of TYPE('a)} \implies (\text{of-bl } x :: 'a :: \text{len word}) < 2 ^ k$
<proof>

11.9 Cardinality, finiteness of set of words

lemmas *card-lessThan'* = *card-lessThan* [*unfolded lessThan-def*]

lemmas *card-eq* = *word-unat.Abs-inj-on* [*THEN card-image*,
unfolded word-unat.image, unfolded unats-def, standard]

lemmas *card-word* = *trans* [*OF card-eq card-lessThan'*, *standard*]

lemma *finite-word-UNIV*: *finite* (*UNIV* :: 'a :: len word set)
 ⟨*proof*⟩

lemma *card-word-size*:
 $\text{card } (UNIV :: 'a :: len \text{ word set}) = (2 \wedge \text{size } (x :: 'a \text{ word}))$
 ⟨*proof*⟩

end

12 WordBitwise: Bitwise Operations on Words

theory *WordBitwise* **imports** *WordArith* **begin**

lemmas *bin-log-bintrs* = *bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or*

lemmas *wils1* = *bin-log-bintrs* [*THEN* *word-ubin.norm-eq-iff* [*THEN* *iffD1*],
folded word-ubin.eq-norm, THEN eq-reflection, standard]

lemmas *word-log-binary-defs* =
word-and-def word-or-def word-xor-def

lemmas *word-no-log-defs* [*simp*] =
word-not-def [**where** *a=number-of a*,
unfolded word-no-wi wils1, folded word-no-wi, standard]
word-log-binary-defs [**where** *a=number-of a and b=number-of b*,
unfolded word-no-wi wils1, folded word-no-wi, standard]

lemmas *word-wi-log-defs* = *word-no-log-defs* [*unfolded word-no-wi*]

lemma *uint-or*: $\text{uint } (x \text{ OR } y) = (\text{uint } x) \text{ OR } (\text{uint } y)$
 ⟨*proof*⟩

lemma *uint-and*: $\text{uint } (x \text{ AND } y) = (\text{uint } x) \text{ AND } (\text{uint } y)$
 ⟨*proof*⟩

lemma *word-ops-nth-size*:
 $n < \text{size } (x :: 'a :: len0 \text{ word}) \implies$
 $(x \text{ OR } y) !! n = (x !! n | y !! n) \&$
 $(x \text{ AND } y) !! n = (x !! n \& y !! n) \&$
 $(x \text{ XOR } y) !! n = (x !! n \sim y !! n) \&$
 $(\text{NOT } x) !! n = (\sim x !! n)$
 ⟨*proof*⟩

lemma *word-ao-nth*:

fixes $x :: 'a::len0$ word
shows $(x \text{ OR } y) !! n = (x !! n | y !! n) \&$
 $(x \text{ AND } y) !! n = (x !! n \& y !! n)$
 $\langle \text{proof} \rangle$

lemmas *bwsimps* =

word-of-int-homs(2)
word-0-wi-Pls
word-m1-wi-Min
word-wi-log-defs

lemma *word-bw-assocs*:

fixes $x :: 'a::len0$ word
shows
 $(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 $\langle \text{proof} \rangle$

lemma *word-bw-comms*:

fixes $x :: 'a::len0$ word
shows
 $x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
 $\langle \text{proof} \rangle$

lemma *word-bw-lcs*:

fixes $x :: 'a::len0$ word
shows
 $y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 $\langle \text{proof} \rangle$

lemma *word-log-esimps* [*simp*]:

fixes $x :: 'a::len0$ word
shows
 $x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$

$0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$
 ⟨proof⟩

lemma *word-not-dist*:

fixes $x :: 'a::\text{len}0 \text{ word}$

shows

$\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$

$\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$

⟨proof⟩

lemma *word-bw-same*:

fixes $x :: 'a::\text{len}0 \text{ word}$

shows

$x \text{ AND } x = x$

$x \text{ OR } x = x$

$x \text{ XOR } x = 0$

⟨proof⟩

lemma *word-ao-absorbs* [*simp*]:

fixes $x :: 'a::\text{len}0 \text{ word}$

shows

$x \text{ AND } (y \text{ OR } x) = x$

$x \text{ OR } y \text{ AND } x = x$

$x \text{ AND } (x \text{ OR } y) = x$

$y \text{ AND } x \text{ OR } x = x$

$(y \text{ OR } x) \text{ AND } x = x$

$x \text{ OR } x \text{ AND } y = x$

$(x \text{ OR } y) \text{ AND } x = x$

$x \text{ AND } y \text{ OR } x = x$

⟨proof⟩

lemma *word-not-not* [*simp*]:

$\text{NOT } \text{NOT } (x :: 'a::\text{len}0 \text{ word}) = x$

⟨proof⟩

lemma *word-ao-dist*:

fixes $x :: 'a::\text{len}0 \text{ word}$

shows $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$

⟨proof⟩

lemma *word-oa-dist*:

fixes $x :: 'a::\text{len}0 \text{ word}$

shows $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$

⟨proof⟩

lemma *word-add-not* [*simp*]:

fixes $x :: 'a::len0$ word
shows $x + NOT\ x = -1$
 $\langle proof \rangle$

lemma *word-plus-and-or* [*simp*]:
fixes $x :: 'a::len0$ word
shows $(x\ AND\ y) + (x\ OR\ y) = x + y$
 $\langle proof \rangle$

lemma *leoa*:
fixes $x :: 'a::len0$ word
shows $(w = (x\ OR\ y)) ==> (y = (w\ AND\ y))$ $\langle proof \rangle$

lemma *leao*:
fixes $x' :: 'a::len0$ word
shows $(w' = (x'\ AND\ y')) ==> (x' = (x'\ OR\ w'))$ $\langle proof \rangle$

lemmas *word-ao-equiv* = *leao* [*COMP* *leoa* [*COMP* *iffI*]]

lemma *le-word-or2*: $x \leq x\ OR\ (y :: 'a::len0$ word)
 $\langle proof \rangle$

lemmas *le-word-or1* = *xtr3* [*OF* *word-bw-comms* (2) *le-word-or2*, *standard*]

lemmas *word-and-le1* =
xtr3 [*OF* *word-ao-absorbs* (4) [*symmetric*] *le-word-or2*, *standard*]

lemmas *word-and-le2* =
xtr3 [*OF* *word-ao-absorbs* (8) [*symmetric*] *le-word-or2*, *standard*]

lemma *bl-word-not*: $to\text{-}bl\ (NOT\ w) = map\ Not\ (to\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *bl-word-xor*: $to\text{-}bl\ (v\ XOR\ w) = app2\ op\ \sim = (to\text{-}bl\ v)\ (to\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *bl-word-or*: $to\text{-}bl\ (v\ OR\ w) = app2\ op\ | = (to\text{-}bl\ v)\ (to\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *bl-word-and*: $to\text{-}bl\ (v\ AND\ w) = app2\ op\ \& = (to\text{-}bl\ v)\ (to\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *word-lsb-alt*: $lsb\ (w :: 'a::len0$ word) = *test-bit* $w\ 0$
 $\langle proof \rangle$

lemma *word-lsb-1-0*: $lsb\ (1 :: 'a::len$ word) $\&\ \sim\ lsb\ (0 :: 'b::len0$ word)
 $\langle proof \rangle$

lemma *word-lsb-last*: $lsb\ (w :: 'a::len$ word) = *last* $(to\text{-}bl\ w)$
 $\langle proof \rangle$

lemma *word-lsb-int*: $lsb\ w = (uint\ w\ mod\ 2 = 1)$

<proof>

lemma *word-msb-sint*: $msb\ w = (sint\ w < 0)$
<proof>

lemma *word-msb-no'*:
 $w = number-of\ bin \implies msb\ (w::'a::len\ word) = bin-nth\ bin\ (size\ w - 1)$
<proof>

lemmas *word-msb-no* = *refl* [*THEN* *word-msb-no'*, *unfolded word-size*]

lemma *word-msb-nth'*: $msb\ (w::'a::len\ word) = bin-nth\ (uint\ w)\ (size\ w - 1)$
<proof>

lemmas *word-msb-nth* = *word-msb-nth'* [*unfolded word-size*]

lemma *word-msb-alt*: $msb\ (w::'a::len\ word) = hd\ (to-bl\ w)$
<proof>

lemma *word-set-nth*:
 $set-bit\ w\ n\ (test-bit\ w\ n) = (w::'a::len0\ word)$
<proof>

lemma *bin-nth-uint'*:
 $bin-nth\ (uint\ w)\ n = (rev\ (bin-to-bl\ (size\ w)\ (uint\ w))\ !\ n\ \&\ n < size\ w)$
<proof>

lemmas *bin-nth-uint* = *bin-nth-uint'* [*unfolded word-size*]

lemma *test-bit-bl*: $w\ !!\ n = (rev\ (to-bl\ w)\ !\ n\ \&\ n < size\ w)$
<proof>

lemma *to-bl-nth*: $n < size\ w \implies to-bl\ w\ !\ n = w\ !!\ (size\ w - Suc\ n)$
<proof>

lemma *test-bit-set*:
fixes $w :: 'a::len0\ word$
shows $(set-bit\ w\ n\ x)\ !!\ n = (n < size\ w\ \&\ x)$
<proof>

lemma *test-bit-set-gen*:
fixes $w :: 'a::len0\ word$
shows $test-bit\ (set-bit\ w\ n\ x)\ m =$
 $(if\ m = n\ then\ n < size\ w\ \&\ x\ else\ test-bit\ w\ m)$
<proof>

lemma *of-bl-rep-False*: $of-bl\ (replicate\ n\ False\ @\ bs) = of-bl\ bs$
<proof>

lemma *msb-nth'*:

fixes $w :: 'a::len\ word$
shows $msb\ w = w\ !!\ (size\ w - 1)$
 $\langle proof \rangle$

lemmas $msb-nth = msb-nth'$ [*unfolded word-size*]

lemmas $msb0 = len-gt-0$ [*THEN diff-Suc-less, THEN word-ops-nth-size [unfolded word-size], standard*]

lemmas $msb1 = msb0$ [**where** $i = 0$]

lemmas $word-ops-msb = msb1$ [*unfolded msb-nth [symmetric, unfolded One-nat-def]*]

lemmas $lsb0 = len-gt-0$ [*THEN word-ops-nth-size [unfolded word-size], standard*]

lemmas $word-ops-lsb = lsb0$ [*unfolded word-lsb-alt*]

lemma *td-ext-nth'*:

$n = size\ (w::'a::len0\ word) ==> ofn = set-bits ==> [w, ofn\ g] = l ==>$
 $td-ext\ test-bit\ ofn\ \{f.\ ALL\ i.\ f\ i\ \longrightarrow\ i < n\}\ (\%h\ i.\ h\ i\ \&\ i < n)$
 $\langle proof \rangle$

lemmas $td-ext-nth = td-ext-nth'$ [*OF refl refl refl, unfolded word-size*]

interpretation *test-bit*:

$td-ext\ [op\ !! :: 'a::len0\ word ==> nat ==> bool$
 $set-bits$
 $\{f.\ \forall i.\ f\ i \longrightarrow i < len-of\ TYPE('a::len0)\}$
 $(\lambda h\ i.\ h\ i \wedge i < len-of\ TYPE('a::len0))$
 $\langle proof \rangle$

declare $test-bit.Rep'$ [*simp del*]

declare $test-bit.Rep'$ [*rule del*]

lemmas $td-nth = test-bit.td-thm$

lemma *word-set-set-same*:

fixes $w :: 'a::len0\ word$
shows $set-bit\ (set-bit\ w\ n\ x)\ n\ y = set-bit\ w\ n\ y$
 $\langle proof \rangle$

lemma *word-set-set-diff*:

fixes $w :: 'a::len0\ word$
assumes $m \sim n$
shows $set-bit\ (set-bit\ w\ m\ x)\ n\ y = set-bit\ (set-bit\ w\ n\ y)\ m\ x$
 $\langle proof \rangle$

lemma *test-bit-no'*:

fixes $w :: 'a::len0\ word$
shows $w = number-of\ bin ==> test-bit\ w\ n = (n < size\ w\ \&\ bin-nth\ bin\ n)$
 $\langle proof \rangle$

lemmas *test-bit-no* =
refl [THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard]

lemma *nth-0*: $\sim (0::'a::len0 \text{ word}) !! n$
<proof>

lemma *nth-sint*:
fixes *w* :: *'a*::*len* *word*
defines *l* \equiv *len-of TYPE ('a)*
shows *bin-nth (sint w) n* = *(if n < l - 1 then w !! n else w !! (l - 1))*
<proof>

lemma *word-lsb-no*:
lsb (number-of bin :: 'a :: len word) = *(bin-last bin = bit.B1)*
<proof>

lemma *word-set-no*:
set-bit (number-of bin::'a::len0 word) n b =
number-of (bin-sc n (if b then bit.B1 else bit.B0) bin)
<proof>

lemmas *setBit-no* = *setBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*
simplified if-simps, THEN eq-reflection, standard]

lemmas *clearBit-no* = *clearBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*
simplified if-simps, THEN eq-reflection, standard]

lemma *to-bl-n1*:
to-bl (-1::'a::len0 word) = *replicate (len-of TYPE ('a)) True*
<proof>

lemma *word-msb-n1*: *msb (-1::'a::len word)*
<proof>

declare *word-set-set-same* [*simp*] *word-set-nth* [*simp*]
test-bit-no [*simp*] *word-set-no* [*simp*] *nth-0* [*simp*]
setBit-no [*simp*] *clearBit-no* [*simp*]
word-lsb-no [*simp*] *word-msb-no* [*simp*] *word-msb-n1* [*simp*] *word-lsb-1-0* [*simp*]

lemma *word-set-nth-iff*:
(set-bit w n b = w) = (w !! n = b | n >= size (w::'a::len0 word))
<proof>

lemma *test-bit-2p'*:
w = word-of-int (2 ^ n) ==>
w !! m = (m = n & m < size (w :: 'a :: len word))
<proof>

lemmas *test-bit-2p* = *refl [THEN test-bit-2p', unfolded word-size]*

lemmas *nth-w2p = test-bit-2p* [*unfolded of-int-number-of-eq*
word-of-int [symmetric] of-int-power]

lemma *uint-2p*:
 $(0 :: 'a :: \text{len } \text{word}) < 2 \wedge n \implies \text{uint } (2 \wedge n :: 'a :: \text{len } \text{word}) = 2 \wedge n$
 ⟨*proof*⟩

lemma *word-of-int-2p*: $(\text{word-of-int } (2 \wedge n) :: 'a :: \text{len } \text{word}) = 2 \wedge n$
 ⟨*proof*⟩

lemma *bang-is-le*: $x \ll m \implies 2 \wedge m \leq (x :: 'a :: \text{len } \text{word})$
 ⟨*proof*⟩

lemma *word-clr-le*:
fixes $w :: 'a :: \text{len } 0 \text{ word}$
shows $w \geq \text{set-bit } w \ n \ \text{False}$
 ⟨*proof*⟩

lemma *word-set-ge*:
fixes $w :: 'a :: \text{len } \text{word}$
shows $w \leq \text{set-bit } w \ n \ \text{True}$
 ⟨*proof*⟩

end

13 WordShift: Shifting, Rotating, and Splitting Words

theory *WordShift* **imports** *WordBitwise* **begin**

13.1 Bit shifting

lemma *shiffl1-number* [*simp*] :
 $\text{shiffl1 } (\text{number-of } w) = \text{number-of } (w \ \text{BIT } \text{bit.B0})$
 ⟨*proof*⟩

lemma *shiffl1-0* [*simp*] : $\text{shiffl1 } 0 = 0$
 ⟨*proof*⟩

lemmas *shiffl1-def-u = shiffl1-def* [*folded word-number-of-def*]

lemma *shiffl1-def-s*: $\text{shiffl1 } w = \text{number-of } (\text{sint } w \ \text{BIT } \text{bit.B0})$
 ⟨*proof*⟩

lemma *shiftr1-0* [*simp*] : $\text{shiftr1 } 0 = 0$

<proof>

lemma *sshiftr1-0* [*simp*] : *sshiftr1 0 = 0*
<proof>

lemma *sshiftr1-n1* [*simp*] : *sshiftr1 -1 = -1*
<proof>

lemma *shiftr1-0* [*simp*] : *(0::'a::len0 word) << n = 0*
<proof>

lemma *shiftr-0* [*simp*] : *(0::'a::len0 word) >> n = 0*
<proof>

lemma *sshiftr-0* [*simp*] : *0 >>> n = 0*
<proof>

lemma *sshiftr-n1* [*simp*] : *-1 >>> n = -1*
<proof>

lemma *nth-shiftr1*: *shiftr1 w !! n = (n < size w & n > 0 & w !! (n - 1))*
<proof>

lemma *nth-shiftr1'* [*rule-format*]:
ALL n. ((w::'a::len0 word) << m) !! n = (n < size w & n >= m & w !! (n - m))
<proof>

lemmas *nth-shiftr1 = nth-shiftr1'* [*unfolded word-size*]

lemma *nth-shiftr1*: *shiftr1 w !! n = w !! Suc n*
<proof>

lemma *nth-shiftr*:
∧ n. ((w::'a::len0 word) >> m) !! n = w !! (n + m)
<proof>

lemma *uint-shiftr1*: *uint (shiftr1 w) = bin-rest (uint w)*
<proof>

lemma *nth-sshiftr1*:
shiftr1 w !! n = (if n = size w - 1 then w !! n else w !! Suc n)
<proof>

lemma *nth-sshiftr* [*rule-format*] :
ALL n. shiftr w m !! n = (n < size w &
(if n + m >= size w then w !! (size w - 1) else w !! (n + m)))

<proof>

lemma *shiftr1-div-2*: $\text{uint } (\text{shiftr1 } w) = \text{uint } w \text{ div } 2$
<proof>

lemma *sshiftr1-div-2*: $\text{sint } (\text{sshiftr1 } w) = \text{sint } w \text{ div } 2$
<proof>

lemma *shiftr-div-2n*: $\text{uint } (\text{shiftr } w \ n) = \text{uint } w \text{ div } 2 \wedge n$
<proof>

lemma *sshiftr-div-2n*: $\text{sint } (\text{sshiftr } w \ n) = \text{sint } w \text{ div } 2 \wedge n$
<proof>

13.1.1 shift functions in terms of lists of bools

lemmas *bshiftr1-no-bin* [*simp*] =
bshiftr1-def [**where** *w=number-of w, unfolded to-bl-no-bin, standard*]

lemma *bshiftr1-bl*: $\text{to-bl } (\text{bshiftr1 } b \ w) = b \ \# \ \text{butlast } (\text{to-bl } w)$
<proof>

lemma *shiftrl1-of-bl*: $\text{shiftrl1 } (\text{of-bl } bl) = \text{of-bl } (bl \ @ \ [\text{False}])$
<proof>

lemma *shiftrl1-bl*: $\text{shiftrl1 } (w :: 'a :: \text{len0 word}) = \text{of-bl } (\text{to-bl } w \ @ \ [\text{False}])$
<proof>

lemma *bl-shiftrl1*:
 $\text{to-bl } (\text{shiftrl1 } (w :: 'a :: \text{len word})) = \text{tl } (\text{to-bl } w) \ @ \ [\text{False}]$
<proof>

lemma *shiftr1-bl*: $\text{shiftr1 } w = \text{of-bl } (\text{butlast } (\text{to-bl } w))$
<proof>

lemma *bl-shiftr1*:
 $\text{to-bl } (\text{shiftr1 } (w :: 'a :: \text{len word})) = \text{False} \ \# \ \text{butlast } (\text{to-bl } w)$
<proof>

lemma *shiftrl1-rev*:
 $\text{shiftrl1 } (w :: 'a :: \text{len word}) = \text{word-reverse } (\text{shiftr1 } (\text{word-reverse } w))$
<proof>

lemma *shiftrl-rev*:
 $\text{shiftrl } (w :: 'a :: \text{len word}) \ n = \text{word-reverse } (\text{shiftr } (\text{word-reverse } w) \ n)$
<proof>

lemmas *rev-shiffl* =
shiffl-rev [where $w = \text{word-reverse } w$, *simplified, standard*]

lemmas *shiftr-rev* = *rev-shiffl* [THEN *word-rev-gal'*, *standard*]
lemmas *rev-shiftr* = *shiftr-rev* [THEN *word-rev-gal'*, *standard*]

