

# Machine Words in Isabelle/HOL

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## Abstract

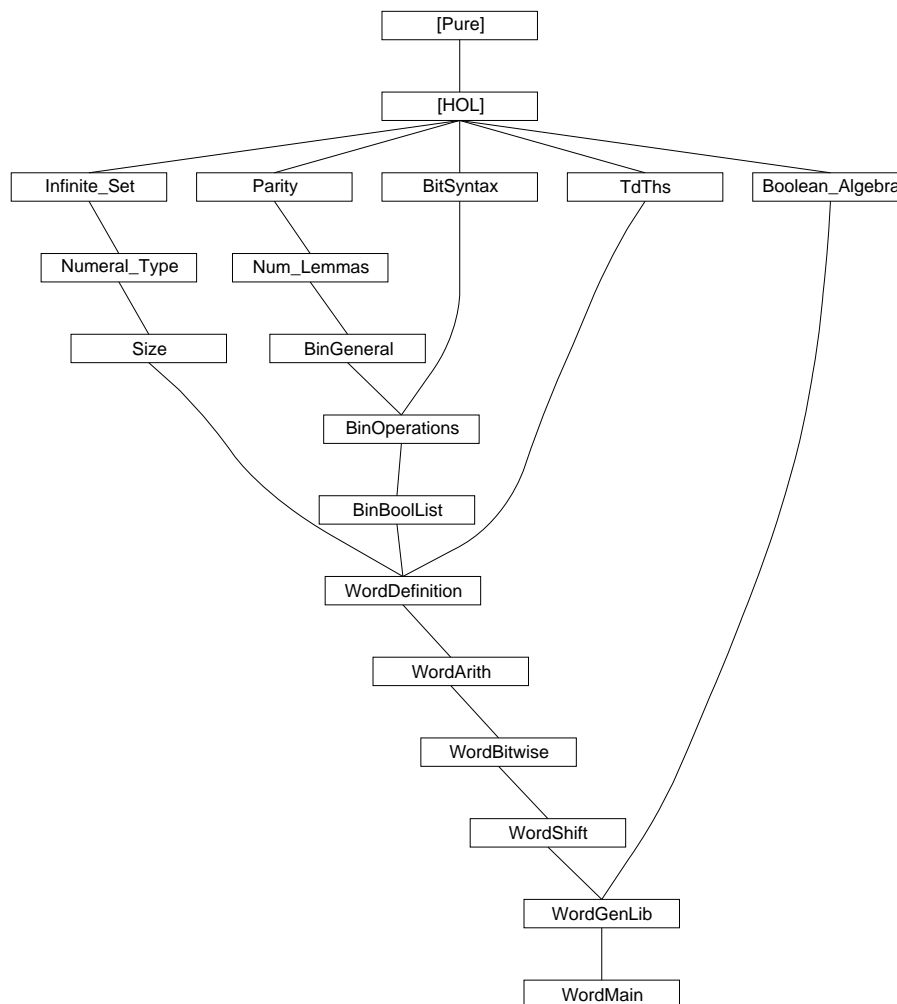
A formalisation of generic, fixed size machine words in Isabelle/HOL.  
An earlier version of this formalisation is described in [1].

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## 1 Numeral-Type: Numeral Syntax for Types

```
theory Numeral-Type
  imports Infinite-Set
begin
```

### 1.1 Preliminary lemmas

```
lemma inj-Inl [simp]: inj-on Inl A
  by (rule inj-onI, simp)
```

```
lemma inj-Inr [simp]: inj-on Inr A
  by (rule inj-onI, simp)
```

```
lemma inj-Some [simp]: inj-on Some A
  by (rule inj-onI, simp)
```

```
lemma card-Plus:
  [| finite A; finite B |] ==> card (A <+> B) = card A + card B
  unfolding Plus-def
  apply (subgoal-tac Inl ‘ A ∩ Inr ‘ B = {})
  apply (simp add: card-Un-disjoint card-image)
  apply fast
done
```

```
lemma (in type-definition) univ:
  UNIV = Abs ‘ A
```

```
proof
  show Abs ‘ A ⊆ UNIV by (rule subset-UNIV)
  show UNIV ⊆ Abs ‘ A
  proof
    fix x :: 'b
    have x = Abs (Rep x) by (rule Rep-inverse [symmetric])
    moreover have Rep x ∈ A by (rule Rep)
    ultimately show x ∈ Abs ‘ A by (rule image-eqI)
  qed
qed
```

```
lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  by (simp add: univ card-image inj-on-def Abs-inject)
```

### 1.2 Cardinalities of types

```
syntax -type-card :: type => nat ((1CARD/(1'(-))))
```

```
translations CARD(t) => card (UNIV::t set)
```

```
typed-print-translation <<
let
  fun card-univ-tr' show-sorts - [Const (@{const-name UNIV}, Type(-,[T]))] =
```

```

    Syntax.const -type-card $ Syntax.term-of-typ show-sorts T;
in [(card, card-univ-tr')]
end
>>

```

```

lemma card-unit: CARD(unit) = 1
  unfolding univ-unit by simp

```

```

lemma card-bool: CARD(bool) = 2
  unfolding univ-bool by simp

```

```

lemma card-prod: CARD('a::finite × 'b::finite) = CARD('a) * CARD('b)
  unfolding univ-prod by (simp only: card-cartesian-product)

```

```

lemma card-sum: CARD('a::finite + 'b::finite) = CARD('a) + CARD('b)
  unfolding univ-sum by (simp only: finite card-Plus)

```

```

lemma card-option: CARD('a::finite option) = Suc CARD('a)
  unfolding univ-option
  apply (subgoal-tac (None::'a option) ∉ range Some)
  apply (simp add: finite card-image)
  apply fast
  done

```

```

lemma card-set: CARD('a::finite set) = 2 ^ CARD('a)
  unfolding univ-set
  by (simp only: card-Pow finite numeral-2-eq-2)

```

### 1.3 Numeral Types

```

typedef (open) num0 = UNIV :: nat set ..
typedef (open) num1 = UNIV :: unit set ..
typedef (open) 'a bit0 = UNIV :: (bool * 'a) set ..
typedef (open) 'a bit1 = UNIV :: (bool * 'a) option set ..

```

```

instance num1 :: finite
proof
  show finite (UNIV::num1 set)
    unfolding type-definition.univ [OF type-definition-num1]
    using finite by (rule finite-imageI)
qed

```

```

instance bit0 :: (finite) finite
proof
  show finite (UNIV::'a bit0 set)
    unfolding type-definition.univ [OF type-definition-bit0]
    using finite by (rule finite-imageI)
qed

```

```

instance bit1 :: (finite) finite
proof
  show finite (UNIV::'a bit1 set)
    unfolding type-definition.univ [OF type-definition-bit1]
    using finite by (rule finite-imageI)
qed

lemma card-num1: CARD(num1) = 1
  unfolding type-definition.card [OF type-definition-num1]
  by (simp only: card-unit)

lemma card-bit0: CARD('a::finite bit0) = 2 * CARD('a)
  unfolding type-definition.card [OF type-definition-bit0]
  by (simp only: card-prod card-bool)

lemma card-bit1: CARD('a::finite bit1) = Suc (2 * CARD('a))
  unfolding type-definition.card [OF type-definition-bit1]
  by (simp only: card-prod card-option card-bool)

lemma card-num0: CARD (num0) = 0
  by (simp add: type-definition.card [OF type-definition-num0])

lemmas card-univ-simps [simp] =
  card-unit
  card-bool
  card-prod
  card-sum
  card-option
  card-set
  card-num1
  card-bit0
  card-bit1
  card-num0

```

## 1.4 Syntax

```

syntax
  -NumeralType :: num-const => type (-)
  -NumeralType0 :: type (0)
  -NumeralType1 :: type (1)

```

```

translations
  -NumeralType1 == (type) num1
  -NumeralType0 == (type) num0

```

```

parse-translation <<
  let

```

```

  val num1-const = Syntax.const Numeral-Type.num1;

```

```

val num0-const = Syntax.const Numeral-Type.num0;
val B0-const = Syntax.const Numeral-Type.bit0;
val B1-const = Syntax.const Numeral-Type.bit1;

fun mk-bintype n =
  let
    fun mk-bit n = if n = 0 then B0-const else B1-const;
    fun bin-of n =
      if n = 1 then num1-const
      else if n = 0 then num0-const
      else if n = ~1 then raise TERM (negative type numeral, [])
      else
        let val (q, r) = Integer.div-mod n 2;
        in mk-bit r $ bin-of q end;
  in bin-of n end;

fun numeral-tr (*-NumeralType*) [Const (str, -)] =
  mk-bintype (valOf (Int.fromString str))
| numeral-tr (*-NumeralType*) ts = raise TERM (numeral-tr, ts);

in [(-NumeralType, numeral-tr)] end;
>>

print-translation <<
let
  fun int-of [] = 0
  | int-of (b :: bs) = b + 2 * int-of bs;

  fun bin-of (Const (num0, -)) = []
  | bin-of (Const (num1, -)) = [1]
  | bin-of (Const (bit0, -) $ bs) = 0 :: bin-of bs
  | bin-of (Const (bit1, -) $ bs) = 1 :: bin-of bs
  | bin-of t = raise TERM (bin-of, [t]);

  fun bit-tr' b [t] =
    let
      val rev-digs = b :: bin-of t handle TERM - => raise Match
      val i = int-of rev-digs;
      val num = string-of-int (abs i);
    in
      Syntax.const -NumeralType $ Syntax.free num
    end
  | bit-tr' b - = raise Match;

in [(bit0, bit-tr' 0), (bit1, bit-tr' 1)] end;
>>

```

### 1.5 Classes with at least 1 and 2

Class `finite` already captures “at least 1”

```

lemma zero-less-card-finite [simp]:
  0 < CARD('a::finite)
proof (cases CARD('a::finite) = 0)
  case False thus ?thesis by (simp del: card-0-eq)
next
  case True
  thus ?thesis by (simp add: finite)
qed

```

```

lemma one-le-card-finite [simp]:
  Suc 0 <= CARD('a::finite)
  by (simp add: less-Suc-eq-le [symmetric] zero-less-card-finite)

```

Class for cardinality “at least 2”

```

class card2 = finite +
  assumes two-le-card: 2 <= CARD('a)

lemma one-less-card: Suc 0 < CARD('a::card2)
  using two-le-card [where 'a='a] by simp

instance bit0 :: (finite) card2
  by intro-classes (simp add: one-le-card-finite)

instance bit1 :: (finite) card2
  by intro-classes (simp add: one-le-card-finite)

```

### 1.6 Examples

```

term TYPE(10)

lemma CARD(0) = 0 by simp
lemma CARD(17) = 17 by simp

end

```

## 2 Size: The size class

```

theory Size
imports Natural_Type
begin

```

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in `Natural_Type`.



**axclass** *len0* < *type*

**consts**

*len-of* :: (*'a* :: *len0* *itself*) ==> *nat*

Some theorems are only true on words with length greater 0.

**axclass** *len* < *len0*

*len-gt-0* [*iff*]:  $0 < \text{len-of } TYPE \text{ } ('a :: \text{len0})$

**instance** *num0* :: *len0* ..

**instance** *num1* :: *len0* ..

**instance** *bit0* :: (*len0*) *len0* ..

**instance** *bit1* :: (*len0*) *len0* ..

**defs** (overloaded)

*len-num0*: *len-of* (*x*::*num0* *itself*) == 0

*len-num1*: *len-of* (*x*::*num1* *itself*) == 1

*len-bit0*: *len-of* (*x*::*'a*::*len0* *bit0* *itself*) == 2 \* *len-of* *TYPE* (*'a*)

*len-bit1*: *len-of* (*x*::*'a*::*len0* *bit1* *itself*) == 2 \* *len-of* *TYPE* (*'a*) + 1

**lemmas** *len-of-numeral-defs* [*simp*] = *len-num0* *len-num1* *len-bit0* *len-bit1*

**instance** *num1* :: *len* **by** (*intro-classes*) *simp*

**instance** *bit0* :: (*len*) *len* **by** (*intro-classes*) *simp*

**instance** *bit1* :: (*len0*) *len* **by** (*intro-classes*) *simp*

— Examples:

**lemma** *len-of* *TYPE*(17) = 17 **by** *simp*

**lemma** *len-of* *TYPE*(0) = 0 **by** *simp*

— not simplified:

**lemma** *len-of* *TYPE*(*'a*::*len0*) = *x*

**oops**

**end**

### 3 Num-Lemmas: Useful Numerical Lemmas

**theory** *Num-Lemmas* **imports** *Parity* **begin**

**lemma** *contentsI*:  $y = \{x\} \implies \text{contents } y = x$

**unfolding** *contents-def* **by** *auto*

**lemma** *prod-case-split*: *prod-case* = *split*

**by** (*rule ext*)+ *auto*

**lemmas** *split-split* = *prod.split* [*unfolded prod-case-split*]

```

lemmas split-split-asm = prod.split-asm [unfolded prod-case-split]
lemmas split.splits = split-split split-split-asm

```

```

lemmas funpow-0 = funpow.simps(1)
lemmas funpow-Suc = funpow.simps(2)

```

```

lemma nonemptyE:  $S \sim = \{\}$   $\implies (!x. x : S \implies R) \implies R$ 
  apply (erule contrapos-np)
  apply (rule equals0I)
  apply auto
done

```

```

lemma gt-or-eq-0:  $0 < y \vee 0 = (y::nat)$  by auto

```

```

constdefs
  mod-alt :: 'a => 'a => 'a :: Divides.div
  mod-alt n m == n mod m

  — alternative way of defining bin-last, bin-rest
  bin-rl :: int => int * bit
  bin-rl w == SOME (r, l). w = r BIT l

```

```

declare iszero-0 [iff]

```

```

lemmas xtr1 = xtrans(1)
lemmas xtr2 = xtrans(2)
lemmas xtr3 = xtrans(3)
lemmas xtr4 = xtrans(4)
lemmas xtr5 = xtrans(5)
lemmas xtr6 = xtrans(6)
lemmas xtr7 = xtrans(7)
lemmas xtr8 = xtrans(8)

```

```

lemma Min-ne-Pls [iff]:
  Numeral.Min  $\sim =$  Numeral.Pls
  unfolding Min-def Pls-def by auto

```

```

lemmas Pls-ne-Min [iff] = Min-ne-Pls [symmetric]

```

```

lemmas PlsMin-defs [intro!] =
  Pls-def Min-def Pls-def [symmetric] Min-def [symmetric]

```

```

lemmas PlsMin-simps [simp] = PlsMin-defs [THEN Eq-TrueI]

```

```

lemma number-of-False-cong:
  False  $\implies$  number-of x = number-of y
  by auto

```

```

lemmas nat-simps = diff-add-inverse2 diff-add-inverse

```

**lemmas** *nat-iffs* = *le-add1 le-add2*

**lemma** *sum-imp-diff*:  $j = k + i \implies j - i = (k :: \text{nat})$   
**by** (*clarsimp simp add: nat-simps*)

**lemma** *nobm1*:  
 $0 < (\text{number-of } w :: \text{nat}) \implies$   
 $\text{number-of } w - (1 :: \text{nat}) = \text{number-of } (\text{Numeral.pred } w)$   
**apply** (*unfold nat-number-of-def One-nat-def nat-1 [symmetric] pred-def*)  
**apply** (*simp add: number-of-eq nat-diff-distrib [symmetric]*)  
**done**

**lemma** *of-int-power*:  
 $\text{of-int } (a \wedge n) = (\text{of-int } a \wedge n :: 'a :: \{\text{recpower, comm-ring-1}\})$   
**by** (*induct n (auto simp add: power-Suc)*)

**lemma** *zless2*:  $0 < (2 :: \text{int})$   
**by** *auto*

**lemmas** *zless2p [simp]* = *zless2 [THEN zero-less-power]*  
**lemmas** *zle2p [simp]* = *zless2p [THEN order-less-imp-le]*

**lemmas** *pos-mod-sign2* = *zless2 [THEN pos-mod-sign [where b = 2::int]]*  
**lemmas** *pos-mod-bound2* = *zless2 [THEN pos-mod-bound [where b = 2::int]]*

— the inverse(s) of *number-of*

**lemma** *nmod2*:  $n \bmod (2 :: \text{int}) = 0 \mid n \bmod 2 = 1$   
**using** *pos-mod-sign2 [of n] pos-mod-bound2 [of n]*  
**unfolding** *mod-alt-def [symmetric]* **by** *auto*

**lemma** *emep1*:  
 $\text{even } n \implies \text{even } d \implies 0 \leq d \implies (n + 1) \bmod (d :: \text{int}) = (n \bmod d) + 1$   
**apply** (*simp add: add-commute*)  
**apply** (*safe dest!: even-equiv-def [THEN iffD1]*)  
**apply** (*subst pos-zmod-mult-2*)  
**apply** *arith*  
**apply** (*simp add: zmod-zmult-zmult1*)  
**done**

**lemmas** *eme1p* = *emep1 [simplified add-commute]*

**lemma** *le-diff-eq'*:  $(a \leq c - b) = (b + a \leq (c :: \text{int}))$   
**by** (*simp add: le-diff-eq add-commute*)

**lemma** *less-diff-eq'*:  $(a < c - b) = (b + a < (c :: \text{int}))$   
**by** (*simp add: less-diff-eq add-commute*)

**lemma** *diff-le-eq'*:  $(a - b \leq c) = (a \leq b + (c :: \text{int}))$

```

by (simp add: diff-le-eq add-commute)

lemma diff-less-eq':  $(a - b < c) = (a < b + (c::int))$ 
by (simp add: diff-less-eq add-commute)

lemmas m1mod2k = zless2p [THEN zmod-minus1]
lemmas m1mod22k = mult-pos-pos [OF zless2 zless2p, THEN zmod-minus1]
lemmas p1mod22k' = zless2p [THEN order-less-imp-le, THEN pos-zmod-mult-2]
lemmas z1pmod2' = zero-le-one [THEN pos-zmod-mult-2, simplified]
lemmas z1pdiv2' = zero-le-one [THEN pos-zdiv-mult-2, simplified]

lemma p1mod22k:
 $(2 * b + 1) \bmod (2 * 2^n) = 2 * (b \bmod 2^n) + (1::int)$ 
by (simp add: p1mod22k' add-commute)

lemma z1pmod2:
 $(2 * b + 1) \bmod 2 = (1::int)$ 
by (simp add: z1pmod2' add-commute)

lemma z1pdiv2:
 $(2 * b + 1) \bmod 2 = (b::int)$ 
by (simp add: z1pdiv2' add-commute)

lemmas zdiv-le-dividend = xtr3 [OF zdiv-1 [symmetric] zdiv-mono2,
  simplified int-one-le-iff-zero-less, simplified, standard]

lemma BIT-eq:  $u \text{ BIT } b = v \text{ BIT } c \implies u = v \ \& \ b = c$ 
apply (unfold Bit-def)
apply (simp (no-asm-use) split: bit.split-asm)
apply simp-all
apply (drule-tac f=even in arg-cong, clarsimp)+
done

lemmas BIT-eqE [elim!] = BIT-eq [THEN conjE, standard]

lemma BIT-eq-iff [simp]:
 $(u \text{ BIT } b = v \text{ BIT } c) = (u = v \ \wedge \ b = c)$ 
by (rule iffI) auto

lemmas BIT-eqI [intro!] = conjI [THEN BIT-eq-iff [THEN iffD2]]

lemma less-Bits:
 $(v \text{ BIT } b < w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ b = \text{bit.B0} \ \& \ c = \text{bit.B1})$ 
unfolding Bit-def by (auto split: bit.split)

lemma le-Bits:

```

$(v \text{ BIT } b \leq w \text{ BIT } c) = (v < w \mid v \leq w \ \& \ (b \sim \text{bit.B1} \mid c \sim \text{bit.B0}))$   
**unfolding** *Bit-def* **by** (*auto split: bit.split*)

**lemma** *neB1E* [*elim!*]:  
**assumes** *ne*:  $y \neq \text{bit.B1}$   
**assumes** *y*:  $y = \text{bit.B0} \implies P$   
**shows** *P*  
**apply** (*rule y*)  
**apply** (*cases y rule: bit.exhaust, simp*)  
**apply** (*simp add: ne*)  
**done**

**lemma** *bin-ex-rl*:  $EX \ w \ b. \ w \text{ BIT } b = \text{bin}$   
**apply** (*unfold Bit-def*)  
**apply** (*cases even bin*)  
**apply** (*clarsimp simp: even-equiv-def*)  
**apply** (*auto simp: odd-equiv-def split: bit.split*)  
**done**

**lemma** *bin-exhaust*:  
**assumes** *Q*:  $\bigwedge x \ b. \text{bin} = x \text{ BIT } b \implies Q$   
**shows** *Q*  
**apply** (*insert bin-ex-rl [of bin]*)  
**apply** (*erule exE*)  
**apply** (*rule Q*)  
**apply** *force*  
**done**

**lemma** *bin-rl-char*:  $(\text{bin-rl } w = (r, l)) = (r \text{ BIT } l = w)$   
**apply** (*unfold bin-rl-def*)  
**apply** *safe*  
**apply** (*cases w rule: bin-exhaust*)  
**apply** *auto*  
**done**

**lemmas** *bin-rl-simps* [*THEN bin-rl-char [THEN iffD2], standard, simp*] =  
*Pls-0-eq Min-1-eq refl*

**lemma** *bin-abs-lem*:  
 $\text{bin} = (w \text{ BIT } b) \implies \sim \text{bin} = \text{Numeral.Min} \dashrightarrow \sim \text{bin} = \text{Numeral.Pls} \dashrightarrow$   
 $\text{nat } (\text{abs } w) < \text{nat } (\text{abs } \text{bin})$   
**apply** (*clarsimp simp add: bin-rl-char*)  
**apply** (*unfold Pls-def Min-def Bit-def*)  
**apply** (*cases b*)  
**apply** (*clarsimp, arith*)  
**apply** (*clarsimp, arith*)  
**done**

**lemma** *bin-induct*:

```

assumes PPls: P Numeral.Pls
and PMin: P Numeral.Min
and PBit: !!bin bit. P bin ==> P (bin BIT bit)
shows P bin
apply (rule-tac P=P and a=bin and f1=nat o abs
        in wf-measure [THEN wf-induct])
apply (simp add: measure-def inv-image-def)
apply (case-tac x rule: bin-exhaust)
apply (frule bin-abs-lem)
apply (auto simp add : PPls PMin PBit)
done

lemma no-no [simp]: number-of (number-of i) = i
unfolding number-of-eq by simp

lemma Bit-B0:
  k BIT bit.B0 = k + k
by (unfold Bit-def) simp

lemma Bit-B1:
  k BIT bit.B1 = k + k + 1
by (unfold Bit-def) simp

lemma Bit-B0-2t: k BIT bit.B0 = 2 * k
by (rule trans, rule Bit-B0) simp

lemma Bit-B1-2t: k BIT bit.B1 = 2 * k + 1
by (rule trans, rule Bit-B1) simp

lemma B-mod-2':
  X = 2 ==> (w BIT bit.B1) mod X = 1 & (w BIT bit.B0) mod X = 0
apply (simp (no-asm) only: Bit-B0 Bit-B1)
apply (simp add: z1pmod2)
done

lemmas B1-mod-2 [simp] = B-mod-2' [OF refl, THEN conjunct1, standard]
lemmas B0-mod-2 [simp] = B-mod-2' [OF refl, THEN conjunct2, standard]

lemma axbby:
  a + m + m = b + n + n ==> (a = 0 | a = 1) ==> (b = 0 | b = 1) ==>
  a = b & m = (n :: int)
apply auto
apply (drule-tac f=%n. n mod 2 in arg-cong)
apply (clarsimp simp: z1pmod2)
apply (drule-tac f=%n. n mod 2 in arg-cong)
apply (clarsimp simp: z1pmod2)
done

lemma axxmod2:

```

$(1 + x + x) \bmod 2 = (1 :: \text{int}) \ \& \ (0 + x + x) \bmod 2 = (0 :: \text{int})$   
**by** *simp* (rule *z1pmod2*)

**lemma** *axxdiv2*:

$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int})$   
**by** *simp* (rule *z1pdiv2*)

**lemmas** *iszero-minus* = *trans* [*THEN trans*,  
*OF iszero-def neg-equal-0-iff-equal iszero-def [symmetric]*, *standard*]

**lemmas** *zadd-diff-inverse* = *trans* [*OF diff-add-cancel [symmetric]* *add-commute*,  
*standard*]

**lemmas** *add-diff-cancel2* = *add-commute* [*THEN diff-eq-eq [THEN iffD2]*, *standard*]

**lemma** *zmod-uminus*:  $-(a :: \text{int}) \bmod b \bmod b = -a \bmod b$   
**by** (*simp add : zmod-zminus1-eq-if*)

**lemma** *zmod-zsub-distrib*:  $((a :: \text{int}) - b) \bmod c = (a \bmod c - b \bmod c) \bmod c$   
**apply** (*unfold diff-int-def*)  
**apply** (*rule trans [OF - zmod-zadd1-eq [symmetric]]*)  
**apply** (*simp add: zmod-uminus zmod-zadd1-eq [symmetric]*)  
**done**

**lemma** *zmod-zsub-right-eq*:  $((a :: \text{int}) - b) \bmod c = (a - b \bmod c) \bmod c$   
**apply** (*unfold diff-int-def*)  
**apply** (*rule trans [OF - zmod-zadd-right-eq [symmetric]]*)  
**apply** (*simp add : zmod-uminus zmod-zadd-right-eq [symmetric]*)  
**done**

**lemma** *zmod-zsub-left-eq*:  $((a :: \text{int}) - b) \bmod c = (a \bmod c - b) \bmod c$   
**by** (*rule zmod-zadd-left-eq [where b = - b, simplified diff-int-def [symmetric]]*)

**lemma** *zmod-zsub-self* [*simp*]:  
 $((b :: \text{int}) - a) \bmod a = b \bmod a$   
**by** (*simp add: zmod-zsub-right-eq*)

**lemma** *zmod-zmult1-eq-rev*:  
 $b * a \bmod c = b \bmod c * a \bmod (c :: \text{int})$   
**apply** (*simp add: mult-commute*)  
**apply** (*subst zmod-zmult1-eq*)  
**apply** *simp*  
**done**

**lemmas** *rdmods* [*symmetric*] = *zmod-uminus* [*symmetric*]  
*zmod-zsub-left-eq zmod-zsub-right-eq zmod-zadd-left-eq*  
*zmod-zadd-right-eq zmod-zmult1-eq zmod-zmult1-eq-rev*

**lemma** *mod-plus-right*:

$((a + x) \bmod m = (b + x) \bmod m) = (a \bmod m = b \bmod (m :: nat))$   
**apply** (*induct* *x*)  
**apply** (*simp-all* *add*: *mod-Suc*)  
**apply** *arith*  
**done**

**lemma** *nat-minus-mod*:  $(n - n \bmod m) \bmod m = (0 :: nat)$

**by** (*induct* *n*) (*simp-all* *add* : *mod-Suc*)

**lemmas** *nat-minus-mod-plus-right* = *trans* [*OF* *nat-minus-mod mod-0* [*symmetric*],  
*THEN* *mod-plus-right* [*THEN* *iffD2*], *standard*, *simplified*]

**lemmas** *push-mods'* = *zmod-zadd1-eq* [*standard*]

*zmod-zmult-distrib* [*standard*] *zmod-zsub-distrib* [*standard*]

*zmod-uminus* [*symmetric*, *standard*]

**lemmas** *push-mods* = *push-mods'* [*THEN* *eq-reflection*, *standard*]

**lemmas** *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN* *eq-reflection*, *standard*]

**lemmas** *mod-simps* =

*zmod-zmult-self1* [*THEN* *eq-reflection*] *zmod-zmult-self2* [*THEN* *eq-reflection*]

*mod-mod-trivial* [*THEN* *eq-reflection*]

**lemma** *nat-mod-eq*:

$!!b. b < n ==> a \bmod n = b \bmod n ==> a \bmod n = (b :: nat)$

**by** (*induct* *a*) *auto*

**lemmas** *nat-mod-eq'* = *refl* [*THEN* [*2*] *nat-mod-eq*]

**lemma** *nat-mod-lem*:

$(0 :: nat) < n ==> b < n = (b \bmod n = b)$

**apply** *safe*

**apply** (*erule* *nat-mod-eq'*)

**apply** (*erule* *subst*)

**apply** (*erule* *mod-less-divisor*)

**done**

**lemma** *mod-nat-add*:

$(x :: nat) < z ==> y < z ==>$

$(x + y) \bmod z = (if\ x + y < z\ then\ x + y\ else\ x + y - z)$

**apply** (*rule* *nat-mod-eq*)

**apply** *auto*

**apply** (*rule* *trans*)

**apply** (*rule* *le-mod-geq*)

**apply** *simp*

**apply** (*rule* *nat-mod-eq'*)

**apply** *arith*

**done**



**lemma** *mod-nat-sub*:

$(x :: \text{nat}) < z \implies (x - y) \bmod z = x - y$   
**by** (*rule nat-mod-eq'*) *arith*

**lemma** *int-mod-lem*:

$(0 :: \text{int}) < n \implies (0 \leq b \ \& \ b < n) = (b \bmod n = b)$   
**apply** *safe*  
**apply** (*erule* (1) *mod-pos-pos-trivial*)  
**apply** (*erule-tac* [!] *subst*)  
**apply** *auto*  
**done**

**lemma** *int-mod-eq*:

$(0 :: \text{int}) \leq b \implies b < n \implies a \bmod n = b \bmod n \implies a \bmod n = b$   
**by** *clarsimp* (*rule mod-pos-pos-trivial*)

**lemmas** *int-mod-eq'* = *refl* [*THEN* [3] *int-mod-eq*]

**lemma** *int-mod-le*:  $0 \leq a \implies 0 < (n :: \text{int}) \implies a \bmod n \leq a$

**apply** (*cases*  $a < n$ )  
**apply** (*auto dest: mod-pos-pos-trivial pos-mod-bound* [**where**  $a=a$ ])  
**done**

**lemma** *int-mod-le'*:  $0 \leq b - n \implies 0 < (n :: \text{int}) \implies b \bmod n \leq b - n$

**by** (*rule int-mod-le* [**where**  $a = b - n$  **and**  $n = n$ , *simplified*])

**lemma** *int-mod-ge*:  $a < n \implies 0 < (n :: \text{int}) \implies a \leq a \bmod n$

**apply** (*cases*  $0 \leq a$ )  
**apply** (*erule* (1) *mod-pos-pos-trivial*)  
**apply** *simp*  
**apply** (*rule order-trans* [*OF* - *pos-mod-sign*])  
**apply** *simp*  
**apply** *assumption*  
**done**

**lemma** *int-mod-ge'*:  $b < 0 \implies 0 < (n :: \text{int}) \implies b + n \leq b \bmod n$

**by** (*rule int-mod-ge* [**where**  $a = b + n$  **and**  $n = n$ , *simplified*])

**lemma** *mod-add-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$   
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$   
**by** (*auto intro: int-mod-eq*)

**lemma** *mod-sub-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$   
 $(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$   
**by** (*auto intro: int-mod-eq*)

**lemmas** *zmde* = *zmod-zdiv-equality* [*THEN diff-eq-eq* [*THEN iffD2*], *symmetric*]  
**lemmas** *mcl* = *mult-cancel-left* [*THEN iffD1*, *THEN make-pos-rule*]

**lemma** *zdiv-mult-self*:  $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$   
**apply** (*rule mcl*)  
**prefer** 2  
**apply** (*erule asm-rl*)  
**apply** (*simp add: zmde ring-distrib*)  
**apply** (*simp add: push-mods*)  
**done**

**lemma** *eqne*:  $\text{equiv } A \ r \implies X : A \ /\ / \ r \implies X \sim = \{\}$   
**unfolding** *equiv-def refl-def quotient-def Image-def* **by** *auto*

**lemmas** *Rep-Integ-ne* = *Integ.Rep-Integ*  
[*THEN equiv-intrel* [*THEN eqne*, *simplified Integ-def* [*symmetric*]], *standard*]

**lemmas** *riq* = *Integ.Rep-Integ* [*simplified Integ-def*]  
**lemmas** *intrel-refl* = *refl* [*THEN equiv-intrel-iff* [*THEN iffD1*], *standard*]  
**lemmas** *Rep-Integ-equiv* = *quotient-eq-iff*  
[*OF equiv-intrel riq riq*, *simplified Integ.Rep-Integ-inject*, *standard*]  
**lemmas** *Rep-Integ-same* =  
*Rep-Integ-equiv* [*THEN intrel-refl* [*THEN rev-iffD2*], *standard*]

**lemma** *RI-int*:  $(a, 0) : \text{Rep-Integ } (\text{int } a)$   
**unfolding** *int-def* **by** *auto*

**lemmas** *RI-intrel* [*simp*] = *UNIV-I* [*THEN quotientI*,  
*THEN Integ.Abs-Integ-inverse* [*simplified Integ-def*], *standard*]

**lemma** *RI-minus*:  $(a, b) : \text{Rep-Integ } x \implies (b, a) : \text{Rep-Integ } (-x)$   
**apply** (*rule-tac z=x in eq-Abs-Integ*)  
**apply** (*clarsimp simp: minus*)  
**done**

**lemma** *RI-add*:  
 $(a, b) : \text{Rep-Integ } x \implies (c, d) : \text{Rep-Integ } y \implies$   
 $(a + c, b + d) : \text{Rep-Integ } (x + y)$   
**apply** (*rule-tac z=x in eq-Abs-Integ*)  
**apply** (*rule-tac z=y in eq-Abs-Integ*)  
**apply** (*clarsimp simp: add*)  
**done**

**lemma** *mem-same*:  $a : S \implies a = b \implies b : S$   
**by** *fast*

```

lemma RI-eq-diff':  $(a, b) : \text{Rep-Integ } (\text{int } a - \text{int } b)$ 
  apply (unfold diff-def)
  apply (rule mem-same)
  apply (rule RI-minus RI-add RI-int)+
  apply simp
done

lemma RI-eq-diff:  $((a, b) : \text{Rep-Integ } x) = (\text{int } a - \text{int } b = x)$ 
  apply safe
  apply (rule Rep-Integ-same)
  prefer 2
  apply (erule asm-rl)
  apply (rule RI-eq-diff')+
done

lemma mod-power-lem:
   $a > 1 \implies a^n \bmod a^m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int})^n)$ 
  apply clarsimp
  apply safe
  apply (simp add: zdvd-iff-zmod-eq-0 [symmetric])
  apply (erule le-iff-add [THEN iffD1])
  apply (force simp: zpower-zadd-distrib)
  apply (rule mod-pos-pos-trivial)
  apply (simp add: zero-le-power)
  apply (rule power-strict-increasing)
  apply auto
done

lemma min-pm [simp]:  $\min a b + (a - b) = (a :: \text{nat})$ 
  by arith

lemmas min-pm1 [simp] = trans [OF add-commute min-pm]

lemma rev-min-pm [simp]:  $\min b a + (a - b) = (a :: \text{nat})$ 
  by simp

lemmas rev-min-pm1 [simp] = trans [OF add-commute rev-min-pm]

lemma pl-pl-rels:
   $a + b = c + d \implies$ 
   $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$ 
  apply (cut-tac n=a and m=c in nat-le-linear)
  apply (safe dest!: le-iff-add [THEN iffD1])
  apply arith+
done

lemmas pl-pl-rels' = add-commute [THEN [2] trans, THEN pl-pl-rels]

lemma minus-eq:  $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$ 

```

```

by arith

lemma pl-pl-mm:  $(a :: \text{nat}) + b = c + d \implies a - c = d - b$ 
by arith

lemmas pl-pl-mm' = add-commute [THEN [2] trans, THEN pl-pl-mm]

lemma min-minus [simp] :  $\min m (m - k) = (m - k :: \text{nat})$ 
by arith

lemmas min-minus' [simp] = trans [OF min-max.inf-commute min-minus]

lemma nat-no-eq-iff:
   $(\text{number-of } b :: \text{int}) \geq 0 \implies (\text{number-of } c :: \text{int}) \geq 0 \implies$ 
   $(\text{number-of } b = (\text{number-of } c :: \text{nat})) = (b = c)$ 
apply (unfold nat-number-of-def)
apply safe
apply (drule (2) eq-nat-nat-iff [THEN iffD1])
apply (simp add: number-of-eq)
done

lemmas dme = box-equals [OF div-mod-equality add-0-right add-0-right]
lemmas dtle = xtr3 [OF dme [symmetric] le-add1]
lemmas th2 = order-trans [OF order-refl [THEN [2] mult-le-mono] dtle]

lemma td-gal:
   $0 < c \implies (a \geq b * c) = (a \text{ div } c \geq (b :: \text{nat}))$ 
apply safe
apply (erule (1) xtr4 [OF div-le-mono div-mult-self-is-m])
apply (erule th2)
done

lemmas td-gal-lt = td-gal [simplified le-def, simplified]

lemma div-mult-le:  $(a :: \text{nat}) \text{ div } b * b \leq a$ 
apply (cases b)
prefer 2
apply (rule order-refl [THEN th2])
apply auto
done

lemmas sdl = split-div-lemma [THEN iffD1, symmetric]

lemma given-quot:  $f > (0 :: \text{nat}) \implies (f * l + (f - 1)) \text{ div } f = l$ 
by (rule sdl, assumption) (simp (no-asm))

lemma given-quot-alt:  $f > (0 :: \text{nat}) \implies (l * f + f - \text{Suc } 0) \text{ div } f = l$ 
apply (frule given-quot)
apply (rule trans)

```

```

    prefer 2
    apply (erule asm-rl)
    apply (rule-tac f=%n. n div f in arg-cong)
    apply (simp add : mult-ac)
    done

lemma diff-mod-le: (a::nat) < d ==> b dvd d ==> a - a mod b <= d - b
  apply (unfold dvd-def)
  apply clarify
  apply (case-tac k)
  apply clarsimp
  apply clarify
  apply (cases b > 0)
  apply (drule mult-commute [THEN xtr1])
  apply (frule (1) td-gal-lt [THEN iffD1])
  apply (clarsimp simp: le-simps)
  apply (rule mult-div-cancel [THEN [2] xtr4])
  apply (rule mult-mono)
  apply auto
  done

lemma less-le-mult':
  w * c < b * c ==> 0 ≤ c ==> (w + 1) * c ≤ b * (c::int)
  apply (rule mult-right-mono)
  apply (rule zless-imp-add1-zle)
  apply (erule (1) mult-right-less-imp-less)
  apply assumption
  done

lemmas less-le-mult = less-le-mult' [simplified left-distrib, simplified]

lemmas less-le-mult-minus = iffD2 [OF le-diff-eq less-le-mult,
  simplified left-diff-distrib, standard]

lemma lrlem':
  assumes d: (i::nat) ≤ j ∨ m < j'
  assumes R1: i * k ≤ j * k ==> R
  assumes R2: Suc m * k' ≤ j' * k' ==> R
  shows R using d
  apply safe
  apply (rule R1, erule mult-le-mono1)
  apply (rule R2, erule Suc-le-eq [THEN iffD2 [THEN mult-le-mono1]])
  done

lemma lrlem: (0::nat) < sc ==>
  (sc - n + (n + lb * n) <= m * n) = (sc + lb * n <= m * n)
  apply safe
  apply arith
  apply (case-tac sc >= n)

```

```

  apply arith
  apply (insert linorder-le-less-linear [of m lb])
  apply (erule-tac k=n and k'=n in brlem')
  apply arith
  apply simp
done

lemma gen-minus:  $0 < n \implies f\ n = f\ (Suc\ (n - 1))$ 
  by auto

lemma mpl-lem:  $j \leq (i :: nat) \implies k < j \implies i - j + k < i$ 
  apply (induct i, clarsimp)
  apply (cases j, clarsimp)
  apply arith
done

lemma nonneg-mod-div:
   $0 \leq a \implies 0 \leq b \implies 0 \leq (a \bmod b :: int) \ \& \ 0 \leq a \div b$ 
  apply (cases b = 0, clarsimp)
  apply (auto intro: pos-imp-zdiv-nonneg-iff [THEN iffD2])
done

end

```

## 4 BinGeneral: Basic Definitions for Binary Integers

**theory** *BinGeneral* **imports** *Num-Lemmas*

**begin**

### 4.1 Recursion combinator for binary integers

**lemma** *brlem*:  $(bin = Numeral.Min) = (\neg bin + Numeral.pred\ 0 = 0)$   
**unfolding** *Min-def pred-def* **by** *arith*

**function**

```

  bin-rec' :: int * 'a * 'a * (int => bit => 'a => 'a) => 'a
  where
    bin-rec' (bin, f1, f2, f3) = (if bin = Numeral.Plus then f1
      else if bin = Numeral.Min then f2
      else case bin-rl bin of (w, b) => f3 w b (bin-rec' (w, f1, f2, f3)))
  by pat-completeness auto

```

**termination**

```

  apply (relation measure (nat o abs o fst))
  apply simp

```

```

apply (simp add: Pls-def brlem)
apply (clarsimp simp: bin-rl-char pred-def)
apply (frule thin-rl [THEN refl [THEN bin-abs-lem [rule-format]]])
  apply (unfold Pls-def Min-def number-of-eq)
  prefer 2
  apply (erule asm-rl)
apply auto
done

constdefs
  bin-rec :: 'a => 'a => (int => bit => 'a => 'a) => int => 'a
  bin-rec f1 f2 f3 bin == bin-rec' (bin, f1, f2, f3)

lemma bin-rec-PM:
  f = bin-rec f1 f2 f3 ==> f Numeral.Pls = f1 & f Numeral.Min = f2
apply safe
apply (unfold bin-rec-def)
apply (auto intro: bin-rec'.simps [THEN trans])
done

lemmas bin-rec-Pls = refl [THEN bin-rec-PM, THEN conjunct1, standard]
lemmas bin-rec-Min = refl [THEN bin-rec-PM, THEN conjunct2, standard]

lemma bin-rec-Bit:
  f = bin-rec f1 f2 f3 ==> f3 Numeral.Pls bit.B0 f1 = f1 ==>
    f3 Numeral.Min bit.B1 f2 = f2 ==> f (w BIT b) = f3 w b (f w)
apply clarify
apply (unfold bin-rec-def)
apply (rule bin-rec'.simps [THEN trans])
apply auto
  apply (unfold Pls-def Min-def Bit-def)
  apply (cases b, auto)+
done

lemmas bin-rec-simps = refl [THEN bin-rec-Bit] bin-rec-Pls bin-rec-Min

```

## 4.2 Destructors for binary integers

### consts

— corresponding operations analysing bins

```

bin-last :: int => bit
bin-rest :: int => int
bin-sign :: int => int
bin-nth :: int => nat => bool

```

### primrec

```

Z : bin-nth w 0 = (bin-last w = bit.B1)
Suc : bin-nth w (Suc n) = bin-nth (bin-rest w) n

```

**defs**

```

bin-rest-def : bin-rest w == fst (bin-rl w)
bin-last-def : bin-last w == snd (bin-rl w)
bin-sign-def : bin-sign == bin-rec Numeral.Pls Numeral.Min (%w b s. s)

```

**lemma** *bin-rl*:  $\text{bin-rl } w = (\text{bin-rest } w, \text{bin-last } w)$   
**unfolding** *bin-rest-def bin-last-def* **by** *auto*

**lemmas** *bin-rl-simp* [*simp*] = *iffD1* [*OF bin-rl-char bin-rl*]

**lemma** *bin-rest-simps* [*simp*]:  
 $\text{bin-rest Numeral.Pls} = \text{Numeral.Pls}$   
 $\text{bin-rest Numeral.Min} = \text{Numeral.Min}$   
 $\text{bin-rest } (w \text{ BIT } b) = w$   
**unfolding** *bin-rest-def* **by** *auto*

**lemma** *bin-last-simps* [*simp*]:  
 $\text{bin-last Numeral.Pls} = \text{bit.B0}$   
 $\text{bin-last Numeral.Min} = \text{bit.B1}$   
 $\text{bin-last } (w \text{ BIT } b) = b$   
**unfolding** *bin-last-def* **by** *auto*

**lemma** *bin-sign-simps* [*simp*]:  
 $\text{bin-sign Numeral.Pls} = \text{Numeral.Pls}$   
 $\text{bin-sign Numeral.Min} = \text{Numeral.Min}$   
 $\text{bin-sign } (w \text{ BIT } b) = \text{bin-sign } w$   
**unfolding** *bin-sign-def* **by** (*auto simp: bin-rec-simps*)

**lemma** *bin-r-l-extras* [*simp*]:  
 $\text{bin-last } 0 = \text{bit.B0}$   
 $\text{bin-last } (-1) = \text{bit.B1}$   
 $\text{bin-last } -1 = \text{bit.B1}$   
 $\text{bin-last } 1 = \text{bit.B1}$   
 $\text{bin-rest } 1 = 0$   
 $\text{bin-rest } 0 = 0$   
 $\text{bin-rest } (-1) = -1$   
 $\text{bin-rest } -1 = -1$   
**apply** (*unfold number-of-Min*)  
**apply** (*unfold Pls-def [symmetric] Min-def [symmetric]*)  
**apply** (*unfold numeral-1-eq-1 [symmetric]*)  
**apply** (*auto simp: number-of-eq*)  
**done**

**lemma** *bin-last-mod*:  
 $\text{bin-last } w = (\text{if } w \bmod 2 = 0 \text{ then bit.B0 else bit.B1})$   
**apply** (*case-tac w rule: bin-exhaust*)  
**apply** (*case-tac b*)  
**apply** *auto*  
**done**



**lemma** *bin-rest-div*:

```

  bin-rest  $w = w \text{ div } 2$ 
  apply (case-tac  $w$  rule: bin-exhaust)
  apply (rule trans)
  apply clarsimp
  apply (rule refl)
  apply (drule trans)
  apply (rule Bit-def)
  apply (simp add: z1pdiv2 split: bit.split)
done

```

**lemma** *Bit-div2* [simp]:  $(w \text{ BIT } b) \text{ div } 2 = w$   
 unfolding *bin-rest-div* [symmetric] by auto

**lemma** *bin-nth-lem* [rule-format]:

```

  ALL  $y$ . bin-nth  $x = \text{bin-nth } y \longrightarrow x = y$ 
  apply (induct  $x$  rule: bin-induct)
  apply safe
  apply (erule rev-mp)
  apply (induct-tac  $y$  rule: bin-induct)
  apply safe
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (erule notE, rule ext,
    drule-tac  $x=\text{Suc } x$  in fun-cong, force)
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (erule rev-mp)
  apply (induct-tac  $y$  rule: bin-induct)
  apply safe
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (erule notE, rule ext,
    drule-tac  $x=\text{Suc } x$  in fun-cong, force)
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (case-tac  $y$  rule: bin-exhaust)
  apply clarify
  apply (erule allE)
  apply (erule impE)
  prefer 2
  apply (erule BIT-eqI)
  apply (drule-tac  $x=0$  in fun-cong, force)
  apply (rule ext)
  apply (drule-tac  $x=\text{Suc } ?x$  in fun-cong, force)
done

```

**lemma** *bin-nth-eq-iff*:  $(\text{bin-nth } x = \text{bin-nth } y) = (x = y)$   
 by (auto elim: bin-nth-lem)

**lemmas** *bin-eqI* = ext [THEN *bin-nth-eq-iff* [THEN *iffD1*], standard]

**lemma** *bin-nth-Pls* [simp]:  $\sim \text{bin-nth Numeral.Pls } n$   
**by** (induct *n*) auto

**lemma** *bin-nth-Min* [simp]:  $\text{bin-nth Numeral.Min } n$   
**by** (induct *n*) auto

**lemma** *bin-nth-0-BIT*:  $\text{bin-nth } (w \text{ BIT } b) \ 0 = (b = \text{bit.B1})$   
**by** auto

**lemma** *bin-nth-Suc-BIT*:  $\text{bin-nth } (w \text{ BIT } b) \ (\text{Suc } n) = \text{bin-nth } w \ n$   
**by** auto

**lemma** *bin-nth-minus* [simp]:  $0 < n \implies \text{bin-nth } (w \text{ BIT } b) \ n = \text{bin-nth } w \ (n - 1)$   
**by** (cases *n*) auto

**lemmas** *bin-nth-0* = *bin-nth.simps*(1)  
**lemmas** *bin-nth-Suc* = *bin-nth.simps*(2)

**lemmas** *bin-nth-simps* =  
*bin-nth-0 bin-nth-Suc bin-nth-Pls bin-nth-Min bin-nth-minus*

**lemma** *bin-sign-rest* [simp]:  
 $\text{bin-sign } (\text{bin-rest } w) = (\text{bin-sign } w)$   
**by** (case-tac *w* rule: *bin-exhaust*) auto

### 4.3 Truncating binary integers

**consts**  
*bintrunc* :: *nat* => *int* => *int*  
**primrec**  
 $Z : \text{bintrunc } 0 \text{ bin} = \text{Numeral.Pls}$   
 $\text{Suc} : \text{bintrunc } (\text{Suc } n) \text{ bin} = \text{bintrunc } n \ (\text{bin-rest bin}) \text{ BIT } (\text{bin-last bin})$

**consts**  
*sbintrunc* :: *nat* => *int* => *int*  
**primrec**  
 $Z : \text{sbintrunc } 0 \text{ bin} =$   
 $(\text{case bin-last bin of bit.B1} \implies \text{Numeral.Min} \mid \text{bit.B0} \implies \text{Numeral.Pls})$   
 $\text{Suc} : \text{sbintrunc } (\text{Suc } n) \text{ bin} = \text{sbintrunc } n \ (\text{bin-rest bin}) \text{ BIT } (\text{bin-last bin})$

**lemma** *sign-bintr*:  
 $!!w. \text{bin-sign } (\text{bintrunc } n \ w) = \text{Numeral.Pls}$   
**by** (induct *n*) auto

**lemma** *bintrunc-mod2p*:  
 $!!w. \text{bintrunc } n \ w = (w \text{ mod } 2^n :: \text{int})$   
**apply** (induct *n*, *clarsimp*)  
**apply** (*simp add: bin-last-mod bin-rest-div Bit-def zmod-zmult2-eq*)

```

      cong: number-of-False-cong)
done

lemma sbintrunc-mod2p:
  !!w. sbintrunc n w = ((w + 2 ^ n) mod 2 ^ (Suc n) - 2 ^ n :: int)
  apply (induct n)
  apply clarsimp
  apply (subst zmod-zadd-left-eq)
  apply (simp add: bin-last-mod)
  apply (simp add: number-of-eq)
  apply clarsimp
  apply (simp add: bin-last-mod bin-rest-div Bit-def
    cong: number-of-False-cong)
  apply (clarsimp simp: zmod-zmult-zmult1 [symmetric]
    zmod-zdiv-equality [THEN diff-eq-eq [THEN iffD2 [THEN sym]]])
  apply (rule trans [symmetric, OF - emep1])
  apply auto
  apply (auto simp: even-def)
done

```

#### 4.4 Simplifications for (s)bintrunc

```

lemma bit-bool:
  (b = (b' = bit.B1)) = (b' = (if b then bit.B1 else bit.B0))
  by (cases b') auto

lemmas bit-bool1 [simp] = refl [THEN bit-bool [THEN iffD1], symmetric]

lemma bin-sign-lem:
  !!bin. (bin-sign (sbintrunc n bin) = Numeral.Min) = bin-nth bin n
  apply (induct n)
  apply (case-tac bin rule: bin-exhaust, case-tac b, auto)+
done

lemma nth-bintr:
  !!w m. bin-nth (bintrunc m w) n = (n < m & bin-nth w n)
  apply (induct n)
  apply (case-tac m, auto)[1]
  apply (case-tac m, auto)[1]
done

lemma nth-sbintr:
  !!w m. bin-nth (sbintrunc m w) n =
    (if n < m then bin-nth w n else bin-nth w m)
  apply (induct n)
  apply (case-tac m, simp-all split: bit.splits)[1]
  apply (case-tac m, simp-all split: bit.splits)[1]
done

```

**lemma** *bin-nth-Bit*:

*bin-nth* (*w BIT b*) *n* = (*n* = 0 & *b* = *bit.B1* | (*EX m. n* = *Suc m* & *bin-nth w m*))  
**by** (*cases n*) *auto*

**lemma** *bintrunc-bintrunc-l*:

*n* <= *m* ==> (*bintrunc m (bintrunc n w)* = *bintrunc n w*)  
**by** (*rule bin-eqI*) (*auto simp add : nth-bintr*)

**lemma** *sbintrunc-sbintrunc-l*:

*n* <= *m* ==> (*sbintrunc m (sbintrunc n w)* = *sbintrunc n w*)  
**by** (*rule bin-eqI*) (*auto simp: nth-sbintr min-def*)

**lemma** *bintrunc-bintrunc-ge*:

*n* <= *m* ==> (*bintrunc n (bintrunc m w)* = *bintrunc n w*)  
**by** (*rule bin-eqI*) (*auto simp: nth-bintr*)

**lemma** *bintrunc-bintrunc-min [simp]*:

*bintrunc m (bintrunc n w)* = *bintrunc (min m n) w*  
**apply** (*unfold min-def*)  
**apply** (*rule bin-eqI*)  
**apply** (*auto simp: nth-bintr*)  
**done**

**lemma** *sbintrunc-sbintrunc-min [simp]*:

*sbintrunc m (sbintrunc n w)* = *sbintrunc (min m n) w*  
**apply** (*unfold min-def*)  
**apply** (*rule bin-eqI*)  
**apply** (*auto simp: nth-sbintr*)  
**done**

**lemmas** *bintrunc-Pls* =

*bintrunc.Suc* [**where** *bin*=*Numeral.Pls*, *simplified bin-last-simps bin-rest-simps*, *standard*]

**lemmas** *bintrunc-Min [simp]* =

*bintrunc.Suc* [**where** *bin*=*Numeral.Min*, *simplified bin-last-simps bin-rest-simps*, *standard*]

**lemmas** *bintrunc-BIT [simp]* =

*bintrunc.Suc* [**where** *bin*=*w BIT b*, *simplified bin-last-simps bin-rest-simps*, *standard*]

**lemmas** *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

**lemmas** *sbintrunc-Suc-Pls* =

*sbintrunc.Suc* [**where** *bin*=*Numeral.Pls*, *simplified bin-last-simps bin-rest-simps*, *standard*]

**lemmas** *sbintrunc-Suc-Min* =  
*sbintrunc.Suc* [where *bin*=Numeral.Min, simplified bin-last-simps bin-rest-simps,  
 standard]

**lemmas** *sbintrunc-Suc-BIT* [simp] =  
*sbintrunc.Suc* [where *bin*=*w* BIT *b*, simplified bin-last-simps bin-rest-simps, stan-  
 dard]

**lemmas** *sbintrunc-Sucs* = *sbintrunc-Suc-Pls* *sbintrunc-Suc-Min* *sbintrunc-Suc-BIT*

**lemmas** *sbintrunc-Pls* =  
*sbintrunc.Z* [where *bin*=Numeral.Pl, simplified bin-last-simps bin-rest-simps bit.simps, standard]

**lemmas** *sbintrunc-Min* =  
*sbintrunc.Z* [where *bin*=Numeral.Min, simplified bin-last-simps bin-rest-simps bit.simps, standard]

**lemmas** *sbintrunc-0-BIT-B0* [simp] =  
*sbintrunc.Z* [where *bin*=*w* BIT bit.B0, simplified bin-last-simps bin-rest-simps bit.simps, standard]

**lemmas** *sbintrunc-0-BIT-B1* [simp] =  
*sbintrunc.Z* [where *bin*=*w* BIT bit.B1, simplified bin-last-simps bin-rest-simps bit.simps, standard]

**lemmas** *sbintrunc-0-simps* =  
*sbintrunc-Pls* *sbintrunc-Min* *sbintrunc-0-BIT-B0* *sbintrunc-0-BIT-B1*

**lemmas** *bintrunc-simps* = *bintrunc.Z* *bintrunc-Sucs*  
**lemmas** *sbintrunc-simps* = *sbintrunc-0-simps* *sbintrunc-Sucs*

**lemma** *bintrunc-minus*:  
 $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) \ w = \text{bintrunc } n \ w$   
 by *auto*

**lemma** *sbintrunc-minus*:  
 $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) \ w = \text{sbintrunc } n \ w$   
 by *auto*

**lemmas** *bintrunc-minus-simps* =  
*bintrunc-Sucs* [THEN [2] *bintrunc-minus* [symmetric, THEN trans], standard]

**lemmas** *sbintrunc-minus-simps* =  
*sbintrunc-Sucs* [THEN [2] *sbintrunc-minus* [symmetric, THEN trans], standard]

**lemma** *bintrunc-n-Pls* [simp]:  
 $\text{bintrunc } n \ \text{Numeral.Pl} = \text{Numeral.Pl}$   
 by (induct *n*) *auto*

**lemma** *sbintrunc-n-PM* [simp]:

*sbintrunc n Numeral.Pls = Numeral.Pls*  
*sbintrunc n Numeral.Min = Numeral.Min*  
**by** (induct n) auto

**lemmas** *thobini1* = arg-cong [where  $f = \%w. w \text{ BIT } b$ , standard]

**lemmas** *bintrunc-BIT-I* = trans [OF *bintrunc-BIT thobini1*]

**lemmas** *bintrunc-Min-I* = trans [OF *bintrunc-Min thobini1*]

**lemmas** *bmsts* = *bintrunc-minus-simps* [THEN *thobini1* [THEN [2] trans], standard]

**lemmas** *bintrunc-Pls-minus-I* = *bmsts*(1)

**lemmas** *bintrunc-Min-minus-I* = *bmsts*(2)

**lemmas** *bintrunc-BIT-minus-I* = *bmsts*(3)

**lemma** *bintrunc-0-Min*: *bintrunc 0 Numeral.Min = Numeral.Pls*

**by** auto

**lemma** *bintrunc-0-BIT*: *bintrunc 0 (w BIT b) = Numeral.Pls*

**by** auto

**lemma** *bintrunc-Suc-lem*:

*bintrunc (Suc n) x = y ==> m = Suc n ==> bintrunc m x = y*

**by** auto

**lemmas** *bintrunc-Suc-Ialts* =

*bintrunc-Min-I bintrunc-BIT-I* [THEN *bintrunc-Suc-lem*, standard]

**lemmas** *sbintrunc-BIT-I* = trans [OF *sbintrunc-Suc-BIT thobini1*]

**lemmas** *sbintrunc-Suc-Is* =

*sbintrunc-Sucs* [THEN *thobini1* [THEN [2] trans], standard]

**lemmas** *sbintrunc-Suc-minus-Is* =

*sbintrunc-minus-simps* [THEN *thobini1* [THEN [2] trans], standard]

**lemma** *sbintrunc-Suc-lem*:

*sbintrunc (Suc n) x = y ==> m = Suc n ==> sbintrunc m x = y*

**by** auto

**lemmas** *sbintrunc-Suc-Ialts* =

*sbintrunc-Suc-Is* [THEN *sbintrunc-Suc-lem*, standard]

**lemma** *sbintrunc-bintrunc-lt*:

$m > n ==> sbintrunc n (bintrunc m w) = sbintrunc n w$

**by** (rule bin-eqI) (auto simp: nth-sbintr nth-bintr)

**lemma** *bintrunc-sbintrunc-le*:

$m \leq Suc n ==> bintrunc m (sbintrunc n w) = bintrunc m w$

```

apply (rule bin-eqI)
apply (auto simp: nth-sbintr nth-bintr)
apply (subgoal-tac x=n, safe, arith+)[1]
apply (subgoal-tac x=n, safe, arith+)[1]
done

lemmas bintrunc-sbintrunc [simp] = order-refl [THEN bintrunc-sbintrunc-le]
lemmas sbintrunc-bintrunc [simp] = lessI [THEN sbintrunc-bintrunc-lt]
lemmas bintrunc-bintrunc [simp] = order-refl [THEN bintrunc-bintrunc-l]
lemmas sbintrunc-sbintrunc [simp] = order-refl [THEN sbintrunc-sbintrunc-l]

lemma bintrunc-sbintrunc' [simp]:
   $0 < n \implies \text{bintrunc } n (\text{sbintrunc } (n - 1) w) = \text{bintrunc } n w$ 
by (cases n) (auto simp del: bintrunc.Suc)

lemma sbintrunc-bintrunc' [simp]:
   $0 < n \implies \text{sbintrunc } (n - 1) (\text{bintrunc } n w) = \text{sbintrunc } (n - 1) w$ 
by (cases n) (auto simp del: bintrunc.Suc)

lemma bin-sbin-eq-iff:
   $\text{bintrunc } (\text{Suc } n) x = \text{bintrunc } (\text{Suc } n) y \iff$ 
   $\text{sbintrunc } n x = \text{sbintrunc } n y$ 
apply (rule iffI)
apply (rule box-equals [OF - sbintrunc-bintrunc sbintrunc-bintrunc])
apply simp
apply (rule box-equals [OF - bintrunc-sbintrunc bintrunc-sbintrunc])
apply simp
done

lemma bin-sbin-eq-iff':
   $0 < n \implies \text{bintrunc } n x = \text{bintrunc } n y \iff$ 
   $\text{sbintrunc } (n - 1) x = \text{sbintrunc } (n - 1) y$ 
by (cases n) (simp-all add: bin-sbin-eq-iff del: bintrunc.Suc)

lemmas bintrunc-sbintruncS0 [simp] = bintrunc-sbintrunc' [unfolded One-nat-def]
lemmas sbintrunc-bintruncS0 [simp] = sbintrunc-bintrunc' [unfolded One-nat-def]

lemmas bintrunc-bintrunc-l' = le-add1 [THEN bintrunc-bintrunc-l]
lemmas sbintrunc-sbintrunc-l' = le-add1 [THEN sbintrunc-sbintrunc-l]

lemmas nat-non0-gr =
  trans [OF iszero-def [THEN Not-eq-iff [THEN iffD2]] refl, standard]

lemmas bintrunc-pred-simps [simp] =
  bintrunc-minus-simps [of number-of bin, simplified nobm1, standard]

lemmas sbintrunc-pred-simps [simp] =

```

*sbintrunc-minus-simps* [of number-of bin, simplified nobm1, standard]

**lemma** *no-bintr-alt*:

*number-of* (bintrunc  $n$   $w$ ) =  $w \bmod 2^n$   
**by** (*simp add*: *number-of-eq bintrunc-mod2p*)

**lemma** *no-bintr-alt1*: bintrunc  $n$  = (% $w$ .  $w \bmod 2^n :: \text{int}$ )  
**by** (*rule ext*) (*rule bintrunc-mod2p*)

**lemma** *range-bintrunc*: range (bintrunc  $n$ ) = { $i$ .  $0 \leq i \ \& \ i < 2^n$ }  
**apply** (*unfold no-bintr-alt1*)  
**apply** (*auto simp add*: *image-iff*)  
**apply** (*rule exI*)  
**apply** (*auto intro*: *int-mod-lem* [*THEN iffD1, symmetric*])  
**done**

**lemma** *no-bintr*:

*number-of* (bintrunc  $n$   $w$ ) = (*number-of*  $w \bmod 2^n :: \text{int}$ )  
**by** (*simp add* : *bintrunc-mod2p number-of-eq*)

**lemma** *no-sbintr-alt2*:

*sbintrunc*  $n$  = (% $w$ .  $(w + 2^n) \bmod 2^{Suc\ n} - 2^n :: \text{int}$ )  
**by** (*rule ext*) (*simp add* : *sbintrunc-mod2p*)

**lemma** *no-sbintr*:

*number-of* (*sbintrunc*  $n$   $w$ ) =  
 ((*number-of*  $w + 2^n \bmod 2^{Suc\ n} - 2^n :: \text{int}$ )  
**by** (*simp add* : *no-sbintr-alt2 number-of-eq*)

**lemma** *range-sbintrunc*:

range (*sbintrunc*  $n$ ) = { $i$ .  $-(2^n) \leq i \ \& \ i < 2^n$ }  
**apply** (*unfold no-sbintr-alt2*)  
**apply** (*auto simp add*: *image-iff eq-diff-eq*)  
**apply** (*rule exI*)  
**apply** (*auto intro*: *int-mod-lem* [*THEN iffD1, symmetric*])  
**done**