lemma *bl-sshiftr1*:
 $\text{to-bl } (\text{sshiftr1 } (w :: 'a :: \text{len word})) = \text{hd } (\text{to-bl } w) \# \text{butlast } (\text{to-bl } w)$
 ⟨*proof*⟩

lemma *drop-shiftr*:
 $\text{drop } n (\text{to-bl } ((w :: 'a :: \text{len word}) \gg n)) = \text{take } (\text{size } w - n) (\text{to-bl } w)$
 ⟨*proof*⟩

lemma *drop-sshiftr*:
 $\text{drop } n (\text{to-bl } ((w :: 'a :: \text{len word}) \ggg n)) = \text{take } (\text{size } w - n) (\text{to-bl } w)$
 ⟨*proof*⟩

lemma *take-shiftr* [rule-format] :
 $n \leq \text{size } (w :: 'a :: \text{len word}) \longrightarrow \text{take } n (\text{to-bl } (w \gg n)) =$
 $\text{replicate } n \text{ False}$
 ⟨*proof*⟩

lemma *take-sshiftr'* [rule-format] :
 $n \leq \text{size } (w :: 'a :: \text{len word}) \longrightarrow \text{hd } (\text{to-bl } (w \ggg n)) = \text{hd } (\text{to-bl } w) \ \&$
 $\text{take } n (\text{to-bl } (w \ggg n)) = \text{replicate } n (\text{hd } (\text{to-bl } w))$
 ⟨*proof*⟩

lemmas *hd-sshiftr* = *take-sshiftr'* [THEN *conjunct1*, *standard*]
lemmas *take-sshiftr* = *take-sshiftr'* [THEN *conjunct2*, *standard*]

lemma *atd-lem*: $\text{take } n \text{ xs} = t \implies \text{drop } n \text{ xs} = d \implies \text{xs} = t @ d$
 ⟨*proof*⟩

lemmas *bl-shiftr* = *atd-lem* [OF *take-shiftr drop-shiftr*]
lemmas *bl-sshiftr* = *atd-lem* [OF *take-sshiftr drop-sshiftr*]

lemma *shiffl-of-bl*: $\text{of-bl } bl \ll n = \text{of-bl } (bl @ \text{replicate } n \text{ False})$
 ⟨*proof*⟩

lemma *shiffl-bl*:
 $(w :: 'a :: \text{len0 word}) \ll (n :: \text{nat}) = \text{of-bl } (\text{to-bl } w @ \text{replicate } n \text{ False})$
 ⟨*proof*⟩

lemmas *shiffl-number* [simp] = *shiffl-def* [where $w = \text{number-of } w$, *standard*]

lemma *bl-shiffl*:
 $\text{to-bl } (w \ll n) = \text{drop } n (\text{to-bl } w) @ \text{replicate } (\text{min } (\text{size } w) \ n) \ \text{False}$
 ⟨*proof*⟩

lemma *shiffl-zero-size*:

fixes $x :: 'a::len0$ word

shows $size\ x \leq n \implies x \ll n = 0$

<proof>

lemma *shiffl1-2t*: $shiffl1\ (w :: 'a :: len\ word) = 2 * w$

<proof>

lemma *shiffl1-p*: $shiffl1\ (w :: 'a :: len\ word) = w + w$

<proof>

lemma *shiffl-t2n*: $shiffl\ (w :: 'a :: len\ word)\ n = 2^{\wedge} n * w$

<proof>

lemma *shiftr1-bintr* [*simp*]:

$(shiftr1\ (number-of\ w) :: 'a :: len0\ word) =$

$number-of\ (bin-rest\ (bintrunc\ (len-of\ TYPE\ ('a))\ w))$

<proof>

lemma *sshiftr1-sbintr* [*simp*] :

$(sshiftr1\ (number-of\ w) :: 'a :: len\ word) =$

$number-of\ (bin-rest\ (sbintrunc\ (len-of\ TYPE\ ('a) - 1)\ w))$

<proof>

lemma *shiftr-no'*:

$w = number-of\ bin \implies$

$(w::'a::len0\ word) \gg n = number-of\ ((bin-rest\ ^\ n)\ (bintrunc\ (size\ w)\ bin))$

<proof>

lemma *sshiftr-no'*:

$w = number-of\ bin \implies w \ggg n = number-of\ ((bin-rest\ ^\ n)$

$(sbintrunc\ (size\ w - 1)\ bin))$

<proof>

lemmas *sshiftr-no* [*simp*] =

shiftr-no' [**where** $w = number-of\ w$, *OF refl, unfolded word-size, standard*]

lemmas *shiftr-no* [*simp*] =

shiftr-no' [**where** $w = number-of\ w$, *OF refl, unfolded word-size, standard*]

lemma *shiftr1-bl-of'*:

$us = shiftr1\ (of-bl\ bl) \implies length\ bl \leq size\ us \implies$

$us = of-bl\ (butlast\ bl)$

<proof>

lemmas *shiftr1-bl-of* = *refl* [*THEN shiftr1-bl-of'*, *unfolded word-size*]

lemma *shiftr-bl-of'* [rule-format]:

$us = of-bl\ bl \gg n \implies length\ bl \leq size\ us \dashrightarrow$
 $us = of-bl\ (take\ (length\ bl - n)\ bl)$
 ⟨proof⟩

lemmas *shiftr-bl-of* = refl [THEN *shiftr-bl-of'*, unfolded word-size]

lemmas *shiftr-bl* = word-bl.Rep' [THEN *eq-imp-le*, THEN *shiftr-bl-of*,
 simplified word-size, simplified, THEN *eq-reflection*, standard]

lemma *msb-shift'*: $msb\ (w::'a::len\ word) \leftrightarrow (w \gg (size\ w - 1)) \approx 0$
 ⟨proof⟩

lemmas *msb-shift* = *msb-shift'* [unfolded word-size]

lemma *align-lem-or* [rule-format] :

$ALL\ x\ m.\ length\ x = n + m \dashrightarrow length\ y = n + m \dashrightarrow$
 $drop\ m\ x = replicate\ n\ False \dashrightarrow take\ m\ y = replicate\ m\ False \dashrightarrow$
 $app2\ op\ | \ x\ y = take\ m\ x\ @\ drop\ m\ y$
 ⟨proof⟩

lemma *align-lem-and* [rule-format] :

$ALL\ x\ m.\ length\ x = n + m \dashrightarrow length\ y = n + m \dashrightarrow$
 $drop\ m\ x = replicate\ n\ False \dashrightarrow take\ m\ y = replicate\ m\ False \dashrightarrow$
 $app2\ op\ \&\ x\ y = replicate\ (n + m)\ False$
 ⟨proof⟩

lemma *aligned-bl-add-size'*:

$size\ x - n = m \implies n \leq size\ x \implies drop\ m\ (to-bl\ x) = replicate\ n\ False$
 \implies
 $take\ m\ (to-bl\ y) = replicate\ m\ False \implies$
 $to-bl\ (x + y) = take\ m\ (to-bl\ x)\ @\ drop\ m\ (to-bl\ y)$
 ⟨proof⟩

lemmas *aligned-bl-add-size* = refl [THEN *aligned-bl-add-size'*]

13.1.2 Mask

lemma *nth-mask'*: $m = mask\ n \implies test-bit\ m\ i = (i < n \ \&\ i < size\ m)$
 ⟨proof⟩

lemmas *nth-mask* [simp] = refl [THEN *nth-mask'*]

lemma *mask-bl*: $mask\ n = of-bl\ (replicate\ n\ True)$
 ⟨proof⟩

lemma *mask-bin*: $mask\ n = number-of\ (bintrunc\ n\ Numeral.Min)$
 ⟨proof⟩

lemma *and-mask-bintr*: $w \text{ AND mask } n = \text{number-of } (\text{bintrunc } n \text{ (uint } w))$
 ⟨proof⟩

lemma *and-mask-no*: $\text{number-of } i \text{ AND mask } n = \text{number-of } (\text{bintrunc } n \text{ } i)$
 ⟨proof⟩

lemmas *and-mask-wi* = *and-mask-no* [unfolded word-number-of-def]

lemma *bl-and-mask*:
 $\text{to-bl } (w \text{ AND mask } n :: 'a :: \text{len word}) =$
 $\text{replicate } (\text{len-of TYPE('a)} - n) \text{ False } @$
 $\text{drop } (\text{len-of TYPE('a)} - n) (\text{to-bl } w)$
 ⟨proof⟩

lemmas *and-mask-mod-2p* =
and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]

lemma *and-mask-lt-2p*: $\text{uint } (w \text{ AND mask } n) < 2 \wedge n$
 ⟨proof⟩

lemmas *eq-mod-iff* = *trans* [symmetric, OF int-mod-lem eq-sym-conv]

lemma *mask-eq-iff*: $(w \text{ AND mask } n) = w \leftrightarrow \text{uint } w < 2 \wedge n$
 ⟨proof⟩

lemma *and-mask-dvd*: $2 \wedge n \text{ dvd uint } w = (w \text{ AND mask } n = 0)$
 ⟨proof⟩

lemma *and-mask-dvd-nat*: $2 \wedge n \text{ dvd unat } w = (w \text{ AND mask } n = 0)$
 ⟨proof⟩

lemma *word-2p-lem*:
 $n < \text{size } w \implies w < 2 \wedge n = (\text{uint } (w :: 'a :: \text{len word}) < 2 \wedge n)$
 ⟨proof⟩

lemma *less-mask-eq*: $x < 2 \wedge n \implies x \text{ AND mask } n = (x :: 'a :: \text{len word})$
 ⟨proof⟩

lemmas *mask-eq-iff-w2p* =
trans [OF mask-eq-iff word-2p-lem [symmetric], standard]

lemmas *and-mask-less'* =
iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]

lemma *and-mask-less-size*: $n < \text{size } x \implies x \text{ AND mask } n < 2 \wedge n$
 ⟨proof⟩

lemma *word-mod-2p-is-mask'*:

$c = 2 \wedge n \implies c > 0 \implies x \bmod c = (x :: 'a :: \text{len word}) \text{ AND mask } n$
 ⟨proof⟩

lemmas *word-mod-2p-is-mask* = refl [THEN *word-mod-2p-is-mask'*]

lemma *mask-egs*:

$(a \text{ AND mask } n) + b \text{ AND mask } n = a + b \text{ AND mask } n$
 $a + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - b \text{ AND mask } n = a - b \text{ AND mask } n$
 $a - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $a * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $(b \text{ AND mask } n) * a \text{ AND mask } n = b * a \text{ AND mask } n$
 $(a \text{ AND mask } n) + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $(a \text{ AND mask } n) * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $-(a \text{ AND mask } n) \text{ AND mask } n = -a \text{ AND mask } n$
 $\text{word-succ } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-succ } a \text{ AND mask } n$
 $\text{word-pred } (a \text{ AND mask } n) \text{ AND mask } n = \text{word-pred } a \text{ AND mask } n$
 ⟨proof⟩

lemma *mask-power-eq*:

$(x \text{ AND mask } n) \wedge k \text{ AND mask } n = x \wedge k \text{ AND mask } n$
 ⟨proof⟩

13.1.3 Recast

lemmas *recast-def'* = *recast-def* [*simplified*]

lemmas *recast-def''* = *recast-def'* [*simplified word-size*]

lemmas *recast-no-def* [*simp*] =

recast-def' [**where** $w = \text{number-of } w$, *unfolded word-size*, *standard*]

lemma *to-bl-recast*:

$\text{to-bl } (\text{recast } w :: 'a :: \text{len0 word}) =$
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$
 ⟨proof⟩

lemma *recast-rev-ucast'*:

$cs = [rc, uc] \implies rc = \text{recast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$
 $rc = \text{word-reverse } uc$
 ⟨proof⟩

lemmas *recast-rev-ucast* = *recast-rev-ucast'* [*OF refl refl refl*]

lemmas *recast-ucast* = *recast-rev-ucast*

[**where** $w = \text{word-reverse } w$, *simplified word-rev-rev*, *standard*]

lemmas *ucast-recast* = *recast-rev-ucast* [*THEN word-rev-gal'*, *standard*]

lemmas *ucast-rev-recast* = *recast-ucast* [*THEN word-rev-gal'*, *standard*]

— linking revcast and cast via shift

lemmas *wsst-TYs = source-size target-size word-size*

lemma *revcast-down-uu'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w \gg n)$
 ⟨proof⟩

lemma *revcast-down-us'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w \ggg n)$
 ⟨proof⟩

lemma *revcast-down-su'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{scast } (w \gg n)$
 ⟨proof⟩

lemma *revcast-down-ss'*:

$rc = \text{revcast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{scast } (w \ggg n)$
 ⟨proof⟩

lemmas *revcast-down-uu = refl [THEN revcast-down-uu']*

lemmas *revcast-down-us = refl [THEN revcast-down-us']*

lemmas *revcast-down-su = refl [THEN revcast-down-su']*

lemmas *revcast-down-ss = refl [THEN revcast-down-ss']*

lemma *cast-down-rev*:

$uc = \text{ucast} \implies \text{source-size } uc = \text{target-size } uc + n \implies$
 $uc w = \text{revcast } ((w :: 'a :: \text{len word}) \ll n)$
 ⟨proof⟩

lemma *revcast-up'*:

$rc = \text{revcast} \implies \text{source-size } rc + n = \text{target-size } rc \implies$
 $rc w = (\text{ucast } w :: 'a :: \text{len word}) \ll n$
 ⟨proof⟩

lemmas *revcast-up = refl [THEN revcast-up']*

lemmas *rc1 = revcast-up [THEN*

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *rc2 = revcast-down-uu [THEN*

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *ucast-up =*

rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]

lemmas *ucast-down* =
rc2 [*simplified rev-shiftr revcast-ucast* [*symmetric*]]

13.1.4 Slices

lemmas *slice1-no-bin* [*simp*] =
slice1-def [**where** *w=number-of w, unfolded to-bl-no-bin, standard*]

lemmas *slice-no-bin* [*simp*] =
trans [*OF slice-def* [*THEN meta-eq-to-obj-eq*]
slice1-no-bin [*THEN meta-eq-to-obj-eq*],
unfolded word-size, standard]

lemma *slice1-0* [*simp*] : *slice1 n 0 = 0*
 ⟨*proof*⟩

lemma *slice-0* [*simp*] : *slice n 0 = 0*
 ⟨*proof*⟩

lemma *slice-take'*: *slice n w = of-bl (take (size w - n) (to-bl w))*
 ⟨*proof*⟩

lemmas *slice-take = slice-take'* [*unfolded word-size*]

— *shiftr* to a word of the same size is just *slice*, *slice* is just *shiftr* then *ucast*

lemmas *shiftr-slice = trans*
 [*OF shiftr-bl* [*THEN meta-eq-to-obj-eq*] *slice-take* [*symmetric*], *standard*]

lemma *slice-shiftr*: *slice n w = ucast (w >> n)*
 ⟨*proof*⟩

lemma *nth-slice*:
 (*slice n w :: 'a :: len0 word*) !! *m* =
 (*w* !! (*m* + *n*) & *m* < *len-of TYPE ('a)*)
 ⟨*proof*⟩

lemma *slice1-down-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>
to-bl sl = takefill False fs (drop k (to-bl w))
 ⟨*proof*⟩

lemma *slice1-up-alt'*:
sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>
to-bl sl = takefill False fs (replicate k False @ (to-bl w))
 ⟨*proof*⟩

lemmas *sd1 = slice1-down-alt'* [*OF refl refl, unfolded word-size*]

lemmas *su1 = slice1-up-alt'* [*OF refl refl, unfolded word-size*]

lemmas *slice1-down-alt = le-add-diff-inverse* [*THEN sd1*]

lemmas *slice1-up-alt*s =

le-add-diff-inverse [*symmetric*, *THEN su1*]

le-add-diff-inverse2 [*symmetric*, *THEN su1*]

lemma *ucast-slice1*: *ucast w = slice1 (size w) w*
 ⟨*proof*⟩

lemma *ucast-slice*: *ucast w = slice 0 w*
 ⟨*proof*⟩

lemmas *slice-id = trans [OF ucast-slice [symmetric] ucast-id]*

lemma *revcast-slice1'*:

rc = revcast w ==> slice1 (size rc) w = rc

⟨*proof*⟩

lemmas *revcast-slice1 = refl [THEN revcast-slice1']*

lemma *slice1-tf-tf'*:

to-bl (slice1 n w :: 'a :: len0 word) =

rev (takefill False (len-of TYPE('a)) (rev (takefill False n (to-bl w))))

⟨*proof*⟩

lemmas *slice1-tf-tf = slice1-tf-tf'*

[*THEN word-bl.Rep-inverse'*, *symmetric*, *standard*]

lemma *rev-slice1*:

n + k = len-of TYPE('a) + len-of TYPE('b) ==>

slice1 n (word-reverse w :: 'b :: len0 word) =

word-reverse (slice1 k w :: 'a :: len0 word)

⟨*proof*⟩

lemma *rev-slice'*:

res = slice n (word-reverse w) ==> n + k + size res = size w ==>

res = word-reverse (slice k w)

⟨*proof*⟩

lemmas *rev-slice = refl [THEN rev-slice', unfolded word-size]*

lemmas *sym-notr =*

not-iff [THEN iffD2, THEN not-sym, THEN not-iff [THEN iffD1]]

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:

(sint (x :: 'a :: len word) + sint y = sint (x + y)) =

((((x+y) XOR x) AND ((x+y) XOR y)) >> (size x - 1) = 0)

⟨*proof*⟩

13.2 Split and cat

lemmas *word-split-bin'* = *word-split-def* [*THEN meta-eq-to-obj-eq, standard*]

lemmas *word-cat-bin'* = *word-cat-def* [*THEN meta-eq-to-obj-eq, standard*]

lemma *word-rsplit-no*:

$$\begin{aligned} & (\text{word-rsplit } (\text{number-of } \text{bin} :: 'b :: \text{len0 } \text{word}) :: 'a \text{ word list}) = \\ & \quad \text{map } \text{number-of } (\text{bin-rsplit } (\text{len-of } \text{TYPE}('a :: \text{len})) \\ & \quad \quad (\text{len-of } \text{TYPE}('b), \text{bintrunc } (\text{len-of } \text{TYPE}('b)) \text{ bin})) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*
[*unfolded bin-rsplittl-def bin-rsplit-l [symmetric]*]

lemma *test-bit-cat*:

$$\begin{aligned} & \text{wc} = \text{word-cat } a \ b \implies \text{wc} !! n = (n < \text{size } \text{wc} \ \& \\ & \quad (\text{if } n < \text{size } b \ \text{then } b !! n \ \text{else } a !! (n - \text{size } b))) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *word-cat-bl*: *word-cat a b = of-bl (to-bl a @ to-bl b)*
<proof>

lemma *of-bl-append*:

$$\begin{aligned} & (\text{of-bl } (xs \ @ \ ys) :: 'a :: \text{len } \text{word}) = \text{of-bl } xs * 2^{(\text{length } ys)} + \text{of-bl } ys \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *of-bl-False* [*simp*]:

$$\begin{aligned} & \text{of-bl } (\text{False}\#xs) = \text{of-bl } xs \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *of-bl-True*:

$$\begin{aligned} & (\text{of-bl } (\text{True}\#xs)::'a::\text{len } \text{word}) = 2^{\text{length } xs} + \text{of-bl } xs \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *of-bl-Cons*:

$$\begin{aligned} & \text{of-bl } (x\#xs) = \text{of-bool } x * 2^{\text{length } xs} + \text{of-bl } xs \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *split-uint-lem*: *bin-split n (uint (w :: 'a :: len0 word)) = (a, b) ==>*
a = bintrunc (len-of TYPE('a) - n) a & b = bintrunc (len-of TYPE('a)) b
<proof>

lemma *word-split-bl'*:

$$\begin{aligned} & \text{std} = \text{size } c - \text{size } b \implies (\text{word-split } c = (a, b)) \implies \\ & \quad (a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \ \& \ b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c))) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *word-split-bl*: *std = size c - size b ==>*

$$\begin{aligned} & (a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \ \& \ b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c))) \iff \\ & \quad \text{word-split } c = (a, b) \end{aligned}$$

<proof>

lemma *word-split-bl-eq*:

$(\text{word-split } (c::'a::\text{len word}) :: ('c :: \text{len0 word} * 'd :: \text{len0 word})) =$
 $(\text{of-bl } (\text{take } (\text{len-of TYPE}('a::\text{len}) - \text{len-of TYPE}('d::\text{len0})) (\text{to-bl } c)),$
 $\text{of-bl } (\text{drop } (\text{len-of TYPE}('a) - \text{len-of TYPE}('d)) (\text{to-bl } c)))$
<proof>

lemma *test-bit-split'*:

$\text{word-split } c = (a, b) \dashrightarrow (\text{ALL } n \ m. b !! n = (n < \text{size } b \ \& \ c !! n) \ \&$
 $a !! m = (m < \text{size } a \ \& \ c !! (m + \text{size } b)))$
<proof>

lemmas *test-bit-split =*

test-bit-split' [THEN mp, simplified all-simps, standard]

lemma *test-bit-split-eq*: $\text{word-split } c = (a, b) \dashrightarrow$

$(\text{ALL } n::\text{nat}. b !! n = (n < \text{size } b \ \& \ c !! n) \ \&$
 $(\text{ALL } m::\text{nat}. a !! m = (m < \text{size } a \ \& \ c !! (m + \text{size } b))))$
<proof>

lemma *word-cat-id*: $\text{word-cat } a \ b = b$

<proof>

lemma *word-cat-hom*:

$\text{len-of TYPE}('a::\text{len0}) \leq \text{len-of TYPE}('b::\text{len0}) + \text{len-of TYPE}('c::\text{len0})$
 \implies
 $(\text{word-cat } (\text{word-of-int } w :: 'b \ \text{word}) (b :: 'c \ \text{word}) :: 'a \ \text{word}) =$
 $\text{word-of-int } (\text{bin-cat } w \ (\text{size } b) \ (\text{uint } b))$
<proof>

lemma *word-cat-split-alt*:

$\text{size } w \leq \text{size } u + \text{size } v \implies \text{word-split } w = (u, v) \implies \text{word-cat } u \ v = w$
<proof>

lemmas *word-cat-split-size =*

sym [THEN [2] *word-cat-split-alt* [symmetric], standard]

13.2.1 Split and slice

lemma *split-slices*:

$\text{word-split } w = (u, v) \implies u = \text{slice } (\text{size } v) \ w \ \& \ v = \text{slice } 0 \ w$
<proof>

lemma *slice-cat1'*:

$wc = \text{word-cat } a \ b \implies \text{size } wc \geq \text{size } a + \text{size } b \implies \text{slice } (\text{size } b) \ wc = a$
<proof>

lemmas *slice-cat1 = refl* [THEN *slice-cat1'*]

lemmas *slice-cat2 = trans* [OF *slice-id word-cat-id*]

lemma *cat-slices*:

$a = \text{slice } n \ c \implies b = \text{slice } 0 \ c \implies n = \text{size } b \implies$
 $\text{size } a + \text{size } b \geq \text{size } c \implies \text{word-cat } a \ b = c$
 ⟨proof⟩

lemma *word-split-cat-alt*:

$w = \text{word-cat } u \ v \implies \text{size } u + \text{size } v \leq \text{size } w \implies \text{word-split } w = (u, v)$
 ⟨proof⟩

lemmas *word-cat-bl-no-bin* [simp] =

word-cat-bl [where $a = \text{number-of } a$
 and $b = \text{number-of } b$,
 unfolded to-bl-no-bin, standard]

lemmas *word-split-bl-no-bin* [simp] =

word-split-bl-eq [where $c = \text{number-of } c$, unfolded to-bl-no-bin, standard]

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

lemma *word-rsplit-same*: $\text{word-rsplit } w = [w]$

⟨proof⟩

lemma *word-rsplit-empty-iff-size*:

$(\text{word-rsplit } w = []) = (\text{size } w = 0)$
 ⟨proof⟩

lemma *test-bit-rsplit*:

$sw = \text{word-rsplit } w \implies m < \text{size } (\text{hd } sw :: 'a :: \text{len } \text{word}) \implies$
 $k < \text{length } sw \implies (\text{rev } sw ! k) !! m = (w !! (k * \text{size } (\text{hd } sw) + m))$
 ⟨proof⟩

lemma *word-rcat-bl*: $\text{word-rcat } wl == \text{of-bl } (\text{concat } (\text{map } \text{to-bl } wl))$

⟨proof⟩

lemma *size-rcat-lem'*:

$\text{size } (\text{concat } (\text{map } \text{to-bl } wl)) = \text{length } wl * \text{size } (\text{hd } wl)$
 ⟨proof⟩

lemmas *size-rcat-lem* = *size-rcat-lem'* [unfolded word-size]

lemmas *td-gal-lt-len* = *len-gt-0* [THEN *td-gal-lt*, standard]

lemma *nth-rcat-lem'* [rule-format] :

$sw = \text{size } (\text{hd } wl :: 'a :: \text{len } \text{word}) \implies (\text{ALL } n. n < \text{size } wl * sw \longrightarrow$
 $\text{rev } (\text{concat } (\text{map } \text{to-bl } wl)) ! n =$
 $\text{rev } (\text{to-bl } (\text{rev } wl ! (n \text{ div } sw))) ! (n \text{ mod } sw))$
 ⟨proof⟩

lemmas *nth-rcat-lem* = refl [THEN *nth-rcat-lem'*, *unfolded word-size*]

lemma *test-bit-rcat*:

$sw = size (hd\ wl :: 'a :: len\ word) ==> rc = word-rcat\ wl ==> rc\ !!\ n =$
 $(n < size\ rc \& n\ div\ sw < size\ wl \& (rev\ wl) ! (n\ div\ sw) !! (n\ mod\ sw))$
 ⟨proof⟩

lemma *foldl-eq-foldr* [rule-format] :

ALL x . $foldl\ op + x\ xs = foldr\ op + (x\ \# \ xs) (0 :: 'a :: comm-monoid-add)$
 ⟨proof⟩

lemmas *test-bit-cong* = arg-cong [where $f = test-bit$, THEN *fun-cong*]

lemmas *test-bit-rsplit-alt* =

trans [OF *nth-rev-alt* [THEN *test-bit-cong*]
test-bit-rsplit [OF *refl asm-rl diff-Suc-less*]]

— lazy way of expressing that u and v , and su and sv , have same types

lemma *word-rsplit-len-indep'*:

$[u,v] = p ==> [su,sv] = q ==> word-rsplit\ u = su ==>$
 $word-rsplit\ v = sv ==> length\ su = length\ sv$
 ⟨proof⟩

lemmas *word-rsplit-len-indep* = *word-rsplit-len-indep'* [OF *refl refl refl refl*]

lemma *length-word-rsplit-size*:

$n = len-of\ TYPE ('a :: len) ==>$
 $(length (word-rsplit\ w :: 'a\ word\ list) <= m) = (size\ w <= m * n)$
 ⟨proof⟩

lemmas *length-word-rsplit-lt-size* =

length-word-rsplit-size [unfolded *Not-eq-iff linorder-not-less* [symmetric]]

lemma *length-word-rsplit-exp-size*:

$n = len-of\ TYPE ('a :: len) ==>$
 $length (word-rsplit\ w :: 'a\ word\ list) = (size\ w + n - 1) div\ n$
 ⟨proof⟩

lemma *length-word-rsplit-even-size*:

$n = len-of\ TYPE ('a :: len) ==> size\ w = m * n ==>$
 $length (word-rsplit\ w :: 'a\ word\ list) = m$
 ⟨proof⟩

lemmas *length-word-rsplit-exp-size'* = refl [THEN *length-word-rsplit-exp-size*]

lemmas *tdle* = *iffD2* [OF *split-div-lemma refl*, THEN *conjunct1*]

lemmas *dtle* = *xtr4* [OF *tdle mult-commute*]

lemma *word-rcat-rsplit*: $\text{word-rcat } (\text{word-rsplit } w) = w$
 ⟨*proof*⟩

lemma *size-word-rsplit-rcat-size'*:
 $\text{word-rcat } (ws :: 'a :: \text{len word list}) = \text{frcw} \implies$
 $\text{size frcw} = \text{length } ws * \text{len-of TYPE } ('a) \implies$
 $\text{size } (\text{hd } [\text{word-rsplit } \text{frcw}, ws]) = \text{size } ws$
 ⟨*proof*⟩

lemmas *size-word-rsplit-rcat-size =*
size-word-rsplit-rcat-size' [simplified]

lemma *msrevs*:
fixes $n::\text{nat}$
shows $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$
and $(k * n + m) \text{ mod } n = m \text{ mod } n$
 ⟨*proof*⟩

lemma *word-rsplit-rcat-size'*:
 $\text{word-rcat } (ws :: 'a :: \text{len word list}) = \text{frcw} \implies$
 $\text{size frcw} = \text{length } ws * \text{len-of TYPE } ('a) \implies \text{word-rsplit } \text{frcw} = ws$
 ⟨*proof*⟩

lemmas *word-rsplit-rcat-size = refl [THEN word-rsplit-rcat-size']*

13.3 Rotation

lemmas *rotater-0' [simp] = rotater-def [where n = 0, simplified]*

lemmas *word-rot-defs = word-roti-def word-rotr-def word-rotl-def*

lemma *rotate-eq-mod*:
 $m \text{ mod length } xs = n \text{ mod length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$
 ⟨*proof*⟩

lemmas *rotate-egs [standard] =*
trans [OF rotate0 [THEN fun-cong] id-apply]
rotate-rotate [symmetric]
rotate-id
rotate-conv-mod
rotate-eq-mod