**lemma** *sb-inc-lem*:

( $a :: \text{int}$ ) +  $2^k < 0 \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$   
**apply** (*erule int-mod-ge'* [**where**  $n = 2^{(Suc\ k)}$  **and**  $b = a + 2^k$ , *simplified zless2p*])  
**apply** (*rule TrueI*)  
**done**

**lemma** *sb-inc-lem'*:

( $a :: \text{int}$ ) <  $-(2^k) \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$   
**by** (*rule iffD1* [*OF less-diff-eq, THEN sb-inc-lem, simplified OrderedGroup.diff-0*])

**lemma** *sbintrunc-inc*:



$x < - (2^n) \implies x + 2^n (Suc\ n) \leq sbintrunc\ n\ x$   
**unfolding** *no-sbintr-alt2* **by** (*drule sb-inc-lem'*) *simp*

**lemma** *sb-dec-lem*:

$(0::int) \leq - (2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$   
**by** (*rule int-mod-le'* [**where**  $n = 2^k (Suc\ k)$  **and**  $b = a + 2^k$ ,  
*simplified zless2p, OF - TrueI, simplified*])

**lemma** *sb-dec-lem'*:

$(2::int)^k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq - (2^k) + a$   
**by** (*rule iffD1* [*OF diff-le-eq', THEN sb-dec-lem, simplified*])

**lemma** *sbintrunc-dec*:

$x \geq (2^n) \implies x - 2^n (Suc\ n) \geq sbintrunc\ n\ x$   
**unfolding** *no-sbintr-alt2* **by** (*drule sb-dec-lem'*) *simp*

**lemmas** *zmod-uminus'* = *zmod-uminus* [**where**  $b=c$ , *standard*]

**lemmas** *zpower-zmod'* = *zpower-zmod* [**where**  $m=c$  **and**  $y=k$ , *standard*]

**lemmas** *brdmod1s'* [*symmetric*] =

*zmod-zadd-left-eq zmod-zadd-right-eq*  
*zmod-zsub-left-eq zmod-zsub-right-eq*  
*zmod-zmult1-eq zmod-zmult1-eq-rev*

**lemmas** *brdmods'* [*symmetric*] =

*zpower-zmod'* [*symmetric*]  
*trans* [*OF zmod-zadd-left-eq zmod-zadd-right-eq*]  
*trans* [*OF zmod-zsub-left-eq zmod-zsub-right-eq*]  
*trans* [*OF zmod-zmult1-eq zmod-zmult1-eq-rev*]  
*zmod-uminus'* [*symmetric*]  
*zmod-zadd-left-eq* [**where**  $b = 1$ ]  
*zmod-zsub-left-eq* [**where**  $b = 1$ ]

**lemmas** *bintr-arith1s* =

*brdmod1s'* [**where**  $c=2^n$ , *folded pred-def succ-def bintrunc-mod2p, standard*]

**lemmas** *bintr-ariths* =

*brdmods'* [**where**  $c=2^n$ , *folded pred-def succ-def bintrunc-mod2p, standard*]

**lemmas** *m2pths* = *pos-mod-sign pos-mod-bound* [*OF zless2p, standard*]

**lemma** *bintr-ge0*:  $(0 :: int) \leq \text{number-of}\ (bintrunc\ n\ w)$

**by** (*simp add : no-bintr m2pths*)

**lemma** *bintr-lt2p*:  $\text{number-of}\ (bintrunc\ n\ w) < (2^n :: int)$

**by** (*simp add : no-bintr m2pths*)

**lemma** *bintr-Min*:

$\text{number-of}\ (bintrunc\ n\ \text{Numeral.Min}) = (2^n :: int) - 1$

**by** (*simp add : no-bintr m1mod2k*)

```

lemma sbintr-ge:  $(-(2^n :: \text{int}) \leq \text{number-of } (\text{sbintrunc } n \ w))$ 
  by (simp add : no-sbintr m2pths)

lemma sbintr-lt:  $\text{number-of } (\text{sbintrunc } n \ w) < (2^n :: \text{int})$ 
  by (simp add : no-sbintr m2pths)

lemma bintrunc-Suc:
   $\text{bintrunc } (\text{Suc } n) \ \text{bin} = \text{bintrunc } n \ (\text{bin-rest } \text{bin}) \ \text{BIT } \text{bin-last } \text{bin}$ 
  by (case-tac bin rule: bin-exhaust) auto

lemma sign-Pls-ge-0:
   $(\text{bin-sign } \text{bin} = \text{Numeral.Pls}) = (\text{number-of } \text{bin} \geq (0 :: \text{int}))$ 
  by (induct bin rule: bin-induct) auto

lemma sign-Min-lt-0:
   $(\text{bin-sign } \text{bin} = \text{Numeral.Min}) = (\text{number-of } \text{bin} < (0 :: \text{int}))$ 
  by (induct bin rule: bin-induct) auto

lemmas sign-Min-neg = trans [OF sign-Min-lt-0 neg-def [symmetric]]

lemma bin-rest-trunc:
   $!!\text{bin. } (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bintrunc } (n - 1) \ (\text{bin-rest } \text{bin})$ 
  by (induct n) auto

lemma bin-rest-power-trunc [rule-format] :
   $(\text{bin-rest } ^k) \ (\text{bintrunc } n \ \text{bin}) =$ 
   $\text{bintrunc } (n - k) \ ((\text{bin-rest } ^k) \ \text{bin})$ 
  by (induct k) (auto simp: bin-rest-trunc)

lemma bin-rest-trunc-i:
   $\text{bintrunc } n \ (\text{bin-rest } \text{bin}) = \text{bin-rest } (\text{bintrunc } (\text{Suc } n) \ \text{bin})$ 
  by auto

lemma bin-rest-strunc:
   $!!\text{bin. } \text{bin-rest } (\text{sbintrunc } (\text{Suc } n) \ \text{bin}) = \text{sbintrunc } n \ (\text{bin-rest } \text{bin})$ 
  by (induct n) auto

lemma bintrunc-rest [simp]:
   $!!\text{bin. } \text{bintrunc } n \ (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bin-rest } (\text{bintrunc } n \ \text{bin})$ 
  apply (induct n, simp)
  apply (case-tac bin rule: bin-exhaust)
  apply (auto simp: bintrunc-bintrunc-l)
  done

lemma sbintrunc-rest [simp]:
   $!!\text{bin. } \text{sbintrunc } n \ (\text{bin-rest } (\text{sbintrunc } n \ \text{bin})) = \text{bin-rest } (\text{sbintrunc } n \ \text{bin})$ 
  apply (induct n, simp)
  apply (case-tac bin rule: bin-exhaust)

```

```

apply (auto simp: bintrunc-bintrunc-l split: bit.splits)
done

```

```

lemma bintrunc-rest':
  bintrunc n o bin-rest o bintrunc n = bin-rest o bintrunc n
by (rule ext) auto

```

```

lemma sbintrunc-rest' :
  sbintrunc n o bin-rest o sbintrunc n = bin-rest o sbintrunc n
by (rule ext) auto

```

```

lemma rco-lem:
  f o g o f = g o f ==> f o (g o f) ^ n = g ^ n o f
apply (rule ext)
apply (induct-tac n)
  apply (simp-all (no-asm))
apply (drule fun-cong)
apply (unfold o-def)
apply (erule trans)
apply simp
done

```

```

lemma rco-alt: (f o g) ^ n o f = f o (g o f) ^ n
apply (rule ext)
apply (induct n)
  apply (simp-all add: o-def)
done

```

```

lemmas rco-bintr = bintrunc-rest'
  [THEN rco-lem [THEN fun-cong], unfolded o-def]
lemmas rco-sbintr = sbintrunc-rest'
  [THEN rco-lem [THEN fun-cong], unfolded o-def]

```

## 4.5 Splitting and concatenation

```

consts
  bin-split :: nat => int => int * int
primrec
  Z : bin-split 0 w = (w, Numeral.Pls)
  Suc : bin-split (Suc n) w = (let (w1, w2) = bin-split n (bin-rest w)
    in (w1, w2 BIT bin-last w))

```

```

consts
  bin-cat :: int => nat => int => int
primrec
  Z : bin-cat w 0 v = w
  Suc : bin-cat w (Suc n) v = bin-cat w n (bin-rest v) BIT bin-last v

```

## 4.6 Miscellaneous lemmas

**lemmas** *funpow-minus-simp* =

*trans [OF gen-minus [where f = power f] funpow-Suc, standard]*

**lemmas** *funpow-pred-simp* [simp] =

*funpow-minus-simp [of number-of bin, simplified nobm1, standard]*

**lemmas** *replicate-minus-simp* =

*trans [OF gen-minus [where f = %n. replicate n x] replicate.replicate-Suc, standard]*

**lemmas** *replicate-pred-simp* [simp] =

*replicate-minus-simp [of number-of bin, simplified nobm1, standard]*

**lemmas** *power-Suc-no* [simp] = *power-Suc [of number-of a, standard]*

**lemmas** *power-minus-simp* =

*trans [OF gen-minus [where f = power f] power-Suc, standard]*

**lemmas** *power-pred-simp* =

*power-minus-simp [of number-of bin, simplified nobm1, standard]*

**lemmas** *power-pred-simp-no* [simp] = *power-pred-simp [where f = number-of f, standard]*

**lemma** *list-exhaust-size-gt0*:

**assumes** *y:  $\bigwedge a \text{ list. } y = a \# \text{ list} \implies P$*

**shows**  *$0 < \text{length } y \implies P$*

**apply** (*cases y, simp*)

**apply** (*rule y*)

**apply** *fastsimp*

**done**

**lemma** *list-exhaust-size-eq0*:

**assumes** *y:  $y = [] \implies P$*

**shows**  *$\text{length } y = 0 \implies P$*

**apply** (*cases y*)

**apply** (*rule y, simp*)

**apply** *simp*

**done**

**lemma** *size-Cons-lem-eq*:

*y = xa # list ==> size y = Suc k ==> size list = k*

**by** *auto*

**lemma** *size-Cons-lem-eq-bin*:

*y = xa # list ==> size y = number-of (Numeral.succ k) ==>*

*size list = number-of k*

**by** (*auto simp: pred-def succ-def split add : split-if-asm*)

```

lemmas ls-splits =
  prod.split split-split prod.split-asm split-split-asm split-if-asm

```

```

lemma not-B1-is-B0:  $y \neq \text{bit}.B1 \implies y = \text{bit}.B0$ 
  by (cases y) auto

```

```

lemma B1-ass-B0:
  assumes  $y: y = \text{bit}.B0 \implies y = \text{bit}.B1$ 
  shows  $y = \text{bit}.B1$ 
  apply (rule classical)
  apply (drule not-B1-is-B0)
  apply (erule y)
  done

```

— simplifications for specific word lengths

```

lemmas n2s-ths [THEN eq-reflection] = add-2-eq-Suc add-2-eq-Suc'

```

```

lemmas s2n-ths = n2s-ths [symmetric]

```

```

end

```

## 5 BitSyntax: Syntactic class for bitwise operations

```

theory BitSyntax
imports Main
begin

```

```

class bit = type +
  fixes bitNOT :: 'a  $\Rightarrow$  'a
    and bitAND :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a
    and bitOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a
    and bitXOR :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a

```

We want the bitwise operations to bind slightly weaker than  $+$  and  $-$ , but  $\sim\sim$  to bind slightly stronger than  $*$ .

```

notation
  bitNOT (NOT - [70] 71) and
  bitAND (infixr AND 64) and
  bitOR (infixr OR 59) and
  bitXOR (infixr XOR 59)

```

Testing and shifting operations.

```

consts
  test-bit :: 'a::bit  $\Rightarrow$  nat  $\Rightarrow$  bool (infixl !! 100)
  lsb      :: 'a::bit  $\Rightarrow$  bool

```

```

msb      :: 'a::bit ⇒ bool
set-bit  :: 'a::bit ⇒ nat ⇒ bool ⇒ 'a
set-bits :: (nat ⇒ bool) ⇒ 'a::bit (binder BITS 10)
shiftrl  :: 'a::bit ⇒ nat ⇒ 'a (infixl << 55)
shiftr   :: 'a::bit ⇒ nat ⇒ 'a (infixl >> 55)

```

### 5.1 Bitwise operations on type *bit*

**instance** *bit* :: *bit*

```

NOT-bit-def: NOT x ≡ case x of bit.B0 ⇒ bit.B1 | bit.B1 ⇒ bit.B0
AND-bit-def: x AND y ≡ case x of bit.B0 ⇒ bit.B0 | bit.B1 ⇒ y
OR-bit-def:  x OR y  :: bit ≡ NOT (NOT x AND NOT y)
XOR-bit-def: x XOR y :: bit ≡ (x AND NOT y) OR (NOT x AND y)
..

```

**lemma** *bit-simps* [*simp*]:

```

NOT bit.B0 = bit.B1
NOT bit.B1 = bit.B0
bit.B0 AND y = bit.B0
bit.B1 AND y = y
bit.B0 OR y = y
bit.B1 OR y = bit.B1
bit.B0 XOR y = y
bit.B1 XOR y = NOT y
by (simp-all add: NOT-bit-def AND-bit-def OR-bit-def
    XOR-bit-def split: bit.split)

```

**lemma** *bit-extra-simps* [*simp*]:

```

x AND bit.B0 = bit.B0
x AND bit.B1 = x
x OR bit.B1 = bit.B1
x OR bit.B0 = x
x XOR bit.B1 = NOT x
x XOR bit.B0 = x
by (cases x, auto)+

```

**lemma** *bit-ops-comm*:

```

(x::bit) AND y = y AND x
(x::bit) OR y = y OR x
(x::bit) XOR y = y XOR x
by (cases y, auto)+

```

**lemma** *bit-ops-same* [*simp*]:

```

(x::bit) AND x = x
(x::bit) OR x = x
(x::bit) XOR x = bit.B0
by (cases x, auto)+

```

**lemma** *bit-not-not* [*simp*]: NOT (NOT (x::bit)) = x

```

  by (cases x) auto
end

```

## 6 BinOperations: Bitwise Operations on Binary Integers

```

theory BinOperations imports BinGeneral BitSyntax

begin

```

### 6.1 Logical operations

bit-wise logical operations on the int type

```

instance int :: bit
  int-not-def: bitNOT  $\equiv$  bin-rec Numeral.Min Numeral.Pls
    ( $\lambda w b s. s \text{ BIT } (\text{NOT } b)$ )
  int-and-def: bitAND  $\equiv$  bin-rec ( $\lambda x. \text{Numeral.Pls}$ ) ( $\lambda y. y$ )
    ( $\lambda w b s y. s \text{ (bin-rest } y) \text{ BIT } (b \text{ AND bin-last } y)$ )
  int-or-def: bitOR  $\equiv$  bin-rec ( $\lambda x. x$ ) ( $\lambda y. \text{Numeral.Min}$ )
    ( $\lambda w b s y. s \text{ (bin-rest } y) \text{ BIT } (b \text{ OR bin-last } y)$ )
  int-xor-def: bitXOR  $\equiv$  bin-rec ( $\lambda x. x$ ) bitNOT
    ( $\lambda w b s y. s \text{ (bin-rest } y) \text{ BIT } (b \text{ XOR bin-last } y)$ )
..

```

```

lemma int-not-simps [simp]:
  NOT Numeral.Pls = Numeral.Min
  NOT Numeral.Min = Numeral.Pls
  NOT (w BIT b) = (NOT w) BIT (NOT b)
  by (unfold int-not-def) (auto intro: bin-rec-simps)

```

```

lemma int-xor-Pls [simp]:
  Numeral.Pls XOR x = x
  unfolding int-xor-def by (simp add: bin-rec-PM)

```

```

lemma int-xor-Min [simp]:
  Numeral.Min XOR x = NOT x
  unfolding int-xor-def by (simp add: bin-rec-PM)

```

```

lemma int-xor-Bits [simp]:
  (x BIT b) XOR (y BIT c) = (x XOR y) BIT (b XOR c)
  apply (unfold int-xor-def)
  apply (rule bin-rec-simps (1) [THEN fun-cong, THEN trans])
  apply (rule ext, simp)
  prefer 2
  apply simp
  apply (rule ext)

```

```

apply (simp add: int-not-simps [symmetric])
done

```

```

lemma int-xor-x-simps':
  w XOR (Numeral.Pls BIT bit.B0) = w
  w XOR (Numeral.Min BIT bit.B1) = NOT w
apply (induct w rule: bin-induct)
  apply simp-all[4]
  apply (unfold int-xor-Bits)
  apply clarsimp+
done

```

```

lemmas int-xor-extra-simps [simp] = int-xor-x-simps' [simplified arith-simps]

```

```

lemma int-or-Pls [simp]:
  Numeral.Pls OR x = x
by (unfold int-or-def) (simp add: bin-rec-PM)

```

```

lemma int-or-Min [simp]:
  Numeral.Min OR x = Numeral.Min
by (unfold int-or-def) (simp add: bin-rec-PM)

```

```

lemma int-or-Bits [simp]:
  (x BIT b) OR (y BIT c) = (x OR y) BIT (b OR c)
unfolding int-or-def by (simp add: bin-rec-simps)

```

```

lemma int-or-x-simps':
  w OR (Numeral.Pls BIT bit.B0) = w
  w OR (Numeral.Min BIT bit.B1) = Numeral.Min
apply (induct w rule: bin-induct)
  apply simp-all[4]
  apply (unfold int-or-Bits)
  apply clarsimp+
done

```

```

lemmas int-or-extra-simps [simp] = int-or-x-simps' [simplified arith-simps]

```

```

lemma int-and-Pls [simp]:
  Numeral.Pls AND x = Numeral.Pls
unfolding int-and-def by (simp add: bin-rec-PM)

```

```

lemma int-and-Min [simp]:
  Numeral.Min AND x = x
unfolding int-and-def by (simp add: bin-rec-PM)

```

```

lemma int-and-Bits [simp]:
  (x BIT b) AND (y BIT c) = (x AND y) BIT (b AND c)
unfolding int-and-def by (simp add: bin-rec-simps)

```



**lemma** *int-and-x-simps'*:

*w AND (Numeral.Pls BIT bit.B0) = Numeral.Pls*  
*w AND (Numeral.Min BIT bit.B1) = w*  
**apply** (*induct w rule: bin-induct*)  
**apply** *simp-all*[4]  
**apply** (*unfold int-and-Bits*)  
**apply** *clarsimp*+  
**done**

**lemmas** *int-and-extra-simps* [*simp*] = *int-and-x-simps'* [*simplified arith-simps*]

**lemma** *bin-ops-comm*:

**shows**  
*int-and-comm: !!y::int. x AND y = y AND x and*  
*int-or-comm: !!y::int. x OR y = y OR x and*  
*int-xor-comm: !!y::int. x XOR y = y XOR x*  
**apply** (*induct x rule: bin-induct*)  
**apply** *simp-all*[6]  
**apply** (*case-tac y rule: bin-exhaust, simp add: bit-ops-comm*) +  
**done**

**lemma** *bin-ops-same* [*simp*]:

*(x::int) AND x = x*  
*(x::int) OR x = x*  
*(x::int) XOR x = Numeral.Pls*  
**by** (*induct x rule: bin-induct*) *auto*

**lemma** *int-not-not* [*simp*]: *NOT (NOT (x::int)) = x*

**by** (*induct x rule: bin-induct*) *auto*

**lemmas** *bin-log-esimps* =

*int-and-extra-simps int-or-extra-simps int-xor-extra-simps*  
*int-and-Pls int-and-Min int-or-Pls int-or-Min int-xor-Pls int-xor-Min*

**lemma** *bbw-ao-absorb*:

*!!y::int. x AND (y OR x) = x & x OR (y AND x) = x*  
**apply** (*induct x rule: bin-induct*)  
**apply** *auto*  
**apply** (*case-tac [!] y rule: bin-exhaust*)  
**apply** *auto*  
**apply** (*case-tac [!] bit*)  
**apply** *auto*  
**done**

**lemma** *bbw-ao-absorbs-other*:

```

 $x \text{ AND } (x \text{ OR } y) = x \wedge (y \text{ AND } x) \text{ OR } x = (x::int)$ 
 $(y \text{ OR } x) \text{ AND } x = x \wedge x \text{ OR } (x \text{ AND } y) = (x::int)$ 
 $(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::int)$ 
apply (auto simp: bbw-ao-absorb int-or-comm)
  apply (subst int-or-comm)
  apply (simp add: bbw-ao-absorb)
  apply (subst int-and-comm)
  apply (subst int-or-comm)
  apply (simp add: bbw-ao-absorb)
  apply (subst int-and-comm)
  apply (simp add: bbw-ao-absorb)
done

```

**lemmas** bbw-ao-absorbs [simp] = bbw-ao-absorb bbw-ao-absorbs-other

**lemma** int-xor-not:

```

!!y::int. (NOT x) XOR y = NOT (x XOR y) &
  x XOR (NOT y) = NOT (x XOR y)
apply (induct x rule: bin-induct)
  apply auto
  apply (case-tac y rule: bin-exhaust, auto,
    case-tac b, auto)+
done

```

**lemma** bbw-assocs':

```

!!y z::int. (x AND y) AND z = x AND (y AND z) &
  (x OR y) OR z = x OR (y OR z) &
  (x XOR y) XOR z = x XOR (y XOR z)
apply (induct x rule: bin-induct)
  apply (auto simp: int-xor-not)
  apply (case-tac [!] y rule: bin-exhaust)
  apply (case-tac [!] z rule: bin-exhaust)
  apply (case-tac [!] bit)
  apply (case-tac [!] b)
  apply auto
done

```

**lemma** int-and-assoc:

```

(x AND y) AND (z::int) = x AND (y AND z)
by (simp add: bbw-assocs')

```

**lemma** int-or-assoc:

```

(x OR y) OR (z::int) = x OR (y OR z)
by (simp add: bbw-assocs')

```

**lemma** int-xor-assoc:

```

(x XOR y) XOR (z::int) = x XOR (y XOR z)
by (simp add: bbw-assocs')

```

**lemmas** *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

**lemma** *bbw-lcs* [*simp*]:  
 $(y::int) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$   
 $(y::int) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$   
 $(y::int) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$   
**apply** (*auto simp: bbw-assocs [symmetric]*)  
**apply** (*auto simp: bin-ops-comm*)  
**done**

**lemma** *bbw-not-dist*:  
 $!!y::int. \text{ NOT } (x \text{ OR } y) = (\text{ NOT } x) \text{ AND } (\text{ NOT } y)$   
 $!!y::int. \text{ NOT } (x \text{ AND } y) = (\text{ NOT } x) \text{ OR } (\text{ NOT } y)$   
**apply** (*induct x rule: bin-induct*)  
**apply** *auto*  
**apply** (*case-tac [!] y rule: bin-exhaust*)  
**apply** (*case-tac [!] bit, auto*)  
**done**

**lemma** *bbw-oa-dist*:  
 $!!y \ z::int. (x \text{ AND } y) \text{ OR } z =$   
 $(x \text{ OR } z) \text{ AND } (y \text{ OR } z)$   
**apply** (*induct x rule: bin-induct*)  
**apply** *auto*  
**apply** (*case-tac y rule: bin-exhaust*)  
**apply** (*case-tac z rule: bin-exhaust*)  
**apply** (*case-tac ba, auto*)  
**done**

**lemma** *bbw-ao-dist*:  
 $!!y \ z::int. (x \text{ OR } y) \text{ AND } z =$   
 $(x \text{ AND } z) \text{ OR } (y \text{ AND } z)$   
**apply** (*induct x rule: bin-induct*)  
**apply** *auto*  
**apply** (*case-tac y rule: bin-exhaust*)  
**apply** (*case-tac z rule: bin-exhaust*)  
**apply** (*case-tac ba, auto*)  
**done**

**lemma** *plus-and-or* [*rule-format*]:  
 $ALL \ y::int. (x \text{ AND } y) + (x \text{ OR } y) = x + y$   
**apply** (*induct x rule: bin-induct*)  
**apply** *clarsimp*  
**apply** *clarsimp*  
**apply** *clarsimp*  
**apply** (*case-tac y rule: bin-exhaust*)  
**apply** *clarsimp*

```

apply (unfold Bit-def)
apply clarsimp
apply (erule-tac  $x = x$  in allE)
apply (simp split: bit.split)
done

```

```

lemma le-int-or:
  !!x. bin-sign y = Numeral.Pls ==>  $x \leq x \text{ OR } y$ 
apply (induct y rule: bin-induct)
  apply clarsimp
  apply clarsimp
apply (case-tac x rule: bin-exhaust)
apply (case-tac b)
  apply (case-tac [!] bit)
  apply (auto simp: less-eq-numeral-code)
done

```

```

lemmas int-and-le =
  xtr3 [OF bbw-ao-absorbs (2) [THEN conjunct2, symmetric] le-int-or]

```

```

lemma bin-nth-ops:
  !!x y. bin-nth (x AND y) n = (bin-nth x n & bin-nth y n)
  !!x y. bin-nth (x OR y) n = (bin-nth x n | bin-nth y n)
  !!x y. bin-nth (x XOR y) n = (bin-nth x n ~ bin-nth y n)
  !!x. bin-nth (NOT x) n = (~ bin-nth x n)
apply (induct n)
  apply safe
  apply (case-tac [!] x rule: bin-exhaust)
  apply simp-all
  apply (case-tac [!] y rule: bin-exhaust)
  apply simp-all
  apply (auto dest: not-B1-is-B0 intro: B1-ass-B0)
done

```

```

lemma bin-add-not:  $x + \text{NOT } x = \text{Numeral.Min}$ 
apply (induct x rule: bin-induct)
  apply clarsimp
  apply clarsimp
apply (case-tac bit, auto)
done

```

```

lemma bin-trunc-ao:
  !!x y. (bintrunc n x) AND (bintrunc n y) = bintrunc n (x AND y)
  !!x y. (bintrunc n x) OR (bintrunc n y) = bintrunc n (x OR y)
apply (induct n)
  apply auto

```

```

    apply (case-tac [!] x rule: bin-exhaust)
    apply (case-tac [!] y rule: bin-exhaust)
    apply auto
done

```

**lemma** *bin-trunc-xor*:

```

!!x y. bintrunc n (bintrunc n x XOR bintrunc n y) =
    bintrunc n (x XOR y)
apply (induct n)
apply auto
apply (case-tac [!] x rule: bin-exhaust)
apply (case-tac [!] y rule: bin-exhaust)
apply auto
done

```

**lemma** *bin-trunc-not*:

```

!!x. bintrunc n (NOT (bintrunc n x)) = bintrunc n (NOT x)
apply (induct n)
apply auto
apply (case-tac [!] x rule: bin-exhaust)
apply auto
done

```

**lemma** *bintr-bintr-i*:

```

x = bintrunc n y ==> bintrunc n x = bintrunc n y
by auto

```

**lemmas** *bin-trunc-and* = *bin-trunc-ao*(1) [*THEN bintr-bintr-i*]

**lemmas** *bin-trunc-or* = *bin-trunc-ao*(2) [*THEN bintr-bintr-i*]

## 6.2 Setting and clearing bits

**consts**

```

bin-sc :: nat => bit => int => int

```

**primrec**

```

Z : bin-sc 0 b w = bin-rest w BIT b

```

```

Suc :

```

```

    bin-sc (Suc n) b w = bin-sc n b (bin-rest w) BIT bin-last w

```

**lemma** *bin-nth-sc* [*simp*]:

```

!!w. bin-nth (bin-sc n b w) n = (b = bit.B1)
by (induct n) auto

```

**lemma** *bin-sc-sc-same* [*simp*]:

```

!!w. bin-sc n c (bin-sc n b w) = bin-sc n c w

```

**by** (*induct n*) *auto*

**lemma** *bin-sc-sc-diff*:

*!!w m. m ~ = n ==>*

*bin-sc m c (bin-sc n b w) = bin-sc n b (bin-sc m c w)*

**apply** (*induct n*)

**apply** (*case-tac* [!] *m*)

**apply** *auto*

**done**

**lemma** *bin-nth-sc-gen*:

*!!w m. bin-nth (bin-sc n b w) m = (if m = n then b = bit.B1 else bin-nth w m)*

**by** (*induct n*) (*case-tac* [!] *m, auto*)

**lemma** *bin-sc-nth [simp]*:

*!!w. (bin-sc n (If (bin-nth w n) bit.B1 bit.B0) w) = w*

**by** (*induct n*) *auto*

**lemma** *bin-sign-sc [simp]*:

*!!w. bin-sign (bin-sc n b w) = bin-sign w*

**by** (*induct n*) *auto*

**lemma** *bin-sc-bintr [simp]*:

*!!w m. bintrunc m (bin-sc n x (bintrunc m (w))) = bintrunc m (bin-sc n x w)*

**apply** (*induct n*)

**apply** (*case-tac* [!] *w rule: bin-exhaust*)

**apply** (*case-tac* [!] *m, auto*)

**done**

**lemma** *bin-clr-le*:

*!!w. bin-sc n bit.B0 w <= w*

**apply** (*induct n*)

**apply** (*case-tac* [!] *w rule: bin-exhaust*)

**apply** *auto*

**apply** (*unfold Bit-def*)

**apply** (*simp-all split: bit.split*)

**done**

**lemma** *bin-set-ge*:

*!!w. bin-sc n bit.B1 w >= w*

**apply** (*induct n*)

**apply** (*case-tac* [!] *w rule: bin-exhaust*)

**apply** *auto*

**apply** (*unfold Bit-def*)

**apply** (*simp-all split: bit.split*)

**done**

**lemma** *bintr-bin-clr-le*:

*!!w m. bintrunc n (bin-sc m bit.B0 w) <= bintrunc n w*

```

apply (induct n)
  apply simp
apply (case-tac w rule: bin-exhaust)
apply (case-tac m)
  apply auto
  apply (unfold Bit-def)
  apply (simp-all split: bit.split)
done

lemma bintr-bin-set-ge:
  !!w m. bintrunc n (bin-sc m bit.B1 w) >= bintrunc n w
  apply (induct n)
  apply simp
  apply (case-tac w rule: bin-exhaust)
  apply (case-tac m)
  apply auto
  apply (unfold Bit-def)
  apply (simp-all split: bit.split)
done

lemma bin-sc-FP [simp]: bin-sc n bit.B0 Numeral.Pls = Numeral.Pls
  by (induct n) auto

lemma bin-sc-TM [simp]: bin-sc n bit.B1 Numeral.Min = Numeral.Min
  by (induct n) auto

lemmas bin-sc-simps = bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP

lemma bin-sc-minus:
  0 < n ==> bin-sc (Suc (n - 1)) b w = bin-sc n b w
  by auto

lemmas bin-sc-Suc-minus =
  trans [OF bin-sc-minus [symmetric] bin-sc.Suc, standard]

lemmas bin-sc-Suc-pred [simp] =
  bin-sc-Suc-minus [of number-of bin, simplified nobm1, standard]

```

### 6.3 Operations on lists of booleans

**consts**

```

bin-to-bl :: nat => int => bool list
bin-to-bl-aux :: nat => int => bool list => bool list
bl-to-bin :: bool list => int
bl-to-bin-aux :: int => bool list => int

bl-of-nth :: nat => (nat => bool) => bool list

```

**primrec**

*Nil* : *bl-to-bin-aux* *w* [] = *w*  
*Cons* : *bl-to-bin-aux* *w* (*b* # *bs*) =  
           *bl-to-bin-aux* (*w BIT* (if *b* then *bit.B1* else *bit.B0*)) *bs*

**primrec**

*Z* : *bin-to-bl-aux* 0 *w bl* = *bl*  
*Suc* : *bin-to-bl-aux* (*Suc n*) *w bl* =  
       *bin-to-bl-aux* *n* (*bin-rest w*) ((*bin-last w* = *bit.B1*) # *bl*)

**defs**

*bin-to-bl-def* : *bin-to-bl n w* == *bin-to-bl-aux n w* []  
*bl-to-bin-def* : *bl-to-bin bs* == *bl-to-bin-aux Numeral.Pls bs*

**primrec**

*Suc* : *bl-of-nth* (*Suc n*) *f* = *f n* # *bl-of-nth n f*  
*Z* : *bl-of-nth* 0 *f* = []

**consts**

*takefill* :: '*a* => *nat* => '*a list* => '*a list*  
*app2* :: ('*a* => '*b* => '*c*) => '*a list* => '*b list* => '*c list*

— *takefill* - like *take* but if argument list too short,  
 — extends result to get requested length

**primrec**

*Z* : *takefill fill* 0 *xs* = []  
*Suc* : *takefill fill* (*Suc n*) *xs* = (  
   *case xs of* [] => *fill* # *takefill fill n xs*  
   | *y* # *ys* => *y* # *takefill fill n ys*)

**defs**

*app2-def* : *app2 f as bs* == *map* (*split f*) (*zip as bs*)

**6.4 Splitting and concatenation**

— *rcat* and *rsplit*

**consts**

*bin-rcat* :: *nat* => *int list* => *int*  
*bin-rsplit-aux* :: *nat* \* *int list* \* *nat* \* *int* => *int list*  
*bin-rsplit* :: *nat* => (*nat* \* *int*) => *int list*  
*bin-rsplittl-aux* :: *nat* \* *int list* \* *nat* \* *int* => *int list*  
*bin-rsplittl* :: *nat* => (*nat* \* *int*) => *int list*

**recdef** *bin-rsplit-aux* *measure* (*fst o snd o snd*)

*bin-rsplit-aux* (*n*, *bs*, (*m*, *c*)) =  
 (if *m* = 0 | *n* = 0 then *bs* else  
   let (*a*, *b*) = *bin-split* *n c*  
   in *bin-rsplit-aux* (*n*, *b* # *bs*, (*m* - *n*, *a*)))

**recdef** *bin-rsplittl-aux* *measure* (*fst o snd o snd*)



```

bin-rsplittl-aux (n, bs, (m, c)) =
  (if m = 0 | n = 0 then bs else
   let (a, b) = bin-split (min m n) c
   in bin-rsplittl-aux (n, b # bs, (m - n, a)))

```

**defs**

```

bin-rcat-def : bin-rcat n bs == foldl (%u v. bin-cat u n v) Numeral.Pls bs
bin-rsplit-def : bin-rsplit n w == bin-rsplit-aux (n, [], w)
bin-rsplittl-def : bin-rsplittl n w == bin-rsplittl-aux (n, [], w)

```

```

declare bin-rsplit-aux.simps [simp del]
declare bin-rsplittl-aux.simps [simp del]

```

**lemma** bin-sign-cat:

```

!!y. bin-sign (bin-cat x n y) = bin-sign x
by (induct n) auto

```

**lemma** bin-cat-Suc-Bit:

```

bin-cat w (Suc n) (v BIT b) = bin-cat w n v BIT b
by auto

```

**lemma** bin-nth-cat:

```

!!n y. bin-nth (bin-cat x k y) n =
  (if n < k then bin-nth y n else bin-nth x (n - k))
apply (induct k)
apply clarsimp
apply (case-tac n, auto)
done

```

**lemma** bin-nth-split:

```

!!b c. bin-split n c = (a, b) ==>
  (ALL k. bin-nth a k = bin-nth c (n + k)) &
  (ALL k. bin-nth b k = (k < n & bin-nth c k))
apply (induct n)
apply clarsimp
apply (clarsimp simp: Let-def split: ls-splits)
apply (case-tac k)
apply auto
done

```

**lemma** bin-cat-assoc:

```

!!z. bin-cat (bin-cat x m y) n z = bin-cat x (m + n) (bin-cat y n z)
by (induct n) auto

```

**lemma** bin-cat-assoc-sym: !!z m.

```

bin-cat x m (bin-cat y n z) = bin-cat (bin-cat x (m - n) y) (min m n) z
apply (induct n, clarsimp)

```

**apply** (*case-tac m, auto*)  
**done**

**lemma** *bin-cat-Pls [simp]*:  
 !!w. *bin-cat Numeral.Pls n w = bintrunc n w*  
**by** (*induct n*) *auto*

**lemma** *bintr-cat1*:  
 !!b. *bintrunc (k + n) (bin-cat a n b) = bin-cat (bintrunc k a) n b*  
**by** (*induct n*) *auto*

**lemma** *bintr-cat*: *bintrunc m (bin-cat a n b) =*  
*bin-cat (bintrunc (m - n) a) n (bintrunc (min m n) b)*  
**by** (*rule bin-eqI*) (*auto simp: bin-nth-cat nth-bintr*)

**lemma** *bintr-cat-same [simp]*:  
*bintrunc n (bin-cat a n b) = bintrunc n b*  
**by** (*auto simp add : bintr-cat*)

**lemma** *cat-bintr [simp]*:  
 !!b. *bin-cat a n (bintrunc n b) = bin-cat a n b*  
**by** (*induct n*) *auto*

**lemma** *split-bintrunc*:  
 !!b c. *bin-split n c = (a, b) ==> b = bintrunc n c*  
**by** (*induct n*) (*auto simp: Let-def split: ls-splits*)

**lemma** *bin-cat-split*:  
 !!v w. *bin-split n w = (u, v) ==> w = bin-cat u n v*  
**by** (*induct n*) (*auto simp: Let-def split: ls-splits*)

**lemma** *bin-split-cat*:  
 !!w. *bin-split n (bin-cat v n w) = (v, bintrunc n w)*  
**by** (*induct n*) *auto*

**lemma** *bin-split-Pls [simp]*:  
*bin-split n Numeral.Pls = (Numeral.Pls, Numeral.Pls)*  
**by** (*induct n*) (*auto simp: Let-def split: ls-splits*)

**lemma** *bin-split-Min [simp]*:  
*bin-split n Numeral.Min = (Numeral.Min, bintrunc n Numeral.Min)*  
**by** (*induct n*) (*auto simp: Let-def split: ls-splits*)

**lemma** *bin-split-trunc*:  
 !!m b c. *bin-split (min m n) c = (a, b) ==>*  
*bin-split n (bintrunc m c) = (bintrunc (m - n) a, b)*  
**apply** (*induct n, clarsimp*)  
**apply** (*simp add: bin-rest-trunc Let-def split: ls-splits*)  
**apply** (*case-tac m*)

```

apply (auto simp: Let-def split: ls-splits)
done

lemma bin-split-trunc1:
  !!m b c. bin-split n c = (a, b) ==>
    bin-split n (bintrunc m c) = (bintrunc (m - n) a, bintrunc m b)
apply (induct n, clarsimp)
apply (simp add: bin-rest-trunc Let-def split: ls-splits)
apply (case-tac m)
apply (auto simp: Let-def split: ls-splits)
done

lemma bin-cat-num:
  !!b. bin-cat a n b = a * 2 ^ n + bintrunc n b
apply (induct n, clarsimp)
apply (simp add: Bit-def cong: number-of-False-cong)
done

lemma bin-split-num:
  !!b. bin-split n b = (b div 2 ^ n, b mod 2 ^ n)
apply (induct n, clarsimp)
apply (simp add: bin-rest-div zdiv-zmult2-eq)
apply (case-tac b rule: bin-exhaust)
apply simp
apply (simp add: Bit-def zmod-zmult-zmult1 p1mod22k
  split: bit.split
  cong: number-of-False-cong)
done

```

## 6.5 Miscellaneous lemmas

```

lemma nth-2p-bin:
  !!m. bin-nth (2 ^ n) m = (m = n)
apply (induct n)
apply clarsimp
apply safe
apply (case-tac m)
apply (auto simp: trans [OF numeral-1-eq-1 [symmetric] number-of-eq])
apply (case-tac m)
apply (auto simp: Bit-B0-2t [symmetric])
done

```

```

lemma ex-eq-or:
  (EX m. n = Suc m & (m = k | P m)) = (n = Suc k | (EX m. n = Suc m & P
m))
by auto

```

end

## 7 BinBoolList: Bool lists and integers

**theory** *BinBoolList* **imports** *BinOperations* **begin**

### 7.1 Arithmetic in terms of bool lists

**consts**

*rbl-succ* :: *bool list* => *bool list*  
*rbl-pred* :: *bool list* => *bool list*  
*rbl-add* :: *bool list* => *bool list* => *bool list*  
*rbl-mult* :: *bool list* => *bool list* => *bool list*

**primrec**

*Nil* : *rbl-succ Nil* = *Nil*  
*Cons* : *rbl-succ* (*x* # *xs*) = (if *x* then *False* # *rbl-succ xs* else *True* # *xs*)

**primrec**

*Nil* : *rbl-pred Nil* = *Nil*  
*Cons* : *rbl-pred* (*x* # *xs*) = (if *x* then *False* # *xs* else *True* # *rbl-pred xs*)

**primrec**

*Nil* : *rbl-add Nil x* = *Nil*  
*Cons* : *rbl-add* (*y* # *ys*) *x* = (let *ws* = *rbl-add ys* (*tl x*) in  
(*y* ~ = *hd x*) # (if *hd x* & *y* then *rbl-succ ws* else *ws*))

**primrec**

*Nil* : *rbl-mult Nil x* = *Nil*  
*Cons* : *rbl-mult* (*y* # *ys*) *x* = (let *ws* = *False* # *rbl-mult ys x* in  
if *y* then *rbl-add ws x* else *ws*)

**lemma** *tl-take*: *tl* (*take n l*) = *take* (*n - 1*) (*tl l*)

**apply** (*cases n, clarsimp*)  
**apply** (*cases l, auto*)  
**done**

**lemma** *take-butlast* [*rule-format*] :

*ALL n. n < length l --> take n* (*butlast l*) = *take n l*  
**apply** (*induct l, clarsimp*)  
**apply** *clarsimp*  
**apply** (*case-tac n*)  
**apply** *auto*  
**done**

**lemma** *butlast-take* [*rule-format*] :

*ALL n. n <= length l --> butlast* (*take n l*) = *take* (*n - 1*) *l*

```

apply (induct l, clarsimp)
apply clarsimp
apply (case-tac n)
apply safe
apply simp-all
apply (case-tac nat)
apply auto
done

```

```

lemma butlast-drop [rule-format] :
  ALL n. butlast (drop n l) = drop n (butlast l)
apply (induct l)
apply clarsimp
apply clarsimp
apply safe
apply ((case-tac n, auto)[1])+
done

```

```

lemma butlast-power:
  (butlast ^ n) bl = take (length bl - n) bl
by (induct n) (auto simp: butlast-take)

```

```

lemma bin-to-bl-aux-Pls-minus-simp:
  0 < n ==> bin-to-bl-aux n Numeral.Pls bl =
    bin-to-bl-aux (n - 1) Numeral.Pls (False # bl)
by (cases n) auto

```

```

lemma bin-to-bl-aux-Min-minus-simp:
  0 < n ==> bin-to-bl-aux n Numeral.Min bl =
    bin-to-bl-aux (n - 1) Numeral.Min (True # bl)
by (cases n) auto

```

```

lemma bin-to-bl-aux-Bit-minus-simp:
  0 < n ==> bin-to-bl-aux n (w BIT b) bl =
    bin-to-bl-aux (n - 1) w ((b = bit.B1) # bl)
by (cases n) auto

```

```

declare bin-to-bl-aux-Pls-minus-simp [simp]
  bin-to-bl-aux-Min-minus-simp [simp]
  bin-to-bl-aux-Bit-minus-simp [simp]

```

```

lemma bl-to-bin-aux-append [rule-format] :
  ALL w. bl-to-bin-aux w (bs @ cs) = bl-to-bin-aux (bl-to-bin-aux w bs) cs
by (induct bs) auto

```

```

lemma bin-to-bl-aux-append [rule-format] :
  ALL w bs. bin-to-bl-aux n w bs @ cs = bin-to-bl-aux n w (bs @ cs)

```

**by** (*induct n*) *auto*

**lemma** *bl-to-bin-append*:

*bl-to-bin (bs @ cs) = bl-to-bin-aux (bl-to-bin bs) cs*

**unfolding** *bl-to-bin-def* **by** (*rule bl-to-bin-aux-append*)

**lemma** *bin-to-bl-aux-alt*:

*bin-to-bl-aux n w bs = bin-to-bl n w @ bs*

**unfolding** *bin-to-bl-def* **by** (*simp add : bin-to-bl-aux-append*)

**lemma** *bin-to-bl-0*: *bin-to-bl 0 bs = []*

**unfolding** *bin-to-bl-def* **by** *auto*

**lemma** *size-bin-to-bl-aux* [*rule-format*] :

*ALL w bs. size (bin-to-bl-aux n w bs) = n + length bs*

**by** (*induct n*) *auto*

**lemma** *size-bin-to-bl*: *size (bin-to-bl n w) = n*

**unfolding** *bin-to-bl-def* **by** (*simp add : size-bin-to-bl-aux*)

**lemma** *bin-bl-bin'* [*rule-format*] :

*ALL w bs. bl-to-bin (bin-to-bl-aux n w bs) =*

*bl-to-bin-aux (bintrunc n w) bs*

**by** (*induct n*) (*auto simp add : bl-to-bin-def*)

**lemma** *bin-bl-bin*: *bl-to-bin (bin-to-bl n w) = bintrunc n w*

**unfolding** *bin-to-bl-def bin-bl-bin'* **by** *auto*

**lemma** *bl-bin-bl'* [*rule-format*] :

*ALL w n. bin-to-bl (n + length bs) (bl-to-bin-aux w bs) =*

*bin-to-bl-aux n w bs*

**apply** (*induct bs*)

**apply** *auto*

**apply** (*simp-all only : add-Suc [symmetric]*)

**apply** (*auto simp add : bin-to-bl-def*)

**done**

**lemma** *bl-bin-bl*: *bin-to-bl (length bs) (bl-to-bin bs) = bs*

**unfolding** *bl-to-bin-def*

**apply** (*rule box-equals*)

**apply** (*rule bl-bin-bl'*)

**prefer** 2

**apply** (*rule bin-to-bl-aux.Z*)

**apply** *simp*

**done**

**declare**

*bin-to-bl-0* [*simp*]

*size-bin-to-bl* [*simp*]

*bin-bl-bin* [simp]  
*bl-bin-bl* [simp]

**lemma** *bl-to-bin-inj*:

*bl-to-bin bs = bl-to-bin cs ==> length bs = length cs ==> bs = cs*

**apply** (*rule-tac box-equals*)

**defer**

**apply** (*rule bl-bin-bl*)

**apply** (*rule bl-bin-bl*)

**apply** *simp*

**done**

**lemma** *bl-to-bin-False*: *bl-to-bin (False # bl) = bl-to-bin bl*

**unfolding** *bl-to-bin-def* **by** *auto*

**lemma** *bl-to-bin-Nil*: *bl-to-bin [] = Numeral.Pls*

**unfolding** *bl-to-bin-def* **by** *auto*

**lemma** *bin-to-bl-Pls-aux* [rule-format] :

*ALL bl. bin-to-bl-aux n Numeral.Pls bl = replicate n False @ bl*

**by** (*induct n*) (*auto simp: replicate-app-Cons-same*)

**lemma** *bin-to-bl-Pls*: *bin-to-bl n Numeral.Pls = replicate n False*

**unfolding** *bin-to-bl-def* **by** (*simp add : bin-to-bl-Pls-aux*)

**lemma** *bin-to-bl-Min-aux* [rule-format] :

*ALL bl. bin-to-bl-aux n Numeral.Min bl = replicate n True @ bl*

**by** (*induct n*) (*auto simp: replicate-app-Cons-same*)

**lemma** *bin-to-bl-Min*: *bin-to-bl n Numeral.Min = replicate n True*

**unfolding** *bin-to-bl-def* **by** (*simp add : bin-to-bl-Min-aux*)

**lemma** *bl-to-bin-rep-F*:

*bl-to-bin (replicate n False @ bl) = bl-to-bin bl*

**apply** (*simp add: bin-to-bl-Pls-aux [symmetric] bin-bl-bin'*)

**apply** (*simp add: bl-to-bin-def*)

**done**

**lemma** *bin-to-bl-trunc*:

*n <= m ==> bin-to-bl n (bintrunc m w) = bin-to-bl n w*

**by** (*auto intro: bl-to-bin-inj*)

**declare**

*bin-to-bl-trunc* [simp]

*bl-to-bin-False* [simp]

*bl-to-bin-Nil* [simp]

**lemma** *bin-to-bl-aux-bintr* [rule-format] :

*ALL m bin bl. bin-to-bl-aux n (bintrunc m bin) bl =*

```

    replicate (n - m) False @ bin-to-bl-aux (min n m) bin bl
  apply (induct-tac n)
  apply clarsimp
  apply clarsimp
  apply (case-tac m)
  apply (clarsimp simp: bin-to-bl-Pls-aux)
  apply (erule thin-rl)
  apply (induct-tac n)
  apply auto
done

lemmas bin-to-bl-bintr =
  bin-to-bl-aux-bintr [where bl = [], folded bin-to-bl-def]

lemma bl-to-bin-rep-False: bl-to-bin (replicate n False) = Numeral.Pls
  by (induct n) auto

lemma len-bin-to-bl-aux [rule-format] :
  ALL w bs. length (bin-to-bl-aux n w bs) = n + length bs
  by (induct n) auto

lemma len-bin-to-bl [simp]: length (bin-to-bl n w) = n
  unfolding bin-to-bl-def len-bin-to-bl-aux by auto

lemma sign-bl-bin' [rule-format] :
  ALL w. bin-sign (bl-to-bin-aux w bs) = bin-sign w
  by (induct bs) auto

lemma sign-bl-bin: bin-sign (bl-to-bin bs) = Numeral.Pls
  unfolding bl-to-bin-def by (simp add: sign-bl-bin')

lemma bl-sbin-sign-aux [rule-format] :
  ALL w bs. hd (bin-to-bl-aux (Suc n) w bs) =
    (bin-sign (sbintrunc n w) = Numeral.Min)
  apply (induct n)
  apply clarsimp
  apply (case-tac w rule: bin-exhaust)
  apply (simp split add: bit.split)
  apply clarsimp
done

lemma bl-sbin-sign:
  hd (bin-to-bl (Suc n) w) = (bin-sign (sbintrunc n w) = Numeral.Min)
  unfolding bin-to-bl-def by (rule bl-sbin-sign-aux)

lemma bin-nth-of-bl-aux [rule-format] :
  ALL w. bin-nth (bl-to-bin-aux w bl) n =
    (n < size bl & rev bl ! n | n >= length bl & bin-nth w (n - size bl))
  apply (induct-tac bl)

```



```

  apply clarsimp
  apply clarsimp
  apply (cut-tac x=n and y=size list in linorder-less-linear)
  apply (erule disjE, simp add: nth-append)+
  apply (simp add: nth-append)
done

```

**lemma** *bin-nth-of-bl*:  $\text{bin-nth } (\text{bl-to-bin } \text{bl}) \ n = (n < \text{length } \text{bl} \ \& \ \text{rev } \text{bl} \ ! \ n)$   
**unfolding** *bl-to-bin-def* **by** (*simp add : bin-nth-of-bl-aux*)

**lemma** *bin-nth-bl* [rule-format] :  $\text{ALL } m \ w. \ n < m \ \longrightarrow$   
 $\text{bin-nth } w \ n = \text{nth } (\text{rev } (\text{bin-to-bl } m \ w)) \ n$   
**apply** (*induct n*)  
**apply** *clarsimp*  
**apply** (*case-tac m, clarsimp*)  
**apply** (*clarsimp simp: bin-to-bl-def*)  
**apply** (*simp add: bin-to-bl-aux-alt*)  
**apply** *clarsimp*  
**apply** (*case-tac m, clarsimp*)  
**apply** (*clarsimp simp: bin-to-bl-def*)  
**apply** (*simp add: bin-to-bl-aux-alt*)  
**done**

**lemma** *nth-rev* [rule-format] :  
 $n < \text{length } \text{xs} \ \longrightarrow \ \text{rev } \text{xs} \ ! \ n = \text{xs} \ ! \ (\text{length } \text{xs} - 1 - n)$   
**apply** (*induct-tac xs*)  
**apply** *simp*  
**apply** (*clarsimp simp add : nth-append nth.simps split add : nat.split*)  
**apply** (*rule-tac f = %n. list ! n in arg-cong*)  
**apply** *arith*  
**done**

**lemmas** *nth-rev-alt = nth-rev* [**where** *xs = rev ys, simplified, standard*]

**lemma** *nth-bin-to-bl-aux* [rule-format] :  
 $\text{ALL } w \ n \ \text{bl}. \ n < m + \text{length } \text{bl} \ \longrightarrow \ (\text{bin-to-bl-aux } m \ w \ \text{bl}) \ ! \ n =$   
 $(\text{if } n < m \ \text{then } \text{bin-nth } w \ (m - 1 - n) \ \text{else } \text{bl} \ ! \ (n - m))$   
**apply** (*induct-tac m*)  
**apply** *clarsimp*  
**apply** *clarsimp*  
**apply** (*case-tac w rule: bin-exhaust*)  
**apply** *clarsimp*  
**apply** (*case-tac na - n*)  
**apply** *arith*  
**apply** *simp*  
**apply** (*rule-tac f = %n. bl ! n in arg-cong*)  
**apply** *arith*  
**done**

**lemma** *nth-bin-to-bl*:  $n < m \implies (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$   
**unfolding** *bin-to-bl-def* **by** (*simp add : nth-bin-to-bl-aux*)

**lemma** *bl-to-bin-lt2p-aux* [*rule-format*] :  
*ALL w. bl-to-bin-aux w bs < (w + 1) \* (2 ^ length bs)*  
**apply** (*induct bs*)  
**apply** *clarsimp*  
**apply** *clarsimp*  
**apply** *safe*  
**apply** (*erule allE, erule xtr8 [rotated]*),  
*simp add: Bit-def ring-simps cong add : number-of-False-cong*)+  
**done**

**lemma** *bl-to-bin-lt2p*: *bl-to-bin bs < (2 ^ length bs)*  
**apply** (*unfold bl-to-bin-def*)  
**apply** (*rule xtr1*)  
**prefer** 2  
**apply** (*rule bl-to-bin-lt2p-aux*)  
**apply** *simp*  
**done**

**lemma** *bl-to-bin-ge2p-aux* [*rule-format*] :  
*ALL w. bl-to-bin-aux w bs >= w \* (2 ^ length bs)*  
**apply** (*induct bs*)  
**apply** *clarsimp*  
**apply** *clarsimp*  
**apply** *safe*  
**apply** (*erule allE, erule less-eq-less.order-trans [rotated]*),  
*simp add: Bit-def ring-simps cong add : number-of-False-cong*)+  
**done**

**lemma** *bl-to-bin-ge0*: *bl-to-bin bs >= 0*  
**apply** (*unfold bl-to-bin-def*)  
**apply** (*rule xtr4*)  
**apply** (*rule bl-to-bin-ge2p-aux*)  
**apply** *simp*  
**done**

**lemma** *butlast-rest-bin*:  
*butlast (bin-to-bl n w) = bin-to-bl (n - 1) (bin-rest w)*  
**apply** (*unfold bin-to-bl-def*)  
**apply** (*cases w rule: bin-exhaust*)  
**apply** (*cases n, clarsimp*)  
**apply** *clarsimp*  
**apply** (*auto simp add: bin-to-bl-aux-alt*)  
**done**

**lemmas** *butlast-bin-rest = butlast-rest-bin*  
[*where w=bl-to-bin bl and n=length bl, simplified, standard*]

**lemma** *butlast-rest-bl2bin-aux* [rule-format] :

*ALL w. bl*  $\sim = [] \dashrightarrow$

*bl-to-bin-aux w (butlast bl) = bin-rest (bl-to-bin-aux w bl)*

**by** (induct bl) auto

**lemma** *butlast-rest-bl2bin*:

*bl-to-bin (butlast bl) = bin-rest (bl-to-bin bl)*

**apply** (unfold bl-to-bin-def)

**apply** (cases bl)

**apply** (auto simp add: butlast-rest-bl2bin-aux)

**done**

**lemma** *trunc-bl2bin-aux* [rule-format] :

*ALL w. bintrunc m (bl-to-bin-aux w bl) =*

*bl-to-bin-aux (bintrunc (m - length bl) w) (drop (length bl - m) bl)*

**apply** (induct-tac bl)

**apply** clarsimp

**apply** clarsimp

**apply** safe

**apply** (case-tac m - size list)

**apply** (simp add: diff-is-0-eq [THEN iffD1, THEN Suc-diff-le])

**apply** simp

**apply** (rule-tac f = %nat. bl-to-bin-aux (bintrunc nat w BIT bit.B1) list  
in arg-cong)

**apply** simp

**apply** (case-tac m - size list)

**apply** (simp add: diff-is-0-eq [THEN iffD1, THEN Suc-diff-le])

**apply** simp

**apply** (rule-tac f = %nat. bl-to-bin-aux (bintrunc nat w BIT bit.B0) list  
in arg-cong)

**apply** simp

**done**

**lemma** *trunc-bl2bin*:

*bintrunc m (bl-to-bin bl) = bl-to-bin (drop (length bl - m) bl)*

**unfolding** bl-to-bin-def **by** (simp add: trunc-bl2bin-aux)

**lemmas** *trunc-bl2bin-len* [simp] =

*trunc-bl2bin [of length bl bl, simplified, standard]*

**lemma** *bl2bin-drop*:

*bl-to-bin (drop k bl) = bintrunc (length bl - k) (bl-to-bin bl)*

**apply** (rule trans)

**prefer** 2

**apply** (rule trunc-bl2bin [symmetric])

**apply** (cases k <= length bl)

**apply** auto

**done**

```

lemma nth-rest-power-bin [rule-format] :
  ALL n. bin-nth ((bin-rest ^ k) w) n = bin-nth w (n + k)
apply (induct k, clarsimp)
apply clarsimp
apply (simp only: bin-nth.Suc [symmetric] add-Suc)
done

lemma take-rest-power-bin:
  m <= n ==> take m (bin-to-bl n w) = bin-to-bl m ((bin-rest ^ (n - m)) w)
apply (rule nth-equalityI)
apply simp
apply (clarsimp simp add: nth-bin-to-bl nth-rest-power-bin)
done

lemma hd-butlast: size xs > 1 ==> hd (butlast xs) = hd xs
by (cases xs) auto

lemma last-bin-last' [rule-format] :
  ALL w. size xs > 0 --> last xs = (bin-last (bl-to-bin-aux w xs) = bit.B1)
by (induct xs) auto

lemma last-bin-last:
  size xs > 0 ==> last xs = (bin-last (bl-to-bin xs) = bit.B1)
unfolding bl-to-bin-def by (erule last-bin-last')

lemma bin-last-last:
  bin-last w = (if last (bin-to-bl (Suc n) w) then bit.B1 else bit.B0)
apply (unfold bin-to-bl-def)
apply simp
apply (auto simp add: bin-to-bl-aux-alt)
done

lemma app2-Nil [simp]: app2 f [] ys = []
unfolding app2-def by auto

lemma app2-Cons [simp]:
  app2 f (x # xs) (y # ys) = f x y # app2 f xs ys
unfolding app2-def by auto

lemma bl-xor-aux-bin [rule-format] : ALL v w bs cs.
  app2 (%x y. x ~ = y) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v XOR w) (app2 (%x y. x ~ = y) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)

```

```

apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

```

```

lemma bl-or-aux-bin [rule-format] : ALL v w bs cs.
  app2 (op | ) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v OR w) (app2 (op | ) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

```

```

lemma bl-and-aux-bin [rule-format] : ALL v w bs cs.
  app2 (op & ) (bin-to-bl-aux n v bs) (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (v AND w) (app2 (op & ) bs cs)
apply (induct-tac n)
apply safe
apply simp
apply (case-tac v rule: bin-exhaust)
apply (case-tac w rule: bin-exhaust)
apply clarsimp
apply (case-tac b)
apply (case-tac ba, safe, simp-all)+
done

```

```

lemma bl-not-aux-bin [rule-format] :
  ALL w cs. map Not (bin-to-bl-aux n w cs) =
  bin-to-bl-aux n (NOT w) (map Not cs)
apply (induct n)
apply clarsimp
apply clarsimp
apply (case-tac w rule: bin-exhaust)
apply (case-tac b)
apply auto
done

```

```

lemmas bl-not-bin = bl-not-aux-bin
  [where cs = [], unfolded bin-to-bl-def [symmetric] map.simps]

```

```

lemmas bl-and-bin = bl-and-aux-bin [where bs=[] and cs=[],
  unfolded app2-Nil, folded bin-to-bl-def]

```

**lemmas** *bl-or-bin* = *bl-or-aux-bin* [**where** *bs*=[] **and** *cs*=[],  
*unfolded app2-Nil, folded bin-to-bl-def*]

**lemmas** *bl-xor-bin* = *bl-xor-aux-bin* [**where** *bs*=[] **and** *cs*=[],  
*unfolded app2-Nil, folded bin-to-bl-def*]

**lemma** *drop-bin2bl-aux* [*rule-format*] :  
*ALL m bin bs. drop m (bin-to-bl-aux n bin bs) =*  
*bin-to-bl-aux (n - m) bin (drop (m - n) bs)*  
**apply** (*induct n, clarsimp*)  
**apply** *clarsimp*  
**apply** (*case-tac bin rule: bin-exhaust*)  
**apply** (*case-tac m <= n, simp*)  
**apply** (*case-tac m - n, simp*)  
**apply** *simp*  
**apply** (*rule-tac f = %nat. drop nat bs in arg-cong*)  
**apply** *simp*  
**done**

**lemma** *drop-bin2bl*: *drop m (bin-to-bl n bin) = bin-to-bl (n - m) bin*  
**unfolding** *bin-to-bl-def* **by** (*simp add : drop-bin2bl-aux*)

**lemma** *take-bin2bl-lem1* [*rule-format*] :  
*ALL w bs. take m (bin-to-bl-aux m w bs) = bin-to-bl m w*  
**apply** (*induct m, clarsimp*)  
**apply** *clarsimp*  
**apply** (*simp add: bin-to-bl-aux-alt*)  
**apply** (*simp add: bin-to-bl-def*)  
**apply** (*simp add: bin-to-bl-aux-alt*)  
**done**

**lemma** *take-bin2bl-lem* [*rule-format*] :  
*ALL w bs. take m (bin-to-bl-aux (m + n) w bs) =*  
*take m (bin-to-bl (m + n) w)*  
**apply** (*induct n*)  
**apply** *clarify*  
**apply** (*simp-all (no-asm) add: bin-to-bl-def take-bin2bl-lem1*)  
**apply** *simp*  
**done**

**lemma** *bin-split-take* [*rule-format*] :  
*ALL b c. bin-split n c = (a, b) -->*  
*bin-to-bl m a = take m (bin-to-bl (m + n) c)*  
**apply** (*induct n*)  
**apply** *clarsimp*  
**apply** (*clarsimp simp: Let-def split: ls-splits*)  
**apply** (*simp add: bin-to-bl-def*)  
**apply** (*simp add: take-bin2bl-lem*)  
**done**

**lemma** *bin-split-take1*:

$k = m + n \implies \text{bin-split } n \ c = (a, b) \implies$   
 $\text{bin-to-bl } m \ a = \text{take } m \ (\text{bin-to-bl } k \ c)$   
**by** (*auto elim: bin-split-take*)

**lemma** *nth-takefill* [*rule-format*] :  $ALL \ m \ l. \ m < n \longrightarrow$

$\text{takefill fill } n \ l \ ! \ m = (\text{if } m < \text{length } l \text{ then } l \ ! \ m \text{ else } \text{fill})$   
**apply** (*induct n, clarsimp*)  
**apply** *clarsimp*  
**apply** (*case-tac m*)  
**apply** (*simp split: list.split*)  
**apply** *clarsimp*  
**apply** (*erule allE*)  
**apply** (*erule (1) impE*)  
**apply** (*simp split: list.split*)  
**done**

**lemma** *takefill-alt* [*rule-format*] :

$ALL \ l. \ \text{takefill fill } n \ l = \text{take } n \ l \ @ \ \text{replicate } (n - \text{length } l) \ \text{fill}$   
**by** (*induct n*) (*auto split: list.split*)