13.3.1 Rotation of list to right

lemma *rotate1-rl'*: $\text{rotater1 } (l @ [a]) = a \# l$
 ⟨*proof*⟩

lemma *rotate1-rl [simp]* : $\text{rotater1 } (\text{rotate1 } l) = l$
 ⟨*proof*⟩

lemma *rotate1-lr* [*simp*] : $\text{rotate1} (\text{rotater1 } l) = l$
 ⟨*proof*⟩

lemma *rotater1-rev'*: $\text{rotater1} (\text{rev } xs) = \text{rev} (\text{rotate1 } xs)$
 ⟨*proof*⟩

lemma *rotater-rev'*: $\text{rotater } n (\text{rev } xs) = \text{rev} (\text{rotate } n xs)$
 ⟨*proof*⟩

lemmas *rotater-rev = rotater-rev'* [**where** $xs = \text{rev } ys$, *simplified*, *standard*]

lemma *rotater-drop-take*:
 $\text{rotater } n xs =$
 $\text{drop} (\text{length } xs - n \text{ mod } \text{length } xs) xs @$
 $\text{take} (\text{length } xs - n \text{ mod } \text{length } xs) xs$
 ⟨*proof*⟩

lemma *rotater-Suc* [*simp*] :
 $\text{rotater} (\text{Suc } n) xs = \text{rotater1} (\text{rotater } n xs)$
 ⟨*proof*⟩

lemma *rotate-inv-plus* [*rule-format*] :
 ALL $k. k = m + n \longrightarrow \text{rotater } k (\text{rotate } n xs) = \text{rotater } m xs \ \&$
 $\text{rotate } k (\text{rotater } n xs) = \text{rotate } m xs \ \&$
 $\text{rotater } n (\text{rotate } k xs) = \text{rotate } m xs \ \&$
 $\text{rotate } n (\text{rotater } k xs) = \text{rotater } m xs$
 ⟨*proof*⟩

lemmas *rotate-inv-rel = le-add-diff-inverse2* [*symmetric*, *THEN rotate-inv-plus*]

lemmas *rotate-inv-eq = order-refl* [*THEN rotate-inv-rel*, *simplified*]

lemmas *rotate-lr* [*simp*] = *rotate-inv-eq* [*THEN conjunct1*, *standard*]

lemmas *rotate-rl* [*simp*] =
rotate-inv-eq [*THEN conjunct2*, *THEN conjunct1*, *standard*]

lemma *rotate-gal*: $(\text{rotater } n xs = ys) = (\text{rotate } n ys = xs)$
 ⟨*proof*⟩

lemma *rotate-gal'*: $(ys = \text{rotater } n xs) = (xs = \text{rotate } n ys)$
 ⟨*proof*⟩

lemma *length-rotater* [*simp*]:
 $\text{length} (\text{rotater } n xs) = \text{length } xs$
 ⟨*proof*⟩

lemmas *rrs0 = rotate-eqs* [*THEN restrict-to-left*,
simplified rotate-gal [*symmetric*] *rotate-gal'* [*symmetric*], *standard*]

lemmas $rrs1 = rrs0$ [THEN refl [THEN rev-iffD1]]
lemmas $rotater-eqs = rrs1$ [simplified length-rotater, standard]
lemmas $rotater-0 = rotater-eqs$ (1)
lemmas $rotater-add = rotater-eqs$ (2)

13.3.2 map, app2, commuting with rotate(r)

lemma $last-map$: $xs \sim = [] \implies last (map f xs) = f (last xs)$
 ⟨proof⟩

lemma $butlast-map$:
 $xs \sim = [] \implies butlast (map f xs) = map f (butlast xs)$
 ⟨proof⟩

lemma $rotater1-map$: $rotater1 (map f xs) = map f (rotater1 xs)$
 ⟨proof⟩

lemma $rotater-map$:
 $rotater n (map f xs) = map f (rotater n xs)$
 ⟨proof⟩

lemma $but-last-zip$ [rule-format] :
 ALL ys . $length xs = length ys \implies xs \sim = [] \implies$
 $last (zip xs ys) = (last xs, last ys) \ \&$
 $butlast (zip xs ys) = zip (butlast xs) (butlast ys)$
 ⟨proof⟩

lemma $but-last-app2$ [rule-format] :
 ALL ys . $length xs = length ys \implies xs \sim = [] \implies$
 $last (app2 f xs ys) = f (last xs) (last ys) \ \&$
 $butlast (app2 f xs ys) = app2 f (butlast xs) (butlast ys)$
 ⟨proof⟩

lemma $rotater1-zip$:
 $length xs = length ys \implies$
 $rotater1 (zip xs ys) = zip (rotater1 xs) (rotater1 ys)$
 ⟨proof⟩

lemma $rotater1-app2$:
 $length xs = length ys \implies$
 $rotater1 (app2 f xs ys) = app2 f (rotater1 xs) (rotater1 ys)$
 ⟨proof⟩

lemmas $lrth =$
 $box-equals$ [OF $asm-rl$ length-rotater [symmetric]
 $length-rotater$ [symmetric],
 THEN $rotater1-app2$]

lemma $rotater-app2$:

$length\ xs = length\ ys ==>$
 $rotater\ n\ (app2\ f\ xs\ ys) = app2\ f\ (rotater\ n\ xs)\ (rotater\ n\ ys)$
 ⟨proof⟩

lemma *rotate1-app2*:
 $length\ xs = length\ ys ==>$
 $rotate1\ (app2\ f\ xs\ ys) = app2\ f\ (rotate1\ xs)\ (rotate1\ ys)$
 ⟨proof⟩

lemmas *lth = box-equals* [OF *asm-rl length-rotate* [symmetric]
length-rotate [symmetric], THEN *rotate1-app2*]

lemma *rotate-app2*:
 $length\ xs = length\ ys ==>$
 $rotate\ n\ (app2\ f\ xs\ ys) = app2\ f\ (rotate\ n\ xs)\ (rotate\ n\ ys)$
 ⟨proof⟩

lemma *to-bl-rotl*:
 $to-bl\ (word-rotl\ n\ w) = rotate\ n\ (to-bl\ w)$
 ⟨proof⟩

lemmas *blrs0 = rotate-eqs* [THEN *to-bl-rotl* [THEN *trans*]]

lemmas *word-rotl-eqs =*
blrs0 [simplified *word-bl.Rep'* *word-bl.Rep-inject to-bl-rotl* [symmetric]]

lemma *to-bl-rotr*:
 $to-bl\ (word-rotr\ n\ w) = rotater\ n\ (to-bl\ w)$
 ⟨proof⟩

lemmas *brrs0 = rotater-eqs* [THEN *to-bl-rotr* [THEN *trans*]]

lemmas *word-rotr-eqs =*
brrs0 [simplified *word-bl.Rep'* *word-bl.Rep-inject to-bl-rotr* [symmetric]]

declare *word-rotr-eqs* (1) [simp]

declare *word-rotl-eqs* (1) [simp]

lemma
word-rot-rl [simp]:
 $word-rotl\ k\ (word-rotr\ k\ v) = v$ **and**
word-rot-lr [simp]:
 $word-rotr\ k\ (word-rotl\ k\ v) = v$
 ⟨proof⟩

lemma
word-rot-gal:
 $(word-rotr\ n\ v = w) = (word-rotl\ n\ w = v)$ **and**
word-rot-gal':

$(w = \text{word-rottr } n \ v) = (v = \text{word-rotl } n \ w)$
 ⟨proof⟩

lemma *word-rottr-rev*:

$\text{word-rottr } n \ w = \text{word-reverse } (\text{word-rotl } n \ (\text{word-reverse } w))$
 ⟨proof⟩

lemma *word-roti-0* [simp]: $\text{word-roti } 0 \ w = w$
 ⟨proof⟩

lemmas *abl-cong* = *arg-cong* [where $f = \text{of-bl}$]

lemma *word-roti-add*:

$\text{word-roti } (m + n) \ w = \text{word-roti } m \ (\text{word-roti } n \ w)$
 ⟨proof⟩

lemma *word-roti-conv-mod'*: $\text{word-roti } n \ w = \text{word-roti } (n \bmod \text{int } (\text{size } w)) \ w$
 ⟨proof⟩

lemmas *word-roti-conv-mod* = *word-roti-conv-mod'* [unfolded *word-size*]

13.3.3 Word rotation commutes with bit-wise operations

locale *word-rotate*

context *word-rotate*

begin

lemmas *word-rot-defs'* = *to-bl-rotl to-bl-rottr*

lemmas *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

lemmas *lbl-lbl* = *trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

lemmas *ths-app2* [OF *lbl-lbl*] = *rotate-app2 rotater-app2*

lemmas *ths-map* [where $xs = \text{to-bl } v$] = *rotate-map rotater-map*

lemmas *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-app2 ths-map*

lemma *word-rot-logs*:

$\text{word-rotl } n \ (\text{NOT } v) = \text{NOT } \text{word-rotl } n \ v$
 $\text{word-rottr } n \ (\text{NOT } v) = \text{NOT } \text{word-rottr } n \ v$
 $\text{word-rotl } n \ (x \ \text{AND } y) = \text{word-rotl } n \ x \ \text{AND } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{AND } y) = \text{word-rottr } n \ x \ \text{AND } \text{word-rottr } n \ y$
 $\text{word-rotl } n \ (x \ \text{OR } y) = \text{word-rotl } n \ x \ \text{OR } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{OR } y) = \text{word-rottr } n \ x \ \text{OR } \text{word-rottr } n \ y$
 $\text{word-rotl } n \ (x \ \text{XOR } y) = \text{word-rotl } n \ x \ \text{XOR } \text{word-rotl } n \ y$
 $\text{word-rottr } n \ (x \ \text{XOR } y) = \text{word-rottr } n \ x \ \text{XOR } \text{word-rottr } n \ y$