**lemma** *takefill-replicate* [*simp*]:

$\text{takefill fill } n \ (\text{replicate } m \ \text{fill}) = \text{replicate } n \ \text{fill}$   
**by** (*simp add : takefill-alt replicate-add [symmetric]*)

**lemma** *takefill-le'* [*rule-format*] :

$ALL \ l \ n. \ n = m + k \longrightarrow \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$   
**by** (*induct m*) (*auto split: list.split*)

**lemma** *length-takefill* [*simp*]:  $\text{length } (\text{takefill fill } n \ l) = n$

**by** (*simp add : takefill-alt*)

**lemma** *take-takefill'*:

$!!w \ n. \ n = k + m \implies \text{take } k \ (\text{takefill fill } n \ w) = \text{takefill fill } k \ w$   
**by** (*induct k*) (*auto split add : list.split*)

**lemma** *drop-takefill*:

$!!w. \ \text{drop } k \ (\text{takefill fill } (m + k) \ w) = \text{takefill fill } m \ (\text{drop } k \ w)$   
**by** (*induct k*) (*auto split add : list.split*)

**lemma** *takefill-le* [*simp*]:

$m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$   
**by** (*auto simp: le-iff-add takefill-le'*)

**lemma** *take-takefill* [*simp*]:

$m \leq n \implies \text{take } m \ (\text{takefill fill } n \ w) = \text{takefill fill } m \ w$   
**by** (*auto simp: le-iff-add take-takefill'*)

**lemma** *takefill-append*:

*takefill fill (m + length xs) (xs @ w) = xs @ (takefill fill m w)*  
**by** (*induct xs*) *auto*

**lemma** *takefill-same'*:

*l = length xs ==> takefill fill l xs = xs*  
**by** *clarify (induct xs, auto)*

**lemmas** *takefill-same [simp] = takefill-same' [OF refl]*

**lemma** *takefill-bintrunc*:

*takefill False n bl = rev (bin-to-bl n (bl-to-bin (rev bl)))*  
**apply** (*rule nth-equalityI*)  
**apply** *simp*  
**apply** (*clarsimp simp: nth-takefill nth-rev nth-bin-to-bl bin-nth-of-bl*)  
**done**

**lemma** *bl-bin-bl-rtf*:

*bin-to-bl n (bl-to-bin bl) = rev (takefill False n (rev bl))*  
**by** (*simp add : takefill-bintrunc*)

**lemmas** *bl-bin-bl-rep-drop =*

*bl-bin-bl-rtf [simplified takefill-alt,*  
*simplified, simplified rev-take, simplified]*

**lemma** *tf-rev*:

*n + k = m + length bl ==> takefill x m (rev (takefill y n bl)) =*  
*rev (takefill y m (rev (takefill x k (rev bl))))*  
**apply** (*rule nth-equalityI*)  
**apply** (*auto simp add: nth-takefill nth-rev*)  
**apply** (*rule-tac f = %n. bl ! n in arg-cong*)  
**apply** *arith*  
**done**

**lemma** *takefill-minus*:

*0 < n ==> takefill fill (Suc (n - 1)) w = takefill fill n w*  
**by** *auto*

**lemmas** *takefill-Suc-cases =*

*list.cases [THEN takefill.Suc [THEN trans], standard]*

**lemmas** *takefill-Suc-Nil = takefill-Suc-cases (1)*

**lemmas** *takefill-Suc-Cons = takefill-Suc-cases (2)*

**lemmas** *takefill-minus-simps = takefill-Suc-cases [THEN [2]*

*takefill-minus [symmetric, THEN trans], standard]*

**lemmas** *takefill-pred-simps [simp] =*

*takefill-minus-simps [where n=number-of bin, simplified nobm1, standard]*



```

lemma bl-to-bin-aux-cat:
  !!nv v. bl-to-bin-aux (bin-cat w nv v) bs =
    bin-cat w (nv + length bs) (bl-to-bin-aux v bs)
apply (induct bs)
apply simp
apply (simp add: bin-cat-Suc-Bit [symmetric] del: bin-cat.simps)
done

```

```

lemma bin-to-bl-aux-cat:
  !!w bs. bin-to-bl-aux (nv + nw) (bin-cat v nw w) bs =
    bin-to-bl-aux nv v (bin-to-bl-aux nw w bs)
by (induct nw) auto

```

```

lemmas bl-to-bin-aux-alt =
  bl-to-bin-aux-cat [where nv = 0 and v = Numeral.Pls,
    simplified bl-to-bin-def [symmetric], simplified]

```

```

lemmas bin-to-bl-cat =
  bin-to-bl-aux-cat [where bs = [], folded bin-to-bl-def]

```

```

lemmas bl-to-bin-aux-app-cat =
  trans [OF bl-to-bin-aux-append bl-to-bin-aux-alt]

```

```

lemmas bin-to-bl-aux-cat-app =
  trans [OF bin-to-bl-aux-cat bin-to-bl-aux-alt]

```

```

lemmas bl-to-bin-app-cat = bl-to-bin-aux-app-cat
  [where w = Numeral.Pls, folded bl-to-bin-def]

```

```

lemmas bin-to-bl-cat-app = bin-to-bl-aux-cat-app
  [where bs = [], folded bin-to-bl-def]

```

```

lemma bl-to-bin-app-cat-alt:
  bin-cat (bl-to-bin cs) n w = bl-to-bin (cs @ bin-to-bl n w)
by (simp add : bl-to-bin-app-cat)

```

```

lemma mask-lem: (bl-to-bin (True # replicate n False)) =
  Numeral.succ (bl-to-bin (replicate n True))
apply (unfold bl-to-bin-def)
apply (induct n)
apply simp
apply (simp only: Suc-eq-add-numeral-1 replicate-add
  append-Cons [symmetric] bl-to-bin-aux-append)
apply simp
done

```

**lemma** *length-bl-of-nth* [*simp*]: *length (bl-of-nth n f) = n*  
**by** (*induct n*) *auto*

**lemma** *nth-bl-of-nth* [*simp*]:  
 $m < n \implies \text{rev } (\text{bl-of-nth } n \ f) \ ! \ m = f \ m$   
**apply** (*induct n*)  
**apply** *simp*  
**apply** (*clarsimp simp add : nth-append*)  
**apply** (*rule-tac f = f in arg-cong*)  
**apply** *simp*  
**done**

**lemma** *bl-of-nth-inj*:  
 $(!!k. \ k < n \implies f \ k = g \ k) \implies \text{bl-of-nth } n \ f = \text{bl-of-nth } n \ g$   
**by** (*induct n*) *auto*

**lemma** *bl-of-nth-nth-le* [*rule-format*] : *ALL xs.*  
 $\text{length } xs \geq n \dashv\vdash \text{bl-of-nth } n \ (\text{nth } (\text{rev } xs)) = \text{drop } (\text{length } xs - n) \ xs$   
**apply** (*induct n, clarsimp*)  
**apply** *clarsimp*  
**apply** (*rule trans [OF - hd-Cons-tl]*)  
**apply** (*frule Suc-le-lessD*)  
**apply** (*simp add: nth-rev trans [OF drop-Suc drop-tl, symmetric]*)  
**apply** (*subst hd-drop-conv-nth*)  
**apply** *force*  
**apply** *simp-all*  
**apply** (*rule-tac f = %n. drop n xs in arg-cong*)  
**apply** *simp*  
**done**

**lemmas** *bl-of-nth-nth* [*simp*] = *order-refl [THEN bl-of-nth-nth-le, simplified]*

**lemma** *size-rbl-pred*: *length (rbl-pred bl) = length bl*  
**by** (*induct bl*) *auto*

**lemma** *size-rbl-succ*: *length (rbl-succ bl) = length bl*  
**by** (*induct bl*) *auto*

**lemma** *size-rbl-add*:  
 $!!cl. \ \text{length } (\text{rbl-add } bl \ cl) = \text{length } bl$   
**by** (*induct bl*) (*auto simp: Let-def size-rbl-succ*)

**lemma** *size-rbl-mult*:  
 $!!cl. \ \text{length } (\text{rbl-mult } bl \ cl) = \text{length } bl$   
**by** (*induct bl*) (*auto simp add : Let-def size-rbl-add*)

**lemmas** *rbl-sizes* [*simp*] =

*size-rbl-pred size-rbl-succ size-rbl-add size-rbl-mult*

**lemmas** *rbl-Nils* =  
*rbl-pred.Nil rbl-succ.Nil rbl-add.Nil rbl-mult.Nil*

**lemma** *rbl-pred*:  
 !!*bin*. *rbl-pred* (*rev* (*bin-to-bl* *n bin*)) = *rev* (*bin-to-bl* *n* (*Numeral.pred bin*))  
**apply** (*induct n, simp*)  
**apply** (*unfold bin-to-bl-def*)  
**apply** *clarsimp*  
**apply** (*case-tac bin rule: bin-exhaust*)  
**apply** (*case-tac b*)  
**apply** (*clarsimp simp: bin-to-bl-aux-alt*) +  
**done**

**lemma** *rbl-succ*:  
 !!*bin*. *rbl-succ* (*rev* (*bin-to-bl* *n bin*)) = *rev* (*bin-to-bl* *n* (*Numeral.succ bin*))  
**apply** (*induct n, simp*)  
**apply** (*unfold bin-to-bl-def*)  
**apply** *clarsimp*  
**apply** (*case-tac bin rule: bin-exhaust*)  
**apply** (*case-tac b*)  
**apply** (*clarsimp simp: bin-to-bl-aux-alt*) +  
**done**

**lemma** *rbl-add*:  
 !!*bina binb*. *rbl-add* (*rev* (*bin-to-bl* *n bina*)) (*rev* (*bin-to-bl* *n binb*)) =  
*rev* (*bin-to-bl* *n* (*bina + binb*))  
**apply** (*induct n, simp*)  
**apply** (*unfold bin-to-bl-def*)  
**apply** *clarsimp*  
**apply** (*case-tac bina rule: bin-exhaust*)  
**apply** (*case-tac binb rule: bin-exhaust*)  
**apply** (*case-tac b*)  
**apply** (*case-tac [!] ba*)  
**apply** (*auto simp: rbl-succ succ-def bin-to-bl-aux-alt Let-def add-ac*)  
**done**

**lemma** *rbl-add-app2*:  
 !!*blb*. *length blb* >= *length bla* ==>  
*rbl-add bla (blb @ blc) = rbl-add bla blb*  
**apply** (*induct bla, simp*)  
**apply** *clarsimp*  
**apply** (*case-tac blb, clarsimp*)  
**apply** (*clarsimp simp: Let-def*)  
**done**

**lemma** *rbl-add-take2*:  
 !!*blb*. *length blb* >= *length bla* ==>

```

    rbl-add bla (take (length bla) bbb) = rbl-add bla bbb
  apply (induct bla, simp)
  apply clarsimp
  apply (case-tac bbb, clarsimp)
  apply (clarsimp simp: Let-def)
done

```

**lemma** *rbl-add-long*:

```

  m >= n ==> rbl-add (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =
    rev (bin-to-bl n (bina + binb))
  apply (rule box-equals [OF - rbl-add-take2 rbl-add])
  apply (rule-tac f = rbl-add (rev (bin-to-bl n bina)) in arg-cong)
  apply (rule rev-swap [THEN iffD1])
  apply (simp add: rev-take drop-bin2bl)
  apply simp
done

```

**lemma** *rbl-mult-app2*:

```

  !!bbb. length bbb >= length bla ==>
    rbl-mult bla (bbb @ bbb) = rbl-mult bla bbb
  apply (induct bla, simp)
  apply clarsimp
  apply (case-tac bbb, clarsimp)
  apply (clarsimp simp: Let-def rbl-add-app2)
done

```

**lemma** *rbl-mult-take2*:

```

  length bbb >= length bla ==>
    rbl-mult bla (take (length bla) bbb) = rbl-mult bla bbb
  apply (rule trans)
  apply (rule rbl-mult-app2 [symmetric])
  apply simp
  apply (rule-tac f = rbl-mult bla in arg-cong)
  apply (rule append-take-drop-id)
done

```

**lemma** *rbl-mult-gt1*:

```

  m >= length bl ==> rbl-mult bl (rev (bin-to-bl m binb)) =
    rbl-mult bl (rev (bin-to-bl (length bl) binb))
  apply (rule trans)
  apply (rule rbl-mult-take2 [symmetric])
  apply simp-all
  apply (rule-tac f = rbl-mult bl in arg-cong)
  apply (rule rev-swap [THEN iffD1])
  apply (simp add: rev-take drop-bin2bl)
done

```

**lemma** *rbl-mult-gt*:

```

  m > n ==> rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl m binb)) =

```

```

    rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb))
  by (auto intro: trans [OF rbl-mult-gt1])

```

**lemmas** *rbl-mult-Suc* = *lessI* [THEN *rbl-mult-gt*]

**lemma** *rbbl-Cons*:

```

  b # rev (bin-to-bl n x) = rev (bin-to-bl (Suc n) (x BIT If b bit.B1 bit.B0))
  apply (unfold bin-to-bl-def)
  apply simp
  apply (simp add: bin-to-bl-aux-alt)
  done

```

**lemma** *rbl-mult*: !!*bina binb*.

```

  rbl-mult (rev (bin-to-bl n bina)) (rev (bin-to-bl n binb)) =
  rev (bin-to-bl n (bina * binb))
  apply (induct n)
  apply simp
  apply (unfold bin-to-bl-def)
  apply clarsimp
  apply (case-tac bina rule: bin-exhaust)
  apply (case-tac binb rule: bin-exhaust)
  apply (case-tac b)
  apply (case-tac [!] ba)
  apply (auto simp: bin-to-bl-aux-alt Let-def)
  apply (auto simp: rbbl-Cons rbl-mult-Suc rbl-add)
  done

```

**lemma** *rbl-add-split*:

```

  P (rbl-add (y # ys) (x # xs)) =
  (ALL ws. length ws = length ys --> ws = rbl-add ys xs -->
  (y --> ((x --> P (False # rbl-succ ws)) & (~ x --> P (True # ws))))
  &
  (~ y --> P (x # ws)))
  apply (auto simp add: Let-def)
  apply (case-tac [!] y)
  apply auto
  done

```

**lemma** *rbl-mult-split*:

```

  P (rbl-mult (y # ys) xs) =
  (ALL ws. length ws = Suc (length ys) --> ws = False # rbl-mult ys xs -->
  (y --> P (rbl-add ws xs)) & (~ y --> P ws))
  by (clarsimp simp add: Let-def)

```

**lemma** *and-len*: *xs = ys ==> xs = ys & length xs = length ys*

**by** *auto*

**lemma** *size-if*: *size (if p then xs else ys) = (if p then size xs else size ys)*

**by** *auto*

**lemma** *tl-if*:  $tl\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ tl\ xs\ else\ tl\ ys)$   
**by** *auto*

**lemma** *hd-if*:  $hd\ (if\ p\ then\ xs\ else\ ys) = (if\ p\ then\ hd\ xs\ else\ hd\ ys)$   
**by** *auto*

**lemma** *if-Not-x*:  $(if\ p\ then\ \sim x\ else\ x) = (p = (\sim x))$   
**by** *auto*

**lemma** *if-x-Not*:  $(if\ p\ then\ x\ else\ \sim x) = (p = x)$   
**by** *auto*

**lemma** *if-same-and*:  $(If\ p\ x\ y\ \&\ If\ p\ u\ v) = (if\ p\ then\ x\ \&\ u\ else\ y\ \&\ v)$   
**by** *auto*

**lemma** *if-same-eq*:  $(If\ p\ x\ y = (If\ p\ u\ v)) = (if\ p\ then\ x = (u)\ else\ y = (v))$   
**by** *auto*

**lemma** *if-same-eq-not*:  
 $(If\ p\ x\ y = (\sim If\ p\ u\ v)) = (if\ p\ then\ x = (\sim u)\ else\ y = (\sim v))$   
**by** *auto*

**lemma** *if-Cons*:  $(if\ p\ then\ x\ \# xs\ else\ y\ \# ys) = If\ p\ x\ y\ \# If\ p\ xs\ ys$   
**by** *auto*

**lemma** *if-single*:  
 $(if\ xc\ then\ [xab]\ else\ [an]) = [if\ xc\ then\ xab\ else\ an]$   
**by** *auto*

**lemma** *if-bool-simps*:  
 $If\ p\ True\ y = (p \mid y) \ \&\ If\ p\ False\ y = (\sim p \ \&\ y) \ \&$   
 $If\ p\ y\ True = (p \dashrightarrow y) \ \&\ If\ p\ y\ False = (p \ \&\ y)$   
**by** *auto*

**lemmas** *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

**lemmas** *seqr* = *eq-reflection* [**where**  $x = size\ w$ , *standard*]

**lemmas** *tl-Nil* = *tl.simps* (1)  
**lemmas** *tl-Cons* = *tl.simps* (2)

## 7.2 Repeated splitting or concatenation

**lemma** *sclem*:  
 $size\ (concat\ (map\ (bin-to-bl\ n)\ xs)) = length\ xs * n$   
**by** (*induct xs*) *auto*

```

lemma bin-cat-foldl-lem [rule-format] :
  ALL x. foldl (%u. bin-cat u n) x xs =
    bin-cat x (size xs * n) (foldl (%u. bin-cat u n) y xs)
apply (induct xs)
apply simp
apply clarify
apply (simp (no-asm))
apply (frule asm-rl)
apply (drule spec)
apply (erule trans)
apply (drule-tac x = bin-cat y n a in spec)
apply (simp add : bin-cat-assoc-sym min-def)
done

lemma bin-rcat-bl:
  (bin-rcat n wl) = bl-to-bin (concat (map (bin-to-bl n) wl))
apply (unfold bin-rcat-def)
apply (rule sym)
apply (induct wl)
apply (auto simp add : bl-to-bin-append)
apply (simp add : bl-to-bin-aux-alt sclem)
apply (simp add : bin-cat-foldl-lem [symmetric])
done

lemmas bin-rsplit-aux-simps = bin-rsplit-aux.simps bin-rsplittl-aux.simps
lemmas rsplit-aux-simps = bin-rsplit-aux-simps

lemmas th-if-simp1 = split-if [where P = op = l,
  THEN iffD1, THEN conjunct1, THEN mp, standard]
lemmas th-if-simp2 = split-if [where P = op = l,
  THEN iffD1, THEN conjunct2, THEN mp, standard]

lemmas rsplit-aux-simp1s = rsplit-aux-simps [THEN th-if-simp1]

lemmas rsplit-aux-simp2ls = rsplit-aux-simps [THEN th-if-simp2]

lemmas bin-rsplit-aux-simp2s [simp] = rsplit-aux-simp2ls [unfolded Let-def]
lemmas rbscl = bin-rsplit-aux-simp2s (2)

lemmas rsplit-aux-0-simps [simp] =
  rsplit-aux-simp1s [OF disjI1] rsplit-aux-simp1s [OF disjI2]

lemma bin-rsplit-aux-append:
  bin-rsplit-aux (n, bs @ cs, m, c) = bin-rsplit-aux (n, bs, m, c) @ cs
apply (rule-tac u=n and v=bs and w=m and x=c in bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (subst bin-rsplit-aux.simps)
apply (clarsimp split: ls-splits)

```

**done**

**lemma** *bin-rsplittl-aux-append*:

*bin-rsplittl-aux* (*n*, *bs* @ *cs*, *m*, *c*) = *bin-rsplittl-aux* (*n*, *bs*, *m*, *c*) @ *cs*  
**apply** (*rule-tac* *u=n* **and** *v=bs* **and** *w=m* **and** *x=c* **in** *bin-rsplittl-aux.induct*)  
**apply** (*subst* *bin-rsplittl-aux.simps*)  
**apply** (*subst* *bin-rsplittl-aux.simps*)  
**apply** (*clarsimp* *split: ls-splits*)  
**done**

**lemmas** *rsplit-aux-apps* [**where** *bs* = []] =  
*bin-rsplit-aux-append* *bin-rsplittl-aux-append*

**lemmas** *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplittl-def*

**lemmas** *rsplit-aux-alts* = *rsplit-aux-apps*  
[*unfolded append-Nil* *rsplit-def-auxs* [*symmetric*]]

**lemma** *bin-split-minus*:  $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) \ w = \text{bin-split } n \ w$   
**by** *auto*

**lemmas** *bin-split-minus-simp* =  
*bin-split.Suc* [*THEN* [2] *bin-split-minus* [*symmetric*, *THEN* *trans*], *standard*]

**lemma** *bin-split-pred-simp* [*simp*]:  
 $(0::\text{nat}) < \text{number-of bin} \implies$   
 $\text{bin-split } (\text{number-of bin}) \ w =$   
 $(\text{let } (w1, w2) = \text{bin-split } (\text{number-of } (\text{Numeral.pred bin})) \ (\text{bin-rest } w)$   
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$   
**by** (*simp* *only: nobm1 bin-split-minus-simp*)

**declare** *bin-split-pred-simp* [*simp*]

**lemma** *bin-rsplit-aux-simp-alt*:

*bin-rsplit-aux* (*n*, *bs*, *m*, *c*) =  
(*if*  $m = 0 \vee n = 0$   
*then* *bs*  
*else* *let* (*a*, *b*) = *bin-split* *n* *c* *in* *bin-rsplit* *n* (*m* - *n*, *a*) @ *b* # *bs*)  
**apply** (*rule* *trans*)  
**apply** (*subst* *bin-rsplit-aux.simps*, *rule* *refl*)  
**apply** (*simp* *add: rsplit-aux-alts*)  
**done**

**lemmas** *bin-rsplit-simp-alt* =  
*trans* [*OF* *bin-rsplit-def* [*THEN* *meta-eq-to-obj-eq*]  
*bin-rsplit-aux-simp-alt*, *standard*]

**lemmas** *bthrs* = *bin-rsplit-simp-alt* [*THEN* [2] *trans*]



**lemma** *bin-rsplit-size-sign'* [rule-format] :  
 $n > 0 \implies (ALL\ nw\ w.\ rev\ sw = bin-rsplit\ n\ (nw, w) \dashv\dashv$   
 $(ALL\ v:\ set\ sw.\ bintrunc\ n\ v = v))$   
**apply** (induct sw)  
**apply** clarsimp  
**apply** clarsimp  
**apply** (drule bthrs)  
**apply** (simp (no-asm-use) add: Let-def split: ls-splits)  
**apply** clarify  
**apply** (erule impE, rule exI, erule exI)  
**apply** (drule split-bintrunc)  
**apply** simp  
**done**

**lemmas** *bin-rsplit-size-sign = bin-rsplit-size-sign'* [OF asm-rl  
rev-rev-ident [THEN trans] set-rev [THEN equalityD2 [THEN subsetD]],  
standard]

**lemma** *bin-nth-rsplit* [rule-format] :  
 $n > 0 \implies m < n \implies (ALL\ w\ k\ nw.\ rev\ sw = bin-rsplit\ n\ (nw, w) \dashv\dashv$   
 $k < size\ sw \dashv\dashv bin-nth\ (sw\ !\ k)\ m = bin-nth\ w\ (k * n + m))$   
**apply** (induct sw)  
**apply** clarsimp  
**apply** clarsimp  
**apply** (drule bthrs)  
**apply** (simp (no-asm-use) add: Let-def split: ls-splits)  
**apply** clarify  
**apply** (erule allE, erule impE, erule exI)  
**apply** (case-tac k)  
**apply** clarsimp  
**prefer** 2  
**apply** clarsimp  
**apply** (erule allE)  
**apply** (erule (1) impE)  
**apply** (drule bin-nth-split, erule conjE, erule allE,  
erule trans, simp add : add-ac)+  
**done**

**lemma** *bin-rsplit-all*:  
 $0 < nw \implies nw \leq n \implies bin-rsplit\ n\ (nw, w) = [bintrunc\ n\ w]$   
**unfolding** bin-rsplit-def  
**by** (clarsimp dest!: split-bintrunc simp: rsplit-aux-simp2ls split: ls-splits)

**lemma** *bin-rsplit-l* [rule-format] :  
 $ALL\ bin.\ bin-rsplitl\ n\ (m, bin) = bin-rsplit\ n\ (m, bintrunc\ m\ bin)$   
**apply** (rule-tac a = m in wf-less-than [THEN wf-induct])  
**apply** (simp (no-asm) add : bin-rsplitl-def bin-rsplit-def)  
**apply** (rule allI)  
**apply** (subst bin-rsplitl-aux.simps)

```

apply (subst bin-rsplit-aux.simps)
apply (clarsimp simp: rsplit-aux-alts Let-def split: ls-splits)
apply (drule bin-split-trunc)
apply (drule sym [THEN trans], assumption)
apply fast
done

```

```

lemma bin-rsplit-rcat [rule-format] :
   $n > 0 \longrightarrow \text{bin-rsplit } n \ (n * \text{size } ws, \text{bin-rcat } n \ ws) = \text{map } (\text{bintrunc } n) \ ws$ 
apply (unfold bin-rsplit-def bin-rcat-def)
apply (rule-tac  $xs = ws$  in rev-induct)
apply clarsimp
apply clarsimp
apply (clarsimp simp add: bin-split-cat rsplit-aux-alts)
done

```

```

lemma bin-rsplit-aux-len-le [rule-format] :
   $\text{ALL } ws \ m. \ n \neq 0 \longrightarrow ws = \text{bin-rsplit-aux } (n, bs, nw, w) \longrightarrow$ 
   $(\text{length } ws \leq m) = (nw + \text{length } bs * n \leq m * n)$ 
apply (rule-tac  $u=n$  and  $v=bs$  and  $w=nw$  and  $x=w$  in bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (simp add: lrlem Let-def split: ls-splits)
done

```

```

lemma bin-rsplit-len-le:
   $n \neq 0 \longrightarrow ws = \text{bin-rsplit } n \ (nw, w) \longrightarrow (\text{length } ws \leq m) = (nw \leq m * n)$ 
unfolding bin-rsplit-def by (clarsimp simp add : bin-rsplit-aux-len-le)

```

```

lemma bin-rsplit-aux-len [rule-format] :
   $n \neq 0 \longrightarrow \text{length } (\text{bin-rsplit-aux } (n, cs, nw, w)) =$ 
   $(nw + n - 1) \text{ div } n + \text{length } cs$ 
apply (rule-tac  $u=n$  and  $v=cs$  and  $w=nw$  and  $x=w$  in bin-rsplit-aux.induct)
apply (subst bin-rsplit-aux.simps)
apply (clarsimp simp: Let-def split: ls-splits)
apply (erule thin-rl)
apply (case-tac  $m \leq n$ )
prefer 2
apply (simp add: div-add-self2 [symmetric])
apply (case-tac  $m$ , clarsimp)
apply (simp add: div-add-self2)
done

```

```

lemma bin-rsplit-len:
   $n \neq 0 \implies \text{length } (\text{bin-rsplit } n \ (nw, w)) = (nw + n - 1) \text{ div } n$ 
unfolding bin-rsplit-def by (clarsimp simp add : bin-rsplit-aux-len)

```

```

lemma bin-rsplit-aux-len-indep [rule-format] :
   $n \neq 0 \implies (\text{ALL } v \ bs. \text{length } bs = \text{length } cs \longrightarrow$ 

```

```

    length (bin-rsplit-aux (n, bs, nw, v)) =
    length (bin-rsplit-aux (n, cs, nw, w))
  apply (rule-tac u=n and v=cs and w=nw and x=w in bin-rsplit-aux.induct)
  apply clarsimp
  apply (simp (no-asm-simp) add: bin-rsplit-aux-simp-alt Let-def
    split: ls-splits)

  apply clarify
  apply (erule allE)+
  apply (erule impE)
  apply (fast elim!: sym)
  apply (simp (no-asm-use) add: rsplit-aux-alts)
  apply (erule impE)
  apply (rule-tac x=ba # bs in exI)
  apply auto
done

lemma bin-rsplit-len-indep:
  n≠0 ==> length (bin-rsplit n (nw, v)) = length (bin-rsplit n (nw, w))
  apply (unfold bin-rsplit-def)
  apply (erule bin-rsplit-aux-len-indep)
  apply (rule refl)
done

end

```

## 8 TdThs: Type Definition Theorems

theory *TdThs* imports *Main* begin

## 9 More lemmas about normal type definitions

```

lemma
  tdD1: type-definition Rep Abs A ==> ∀x. Rep x ∈ A and
  tdD2: type-definition Rep Abs A ==> ∀x. Abs (Rep x) = x and
  tdD3: type-definition Rep Abs A ==> ∀y. y ∈ A ⟶ Rep (Abs y) = y
  by (auto simp: type-definition-def)

```

```

lemma td-nat-int:
  type-definition int nat (Collect (op <= 0))
  unfolding type-definition-def by auto

```

```

context type-definition
begin

```

```

lemmas Rep' [iff] = Rep [simplified]

```

**declare** *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

**lemma** *Abs-eqD*:  $Abs\ x = Abs\ y \implies x \in A \implies y \in A \implies x = y$   
**by** (*simp add: Abs-inject*)

**lemma** *Abs-inverse'*:  
 $r : A \implies Abs\ r = a \implies Rep\ a = r$   
**by** (*safe elim!: Abs-inverse*)

**lemma** *Rep-comp-inverse*:  
 $Rep\ o\ f = g \implies Abs\ o\ g = f$   
**using** *Rep-inverse* **by** (*auto intro: ext*)

**lemma** *Rep-eqD* [*elim!*]:  $Rep\ x = Rep\ y \implies x = y$   
**by** *simp*

**lemma** *Rep-inverse'*:  $Rep\ a = r \implies Abs\ r = a$   
**by** (*safe intro!: Rep-inverse*)

**lemma** *comp-Abs-inverse*:  
 $f\ o\ Abs = g \implies g\ o\ Rep = f$   
**using** *Rep-inverse* **by** (*auto intro: ext*)

**lemma** *set-Rep*:  
 $A = range\ Rep$   
**proof** (*rule set-ext*)  
**fix**  $x$   
**show**  $(x \in A) = (x \in range\ Rep)$   
**by** (*auto dest: Abs-inverse [of x, symmetric]*)  
**qed**

**lemma** *set-Rep-Abs*:  $A = range\ (Rep\ o\ Abs)$   
**proof** (*rule set-ext*)  
**fix**  $x$   
**show**  $(x \in A) = (x \in range\ (Rep\ o\ Abs))$   
**by** (*auto dest: Abs-inverse [of x, symmetric]*)  
**qed**

**lemma** *Abs-inj-on*: *inj-on*  $Abs\ A$   
**unfolding** *inj-on-def*  
**by** (*auto dest: Abs-inject [THEN iffD1]*)

**lemma** *image*:  $Abs\ `A = UNIV$   
**by** (*auto intro!: image-eqI*)

**lemmas** *td-thm* = *type-definition-axioms*

**lemma** *fns1*:  
 $Rep\ o\ fa = fr\ o\ Rep \mid fa\ o\ Abs = Abs\ o\ fr \implies Abs\ o\ fr\ o\ Rep = fa$

**by** (*auto dest: Rep-comp-inverse elim: comp-Abs-inverse simp: o-assoc*)

**lemmas** *fns1a* = *disjI1* [*THEN fns1*]

**lemmas** *fns1b* = *disjI2* [*THEN fns1*]

**lemma** *fns4*:

*Rep o fa o Abs = fr ==>*

*Rep o fa = fr o Rep & fa o Abs = Abs o fr*

**by** (*auto intro!: ext*)

**end**

**interpretation** *nat-int*: *type-definition* [*int nat Collect (op <= 0)*]

**by** (*rule td-nat-int*)

— resetting to the default nat induct and cases rules

**declare** *Nat.induct* [*case-names 0 Suc, induct type*]

**declare** *Nat.exhaust* [*case-names 0 Suc, cases type*]

## 9.1 Extended form of type definition predicate

**lemma** *td-conds*:

*norm o norm = norm ==> (fr o norm = norm o fr) =*

*(norm o fr o norm = fr o norm & norm o fr o norm = norm o fr)*

**apply** *safe*

**apply** (*simp-all add: o-assoc [symmetric]*)

**apply** (*simp-all add: o-assoc*)

**done**

**lemma** *fn-comm-power*:

*fa o tr = tr o fr ==> fa ^ n o tr = tr o fr ^ n*

**apply** (*rule ext*)

**apply** (*induct n*)

**apply** (*auto dest: fun-cong*)

**done**

**lemmas** *fn-comm-power'* =

*ext [THEN fn-comm-power, THEN fun-cong, unfolded o-def, standard]*

**locale** *td-ext* = *type-definition* +

**fixes** *norm*

**assumes** *eq-norm*:  $\bigwedge x. \text{Rep } (\text{Abs } x) = \text{norm } x$

**begin**

**lemma** *Abs-norm* [*simp*]:

*Abs (norm x) = Abs x*

**using** *eq-norm* [*of x*] **by** (*auto elim: Rep-inverse'*)

```

lemma td-th:
   $g \circ \text{Abs} = f \implies f (\text{Rep } x) = g \ x$ 
  by (drule comp-Abs-inverse [symmetric]) simp

lemma eq-norm':  $\text{Rep } o \ \text{Abs} = \text{norm}$ 
  by (auto simp: eq-norm intro!: ext)

lemma norm-Rep [simp]:  $\text{norm } (\text{Rep } x) = \text{Rep } x$ 
  by (auto simp: eq-norm' intro: td-th)

lemmas td = td-thm

lemma set-iff-norm:  $w : A \longleftrightarrow w = \text{norm } w$ 
  by (auto simp: set-Rep-Abs eq-norm' eq-norm [symmetric])

lemma inverse-norm:
   $(\text{Abs } n = w) = (\text{Rep } w = \text{norm } n)$ 
  apply (rule iffI)
  apply (clarsimp simp add: eq-norm)
  apply (simp add: eq-norm' [symmetric])
  done

lemma norm-eq-iff:
   $(\text{norm } x = \text{norm } y) = (\text{Abs } x = \text{Abs } y)$ 
  by (simp add: eq-norm' [symmetric])

lemma norm-comps:
   $\text{Abs } o \ \text{norm} = \text{Abs}$ 
   $\text{norm } o \ \text{Rep} = \text{Rep}$ 
   $\text{norm } o \ \text{norm} = \text{norm}$ 
  by (auto simp: eq-norm' [symmetric] o-def)

lemmas norm-norm [simp] = norm-comps

lemma fns5:
   $\text{Rep } o \ \text{fa } o \ \text{Abs} = \text{fr} \implies$ 
   $\text{fr } o \ \text{norm} = \text{fr} \ \& \ \text{norm } o \ \text{fr} = \text{fr}$ 
  by (fold eq-norm') (auto intro!: ext)

lemma fns2:
   $\text{Abs } o \ \text{fr } o \ \text{Rep} = \text{fa} \implies$ 
   $(\text{norm } o \ \text{fr } o \ \text{norm} = \text{fr } o \ \text{norm}) = (\text{Rep } o \ \text{fa} = \text{fr } o \ \text{Rep})$ 
  apply (fold eq-norm')
  apply safe
  prefer 2
  apply (simp add: o-assoc)
  apply (rule ext)
  apply (drule-tac x=Rep x in fun-cong)

```

```

apply auto
done

lemma fns3:
   $Abs \circ fr \circ Rep = fa ==>$ 
   $(norm \circ fr \circ norm = norm \circ fr) = (fa \circ Abs = Abs \circ fr)$ 
apply (fold eq-norm^)
apply safe
prefer 2
apply (simp add: o-assoc [symmetric])
apply (rule ext)
apply (drule fun-cong)
apply simp
done

lemma fns:
   $fr \circ norm = norm \circ fr ==>$ 
   $(fa \circ Abs = Abs \circ fr) = (Rep \circ fa = fr \circ Rep)$ 
apply safe
apply (frule fns1b)
prefer 2
apply (frule fns1a)
apply (rule fns3 [THEN iffD1])
prefer 3
apply (rule fns2 [THEN iffD1])
apply (simp-all add: o-assoc [symmetric])
apply (simp-all add: o-assoc)
done

lemma range-norm:
   $range (Rep \circ Abs) = A$ 
by (simp add: set-Rep-Abs)

end

lemmas td-ext-def' =
  td-ext-def [unfolded type-definition-def td-ext-axioms-def]

end

```

## 10 WordDefinition: Definition of Word Type

```

theory WordDefinition imports Size BinBoolList TdThs begin

typedef (open word) 'a word'
  =  $\{(0::int) ..< 2^{len-of\ TYPE('a::len0)}\}$  by auto

```

```

instance word :: (len0) number ..
instance word :: (type) minus ..
instance word :: (type) plus ..
instance word :: (type) one ..
instance word :: (type) zero ..
instance word :: (type) times ..
instance word :: (type) Divides.div ..
instance word :: (type) power ..
instance word :: (type) ord ..
instance word :: (type) size ..
instance word :: (type) inverse ..
instance word :: (type) bit ..