<proof>
end

lemmas *word-rot-logs* = *word-rotate.word-rot-logs*

lemmas *bl-word-rotl-dt* = *trans* [*OF to-bl-rotl rotate-drop-take*,
simplified word-bl.Rep', *standard*]

lemmas *bl-word-rotr-dt* = *trans* [*OF to-bl-rotr rotater-drop-take*,
simplified word-bl.Rep', *standard*]

lemma *bl-word-roti-dt'*:
 $n = \text{nat } ((- i) \text{ mod int } (\text{size } (w :: 'a :: \text{len word}))) \implies$
 $\text{to-bl } (\text{word-roti } i w) = \text{drop } n (\text{to-bl } w) @ \text{take } n (\text{to-bl } w)$
<proof>

lemmas *bl-word-roti-dt* = *bl-word-roti-dt'* [*unfolded word-size*]

lemmas *word-rotl-dt* = *bl-word-rotl-dt*
 [*THEN word-bl.Rep-inverse'* [*symmetric*], *standard*]

lemmas *word-rotr-dt* = *bl-word-rotr-dt*
 [*THEN word-bl.Rep-inverse'* [*symmetric*], *standard*]

lemmas *word-roti-dt* = *bl-word-roti-dt*
 [*THEN word-bl.Rep-inverse'* [*symmetric*], *standard*]

lemma *word-rotx-0* [*simp*] : *word-rotr* *i* *0* = *0* & *word-rotl* *i* *0* = *0*
<proof>

lemma *word-roti-0'* [*simp*] : *word-roti* *n* *0* = *0*
<proof>

lemmas *word-rotr-dt-no-bin'* [*simp*] =
word-rotr-dt [**where** *w=number-of* *w*, *unfolded to-bl-no-bin*, *standard*]

lemmas *word-rotl-dt-no-bin'* [*simp*] =
word-rotl-dt [**where** *w=number-of* *w*, *unfolded to-bl-no-bin*, *standard*]

declare *word-roti-def* [*simp*]

end

14 Boolean-Algebra: Boolean Algebras

theory *Boolean-Algebra*
imports *Main*
begin

locale *boolean* =
fixes *conj* :: 'a \Rightarrow 'a \Rightarrow 'a (**infixr** \sqcap 70)
fixes *disj* :: 'a \Rightarrow 'a \Rightarrow 'a (**infixr** \sqcup 65)
fixes *compl* :: 'a \Rightarrow 'a (\sim - [81] 80)
fixes *zero* :: 'a (**0**)
fixes *one* :: 'a (**1**)
assumes *conj-assoc*: $(x \sqcap y) \sqcap z = x \sqcap (y \sqcap z)$
assumes *disj-assoc*: $(x \sqcup y) \sqcup z = x \sqcup (y \sqcup z)$
assumes *conj-commute*: $x \sqcap y = y \sqcap x$
assumes *disj-commute*: $x \sqcup y = y \sqcup x$
assumes *conj-disj-distrib*: $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$
assumes *disj-conj-distrib*: $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$
assumes *conj-one-right* [*simp*]: $x \sqcap \mathbf{1} = x$
assumes *disj-zero-right* [*simp*]: $x \sqcup \mathbf{0} = x$
assumes *conj-cancel-right* [*simp*]: $x \sqcap \sim x = \mathbf{0}$
assumes *disj-cancel-right* [*simp*]: $x \sqcup \sim x = \mathbf{1}$
begin

lemmas *disj-ac* =
disj-assoc disj-commute
mk-left-commute [**where** 'a = 'a, of *disj*, OF *disj-assoc disj-commute*]

lemmas *conj-ac* =
conj-assoc conj-commute
mk-left-commute [**where** 'a = 'a, of *conj*, OF *conj-assoc conj-commute*]

lemma *dual*: *boolean disj conj compl one zero*
<proof>

14.1 Complement

lemma *complement-unique*:

assumes 1: $a \sqcap x = \mathbf{0}$
assumes 2: $a \sqcup x = \mathbf{1}$
assumes 3: $a \sqcap y = \mathbf{0}$
assumes 4: $a \sqcup y = \mathbf{1}$
shows $x = y$

<proof>

lemma *compl-unique*: $\llbracket x \sqcap y = \mathbf{0}; x \sqcup y = \mathbf{1} \rrbracket \Longrightarrow \sim x = y$
<proof>

lemma *double-compl* [*simp*]: $\sim(\sim x) = x$
<proof>

lemma *compl-eq-compl-iff* [*simp*]: $(\sim x = \sim y) = (x = y)$
<proof>

14.2 Conjunction

lemma *conj-absorb* [*simp*]: $x \sqcap x = x$
 ⟨*proof*⟩

lemma *conj-zero-right* [*simp*]: $x \sqcap \mathbf{0} = \mathbf{0}$
 ⟨*proof*⟩

lemma *compl-one* [*simp*]: $\sim \mathbf{1} = \mathbf{0}$
 ⟨*proof*⟩

lemma *conj-zero-left* [*simp*]: $\mathbf{0} \sqcap x = \mathbf{0}$
 ⟨*proof*⟩

lemma *conj-one-left* [*simp*]: $\mathbf{1} \sqcap x = x$
 ⟨*proof*⟩

lemma *conj-cancel-left* [*simp*]: $\sim x \sqcap x = \mathbf{0}$
 ⟨*proof*⟩

lemma *conj-left-absorb* [*simp*]: $x \sqcap (x \sqcap y) = x \sqcap y$
 ⟨*proof*⟩

lemma *conj-disj-distrib2*:
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$
 ⟨*proof*⟩

lemmas *conj-disj-distrib* =
conj-disj-distrib conj-disj-distrib2

14.3 Disjunction

lemma *disj-absorb* [*simp*]: $x \sqcup x = x$
 ⟨*proof*⟩

lemma *disj-one-right* [*simp*]: $x \sqcup \mathbf{1} = \mathbf{1}$
 ⟨*proof*⟩

lemma *compl-zero* [*simp*]: $\sim \mathbf{0} = \mathbf{1}$
 ⟨*proof*⟩

lemma *disj-zero-left* [*simp*]: $\mathbf{0} \sqcup x = x$
 ⟨*proof*⟩

lemma *disj-one-left* [*simp*]: $\mathbf{1} \sqcup x = \mathbf{1}$
 ⟨*proof*⟩

lemma *disj-cancel-left* [*simp*]: $\sim x \sqcup x = \mathbf{1}$
 ⟨*proof*⟩

lemma *disj-left-absorb* [*simp*]: $x \sqcup (x \sqcup y) = x \sqcup y$
 ⟨*proof*⟩

lemma *disj-conj-distrib2*:
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$
 ⟨*proof*⟩

lemmas *disj-conj-distrib* =
disj-conj-distrib disj-conj-distrib2

14.4 De Morgan’s Laws

lemma *de-Morgan-conj* [*simp*]: $\sim (x \sqcap y) = \sim x \sqcup \sim y$
 ⟨*proof*⟩

lemma *de-Morgan-disj* [*simp*]: $\sim (x \sqcup y) = \sim x \sqcap \sim y$
 ⟨*proof*⟩

end

14.5 Symmetric Difference

locale *boolean-xor* = *boolean* +
fixes *xor* :: 'a => 'a => 'a (**infixr** \oplus 65)
assumes *xor-def*: $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$
begin

lemma *xor-def2*:
 $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$
 ⟨*proof*⟩

lemma *xor-commute*: $x \oplus y = y \oplus x$
 ⟨*proof*⟩

lemma *xor-assoc*: $(x \oplus y) \oplus z = x \oplus (y \oplus z)$
 ⟨*proof*⟩

lemmas *xor-ac* =
xor-assoc xor-commute
mk-left-commute [**where** 'a = 'a, of *xor*, OF *xor-assoc xor-commute*]

lemma *xor-zero-right* [*simp*]: $x \oplus \mathbf{0} = x$
 ⟨*proof*⟩

lemma *xor-zero-left* [*simp*]: $\mathbf{0} \oplus x = x$
 ⟨*proof*⟩

lemma *xor-one-right* [*simp*]: $x \oplus \mathbf{1} = \sim x$
 ⟨*proof*⟩

lemma *xor-one-left* [*simp*]: $\mathbf{1} \oplus x = \sim x$
 ⟨*proof*⟩

lemma *xor-self* [*simp*]: $x \oplus x = \mathbf{0}$
 ⟨*proof*⟩

lemma *xor-left-self* [*simp*]: $x \oplus (x \oplus y) = y$
 ⟨*proof*⟩

lemma *xor-compl-left*: $\sim x \oplus y = \sim (x \oplus y)$
 ⟨*proof*⟩

lemma *xor-compl-right*: $x \oplus \sim y = \sim (x \oplus y)$
 ⟨*proof*⟩

lemma *xor-cancel-right* [*simp*]: $x \oplus \sim x = \mathbf{1}$
 ⟨*proof*⟩

lemma *xor-cancel-left* [*simp*]: $\sim x \oplus x = \mathbf{1}$
 ⟨*proof*⟩

lemma *conj-xor-distrib*: $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$
 ⟨*proof*⟩

lemma *conj-xor-distrib2*:
 $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$
 ⟨*proof*⟩

lemmas *conj-xor-distrib* =
conj-xor-distrib conj-xor-distrib2

end

end

15 WordGenLib: Miscellaneous Library for Words

theory *WordGenLib* **imports** *WordShift Boolean-Algebra*
begin

declare *of-nat-2p* [*simp*]

lemma *word-int-cases*:
 $\llbracket \bigwedge n. \llbracket (x :: 'a::\text{len } 0 \text{ word}) = \text{word-of-int } n; 0 \leq n; n < 2^{\text{len-of TYPE('a)}} \rrbracket \implies P \rrbracket$
 $\implies P$
 ⟨*proof*⟩

lemma *word-nat-cases* [*cases type: word*]:

$$\llbracket \bigwedge n. \llbracket (x :: 'a::len\ word) = of\ nat\ n; n < 2^{len\ of\ TYPE('a)} \rrbracket \implies P \rrbracket$$

$$\implies P$$
<proof>

lemma *max-word-eq*:

$$(max\ word :: 'a::len\ word) = 2^{len\ of\ TYPE('a)} - 1$$
<proof>

lemma *max-word-max* [*simp,intro!*]:

$$n \leq max\ word$$
<proof>

lemma *word-of-int-2p-len*:

$$word\ of\ int\ (2^{len\ of\ TYPE('a)}) = (0 :: 'a::len0\ word)$$
<proof>

lemma *word-pow-0*:

$$(2 :: 'a::len\ word)^{len\ of\ TYPE('a)} = 0$$
<proof>

lemma *max-word-wrap*: $x + 1 = 0 \implies x = max\ word$
<proof>

lemma *max-word-minus*:

$$max\ word = (-1 :: 'a::len\ word)$$
<proof>

lemma *max-word-bl* [*simp*]:

$$to\ bl\ (max\ word :: 'a::len\ word) = replicate\ (len\ of\ TYPE('a))\ True$$
<proof>

lemma *max-test-bit* [*simp*]:

$$(max\ word :: 'a::len\ word) !! n = (n < len\ of\ TYPE('a))$$
<proof>

lemma *word-and-max* [*simp*]:

$$x\ AND\ max\ word = x$$
<proof>

lemma *word-or-max* [*simp*]:

$$x\ OR\ max\ word = max\ word$$
<proof>

lemma *word-ao-dist2*:

$$x\ AND\ (y\ OR\ z) = x\ AND\ y\ OR\ x\ AND\ (z :: 'a::len0\ word)$$
<proof>

lemma *word-oa-dist2*:

$x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z::'a::len0 \text{ word}))$
 ⟨proof⟩

lemma *word-and-not* [simp]:
 $x \text{ AND } \text{NOT } x = (0::'a::len0 \text{ word})$
 ⟨proof⟩

lemma *word-or-not* [simp]:
 $x \text{ OR } \text{NOT } x = \text{max-word}$
 ⟨proof⟩

lemma *word-boolean*:
 $\text{boolean } (op \text{ AND}) (op \text{ OR}) \text{ bitNOT } 0 \text{ max-word}$
 ⟨proof⟩

interpretation *word-bool-alg*:
 $\text{boolean } [op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word}]$
 ⟨proof⟩

lemma *word-xor-and-or*:
 $x \text{ XOR } y = x \text{ AND } \text{NOT } y \text{ OR } \text{NOT } x \text{ AND } (y::'a::len0 \text{ word})$
 ⟨proof⟩

interpretation *word-bool-alg*:
 $\text{boolean-xor } [op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word } op \text{ XOR}]$
 ⟨proof⟩

lemma *shiftr-0* [iff]:
 $(x::'a::len0 \text{ word}) \gg 0 = x$
 ⟨proof⟩

lemma *shiftr-0* [simp]:
 $(x :: 'a :: len \text{ word}) \ll 0 = x$
 ⟨proof⟩

lemma *shiftr-1* [simp]:
 $(1::'a::len \text{ word}) \ll n = 2^n$
 ⟨proof⟩

lemma *uint-lt-0* [simp]:
 $\text{uint } x < 0 = \text{False}$
 ⟨proof⟩

lemma *shiftr1-1* [simp]:
 $\text{shiftr1 } (1::'a::len \text{ word}) = 0$
 ⟨proof⟩

lemma *shiftr-1* [simp]:
 $(1::'a::len \text{ word}) \gg n = (\text{if } n = 0 \text{ then } 1 \text{ else } 0)$

<proof>

lemma *word-less-1* [*simp*]:
 $((x::'a::len\ word) < 1) = (x = 0)$
<proof>

lemma *to-bl-mask*:
 $to-bl\ (mask\ n\ ::\ 'a::len\ word) =$
 $replicate\ (len-of\ TYPE('a) - n)\ False\ @$
 $replicate\ (min\ (len-of\ TYPE('a))\ n)\ True$
<proof>

lemma *map-replicate-True*:
 $n = length\ xs ==>$
 $map\ (\lambda(x,y).\ x\ \&\ y)\ (zip\ xs\ (replicate\ n\ True)) = xs$
<proof>

lemma *map-replicate-False*:
 $n = length\ xs ==> map\ (\lambda(x,y).\ x\ \&\ y)$
 $(zip\ xs\ (replicate\ n\ False)) = replicate\ n\ False$
<proof>

lemma *bl-and-mask*:
fixes $w\ ::\ 'a::len\ word$
fixes n
defines $n' \equiv len-of\ TYPE('a) - n$
shows $to-bl\ (w\ AND\ mask\ n) = replicate\ n'\ False\ @\ drop\ n'\ (to-bl\ w)$
<proof>

lemma *drop-rev-takefill*:
 $length\ xs \leq n ==>$
 $drop\ (n - length\ xs)\ (rev\ (takefill\ False\ n\ (rev\ xs))) = xs$
<proof>

lemma *map-nth-0* [*simp*]:
 $map\ (op\ !!\ (0::'a::len0\ word))\ xs = replicate\ (length\ xs)\ False$
<proof>

lemma *uint-plus-if-size*:
 $uint\ (x + y) =$
 $(if\ uint\ x + uint\ y < 2^{size\ x}\ then$
 $uint\ x + uint\ y$
 $else$
 $uint\ x + uint\ y - 2^{size\ x})$
<proof>

lemma *unat-plus-if-size*:
 $unat\ (x + (y::'a::len\ word)) =$
 $(if\ unat\ x + unat\ y < 2^{size\ x}\ then$

```

      unat x + unat y
    else
      unat x + unat y - 2size x
  ⟨proof⟩

```

lemma *word-neq-0-conv* [simp]:
fixes $w :: 'a :: \text{len word}$
shows $(w \neq 0) = (0 < w)$
 ⟨proof⟩

lemma *max-lt*:
 $\text{unat } (\max a b \text{ div } c) = \text{unat } (\max a b) \text{ div unat } (c :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *uint-sub-if-size*:
 $\text{uint } (x - y) =$
(if $\text{uint } y \leq \text{uint } x$ *then*
 $\text{uint } x - \text{uint } y$
else
 $\text{uint } x - \text{uint } y + 2^{\text{size } x}$)
 ⟨proof⟩

lemma *unat-sub-simple*:
 $x \leq y \implies \text{unat } (y - x) = \text{unat } y - \text{unat } x$
 ⟨proof⟩

lemmas *unat-sub = unat-sub-simple*

lemma *word-less-sub1*:
fixes $x :: 'a :: \text{len word}$
shows $x \neq 0 \implies 1 < x = (0 < x - 1)$
 ⟨proof⟩

lemma *word-le-sub1*:
fixes $x :: 'a :: \text{len word}$
shows $x \neq 0 \implies 1 \leq x = (0 \leq x - 1)$
 ⟨proof⟩

lemmas *word-less-sub1-numberof* [simp] =
word-less-sub1 [of number-of w, standard]

lemmas *word-le-sub1-numberof* [simp] =
word-le-sub1 [of number-of w, standard]

lemma *word-of-int-minus*:
 $\text{word-of-int } (2^{\text{len-of TYPE('a)} - i}) = (\text{word-of-int } (-i) :: 'a :: \text{len word})$
 ⟨proof⟩

lemmas *word-of-int-inj* =
word-uint.Abs-inject [unfolded uints-num, simplified]

lemma *word-le-less-eq*:

$(x :: 'z :: \text{len word}) \leq y = (x = y \vee x < y)$
 ⟨proof⟩

lemma *mod-plus-cong*:

assumes 1: $(b :: \text{int}) = b'$
and 2: $x \bmod b' = x' \bmod b'$
and 3: $y \bmod b' = y' \bmod b'$
and 4: $x' + y' = z'$
shows $(x + y) \bmod b = z' \bmod b'$
 ⟨proof⟩

lemma *mod-minus-cong*:

assumes 1: $(b :: \text{int}) = b'$
and 2: $x \bmod b' = x' \bmod b'$
and 3: $y \bmod b' = y' \bmod b'$
and 4: $x' - y' = z'$
shows $(x - y) \bmod b = z' \bmod b'$
 ⟨proof⟩

lemma *word-induct-less*:

$\llbracket P (0 :: 'a :: \text{len word}); \bigwedge n. \llbracket n < m; P n \rrbracket \implies P (1 + n) \rrbracket \implies P m$
 ⟨proof⟩

lemma *word-induct*:

$\llbracket P (0 :: 'a :: \text{len word}); \bigwedge n. P n \implies P (1 + n) \rrbracket \implies P m$
 ⟨proof⟩

lemma *word-induct2* [*induct type*]:

$\llbracket P 0; \bigwedge n. \llbracket 1 + n \neq 0; P n \rrbracket \implies P (1 + n) \rrbracket \implies P (n :: 'b :: \text{len word})$
 ⟨proof⟩

constdefs

$\text{word-rec} :: 'a \Rightarrow ('b :: \text{len word} \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b \text{ word} \Rightarrow 'a$
 $\text{word-rec forZero forSuc } n \equiv \text{nat-rec forZero (forSuc } \circ \text{ of-nat)} (\text{unat } n)$

lemma *word-rec-0*: $\text{word-rec } z \text{ } s \text{ } 0 = z$

⟨proof⟩

lemma *word-rec-Suc*:

$1 + n \neq (0 :: 'a :: \text{len word}) \implies \text{word-rec } z \text{ } s \text{ } (1 + n) = s \text{ } n (\text{word-rec } z \text{ } s \text{ } n)$
 ⟨proof⟩

lemma *word-rec-Pred*:

$n \neq 0 \implies \text{word-rec } z \text{ } s \text{ } n = s \text{ } (n - 1) (\text{word-rec } z \text{ } s \text{ } (n - 1))$
 ⟨proof⟩

lemma *word-rec-in*:

$f \text{ (word-rec } z \text{ (}\lambda\text{-. } f \text{) } n) = \text{word-rec (} f \text{ } z \text{) (}\lambda\text{-. } f \text{) } n$
 ⟨proof⟩

lemma *word-rec-in2*:

$f \text{ } n \text{ (word-rec } z \text{ } f \text{ } n) = \text{word-rec (} f \text{ } 0 \text{ } z \text{) (} f \text{ } \circ \text{ } op \text{ } + \text{ } 1 \text{) } n$
 ⟨proof⟩

lemma *word-rec-twice*:

$m \leq n \implies \text{word-rec } z \text{ } f \text{ } n = \text{word-rec (word-rec } z \text{ } f \text{ (} n - m \text{)) (} f \text{ } \circ \text{ } op \text{ } + \text{ (} n - m \text{)) } m$
 ⟨proof⟩

lemma *word-rec-id*: $\text{word-rec } z \text{ (}\lambda\text{-. } id \text{) } n = z$

⟨proof⟩

lemma *word-rec-id-eq*: $\forall m < n. f \text{ } m = id \implies \text{word-rec } z \text{ } f \text{ } n = z$

⟨proof⟩

lemma *word-rec-max*:

$\forall m \geq n. m \neq -1 \longrightarrow f \text{ } m = id \implies \text{word-rec } z \text{ } f \text{ } -1 = \text{word-rec } z \text{ } f \text{ } n$
 ⟨proof⟩

lemma *unatSuc*:

$1 + n \neq (0::'a::len \text{ word}) \implies \text{unat (} 1 + n \text{) = Suc (unat } n)$
 ⟨proof⟩

end

16 WordMain: Main Word Library

theory *WordMain* **imports** *WordGenLib*

begin

lemmas *word-no-1* [*simp*] = *word-1-no* [*symmetric*]

lemmas *word-no-0* [*simp*] = *word-0-no* [*symmetric*]

declare *word-0-bl* [*simp*]

declare *bin-to-bl-def* [*simp*]

declare *to-bl-0* [*simp*]

declare *of-bl-True* [*simp*]

Examples

types *word32* = 32 *word*

types *word8* = 8 *word*

types *byte* = *word8*

for more see *WordExampes.thy*

end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.