```

## 10.1 Type conversions and casting

### constdefs

— representation of words using unsigned or signed bins, only difference in these is the type class

```

word-of-int :: int => 'a :: len0 word
word-of-int w == Abs-word (bintrunc (len-of TYPE ('a)) w)

```

— uint and sint cast a word to an integer, uint treats the word as unsigned, sint treats the most-significant-bit as a sign bit

```

uint :: 'a :: len0 word => int
uint w == Rep-word w
sint :: 'a :: len word => int
sint-uint: sint w == sbintrunc (len-of TYPE ('a) - 1) (uint w)
unat :: 'a :: len0 word => nat
unat w == nat (uint w)

```

— the sets of integers representing the words

```

uints :: nat => int set
uints n == range (bintrunc n)
sints :: nat => int set
sints n == range (sbintrunc (n - 1))
unats :: nat => nat set
unats n == {i. i < 2 ^ n}
norm-sint :: nat => int => int
norm-sint n w == (w + 2 ^ (n - 1)) mod 2 ^ n - 2 ^ (n - 1)

```

— cast a word to a different length

```

scast :: 'a :: len word => 'b :: len word
scast w == word-of-int (sint w)
ucast :: 'a :: len0 word => 'b :: len0 word
ucast w == word-of-int (uint w)

```

— whether a cast (or other) function is to a longer or shorter length

```

source-size :: ('a :: len0 word => 'b) => nat
source-size c == let arb = arbitrary ; x = c arb in size arb

```



```

target-size :: ('a => 'b :: len0 word) => nat
target-size c == size (c arbitrary)
is-up :: ('a :: len0 word => 'b :: len0 word) => bool
is-up c == source-size c <= target-size c
is-down :: ('a :: len0 word => 'b :: len0 word) => bool
is-down c == target-size c <= source-size c

```

**constdefs**

```

of-bl :: bool list => 'a :: len0 word
of-bl bl == word-of-int (bl-to-bin bl)
to-bl :: 'a :: len0 word => bool list
to-bl w ==
bin-to-bl (len-of TYPE ('a)) (uint w)

word-reverse :: 'a :: len0 word => 'a word
word-reverse w == of-bl (rev (to-bl w))

```

**defs (overloaded)**

```

word-size: size (w :: 'a :: len0 word) == len-of TYPE('a)
word-number-of-def: number-of w == word-of-int w

```

**constdefs**

```

word-int-case :: (int => 'b) => ('a :: len0 word) => 'b
word-int-case f w == f (uint w)

```

**syntax**

```

of-int :: int => 'a

```

**translations**

```

case x of of-int y => b == word-int-case (%y. b) x

```

**10.2 Arithmetic operations****defs (overloaded)**

```

word-1-wi: (1 :: ('a :: len0) word) == word-of-int 1
word-0-wi: (0 :: ('a :: len0) word) == word-of-int 0

word-le-def: a <= b == uint a <= uint b
word-less-def: x < y == x <= y & x ~ = (y :: 'a :: len0 word)

```

**constdefs**

```

word-succ :: 'a :: len0 word => 'a word
word-succ a == word-of-int (Numeral.succ (uint a))

word-pred :: 'a :: len0 word => 'a word
word-pred a == word-of-int (Numeral.pred (uint a))

udvd :: 'a::len word => 'a::len word => bool (infixl udvd 50)
a udvd b == EX n>=0. uint b = n * uint a

```

*word-sle* :: 'a :: len word => 'a word => bool ((-/ <=s -) [50, 51] 50)  
*a <=s b* == *sint a <= sint b*

*word-sless* :: 'a :: len word => 'a word => bool ((-/ <s -) [50, 51] 50)  
*(x <s y)* == *(x <=s y & x ~ = y)*

**consts**

*word-power* :: 'a :: len0 word => nat => 'a word

**primrec**

*word-power a 0* = 1  
*word-power a (Suc n)* = *a \* word-power a n*

**defs (overloaded)**

*word-pow*: *power* == *word-power*  
*word-add-def*: *a + b* == *word-of-int (uint a + uint b)*  
*word-sub-wi*: *a - b* == *word-of-int (uint a - uint b)*  
*word-minus-def*: *- a* == *word-of-int (- uint a)*  
*word-mult-def*: *a \* b* == *word-of-int (uint a \* uint b)*  
*word-div-def*: *a div b* == *word-of-int (uint a div uint b)*  
*word-mod-def*: *a mod b* == *word-of-int (uint a mod uint b)*

**10.3 Bit-wise operations****defs (overloaded)**

*word-and-def*:  
*(a::'a::len0 word) AND b* == *word-of-int (uint a AND uint b)*

*word-or-def*:  
*(a::'a::len0 word) OR b* == *word-of-int (uint a OR uint b)*

*word-xor-def*:  
*(a::'a::len0 word) XOR b* == *word-of-int (uint a XOR uint b)*

*word-not-def*:  
*NOT (a::'a::len0 word)* == *word-of-int (NOT (uint a))*

*word-test-bit-def*:  
*test-bit (a::'a::len0 word)* == *bin-nth (uint a)*

*word-set-bit-def*:  
*set-bit (a::'a::len0 word) n x* ==  
*word-of-int (bin-sc n (If x bit.B1 bit.B0) (uint a))*

*word-set-bits-def*:  
*(BITS n. f n)::'a::len0 word* == *of-bl (bl-of-nth (len-of TYPE ('a)) f)*

*word-lsb-def*:  
*lsb (a::'a::len0 word)* == *bin-last (uint a) = bit.B1*

*word-msb-def:*  
*msb* (*a*::*'a*::*len* *word*) == *bin-sign* (*sint* *a*) = *Numeral.Min*

**constdefs**

*setBit* :: *'a* :: *len0* *word* => *nat* => *'a* *word*  
*setBit* *w n* == *set-bit* *w n True*

*clearBit* :: *'a* :: *len0* *word* => *nat* => *'a* *word*  
*clearBit* *w n* == *set-bit* *w n False*

**10.4 Shift operations****constdefs**

*shiftr1* :: *'a* :: *len0* *word* => *'a* *word*  
*shiftr1* *w* == *word-of-int* (*uint* *w* *BIT* *bit.B0*)

— shift right as unsigned or as signed, ie logical or arithmetic

*shiftr1* :: *'a* :: *len0* *word* => *'a* *word*  
*shiftr1* *w* == *word-of-int* (*bin-rest* (*uint* *w*))

*sshiftr1* :: *'a* :: *len* *word* => *'a* *word*  
*sshiftr1* *w* == *word-of-int* (*bin-rest* (*sint* *w*))

*bshiftr1* :: *bool* => *'a* :: *len* *word* => *'a* *word*  
*bshiftr1* *b w* == *of-bl* (*b* # *butlast* (*to-bl* *w*))

*sshiftr* :: *'a* :: *len* *word* => *nat* => *'a* *word* (**infixl** >>> 55)  
*w* >>> *n* == (*sshiftr1* ^ *n*) *w*

*mask* :: *nat* => *'a*::*len* *word*  
*mask* *n* == (*1* << *n*) - *1*

*revcast* :: *'a* :: *len0* *word* => *'b* :: *len0* *word*  
*revcast* *w* == *of-bl* (*takefill* *False* (*len-of* *TYPE*('b)) (*to-bl* *w*))

*slice1* :: *nat* => *'a* :: *len0* *word* => *'b* :: *len0* *word*  
*slice1* *n w* == *of-bl* (*takefill* *False* *n* (*to-bl* *w*))

*slice* :: *nat* => *'a* :: *len0* *word* => *'b* :: *len0* *word*  
*slice* *n w* == *slice1* (*size* *w* - *n*) *w*

**defs (overloaded)**

*shiftr1-def*: (*w*::*'a*::*len0* *word*) << *n* == (*shiftr1* ^ *n*) *w*  
*shiftr1-def*: (*w*::*'a*::*len0* *word*) >> *n* == (*shiftr1* ^ *n*) *w*

**10.5 Rotation****constdefs**

```

rotater1 :: 'a list => 'a list
rotater1 ys ==
  case ys of [] => [] | x # xs => last ys # butlast ys

rotater :: nat => 'a list => 'a list
rotater n == rotater1 ^ n

word-rotr :: nat => 'a :: len0 word => 'a :: len0 word
word-rotr n w == of-bl (rotater n (to-bl w))

word-rotl :: nat => 'a :: len0 word => 'a :: len0 word
word-rotl n w == of-bl (rotate n (to-bl w))

word-roti :: int => 'a :: len0 word => 'a :: len0 word
word-roti i w == if i >= 0 then word-rotr (nat i) w
                  else word-rotl (nat (- i)) w

```

## 10.6 Split and cat operations

### constdefs

```

word-cat :: 'a :: len0 word => 'b :: len0 word => 'c :: len0 word
word-cat a b == word-of-int (bin-cat (uint a) (len-of TYPE ('b)) (uint b))

word-split :: 'a :: len0 word => ('b :: len0 word) * ('c :: len0 word)
word-split a ==
  case bin-split (len-of TYPE ('c)) (uint a) of
    (u, v) => (word-of-int u, word-of-int v)

word-rcat :: 'a :: len0 word list => 'b :: len0 word
word-rcat ws ==
  word-of-int (bin-rcat (len-of TYPE ('a)) (map uint ws))

word-rsplit :: 'a :: len0 word => 'b :: len word list
word-rsplit w ==
  map word-of-int (bin-rsplit (len-of TYPE ('b)) (len-of TYPE ('a), uint w))

```

### constdefs

— Largest representable machine integer.

```

max-word :: 'a::len word
max-word ≡ word-of-int (2^len-of TYPE('a) - 1)

```

### consts

```

of-bool :: bool ⇒ 'a::len word

```

### primrec

```

of-bool False = 0
of-bool True = 1

```

**lemmas** *of-nth-def* = *word-set-bits-def*

**lemmas** *word-size-gt-0* [iff] =

*xtr1* [OF *word-size* [THEN *meta-eq-to-obj-eq*] *len-gt-0*, *standard*]

**lemmas** *lens-gt-0* = *word-size-gt-0 len-gt-0*

**lemmas** *lens-not-0* [iff] = *lens-gt-0* [THEN *gr-implies-not0*, *standard*]

**lemma** *uints-num*: *uints* *n* = {*i*.  $0 \leq i \wedge i < 2^n$ }

**by** (*simp add: uints-def range-bintrunc*)

**lemma** *sints-num*: *sints* *n* = {*i*.  $-(2^{n-1}) \leq i \wedge i < 2^{n-1}$ }

**by** (*simp add: sints-def range-sbintrunc*)

**lemmas** *atLeastLessThan-alt* = *atLeastLessThan-def* [unfolded

*atLeast-def lessThan-def Collect-conj-eq* [symmetric]]

**lemma** *mod-in-reps*:  $m > 0 \implies y \bmod m : \{0::\text{int} \mid ..< m\}$

**unfolding** *atLeastLessThan-alt* **by** *auto*

**lemma**

*Rep-word-0:0* <= *Rep-word* *x* **and**

*Rep-word-lt*: *Rep-word* (*x*::*a*::*len0* *word*) <  $2^{\text{len-of TYPE('a)}}$

**by** (*auto simp: Rep-word* [simplified])

**lemma** *Rep-word-mod-same*:

*Rep-word* *x* mod  $2^{\text{len-of TYPE('a)}}$  = *Rep-word* (*x*::*a*::*len0* *word*)

**by** (*simp add: int-mod-eq Rep-word-lt Rep-word-0*)

**lemma** *td-ext-uint*:

*td-ext* (*uint* :: *a* *word* => *int*) *word-of-int* (*uints* (*len-of TYPE('a::len0)*))

( $\%w::\text{int}. w \bmod 2^{\text{len-of TYPE('a)}}$ )

**apply** (*unfold td-ext-def*)

**apply** (*simp add: uints-num uint-def word-of-int-def bintrunc-mod2p*)

**apply** (*simp add: Rep-word-mod-same Rep-word-0 Rep-word-lt*

*word.Rep-word-inverse word.Abs-word-inverse int-mod-lem*)

**done**

**lemmas** *int-word-uint* = *td-ext-uint* [THEN *td-ext.eq-norm*, *standard*]

**interpretation** *word-uint*:

*td-ext* [*uint*::*a*::*len0* *word* => *int*

*word-of-int*

*uints* (*len-of TYPE('a::len0)*)

$\lambda w. w \bmod 2^{\text{len-of TYPE('a::len0)}}$ ]

**by** (*rule td-ext-uint*)

**lemmas** *td-uint* = *word-uint.td-thm*

**lemmas** *td-ext-ubin* = *td-ext-uint*

[*simplified len-gt-0 no-bintr-alt1 [symmetric]*]

**interpretation** *word-ubin*:

*td-ext* [*uint*::'*a*::*len0* *word*  $\Rightarrow$  *int*  
*word-of-int*  
*uints* (*len-of TYPE*('a::*len0*))  
*bintrunc* (*len-of TYPE*('a::*len0*))]  
**by** (*rule td-ext-ubin*)

**lemma** *sint-sbintrunc'*:

*sint* (*word-of-int bin* :: '*a word*) =  
(*sbintrunc* (*len-of TYPE* ('a :: *len*) - 1) *bin*)  
**unfolding** *sint-uint*  
**by** (*auto simp: word-ubin.eq-norm sbintrunc-bintrunc-lt*)

**lemma** *uint-sint*:

*uint w* = *bintrunc* (*len-of TYPE*('a)) (*sint* (*w* :: '*a* :: *len word*))  
**unfolding** *sint-uint* **by** (*auto simp: bintrunc-sbintrunc-le*)

**lemma** *bintr-uint'*:

*n* >= *size w* ==> *bintrunc n* (*uint w*) = *uint w*  
**apply** (*unfold word-size*)  
**apply** (*subst word-ubin.norm-Rep [symmetric]*)  
**apply** (*simp only: bintrunc-bintrunc-min word-size min-def*)  
**apply** *simp*  
**done**

**lemma** *wi-bintr'*:

*wb* = *word-of-int bin* ==> *n* >= *size wb* ==>  
*word-of-int* (*bintrunc n bin*) = *wb*  
**unfolding** *word-size*  
**by** (*clarsimp simp add : word-ubin.norm-eq-iff [symmetric] min-def*)

**lemmas** *bintr-uint* = *bintr-uint'* [*unfolded word-size*]

**lemmas** *wi-bintr* = *wi-bintr'* [*unfolded word-size*]

**lemma** *td-ext-sbin*:

*td-ext* (*sint* :: '*a word* ==> *int*) *word-of-int* (*sints* (*len-of TYPE*('a::*len*)))  
(*sbintrunc* (*len-of TYPE*('a) - 1))  
**apply** (*unfold td-ext-def' sint-uint*)  
**apply** (*simp add : word-ubin.eq-norm*)  
**apply** (*cases len-of TYPE*('a))  
**apply** (*auto simp add : sints-def*)  
**apply** (*rule sym [THEN trans]*)  
**apply** (*rule word-ubin.Abs-norm*)  
**apply** (*simp only: bintrunc-sbintrunc*)  
**apply** (*drule sym*)  
**apply** *simp*  
**done**

**lemmas** *td-ext-sint* = *td-ext-sbin*  
 [simplified len-gt-0 no-sbintr-alt2 Suc-pred' [symmetric]]

**interpretation** *word-sint*:  
*td-ext* [*sint* :: 'a::len word => int  
   *word-of-int*  
   *sints* (len-of TYPE('a::len))  
   %w. (w + 2<sup>(len-of TYPE('a::len) - 1)</sup>) mod 2<sup>len-of TYPE('a::len) - 1</sup>]  
**by** (rule *td-ext-sint*)

**interpretation** *word-sbin*:  
*td-ext* [*sint* :: 'a::len word => int  
   *word-of-int*  
   *sints* (len-of TYPE('a::len))  
   *sbintrunc* (len-of TYPE('a::len) - 1)]  
**by** (rule *td-ext-sbin*)

**lemmas** *int-word-sint* = *td-ext-sint* [THEN *td-ext.eq-norm*, standard]

**lemmas** *td-sint* = *word-sint.td*

**lemma** *word-number-of-alt*: *number-of* b == *word-of-int* (*number-of* b)  
**unfolding** *word-number-of-def* **by** (*simp add: number-of-eq*)

**lemma** *word-no-wi*: *number-of* = *word-of-int*  
**by** (*auto simp: word-number-of-def intro: ext*)

**lemma** *to-bl-def*:  
 (*to-bl* :: 'a :: len0 word => bool list) =  
   *bin-to-bl* (len-of TYPE('a)) o *uint*  
**by** (*auto simp: to-bl-def intro: ext*)

**lemmas** *word-reverse-no-def* [*simp*] = *word-reverse-def* [*of number-of w*, standard]

**lemmas** *uints-mod* = *uints-def* [*unfolded no-bintr-alt1*]

**lemma** *uint-bintrunc*: *uint* (*number-of bin* :: 'a word) =  
   *number-of* (*bintrunc* (len-of TYPE('a :: len0)) *bin*)  
**unfolding** *word-number-of-def number-of-eq*  
**by** (*auto intro: word-ubin.eq-norm*)

**lemma** *sint-sbintrunc*: *sint* (*number-of bin* :: 'a word) =  
   *number-of* (*sbintrunc* (len-of TYPE('a :: len) - 1) *bin*)  
**unfolding** *word-number-of-def number-of-eq*  
**by** (*subst word-sbin.eq-norm*) *simp*

**lemma** *unat-bintrunc*:

*unat* (*number-of* *bin* :: 'a :: len0 word) =  
*number-of* (*bintrunc* (*len-of* *TYPE*('a)) *bin*)  
**unfolding** *unat-def* *nat-number-of-def*  
**by** (*simp* *only*: *uint-bintrunc*)

**declare**

*uint-bintrunc* [*simp*]  
*sint-sbintrunc* [*simp*]  
*unat-bintrunc* [*simp*]

**lemma** *size-0-eq*: *size* (*w* :: 'a :: len0 word) = 0 ==> *v* = *w*

**apply** (*unfold* *word-size*)  
**apply** (*rule* *word-uint.Rep-eqD*)  
**apply** (*rule* *box-equals*)  
**defer**  
**apply** (*rule* *word-ubin.norm-Rep*)+  
**apply** *simp*  
**done**

**lemmas** *uint-lem* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]

**lemmas** *sint-lem* = *word-sint.Rep* [*unfolded sints-num mem-Collect-eq*]

**lemmas** *uint-ge-0* [*iff*] = *uint-lem* [*THEN conjunct1, standard*]

**lemmas** *uint-lt2p* [*iff*] = *uint-lem* [*THEN conjunct2, standard*]

**lemmas** *sint-ge* = *sint-lem* [*THEN conjunct1, standard*]

**lemmas** *sint-lt* = *sint-lem* [*THEN conjunct2, standard*]

**lemma** *sign-uint-Pls* [*simp*]:

*bin-sign* (*uint* *x*) = *Numeral.Pl*s  
**by** (*simp* *add*: *sign-Pls-ge-0 number-of-eq*)

**lemmas** *uint-m2p-neg* = *iffD2* [*OF diff-less-0-iff-less uint-lt2p, standard*]

**lemmas** *uint-m2p-not-non-neg* =  
*iffD2* [*OF linorder-not-le uint-m2p-neg, standard*]

**lemma** *lt2p-lem*:

*len-of* *TYPE*('a) <= *n* ==> *uint* (*w* :: 'a :: len0 word) < 2 ^ *n*  
**by** (*rule* *xtr8* [*OF - uint-lt2p*]) *simp*

**lemmas** *uint-le-0-iff* [*simp*] =

*uint-ge-0* [*THEN leD, THEN linorder-antisym-conv1, standard*]

**lemma** *uint-nat*: *uint* *w* == *int* (*unat* *w*)

**unfolding** *unat-def* **by** *auto*

**lemma** *uint-number-of*:

*uint* (*number-of* *b* :: 'a :: len0 word) = *number-of* *b* mod 2 ^ *len-of* *TYPE*('a)  
**unfolding** *word-number-of-alt*



**by** (*simp only: int-word-uint*)

**lemma** *unat-number-of*:

*bin-sign b = Numeral.Pls ==>*

*unat (number-of b :: 'a :: len0 word) = number-of b mod 2 ^ len-of TYPE ('a)*

**apply** (*unfold unat-def*)

**apply** (*clarsimp simp only: uint-number-of*)

**apply** (*rule nat-mod-distrib [THEN trans]*)

**apply** (*erule sign-Pls-ge-0 [THEN iffD1]*)

**apply** (*simp-all add: nat-power-eq*)

**done**

**lemma** *sint-number-of*: *sint (number-of b :: 'a :: len word) = (number-of b + 2 ^ (len-of TYPE('a) - 1)) mod 2 ^ len-of TYPE('a) - 2 ^ (len-of TYPE('a) - 1)*

**unfolding** *word-number-of-alt* **by** (*rule int-word-sint*)

**lemma** *word-of-int-bin* [*simp*] :

*(word-of-int (number-of bin) :: 'a :: len0 word) = (number-of bin)*

**unfolding** *word-number-of-alt* **by** *auto*

**lemma** *word-int-case-wi*:

*word-int-case f (word-of-int i :: 'b word) =*

*f (i mod 2 ^ len-of TYPE('b::len0))*

**unfolding** *word-int-case-def* **by** (*simp add: word-uint.eq-norm*)

**lemma** *word-int-split*:

*P (word-int-case f x) =*

*(ALL i. x = (word-of-int i :: 'b :: len0 word) &*

*0 <= i & i < 2 ^ len-of TYPE('b) --> P (f i))*

**unfolding** *word-int-case-def*

**by** (*auto simp: word-uint.eq-norm int-mod-eq'*)

**lemma** *word-int-split-asm*:

*P (word-int-case f x) =*

*(~ (EX n. x = (word-of-int n :: 'b::len0 word) &*

*0 <= n & n < 2 ^ len-of TYPE('b::len0) & ~ P (f n)))*

**unfolding** *word-int-case-def*

**by** (*auto simp: word-uint.eq-norm int-mod-eq'*)

**lemmas** *uint-range'* =

*word-uint.Rep [unfolded uints-num mem-Collect-eq, standard]*

**lemmas** *sint-range'* = *word-sint.Rep [unfolded One-nat-def*

*sints-num mem-Collect-eq, standard]*

**lemma** *uint-range-size*: *0 <= uint w & uint w < 2 ^ size w*

**unfolding** *word-size* **by** (*rule uint-range'*)

**lemma** *sint-range-size*:

–  $(2 \wedge (\text{size } w - \text{Suc } 0)) \leq \text{sint } w \ \& \ \text{sint } w < 2 \wedge (\text{size } w - \text{Suc } 0)$   
**unfolding** *word-size* **by** (*rule sint-range'*)

**lemmas** *sint-above-size = sint-range-size*  
 [THEN *conjunct2*, THEN [2] *xtr8*, folded *One-nat-def*, standard]

**lemmas** *sint-below-size = sint-range-size*  
 [THEN *conjunct1*, THEN [2] *order-trans*, folded *One-nat-def*, standard]

**lemma** *test-bit-eq-iff*:  $(\text{test-bit } (u::'a::\text{len } 0 \text{ word}) = \text{test-bit } v) = (u = v)$   
**unfolding** *word-test-bit-def* **by** (*simp add: bin-nth-eq-iff*)

**lemma** *test-bit-size* [rule-format] :  $(w::'a::\text{len } 0 \text{ word}) !! n \longrightarrow n < \text{size } w$   
**apply** (*unfold word-test-bit-def*)  
**apply** (*subst word-ubin.norm-Rep [symmetric]*)  
**apply** (*simp only: nth-bintr word-size*)  
**apply** *fast*  
**done**

**lemma** *word-eqI* [rule-format] :  
**fixes**  $u :: 'a::\text{len } 0 \text{ word}$   
**shows**  $(\text{ALL } n. n < \text{size } u \longrightarrow u !! n = v !! n) \implies u = v$   
**apply** (*rule test-bit-eq-iff [THEN iffD1]*)  
**apply** (*rule ext*)  
**apply** (*erule allE*)  
**apply** (*erule impCE*)  
**prefer** 2  
**apply** *assumption*  
**apply** (*auto dest!: test-bit-size simp add: word-size*)  
**done**

**lemmas** *word-eqD = test-bit-eq-iff* [THEN *iffD2*, THEN *fun-cong*, standard]

**lemma** *test-bit-bin'*:  $w !! n = (n < \text{size } w \ \& \ \text{bin-nth } (\text{uint } w) \ n)$   
**unfolding** *word-test-bit-def word-size*  
**by** (*simp add: nth-bintr [symmetric]*)

**lemmas** *test-bit-bin = test-bit-bin'* [unfolded *word-size*]

**lemma** *bin-nth-uint-imp'*:  $\text{bin-nth } (\text{uint } w) \ n \longrightarrow n < \text{size } w$   
**apply** (*unfold word-size*)  
**apply** (*rule impI*)  
**apply** (*rule nth-bintr [THEN iffD1, THEN conjunct1]*)  
**apply** (*subst word-ubin.norm-Rep*)  
**apply** *assumption*  
**done**

**lemma** *bin-nth-sint'*:  
 $n \geq \text{size } w \longrightarrow \text{bin-nth } (\text{sint } w) \ n = \text{bin-nth } (\text{sint } w) \ (\text{size } w - 1)$

```

apply (rule impI)
apply (subst word-sbin.norm-Rep [symmetric])
apply (simp add : nth-sbintr word-size)
apply auto
done

```

```

lemmas bin-nth-uint-imp = bin-nth-uint-imp' [rule-format, unfolded word-size]
lemmas bin-nth-sint = bin-nth-sint' [rule-format, unfolded word-size]

```

```

lemma td-bl:
  type-definition (to-bl :: 'a::len0 word => bool list)
    of-bl
    {bl. length bl = len-of TYPE('a)}
apply (unfold type-definition-def of-bl-def to-bl-def)
apply (simp add: word-ubin.eq-norm)
apply safe
apply (drule sym)
apply simp
done

```

```

interpretation word-bl:
  type-definition [to-bl :: 'a::len0 word => bool list]
    of-bl
    {bl. length bl = len-of TYPE('a::len0)}
  by (rule td-bl)

```

```

lemma word-size-bl: size w == size (to-bl w)
  unfolding word-size by auto

```

```

lemma to-bl-use-of-bl:
  (to-bl w = bl) = (w = of-bl bl  $\wedge$  length bl = length (to-bl w))
  by (fastsimp elim!: word-bl.Abs-inverse [simplified])

```

```

lemma to-bl-word-rev: to-bl (word-reverse w) = rev (to-bl w)
  unfolding word-reverse-def by (simp add: word-bl.Abs-inverse)

```

```

lemma word-rev-rev [simp] : word-reverse (word-reverse w) = w
  unfolding word-reverse-def by (simp add : word-bl.Abs-inverse)

```

```

lemma word-rev-gal: word-reverse w = u ==> word-reverse u = w
  by auto

```

```

lemmas word-rev-gal' = sym [THEN word-rev-gal, symmetric, standard]

```

```

lemmas length-bl-gt-0 [iff] = xtr1 [OF word-bl.Rep' len-gt-0, standard]
lemmas bl-not-Nil [iff] =
  length-bl-gt-0 [THEN length-greater-0-conv [THEN iffD1], standard]
lemmas length-bl-neq-0 [iff] = length-bl-gt-0 [THEN gr-implies-not0]

```

```

lemma hd-bl-sign-sint: hd (to-bl w) = (bin-sign (sint w) = Numeral.Min)
  apply (unfold to-bl-def sint-uint)
  apply (rule trans [OF - bl-sbin-sign])
  apply simp
  done

```

```

lemma of-bl-drop':
  lend = length bl - len-of TYPE ('a :: len0) ==>
    of-bl (drop lend bl) = (of-bl bl :: 'a word)
  apply (unfold of-bl-def)
  apply (clarsimp simp add : trunc-bl2bin [symmetric])
  done

```

```

lemmas of-bl-no = of-bl-def [folded word-number-of-def]

```

```

lemma test-bit-of-bl:
  (of-bl bl :: 'a :: len0 word) !! n = (rev bl ! n ∧ n < len-of TYPE('a) ∧ n < length
bl)
  apply (unfold of-bl-def word-test-bit-def)
  apply (auto simp add : word-size word-ubin.eq-norm nth-bintr bin-nth-of-bl)
  done

```

```

lemma no-of-bl:
  (number-of bin :: 'a :: len0 word) = of-bl (bin-to-bl (len-of TYPE('a)) bin)
  unfolding word-size of-bl-no by (simp add : word-number-of-def)

```

```

lemma uint-bl: to-bl w == bin-to-bl (size w) (uint w)
  unfolding word-size to-bl-def by auto

```

```

lemma to-bl-bin: bl-to-bin (to-bl w) = uint w
  unfolding uint-bl by (simp add : word-size)

```

```

lemma to-bl-of-bin:
  to-bl (word-of-int bin :: 'a :: len0 word) = bin-to-bl (len-of TYPE('a)) bin
  unfolding uint-bl by (clarsimp simp add : word-ubin.eq-norm word-size)

```

```

lemmas to-bl-no-bin [simp] = to-bl-of-bin [folded word-number-of-def]

```

```

lemma to-bl-to-bin [simp] : bl-to-bin (to-bl w) = uint w
  unfolding uint-bl by (simp add : word-size)

```

```

lemmas uint-bl-bin [simp] = trans [OF bin-bl-bin word-ubin.norm-Rep, standard]

```

```

lemmas num-AB-u [simp] = word-uint.Rep-inverse
  [unfolded o-def word-number-of-def [symmetric], standard]

```

```

lemmas num-AB-s [simp] = word-sint.Rep-inverse
  [unfolded o-def word-number-of-def [symmetric], standard]

```

```

lemma uints-unats: uints n = int ‘ unats n
  apply (unfold unats-def uints-num)
  apply safe
  apply (rule-tac image-eqI)
  apply (erule-tac nat-0-le [symmetric])
  apply auto
  apply (erule-tac nat-less-iff [THEN iffD2])
  apply (rule-tac [2] zless-nat-eq-int-zless [THEN iffD1])
  apply (auto simp add : nat-power-eq int-power)
  done

```

```

lemma unats-uints: unats n = nat ‘ uints n
  by (auto simp add : uints-unats image-iff)

```

```

lemmas bintr-num = word-ubin.norm-eq-iff
  [symmetric, folded word-number-of-def, standard]
lemmas sbintr-num = word-sbin.norm-eq-iff
  [symmetric, folded word-number-of-def, standard]

```

```

lemmas num-of-bintr = word-ubin.Abs-norm [folded word-number-of-def, stan-
dard]
lemmas num-of-sbintr = word-sbin.Abs-norm [folded word-number-of-def, stan-
dard]

```

```

lemma num-of-bintr':
  bintrunc (len-of TYPE('a :: len0)) a = b ==>
    number-of a = (number-of b :: 'a word)
  apply safe
  apply (rule-tac num-of-bintr [symmetric])
  done

```

```

lemma num-of-sbintr':
  sbintrunc (len-of TYPE('a :: len) - 1) a = b ==>
    number-of a = (number-of b :: 'a word)
  apply safe
  apply (rule-tac num-of-sbintr [symmetric])
  done

```

```

lemmas num-abs-bintr = sym [THEN trans,
  OF num-of-bintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]
lemmas num-abs-sbintr = sym [THEN trans,
  OF num-of-sbintr word-number-of-def [THEN meta-eq-to-obj-eq], standard]

```

```

lemma ucast-id: ucast w = w

```

```

unfolding ucast-def by auto

lemma scast-id: scast w = w
unfolding scast-def by auto

lemma ucast-bl: ucast w == of-bl (to-bl w)
unfolding ucast-def of-bl-def uint-bl
by (auto simp add : word-size)

lemma nth-ucast:
  (ucast w :: 'a::len0 word) !! n = (w !! n & n < len-of TYPE('a))
apply (unfold ucast-def test-bit-bin)
apply (simp add: word-ubin.eq-norm nth-bintr word-size)
apply (fast elim!: bin-nth-uint-imp)
done

lemma ucast-bintr [simp]:
  ucast (number-of w :: 'a::len0 word) =
    number-of (bintrunc (len-of TYPE('a)) w)
unfolding ucast-def by simp

lemma scast-sbintr [simp]:
  scast (number-of w :: 'a::len word) =
    number-of (sbintrunc (len-of TYPE('a) - Suc 0) w)
unfolding scast-def by simp

lemmas source-size = source-size-def [unfolded Let-def word-size]
lemmas target-size = target-size-def [unfolded Let-def word-size]
lemmas is-down = is-down-def [unfolded source-size target-size]
lemmas is-up = is-up-def [unfolded source-size target-size]

lemmas is-up-down =
  trans [OF is-up [THEN meta-eq-to-obj-eq]
    is-down [THEN meta-eq-to-obj-eq, symmetric],
    standard]

lemma down-cast-same': uc = ucast ==> is-down uc ==> uc = scast
apply (unfold is-down)
apply safe
apply (rule ext)
apply (unfold ucast-def scast-def uint-sint)
apply (rule word-ubin.norm-eq-iff [THEN iffD1])
apply simp
done

lemma word-rev-tf':
  r = to-bl (of-bl bl) ==> r = rev (takefill False (length r) (rev bl))

```

```

unfolding of-bl-def uint-bl
by (clarsimp simp add: bl-bin-bl-rtf word-ubin.eq-norm word-size)

lemmas word-rev-tf = refl [THEN word-rev-tf', unfolded word-bl.Rep', standard]

lemmas word-rep-drop = word-rev-tf [simplified takefill-alt,
  simplified, simplified rev-take, simplified]

lemma to-bl-ucast:
  to-bl (ucast (w::'b::len0 word) :: 'a::len0 word) =
    replicate (len-of TYPE('a) - len-of TYPE('b)) False @
    drop (len-of TYPE('b) - len-of TYPE('a)) (to-bl w)
apply (unfold ucast-bl)
apply (rule trans)
apply (rule word-rep-drop)
apply simp
done

lemma ucast-up-app':
  uc = ucast ==> source-size uc + n = target-size uc ==>
    to-bl (uc w) = replicate n False @ (to-bl w)
apply (auto simp add : source-size target-size to-bl-ucast)
apply (rule-tac f = %n. replicate n False in arg-cong)
apply simp
done

lemma ucast-down-drop':
  uc = ucast ==> source-size uc = target-size uc + n ==>
    to-bl (uc w) = drop n (to-bl w)
by (auto simp add : source-size target-size to-bl-ucast)

lemma scast-down-drop':
  sc = scast ==> source-size sc = target-size sc + n ==>
    to-bl (sc w) = drop n (to-bl w)
apply (subgoal-tac sc = ucast)
apply safe
apply simp
apply (erule refl [THEN ucast-down-drop'])
apply (rule refl [THEN down-cast-same', symmetric])
apply (simp add : source-size target-size is-down)
done

lemma sint-up-scast':
  sc = scast ==> is-up sc ==> sint (sc w) = sint w
apply (unfold is-up)
apply safe
apply (simp add: scast-def word-sbin.eq-norm)
apply (rule box-equals)
prefer 3

```

```

    apply (rule word-sbin.norm-Rep)
    apply (rule sbintrunc-sbintrunc-l)
    defer
    apply (subst word-sbin.norm-Rep)
    apply (rule refl)
    apply simp
    done

lemma uint-up-ucast':
  uc = ucast ==> is-up uc ==> uint (uc w) = uint w
  apply (unfold is-up)
  apply safe
  apply (rule bin-eqI)
  apply (fold word-test-bit-def)
  apply (auto simp add: nth-ucast)
  apply (auto simp add: test-bit-bin)
  done

lemmas down-cast-same = refl [THEN down-cast-same]
lemmas ucast-up-app = refl [THEN ucast-up-app]
lemmas ucast-down-drop = refl [THEN ucast-down-drop]
lemmas scast-down-drop = refl [THEN scast-down-drop]
lemmas uint-up-ucast = refl [THEN uint-up-ucast]
lemmas sint-up-scast = refl [THEN sint-up-scast]

lemma ucast-up-ucast': uc = ucast ==> is-up uc ==> ucast (uc w) = ucast w
  apply (simp (no-asm) add: ucast-def)
  apply (clarsimp simp add: uint-up-ucast)
  done

lemma scast-up-scast': sc = scast ==> is-up sc ==> scast (sc w) = scast w
  apply (simp (no-asm) add: scast-def)
  apply (clarsimp simp add: sint-up-scast)
  done

lemma ucast-of-bl-up':
  w = of-bl bl ==> size bl <= size w ==> ucast w = of-bl bl
  by (auto simp add : nth-ucast word-size test-bit-of-bl intro!: word-eqI)

lemmas ucast-up-ucast = refl [THEN ucast-up-ucast]
lemmas scast-up-scast = refl [THEN scast-up-scast]
lemmas ucast-of-bl-up = refl [THEN ucast-of-bl-up]

lemmas ucast-up-ucast-id = trans [OF ucast-up-ucast ucast-id]
lemmas scast-up-scast-id = trans [OF scast-up-scast scast-id]

lemmas isduu = is-up-down [where c = ucast, THEN iffD2]
lemmas isdus = is-up-down [where c = scast, THEN iffD2]
lemmas ucast-down-ucast-id = isduu [THEN ucast-up-ucast-id]

```



**lemmas** *scast-down-scast-id* = *isdus* [*THEN* *ucast-up-ucast-id*]

**lemma** *up-ucast-surj*:

*is-up* (*ucast* :: 'b::len0 word => 'a::len0 word) ==>  
*surj* (*ucast* :: 'a word => 'b word)  
**by** (*rule* *surjI*, *erule* *ucast-up-ucast-id*)

**lemma** *up-scast-surj*:

*is-up* (*scast* :: 'b::len word => 'a::len word) ==>  
*surj* (*scast* :: 'a word => 'b word)  
**by** (*rule* *surjI*, *erule* *scast-up-scast-id*)

**lemma** *down-scast-inj*:

*is-down* (*scast* :: 'b::len word => 'a::len word) ==>  
*inj-on* (*ucast* :: 'a word => 'b word) *A*  
**by** (*rule* *inj-on-inverseI*, *erule* *scast-down-scast-id*)

**lemma** *down-ucast-inj*:

*is-down* (*ucast* :: 'b::len0 word => 'a::len0 word) ==>  
*inj-on* (*ucast* :: 'a word => 'b word) *A*  
**by** (*rule* *inj-on-inverseI*, *erule* *ucast-down-ucast-id*)

**lemma** *of-bl-append-same*: *of-bl* (*X* @ *to-bl* *w*) = *w*

**by** (*rule* *word-bl.Rep-eqD*) (*simp* *add*: *word-rep-drop*)

**lemma** *ucast-down-no'*:

*uc* = *ucast* ==> *is-down* *uc* ==> *uc* (*number-of* *bin*) = *number-of* *bin*  
**apply** (*unfold* *word-number-of-def* *is-down*)  
**apply** (*clarsimp* *simp* *add*: *ucast-def* *word-ubin.eq-norm*)  
**apply** (*rule* *word-ubin.norm-eq-iff* [*THEN* *iffD1*])  
**apply** (*erule* *bintrunc-bintrunc-ge*)  
**done**

**lemmas** *ucast-down-no* = *ucast-down-no'* [*OF* *refl*]

**lemma** *ucast-down-bl'*: *uc* = *ucast* ==> *is-down* *uc* ==> *uc* (*of-bl* *bl*) = *of-bl* *bl*

**unfolding** *of-bl-no* **by** *clarify* (*erule* *ucast-down-no*)

**lemmas** *ucast-down-bl* = *ucast-down-bl'* [*OF* *refl*]

**lemmas** *slice-def'* = *slice-def* [*unfolded* *word-size*]

**lemmas** *test-bit-def'* = *word-test-bit-def* [*THEN* *meta-eq-to-obj-eq*, *THEN* *fun-cong*]

**lemmas** *word-log-defs* = *word-and-def* *word-or-def* *word-xor-def* *word-not-def*

**lemmas** *word-log-bin-defs* = *word-log-defs*

**end**

## 11 WordArith: Word Arithmetic

**theory** *WordArith* **imports** *WordDefinition* **begin**

**lemma** *word-less-alt*:  $(a < b) = (uint\ a < uint\ b)$   
**unfolding** *word-less-def word-le-def*  
**by** (*auto simp del: word-uint.Rep-inject*  
*simp: word-uint.Rep-inject [symmetric]*)

**lemma** *signed-linorder*: *linorder word-sle word-sless*  
**apply** *unfold-locales*  
**apply** (*unfold word-sle-def word-sless-def*)  
**by** *auto*

**interpretation** *signed*: *linorder* [*word-sle word-sless*]  
**by** (*rule signed-linorder*)

**lemmas** *word-arith-wis* [*THEN meta-eq-to-obj-eq*] =  
*word-add-def word-mult-def word-minus-def*  
*word-succ-def word-pred-def word-0-wi word-1-wi*

**lemma** *udvdI*:  
 $0 \leq n \implies uint\ b = n * uint\ a \implies a\ udvd\ b$   
**by** (*auto simp: udvd-def*)

**lemmas** *word-div-no* [*simp*] =  
*word-div-def* [*of number-of a number-of b, standard*]

**lemmas** *word-mod-no* [*simp*] =  
*word-mod-def* [*of number-of a number-of b, standard*]

**lemmas** *word-less-no* [*simp*] =  
*word-less-def* [*of number-of a number-of b, standard*]

**lemmas** *word-le-no* [*simp*] =  
*word-le-def* [*of number-of a number-of b, standard*]

**lemmas** *word-sless-no* [*simp*] =  
*word-sless-def* [*of number-of a number-of b, standard*]

**lemmas** *word-sle-no* [*simp*] =  
*word-sle-def* [*of number-of a number-of b, standard*]

**lemmas** *word-0-wi-Pls* = *word-0-wi* [*folded Pls-def*]  
**lemmas** *word-0-no* = *word-0-wi-Pls* [*folded word-no-wi*]

**lemma** *int-one-bin*:  $(1 :: int) == (Numeral.Pls\ BIT\ bit.B1)$

```

unfolding Pls-def Bit-def by auto

lemma word-1-no:
   $(1 :: 'a :: \text{len0 } \text{word}) == \text{number-of } (\text{Numeral.Pls BIT bit.B1})$ 
unfolding word-1-wi word-number-of-def int-one-bin by auto

lemma word-m1-wi:  $-1 == \text{word-of-int } -1$ 
by  $(\text{rule word-number-of-alt})$ 

lemma word-m1-wi-Min:  $-1 = \text{word-of-int Numeral.Min}$ 
by  $(\text{simp add: word-m1-wi number-of-eq})$ 

lemma word-0-bl:  $\text{of-bl } [] = 0$ 
unfolding word-0-wi of-bl-def by  $(\text{simp add : Pls-def})$ 

lemma word-1-bl:  $\text{of-bl } [\text{True}] = 1$ 
unfolding word-1-wi of-bl-def
by  $(\text{simp add : bl-to-bin-def Bit-def Pls-def})$ 

lemma uint-0 [simp]:  $(\text{uint } 0 = 0)$ 
unfolding word-0-wi
by  $(\text{simp add: word-ubin.eq-norm Pls-def [symmetric]})$ 

lemma of-bl-0 [simp]:  $\text{of-bl } (\text{replicate } n \text{ False}) = 0$ 
by  $(\text{simp add : word-0-wi of-bl-def bl-to-bin-rep-False Pls-def})$ 

lemma to-bl-0:
   $\text{to-bl } (0 :: 'a :: \text{len0 } \text{word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ False}$ 
unfolding uint-bl
by  $(\text{simp add : word-size bin-to-bl-Pls Pls-def [symmetric]})$ 

lemma uint-0-iff:  $(\text{uint } x = 0) = (x = 0)$ 
by  $(\text{auto intro!: word-uint.Rep-eqD})$ 

lemma unat-0-iff:  $(\text{unat } x = 0) = (x = 0)$ 
unfolding unat-def by  $(\text{auto simp add : nat-eq-iff uint-0-iff})$ 

lemma unat-0 [simp]:  $\text{unat } 0 = 0$ 
unfolding unat-def by auto

lemma size-0-same':  $\text{size } w = 0 ==> w = (v :: 'a :: \text{len0 } \text{word})$ 
apply  $(\text{unfold word-size})$ 
apply  $(\text{rule box-equals})$ 
defer
apply  $(\text{rule word-uint.Rep-inverse})+$ 
apply  $(\text{rule word-ubin.norm-eq-iff [THEN iffD1]})$ 
apply simp
done

```

```

lemmas size-0-same = size-0-same' [folded word-size]

lemmas unat-eq-0 = unat-0-iff
lemmas unat-eq-zero = unat-0-iff

lemma unat-gt-0: (0 < unat x) = (x ~ = 0)
by (auto simp: unat-0-iff [symmetric])

lemma ucast-0 [simp] : ucast 0 = 0
unfolding ucast-def
by simp (simp add: word-0-wi)

lemma sint-0 [simp] : sint 0 = 0
unfolding sint-uint
by (simp add: Pls-def [symmetric])

lemma scast-0 [simp] : scast 0 = 0
apply (unfold scast-def)
apply simp
apply (simp add: word-0-wi)
done

lemma sint-n1 [simp] : sint -1 = -1
apply (unfold word-m1-wi-Min)
apply (simp add: word-sbin.eq-norm)
apply (unfold Min-def number-of-eq)
apply simp
done

lemma scast-n1 [simp] : scast -1 = -1
  apply (unfold scast-def sint-n1)
  apply (unfold word-number-of-alt)
  apply (rule refl)
done

lemma uint-1 [simp] : uint (1 :: 'a :: len word) = 1
  unfolding word-1-wi
  by (simp add: word-ubin.eq-norm int-one-bin bintrunc-minus-simps)

lemma unat-1 [simp] : unat (1 :: 'a :: len word) = 1
  by (unfold unat-def uint-1) auto

lemma ucast-1 [simp] : ucast (1 :: 'a :: len word) = 1
  unfolding ucast-def word-1-wi
  by (simp add: word-ubin.eq-norm int-one-bin bintrunc-minus-simps)

lemmas arths =

```

*bintr-ariths* [*THEN* *word-ubin.norm-eq-iff* [*THEN iffD1*],  
*folded word-ubin.eq-norm, standard*]

**lemma** *wi-homs*:

**shows**

*wi-hom-add*: *word-of-int a + word-of-int b = word-of-int (a + b)* **and**  
*wi-hom-mult*: *word-of-int a \* word-of-int b = word-of-int (a \* b)* **and**  
*wi-hom-neg*: *− word-of-int a = word-of-int (− a)* **and**  
*wi-hom-succ*: *word-succ (word-of-int a) = word-of-int (Numeral.succ a)* **and**  
*wi-hom-pred*: *word-pred (word-of-int a) = word-of-int (Numeral.pred a)*  
**by** (*auto simp: word-arith-wis arths*)

**lemmas** *wi-hom-syms = wi-homs* [*symmetric*]

**lemma** *word-sub-def*: *a − b == a + − (b :: 'a :: len0 word)*

**unfolding** *word-sub-wi diff-def*

**by** (*simp only : word-uint.Rep-inverse wi-hom-syms*)

**lemmas** *word-diff-minus = word-sub-def* [*THEN meta-eq-to-obj-eq, standard*]

**lemma** *word-of-int-sub-hom*:

*(word-of-int a) − word-of-int b = word-of-int (a − b)*

**unfolding** *word-sub-def diff-def* **by** (*simp only : wi-homs*)

**lemmas** *new-word-of-int-homs =*

*word-of-int-sub-hom wi-homs word-0-wi word-1-wi*

**lemmas** *new-word-of-int-hom-syms = new-word-of-int-homs* [*symmetric, standard*]

**lemmas** *word-of-int-hom-syms =*

*new-word-of-int-hom-syms* [*unfolded succ-def pred-def*]

**lemmas** *word-of-int-homs =*

*new-word-of-int-homs* [*unfolded succ-def pred-def*]

**lemmas** *word-of-int-add-hom = word-of-int-homs* (2)

**lemmas** *word-of-int-mult-hom = word-of-int-homs* (3)

**lemmas** *word-of-int-minus-hom = word-of-int-homs* (4)

**lemmas** *word-of-int-succ-hom = word-of-int-homs* (5)

**lemmas** *word-of-int-pred-hom = word-of-int-homs* (6)

**lemmas** *word-of-int-0-hom = word-of-int-homs* (7)

**lemmas** *word-of-int-1-hom = word-of-int-homs* (8)

**lemmas** *word-arith-alts =*

*word-sub-wi* [*unfolded succ-def pred-def, THEN meta-eq-to-obj-eq, standard*]

*word-arith-wis* [*unfolded succ-def pred-def, standard*]

**lemmas** *word-sub-alt* = *word-arith-alts* (1)  
**lemmas** *word-add-alt* = *word-arith-alts* (2)  
**lemmas** *word-mult-alt* = *word-arith-alts* (3)  
**lemmas** *word-minus-alt* = *word-arith-alts* (4)  
**lemmas** *word-succ-alt* = *word-arith-alts* (5)  
**lemmas** *word-pred-alt* = *word-arith-alts* (6)  
**lemmas** *word-0-alt* = *word-arith-alts* (7)  
**lemmas** *word-1-alt* = *word-arith-alts* (8)

### 11.1 Transferring goals from words to ints

**lemma** *word-ths*:

**shows**

*word-succ-p1*:  $\text{word-succ } a = a + 1$  **and**  
*word-pred-m1*:  $\text{word-pred } a = a - 1$  **and**  
*word-pred-succ*:  $\text{word-pred } (\text{word-succ } a) = a$  **and**  
*word-succ-pred*:  $\text{word-succ } (\text{word-pred } a) = a$  **and**  
*word-mult-succ*:  $\text{word-succ } a * b = b + a * b$   
**by** (*rule* *word-uint.Abs-cases* [of *b*],  
*rule* *word-uint.Abs-cases* [of *a*],  
*simp* *add*: *pred-def succ-def add-commute mult-commute*  
*ring-distrib new-word-of-int-homs*)**+**

**lemmas** *uint-cong* = *arg-cong* [**where** *f* = *uint*]

**lemmas** *uint-word-ariths* =

*word-arith-alts* [*THEN* *trans* [*OF* *uint-cong int-word-uint*], *standard*]

**lemmas** *uint-word-arith-bintrs* = *uint-word-ariths* [*folded bintrunc-mod2p*]

**lemmas** *sint-word-ariths* = *uint-word-arith-bintrs*

*[THEN uint-sint* [*symmetric*, *THEN trans*],  
*unfolded uint-sint bintr-arith1s bintr-ariths*  
*len-gt-0* [*THEN bin-sbin-eq-iff*] *word-sbin.norm-Rep*, *standard*]

**lemmas** *uint-div-alt* = *word-div-def*

*[THEN meta-eq-to-obj-eq* [*THEN trans* [*OF uint-cong int-word-uint*]], *standard*]

**lemmas** *uint-mod-alt* = *word-mod-def*

*[THEN meta-eq-to-obj-eq* [*THEN trans* [*OF uint-cong int-word-uint*]], *standard*]

**lemma** *word-pred-0-n1*:  $\text{word-pred } 0 = \text{word-of-int } -1$

**unfolding** *word-pred-def number-of-eq*

**by** (*simp* *add* : *pred-def word-no-wi*)

**lemma** *word-pred-0-Min*:  $\text{word-pred } 0 = \text{word-of-int } \text{Numeral.Min}$

**by** (*simp* *add*: *word-pred-0-n1 number-of-eq*)

**lemma** *word-m1-Min*:  $-1 = \text{word-of-int } \text{Numeral.Min}$

**unfolding** *Min-def* **by** (*simp only: word-of-int-hom-syms*)

**lemma** *succ-pred-no* [*simp*]:

*word-succ* (*number-of bin*) = *number-of* (*Numeral.succ bin*) &

*word-pred* (*number-of bin*) = *number-of* (*Numeral.pred bin*)

**unfolding** *word-number-of-def* **by** (*simp add : new-word-of-int-homs*)

**lemma** *word-sp-01* [*simp*] :

*word-succ -1 = 0* & *word-succ 0 = 1* & *word-pred 0 = -1* & *word-pred 1 = 0*

**by** (*unfold word-0-no word-1-no*) *auto*

**lemma** *word-of-int-Ex*:

$\exists y. x = \text{word-of-int } y$

**by** (*rule-tac x=uint x in exI*) *simp*

**lemma** *word-arith-egs*:

**fixes** *a* :: '*a::len0 word*

**fixes** *b* :: '*a::len0 word*

**shows**

*word-add-0*:  $0 + a = a$  **and**

*word-add-0-right*:  $a + 0 = a$  **and**

*word-mult-1*:  $1 * a = a$  **and**

*word-mult-1-right*:  $a * 1 = a$  **and**

*word-add-commute*:  $a + b = b + a$  **and**

*word-add-assoc*:  $a + b + c = a + (b + c)$  **and**

*word-add-left-commute*:  $a + (b + c) = b + (a + c)$  **and**

*word-mult-commute*:  $a * b = b * a$  **and**

*word-mult-assoc*:  $a * b * c = a * (b * c)$  **and**

*word-mult-left-commute*:  $a * (b * c) = b * (a * c)$  **and**

*word-left-distrib*:  $(a + b) * c = a * c + b * c$  **and**

*word-right-distrib*:  $a * (b + c) = a * b + a * c$  **and**

*word-left-minus*:  $- a + a = 0$  **and**

*word-diff-0-right*:  $a - 0 = a$  **and**

*word-diff-self*:  $a - a = 0$

**using** *word-of-int-Ex* [*of a*]

*word-of-int-Ex* [*of b*]

*word-of-int-Ex* [*of c*]

**by** (*auto simp: word-of-int-hom-syms* [*symmetric*]

*zadd-0-right add-commute add-assoc add-left-commute*

*mult-commute mult-assoc mult-left-commute*

*plus-times.left-distrib plus-times.right-distrib*)

**lemmas** *word-add-ac* = *word-add-commute word-add-assoc word-add-left-commute*

**lemmas** *word-mult-ac* = *word-mult-commute word-mult-assoc word-mult-left-commute*

**lemmas** *word-plus-ac0* = *word-add-0 word-add-0-right word-add-ac*

**lemmas** *word-times-ac1* = *word-mult-1 word-mult-1-right word-mult-ac*

## 11.2 Order on fixed-length words

**lemma** *word-order-trans*:  $x \leq y \implies y \leq z \implies x \leq (z :: 'a :: \text{len0 word})$   
**unfolding** *word-le-def* **by** *auto*

**lemma** *word-order-refl*:  $z \leq (z :: 'a :: \text{len0 word})$   
**unfolding** *word-le-def* **by** *auto*

**lemma** *word-order-antisym*:  $x \leq y \implies y \leq x \implies x = (y :: 'a :: \text{len0 word})$   
**unfolding** *word-le-def* **by** (*auto intro!*: *word-uint.Rep-eqD*)

**lemma** *word-order-linear*:  
 $y \leq x \mid x \leq (y :: 'a :: \text{len0 word})$   
**unfolding** *word-le-def* **by** *auto*

**lemma** *word-zero-le* [*simp*] :  
 $0 \leq (y :: 'a :: \text{len0 word})$   
**unfolding** *word-le-def* **by** *auto*

**instance** *word* :: (*len0*) *semigroup-add*  
**by** *intro-classes* (*simp add*: *word-add-assoc*)

**instance** *word* :: (*len0*) *linorder*  
**by** *intro-classes* (*auto simp*: *word-less-def word-le-def*)

**instance** *word* :: (*len0*) *ring*  
**by** *intro-classes*  
(*auto simp*: *word-arith-egs word-diff-minus*  
*word-diff-self* [*unfolded word-diff-minus*])

**lemma** *word-m1-ge* [*simp*] : *word-pred* 0  $\geq y$   
**unfolding** *word-le-def*  
**by** (*simp only* : *word-pred-0-n1 word-uint.eq-norm m1mod2k*) *auto*

**lemmas** *word-n1-ge* [*simp*] = *word-m1-ge* [*simplified word-sp-01*]

**lemmas** *word-not-simps* [*simp*] =  
*word-zero-le* [*THEN leD*] *word-m1-ge* [*THEN leD*] *word-n1-ge* [*THEN leD*]

**lemma** *word-gt-0*:  $0 < y = (0 \sim (y :: 'a :: \text{len0 word}))$   
**unfolding** *word-less-def* **by** *auto*

**lemmas** *word-gt-0-no* [*simp*] = *word-gt-0* [*of number-of y, standard*]

**lemma** *word-sless-alt*:  $(a <_s b) == (\text{sint } a < \text{sint } b)$   
**unfolding** *word-sle-def word-sless-def*  
**by** (*auto simp add* : *less-eq-less.less-le*)

**lemma** *word-le-nat-alt*:  $(a \leq b) = (\text{unat } a \leq \text{unat } b)$   
**unfolding** *unat-def word-le-def*



```

  by (rule nat-le-eq-zle [symmetric]) simp

lemma word-less-nat-alt: (a < b) = (unat a < unat b)
  unfolding unat-def word-less-alt
  by (rule nat-less-eq-zless [symmetric]) simp

lemma wi-less:
  (word-of-int n < (word-of-int m :: 'a :: len0 word)) =
    (n mod 2 ^ len-of TYPE('a) < m mod 2 ^ len-of TYPE('a))
  unfolding word-less-alt by (simp add: word-uint.eq-norm)

lemma wi-le:
  (word-of-int n <= (word-of-int m :: 'a :: len0 word)) =
    (n mod 2 ^ len-of TYPE('a) <= m mod 2 ^ len-of TYPE('a))
  unfolding word-le-def by (simp add: word-uint.eq-norm)

lemma udvd-nat-alt: a udvd b = (EX n>=0. unat b = n * unat a)
  apply (unfold udvd-def)
  apply safe
  apply (simp add: unat-def nat-mult-distrib)
  apply (simp add: uint-nat int-mult)
  apply (rule exI)
  apply safe
  prefer 2
  apply (erule notE)
  apply (rule refl)
  apply force
  done

lemma udvd-iff-dvd: x udvd y <-> unat x dvd unat y
  unfolding dvd-def udvd-nat-alt by force

lemmas unat-mono = word-less-nat-alt [THEN iffD1, standard]

lemma word-zero-neq-one: 0 < len-of TYPE ('a :: len0) ==> (0 :: 'a word) ~=
1
  unfolding word-arith-wis
  by (auto simp add: word-ubin.norm-eq-iff [symmetric] gr0-conv-Suc)

lemmas lenw1-zero-neq-one = len-gt-0 [THEN word-zero-neq-one]

lemma no-no [simp] : number-of (number-of b) = number-of b
  by (simp add: number-of-eq)

lemma unat-minus-one: x ~= 0 ==> unat (x - 1) = unat x - 1
  apply (unfold unat-def)
  apply (simp only: int-word-uint word-arith-alts rdmods)
  apply (subgoal-tac uint x >= 1)
  prefer 2

```

```

apply (drule contrapos-nn)
apply (erule word-uint.Rep-inverse' [symmetric])
apply (insert uint-ge-0 [of x])[1]
apply arith
apply (rule box-equals)
apply (rule nat-diff-distrib)
prefer 2
apply assumption
apply simp
apply (subst mod-pos-pos-trivial)
apply arith
apply (insert uint-lt2p [of x])[1]
apply arith
apply (rule refl)
apply simp
done

```

```

lemma measure-unat:  $p \sim = 0 \implies \text{unat } (p - 1) < \text{unat } p$ 
by (simp add: unat-minus-one) (simp add: unat-0-iff [symmetric])

```

```

lemmas uint-add-ge0 [simp] =
  add-nonneg-nonneg [OF uint-ge-0 uint-ge-0, standard]
lemmas uint-mult-ge0 [simp] =
  mult-nonneg-nonneg [OF uint-ge-0 uint-ge-0, standard]

```

```

lemma uint-sub-lt2p [simp]:
   $\text{uint } (x :: 'a :: \text{len0 word}) - \text{uint } (y :: 'b :: \text{len0 word}) <$ 
   $2 \wedge \text{len-of TYPE('a)}$ 
using uint-ge-0 [of y] uint-lt2p [of x] by arith

```

### 11.3 Conditions for the addition (etc) of two words to overflow

```

lemma uint-add-lem:
   $(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$ 
   $(\text{uint } (x + y :: 'a :: \text{len0 word}) = \text{uint } x + \text{uint } y)$ 
by (unfold uint-word-ariths) (auto intro!: trans [OF - int-mod-lem])

```

```

lemma uint-mult-lem:
   $(\text{uint } x * \text{uint } y < 2 \wedge \text{len-of TYPE('a)}) =$ 
   $(\text{uint } (x * y :: 'a :: \text{len0 word}) = \text{uint } x * \text{uint } y)$ 
by (unfold uint-word-ariths) (auto intro!: trans [OF - int-mod-lem])

```

```

lemma uint-sub-lem:
   $(\text{uint } x \geq \text{uint } y) = (\text{uint } (x - y) = \text{uint } x - \text{uint } y)$ 
by (unfold uint-word-ariths) (auto intro!: trans [OF - int-mod-lem])

```

```

lemma uint-add-le:  $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$ 
unfolding uint-word-ariths by (auto simp: mod-add-if-z)

```

**lemma** *uint-sub-ge*:  $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$   
**unfolding** *uint-word-ariths* **by** (*auto simp: mod-sub-if-z*)

**lemmas** *uint-sub-if'* =  
*trans* [*OF uint-word-ariths*(1) *mod-sub-if-z, simplified, standard*]

**lemmas** *uint-plus-if'* =  
*trans* [*OF uint-word-ariths*(2) *mod-add-if-z, simplified, standard*]

## 11.4 Definition of uint\_arith

**lemma** *word-of-int-inverse*:  
 $\text{word-of-int } r = a \implies 0 \leq r \implies r < 2^{\text{len-of TYPE('a)}} \implies$   
 $\text{uint } (a::'a::\text{len0 word}) = r$   
**apply** (*erule word-uint.Abs-inverse'* [*rotated*])  
**apply** (*simp add: uints-num*)  
**done**

**lemma** *uint-split*:  
**fixes** *x::'a::len0 word*  
**shows**  $P (\text{uint } x) =$   
 $(\text{ALL } i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \longrightarrow P \ i)$   
**apply** (*fold word-int-case-def*)  
**apply** (*auto dest!: word-of-int-inverse simp: int-word-uint int-mod-eq'*  
*split: word-int-split*)  
**done**

**lemma** *uint-split-asm*:  
**fixes** *x::'a::len0 word*  
**shows**  $P (\text{uint } x) =$   
 $(\sim (EX \ i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ i))$   
**by** (*auto dest!: word-of-int-inverse*  
*simp: int-word-uint int-mod-eq'*  
*split: uint-split*)

**lemmas** *uint-splits* = *uint-split uint-split-asm*

**lemmas** *uint-arith-simps* =  
*word-le-def word-less-alt*  
*word-uint.Rep-inject* [*symmetric*]  
*uint-sub-if' uint-plus-if'*

**lemma** *power-False-cong*:  $\text{False} \implies a \wedge b = c \wedge d$   
**by** *auto*

**ML**  $\ll$   
*fun uint-arith-ss-of ss =*

```

    ss addsimps @{thms uint-arith-simps}
    delsimps @{thms word-uint.Rep-inject}
    addsplits @{thms split-if-asm}
    addcongs @{thms power-False-cong}

fun uint-arith-tacs ctxt =
  let fun arith-tac' n t = arith-tac ctxt n t handle COOPER => Seq.empty
  in
    [ CLASET' clarify-tac 1,
      SIMPSET' (full-simp-tac o uint-arith-ss-of) 1,
      ALLGOALS (full-simp-tac (HOL-ss addsplits @{thms uint-splits}
                               addcongs @{thms power-False-cong})),
      rewrite-goals-tac @{thms word-size},
      ALLGOALS (fn n => REPEAT (resolve-tac [allI, impI] n) THEN
        REPEAT (etac conjE n) THEN
        REPEAT (dtac @{thm word-of-int-inverse} n
          THEN atac n
          THEN atac n)),
      TRYALL arith-tac' ]
  end

fun uint-arith-tac ctxt = SELECT-GOAL (EVERY (uint-arith-tacs ctxt))
>>

method-setup uint-arith =
  Method.ctxt-args (fn ctxt => Method.SIMPLE-METHOD (uint-arith-tac ctxt 1))

  solving word arithmetic via integers and arith

11.5 More on overflows and monotonicity

lemma no-plus-overflow-uint-size:
  ((x :: 'a :: len0 word) <= x + y) = (uint x + uint y < 2 ^ size x)
  unfolding word-size by uint-arith

lemmas no-olen-add = no-plus-overflow-uint-size [unfolded word-size]

lemma no-ulen-sub: ((x :: 'a :: len0 word) >= x - y) = (uint y <= uint x)
  by uint-arith

lemma no-olen-add':
  fixes x :: 'a::len0 word
  shows (x ≤ y + x) = (uint y + uint x < 2 ^ len-of TYPE('a))
  by (simp add: word-add-ac add-ac no-olen-add)

lemmas olen-add-eqv = trans [OF no-olen-add no-olen-add' [symmetric], stan-
  dard]

lemmas uint-plus-simple-iff = trans [OF no-olen-add uint-add-lem, standard]

```

**lemmas** *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1, standard*]  
**lemmas** *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem, standard*]  
**lemmas** *uint-minus-simple-alt* = *uint-sub-lem* [*folded word-le-def*]  
**lemmas** *word-sub-le-iff* = *no-ulen-sub* [*folded word-le-def*]  
**lemmas** *word-sub-le* = *word-sub-le-iff* [*THEN iffD2, standard*]

**lemma** *word-less-sub1*:  
 $(x :: 'a :: \text{len } \text{word}) \sim = 0 \implies (1 < x) = (0 < x - 1)$   
**by** *uint-arith*

**lemma** *word-le-sub1*:  
 $(x :: 'a :: \text{len } \text{word}) \sim = 0 \implies (1 \leq x) = (0 \leq x - 1)$   
**by** *uint-arith*

**lemma** *sub-wrap-lt*:  
 $((x :: 'a :: \text{len0 } \text{word}) < x - z) = (x < z)$   
**by** *uint-arith*

**lemma** *sub-wrap*:  
 $((x :: 'a :: \text{len0 } \text{word}) \leq x - z) = (z = 0 \mid x < z)$   
**by** *uint-arith*

**lemma** *plus-minus-not-NULL-ab*:  
 $(x :: 'a :: \text{len0 } \text{word}) \leq ab - c \implies c \leq ab \implies c \sim = 0 \implies x + c \sim = 0$   
**by** *uint-arith*

**lemma** *plus-minus-no-overflow-ab*:  
 $(x :: 'a :: \text{len0 } \text{word}) \leq ab - c \implies c \leq ab \implies x \leq x + c$   
**by** *uint-arith*

**lemma** *le-minus'*:  
 $(a :: 'a :: \text{len0 } \text{word}) + c \leq b \implies a \leq a + c \implies c \leq b - a$   
**by** *uint-arith*

**lemma** *le-plus'*:  
 $(a :: 'a :: \text{len0 } \text{word}) \leq b \implies c \leq b - a \implies a + c \leq b$   
**by** *uint-arith*

**lemmas** *le-plus* = *le-plus'* [*rotated*]

**lemmas** *le-minus* = *leD* [*THEN thin-rl, THEN le-minus', standard*]

**lemma** *word-plus-mono-right*:  
 $(y :: 'a :: \text{len0 } \text{word}) \leq z \implies x \leq x + z \implies x + y \leq x + z$   
**by** *uint-arith*

**lemma** *word-less-minus-cancel*:  
 $y - x < z - x \implies x \leq z \implies (y :: 'a :: \text{len0 } \text{word}) < z$   
**by** *uint-arith*

**lemma** *word-less-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) < z \implies x \leq y \implies y - x < z - x$   
**by** *uint-arith*

**lemma** *word-less-minus-mono*:

$a < c \implies d < b \implies a - b < a \implies c - d < c$   
 $\implies a - b < c - (d :: 'a :: \text{len word})$   
**by** *uint-arith*

**lemma** *word-le-minus-cancel*:

$y - x \leq z - x \implies x \leq z \implies (y :: 'a :: \text{len0 word}) \leq z$   
**by** *uint-arith*

**lemma** *word-le-minus-mono-left*:

$(y :: 'a :: \text{len0 word}) \leq z \implies x \leq y \implies y - x \leq z - x$   
**by** *uint-arith*

**lemma** *word-le-minus-mono*:

$a \leq c \implies d \leq b \implies a - b \leq a \implies c - d \leq c$   
 $\implies a - b \leq c - (d :: 'a :: \text{len word})$   
**by** *uint-arith*

**lemma** *plus-le-left-cancel-wrap*:

$(x :: 'a :: \text{len0 word}) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$   
**by** *uint-arith*

**lemma** *plus-le-left-cancel-nowrap*:

$(x :: 'a :: \text{len0 word}) \leq x + y' \implies x \leq x + y \implies$   
 $(x + y' < x + y) = (y' < y)$   
**by** *uint-arith*

**lemma** *word-plus-mono-right2*:

$(a :: 'a :: \text{len0 word}) \leq a + b \implies c \leq b \implies a \leq a + c$   
**by** *uint-arith*

**lemma** *word-less-add-right*:

$(x :: 'a :: \text{len0 word}) < y - z \implies z \leq y \implies x + z < y$   
**by** *uint-arith*

**lemma** *word-less-sub-right*:

$(x :: 'a :: \text{len0 word}) < y + z \implies y \leq x \implies x - y < z$   
**by** *uint-arith*

**lemma** *word-le-plus-either*:

$(x :: 'a :: \text{len0 word}) \leq y \mid x \leq z \implies y \leq y + z \implies x \leq y + z$   
**by** *uint-arith*

```

lemma word-less-nowrapI:
  (x :: 'a :: len0 word) < z - k ==> k <= z ==> 0 < k ==> x < x + k
  by uint-arith

lemma inc-le: (i :: 'a :: len word) < m ==> i + 1 <= m
  by uint-arith

lemma inc-i:
  (1 :: 'a :: len word) <= i ==> i < m ==> 1 <= (i + 1) & i + 1 <= m
  by uint-arith

lemma udvd-incr-lem:
  up < uq ==> up = ua + n * uint K ==>
  uq = ua + n' * uint K ==> up + uint K <= uq
  apply clarsimp
  apply (drule less-le-mult)
  apply safe
  done

lemma udvd-incr':
  p < q ==> uint p = ua + n * uint K ==>
  uint q = ua + n' * uint K ==> p + K <= q
  apply (unfold word-less-alt word-le-def)
  apply (drule (2) udvd-incr-lem)
  apply (erule uint-add-le [THEN order-trans])
  done

lemma udvd-decr':
  p < q ==> uint p = ua + n * uint K ==>
  uint q = ua + n' * uint K ==> p <= q - K
  apply (unfold word-less-alt word-le-def)
  apply (drule (2) udvd-incr-lem)
  apply (drule le-diff-eq [THEN iffD2])
  apply (erule order-trans)
  apply (rule uint-sub-ge)
  done

lemmas udvd-incr-lem0 = udvd-incr-lem [where ua=0, simplified]
lemmas udvd-incr0 = udvd-incr' [where ua=0, simplified]
lemmas udvd-decr0 = udvd-decr' [where ua=0, simplified]

lemma udvd-minus-le':
  xy < k ==> z udvd xy ==> z udvd k ==> xy <= k - z
  apply (unfold udvd-def)
  apply clarify
  apply (erule (2) udvd-decr0)
  done

lemma udvd-incr2-K:

```

```

   $p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq$ 
 $p \implies$ 
   $0 < K \implies p \leq p + K \ \& \ p + K \leq a + s$ 
  apply (unfold udvd-def)
  apply clarify
  apply (simp add: uint-arith-simps split: split-if-asm)
  prefer 2
  apply (insert uint-range' [of s])[1]
  apply arith
  apply (drule add-commute [THEN xtr1])
  apply (simp add: diff-less-eq [symmetric])
  apply (drule less-le-mult)
  apply arith
  apply simp
  done

```

```

lemma word-succ-rbl:
   $to\_bl \ w = bl \implies to\_bl \ (word\_succ \ w) = (rev \ (rbl\_succ \ (rev \ bl)))$ 
  apply (unfold word-succ-def)
  apply clarify
  apply (simp add: to-bl-of-bin)
  apply (simp add: to-bl-def rbl-succ)
  done

```

```

lemma word-pred-rbl:
   $to\_bl \ w = bl \implies to\_bl \ (word\_pred \ w) = (rev \ (rbl\_pred \ (rev \ bl)))$ 
  apply (unfold word-pred-def)
  apply clarify
  apply (simp add: to-bl-of-bin)
  apply (simp add: to-bl-def rbl-pred)
  done

```

```

lemma word-add-rbl:
   $to\_bl \ v = vbl \implies to\_bl \ w = wbl \implies$ 
   $to\_bl \ (v + w) = (rev \ (rbl\_add \ (rev \ vbl) \ (rev \ wbl)))$ 
  apply (unfold word-add-def)
  apply clarify
  apply (simp add: to-bl-of-bin)
  apply (simp add: to-bl-def rbl-add)
  done

```

```

lemma word-mult-rbl:
   $to\_bl \ v = vbl \implies to\_bl \ w = wbl \implies$ 
   $to\_bl \ (v * w) = (rev \ (rbl\_mult \ (rev \ vbl) \ (rev \ wbl)))$ 
  apply (unfold word-mult-def)
  apply clarify
  apply (simp add: to-bl-of-bin)
  apply (simp add: to-bl-def rbl-mult)

```



**done**

**lemma** *rtb-rbl-ariths*:

$rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-succ\ w)) = rbl-succ\ ys$

$rev\ (to-bl\ w) = ys \implies rev\ (to-bl\ (word-pred\ w)) = rbl-pred\ ys$

$[| rev\ (to-bl\ v) = ys; rev\ (to-bl\ w) = xs |]$   
 $\implies rev\ (to-bl\ (v * w)) = rbl-mult\ ys\ xs$

$[| rev\ (to-bl\ v) = ys; rev\ (to-bl\ w) = xs |]$   
 $\implies rev\ (to-bl\ (v + w)) = rbl-add\ ys\ xs$

**by** (*auto simp: rev-swap [symmetric] word-succ-rbl*  
*word-pred-rbl word-mult-rbl word-add-rbl*)

## 11.6 Arithmetic type class instantiations

**instance** *word* :: (*len0*) *comm-monoid-add* ..

**instance** *word* :: (*len0*) *comm-monoid-mult*  
**apply** (*intro-classes*)  
**apply** (*simp add: word-mult-commute*)  
**apply** (*simp add: word-mult-1*)  
**done**

**instance** *word* :: (*len0*) *comm-semiring*  
**by** (*intro-classes*) (*simp add: word-left-distrib*)

**instance** *word* :: (*len0*) *ab-group-add* ..

**instance** *word* :: (*len0*) *comm-ring* ..

**instance** *word* :: (*len*) *comm-semiring-1*  
**by** (*intro-classes*) (*simp add: lenw1-zero-neq-one*)

**instance** *word* :: (*len*) *comm-ring-1* ..

**instance** *word* :: (*len0*) *comm-semiring-0* ..

**instance** *word* :: (*len0*) *order* ..

**instance** *word* :: (*len*) *recpower*  
**by** (*intro-classes*) (*simp-all add: word-pow*)

**lemma** *zero-bintrunc*:

*iszero* (*number-of* *x* :: 'a :: *len word*) =  
(*bintrunc* (*len-of TYPE*('a')) *x* = *Numeral.Pls*)  
**apply** (*unfold iszero-def word-0-wi word-no-wi*)

```

apply (rule word-ubin.norm-eq-iff [symmetric, THEN trans])
apply (simp add : Pls-def [symmetric])
done

lemmas word-le-0-iff [simp] =
  word-zero-le [THEN leD, THEN linorder-antisym-conv1]

lemma word-of-nat: of-nat n = word-of-int (int n)
by (induct n) (auto simp add : word-of-int-hom-syms)

lemma word-of-int: of-int = word-of-int
apply (rule ext)
apply (unfold of-int-def)
apply (rule contentsI)
apply safe
apply (simp-all add: word-of-nat word-of-int-homs)
defer
apply (rule Rep-Integ-ne [THEN nonemptyE])
apply (rule bexI)
prefer 2
apply assumption
apply (auto simp add: RI-eq-diff)
done

lemma word-of-int-nat:
  0 <= x ==> word-of-int x = of-nat (nat x)
by (simp add: of-nat-nat word-of-int)

lemma word-number-of-eq:
  number-of w = (of-int w :: 'a :: len word)
unfolding word-number-of-def word-of-int by auto

instance word :: (len) number-ring
by (intro-classes) (simp add : word-number-of-eq)

lemma iszero-word-no [simp] :
  iszero (number-of bin :: 'a :: len word) =
    iszero (number-of (bintrunc (len-of TYPE('a)) bin) :: int)
apply (simp add: zero-bintrunc number-of-is-id)
apply (unfold iszero-def Pls-def)
apply (rule refl)
done

```

## 11.7 Word and nat

```

lemma td-ext-unat':
  n = len-of TYPE ('a :: len) ==>
    td-ext (unat :: 'a word => nat) of-nat
    (unats n) (%i. i mod 2 ^ n)

```

```

apply (unfold td-ext-def' unat-def word-of-nat unats-uints)
apply (auto intro!: imageI simp add : word-of-int-hom-syms)
apply (erule word-uint.Abs-inverse [THEN arg-cong])
apply (simp add: int-word-uint nat-mod-distrib nat-power-eq)
done

lemmas td-ext-unat = refl [THEN td-ext-unat]
lemmas unat-of-nat = td-ext-unat [THEN td-ext.eq-norm, standard]

interpretation word-unat:
  td-ext [unat::'a::len word ==> nat
    of-nat
    unats (len-of TYPE('a::len))
    %i. i mod 2 ^ len-of TYPE('a::len)]
  by (rule td-ext-unat)

lemmas td-unat = word-unat.td-thm

lemmas unat-lt2p [iff] = word-unat.Rep [unfolded unats-def mem-Collect-eq]

lemma unat-le:  $y \leq \text{unat } (z :: 'a :: \text{len word}) \implies y : \text{unats } (\text{len-of TYPE } ('a))$ 
apply (unfold unats-def)
apply clarsimp
apply (rule xtrans, rule unat-lt2p, assumption)
done

lemma word-nchotomy:
  ALL  $w$ . EX  $n$ .  $(w :: 'a :: \text{len word}) = \text{of-nat } n \ \& \ n < 2 ^ \text{len-of TYPE } ('a)$ 
apply (rule allI)
apply (rule word-unat.Abs-cases)
apply (unfold unats-def)
apply auto
done

lemma of-nat-eq:
  fixes  $w :: 'a :: \text{len word}$ 
  shows  $(\text{of-nat } n = w) = (\exists q. n = \text{unat } w + q * 2 ^ \text{len-of TYPE } ('a))$ 
apply (rule trans)
apply (rule word-unat.inverse-norm)
apply (rule iffI)
apply (rule mod-eqD)
apply simp
apply clarsimp
done

lemma of-nat-eq-size:
   $(\text{of-nat } n = w) = (\text{EX } q. n = \text{unat } w + q * 2 ^ \text{size } w)$ 
  unfolding word-size by (rule of-nat-eq)

```

**lemma** *of-nat-0*:

$(\text{of-nat } m = (0 :: 'a :: \text{len word})) = (\exists q. m = q * 2 \wedge \text{len-of TYPE}('a))$   
**by** (*simp add: of-nat-eq*)

**lemmas** *of-nat-2p = mult-1 [symmetric, THEN iffD2 [OF of-nat-0 exI]]*

**lemma** *of-nat-gt-0*:  $\text{of-nat } k \sim = 0 \implies 0 < k$

**by** (*cases k*) *auto*

**lemma** *of-nat-neg-0*:

$0 < k \implies k < 2 \wedge \text{len-of TYPE}('a :: \text{len}) \implies \text{of-nat } k \sim = (0 :: 'a \text{ word})$   
**by** (*clarsimp simp add : of-nat-0*)

**lemma** *Abs-fnat-hom-add*:

$\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$   
**by** *simp*

**lemma** *Abs-fnat-hom-mult*:

$\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$   
**by** (*simp add: word-of-nat word-of-int-mult-hom zmultip-int*)

**lemma** *Abs-fnat-hom-Suc*:

$\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$   
**by** (*simp add: word-of-nat word-of-int-succ-hom add-ac*)

**lemma** *Abs-fnat-hom-0*:  $(0 :: 'a :: \text{len word}) = \text{of-nat } 0$

**by** (*simp add: word-of-nat word-0-wi*)

**lemma** *Abs-fnat-hom-1*:  $(1 :: 'a :: \text{len word}) = \text{of-nat } (\text{Suc } 0)$

**by** (*simp add: word-of-nat word-1-wi*)

**lemmas** *Abs-fnat-homs =*

*Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc*  
*Abs-fnat-hom-0 Abs-fnat-hom-1*

**lemma** *word-arith-nat-add*:

$a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$   
**by** *simp*

**lemma** *word-arith-nat-mult*:

$a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$   
**by** (*simp add: Abs-fnat-hom-mult [symmetric]*)

**lemma** *word-arith-nat-Suc*:

$\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$   
**by** (*subst Abs-fnat-hom-Suc [symmetric]*) *simp*

**lemma** *word-arith-nat-div*:

$a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div unat } b)$   
**by** (*simp add: word-div-def word-of-nat zdiv-int uint-nat*)

**lemma** *word-arith-nat-mod*:  
 $a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod unat } b)$   
**by** (*simp add: word-mod-def word-of-nat zmod-int uint-nat*)

**lemmas** *word-arith-nat-defs* =  
*word-arith-nat-add word-arith-nat-mult*  
*word-arith-nat-Suc Abs-fnat-hom-0*  
*Abs-fnat-hom-1 word-arith-nat-div*  
*word-arith-nat-mod*

**lemmas** *unat-cong* = *arg-cong* [**where**  $f = \text{unat}$ ]

**lemmas** *unat-word-ariths* = *word-arith-nat-defs*  
[*THEN trans* [*OF unat-cong unat-of-nat*], *standard*]

**lemmas** *word-sub-less-iff* = *word-sub-le-iff*  
[*simplified linorder-not-less* [*symmetric*], *simplified*]

**lemma** *unat-add-lem*:  
 $(\text{unat } x + \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$   
 $(\text{unat } (x + y :: 'a :: \text{len word}) = \text{unat } x + \text{unat } y)$   
**unfolding** *unat-word-ariths*  
**by** (*auto intro!: trans* [*OF - nat-mod-lem*])

**lemma** *unat-mult-lem*:  
 $(\text{unat } x * \text{unat } y < 2 \wedge \text{len-of TYPE('a)}) =$   
 $(\text{unat } (x * y :: 'a :: \text{len word}) = \text{unat } x * \text{unat } y)$   
**unfolding** *unat-word-ariths*  
**by** (*auto intro!: trans* [*OF - nat-mod-lem*])

**lemmas** *unat-plus-if'* =  
*trans* [*OF unat-word-ariths(1) mod-nat-add, simplified, standard*]

**lemma** *le-no-overflow*:  
 $x \leq b \implies a \leq a + b \implies x \leq a + (b :: 'a :: \text{len0 word})$   
**apply** (*erule order-trans*)  
**apply** (*erule olen-add-eqv* [*THEN iffD1*])  
**done**

**lemmas** *un-ui-le* = *trans*  
[*OF word-le-nat-alt* [*symmetric*]  
*word-le-def* [*THEN meta-eq-to-obj-eq*],  
*standard*]

**lemma** *unat-sub-if-size*:  
 $\text{unat } (x - y) = (\text{if } \text{unat } y \leq \text{unat } x$

```

    then unat x - unat y
  else unat x + 2 ^ size x - unat y)
apply (unfold word-size)
apply (simp add: un-ui-le)
apply (auto simp add: unat-def uint-sub-if')
apply (rule nat-diff-distrib)
prefer 3
apply (simp add: group-simps)
apply (rule nat-diff-distrib [THEN trans])
prefer 3
apply (subst nat-add-distrib)
prefer 3
apply (simp add: nat-power-eq)
apply auto
apply uint-arith
done

```

**lemmas** unat-sub-if' = unat-sub-if-size [unfolded word-size]

```

lemma unat-div: unat ((x :: 'a :: len word) div y) = unat x div unat y
apply (simp add : unat-word-ariths)
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq])
apply (rule div-le-dividend)
done

```

```

lemma unat-mod: unat ((x :: 'a :: len word) mod y) = unat x mod unat y
apply (clarsimp simp add : unat-word-ariths)
apply (cases unat y)
prefer 2
apply (rule unat-lt2p [THEN xtr7, THEN nat-mod-eq])
apply (rule mod-le-divisor)
apply auto
done

```

```

lemma uint-div: uint ((x :: 'a :: len word) div y) = uint x div uint y
unfolding uint-nat by (simp add : unat-div zdiv-int)

```

```

lemma uint-mod: uint ((x :: 'a :: len word) mod y) = uint x mod uint y
unfolding uint-nat by (simp add : unat-mod zmod-int)

```

## 11.8 Definition of unat\_arith tactic

```

lemma unat-split:
  fixes x::'a::len word
  shows P (unat x) =
    (ALL n. of-nat n = x & n < 2^len-of TYPE('a) --> P n)
  by (auto simp: unat-of-nat)

```

**lemma** unat-split-asm:

```

fixes  $x :: 'a :: \text{len word}$ 
shows  $P (\text{unat } x) =$ 
   $(\sim (EX n. \text{of-nat } n = x \ \& \ n < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ n))$ 
by  $(\text{auto simp: unat-of-nat})$ 

```

```

lemmas  $\text{of-nat-inverse} =$ 
   $\text{word-unat.Abs-inverse' [rotated, unfolded unats-def, simplified]}$ 

```

```

lemmas  $\text{unat-splits} = \text{unat-split unat-split-asm}$ 

```

```

lemmas  $\text{unat-arith-simps} =$ 
   $\text{word-le-nat-alt word-less-nat-alt}$ 
   $\text{word-unat.Rep-inject [symmetric]}$ 
   $\text{unat-sub-if' unat-plus-if' unat-div unat-mod}$ 

```

```

ML <<
  fun unat-arith-ss-of ss =
    ss addsimps @{thms unat-arith-simps}
      delsimps @{thms word-unat.Rep-inject}
      addsplits @{thms split-if-asm}
      addcongs @{thms power-False-cong}

  fun unat-arith-tacs ctxt =
    let fun arith-tac' n t = arith-tac ctxt n t handle COOPER => Seq.empty
    in
      [ CLASET' clarify-tac 1,
        SIMPSET' (full-simp-tac o unat-arith-ss-of) 1,
        ALLGOALS (full-simp-tac (HOL-ss addsplits @{thms unat-splits}
                                addcongs @{thms power-False-cong})),
        rewrite-goals-tac @{thms word-size},
        ALLGOALS (fn n => REPEAT (resolve-tac [allI, impI] n) THEN
          REPEAT (etac conjE n) THEN
          REPEAT (dtac @{thm of-nat-inverse} n THEN atac n)),
        TRYALL arith-tac' ]
    end

  fun unat-arith-tac ctxt = SELECT-GOAL (EVERY (unat-arith-tacs ctxt))
  >>

```

```

method-setup unat-arith =
  Method.ctxt-args (fn ctxt => Method.SIMPLE-METHOD (unat-arith-tac ctxt
1))
  solving word arithmetic via natural numbers and arith

```

```

lemma  $\text{no-plus-overflow-unat-size:}$ 
   $((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2^{\text{size } x})$ 
unfolding  $\text{word-size by unat-arith}$ 

```

**lemma** *unat-sub*:  $b \leq a \implies \text{unat } (a - b) = \text{unat } a - \text{unat } (b :: 'a :: \text{len word})$   
**by** *unat-arith*

**lemmas** *no-olen-add-nat* = *no-plus-overflow-unat-size* [*unfolded word-size*]

**lemmas** *unat-plus-simple* = *trans* [*OF no-olen-add-nat unat-add-lem, standard*]

**lemma** *word-div-mult*:  
 $(0 :: 'a :: \text{len word}) < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies$   
 $x * y \text{ div } y = x$   
**apply** *unat-arith*  
**apply** *clarsimp*  
**apply** (*subst unat-mult-lem* [*THEN iffD1*])  
**apply** *auto*  
**done**

**lemma** *div-lt'*:  $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies$   
 $\text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$   
**apply** *unat-arith*  
**apply** *clarsimp*  
**apply** (*drule mult-le-mono1*)  
**apply** (*erule order-le-less-trans*)  
**apply** (*rule xtr7* [*OF unat-lt2p div-mult-le*])  
**done**

**lemmas** *div-lt''* = *order-less-imp-le* [*THEN div-lt'*]

**lemma** *div-lt-mult*:  $(i :: 'a :: \text{len word}) < k \text{ div } x \implies 0 < x \implies i * x < k$   
**apply** (*frule div-lt''* [*THEN unat-mult-lem* [*THEN iffD1*]])  
**apply** (*simp add: unat-arith-simps*)  
**apply** (*drule* (1) *mult-less-mono1*)  
**apply** (*erule order-less-le-trans*)  
**apply** (*rule div-mult-le*)  
**done**

**lemma** *div-le-mult*:  
 $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$   
**apply** (*frule div-lt'* [*THEN unat-mult-lem* [*THEN iffD1*]])  
**apply** (*simp add: unat-arith-simps*)  
**apply** (*drule mult-le-mono1*)  
**apply** (*erule order-trans*)  
**apply** (*rule div-mult-le*)  
**done**

**lemma** *div-lt-uint'*:  
 $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2^{\text{len-of TYPE('a)}}$   
**apply** (*unfold uint-nat*)  
**apply** (*drule div-lt'*)  
**apply** (*simp add: zmult-int zless-nat-eq-int-zless* [*symmetric*])



```

      nat-power-eq)
done

lemmas div-lt-uint'' = order-less-imp-le [THEN div-lt-uint']

lemma word-le-exists':
  (x :: 'a :: len0 word) <= y ==>
    (EX z. y = x + z & uint x + uint z < 2 ^ len-of TYPE('a))
  apply (rule exI)
  apply (rule conjI)
  apply (rule zadd-diff-inverse)
  apply uint-arith
done

lemmas plus-minus-not-NULL = order-less-imp-le [THEN plus-minus-not-NULL-ab]

lemmas plus-minus-no-overflow =
  order-less-imp-le [THEN plus-minus-no-overflow-ab]

lemmas mcs = word-less-minus-cancel word-less-minus-mono-left
  word-le-minus-cancel word-le-minus-mono-left

lemmas word-l-diffs = mcs [where y = w + x, unfolded add-diff-cancel, standard]
lemmas word-diff-ls = mcs [where z = w + x, unfolded add-diff-cancel, standard]
lemmas word-plus-mcs = word-diff-ls
  [where y = v + x, unfolded add-diff-cancel, standard]

lemmas le-unat-voi = unat-le [THEN word-unat.Abs-inverse]

lemmas thd = refl [THEN [2] split-div-lemma [THEN iffD2], THEN conjunct1]

lemma thd1:
  a div b * b ≤ (a::nat)
  using gt-or-eq-0 [of b]
  apply (rule disjE)
  apply (erule xtr4 [OF thd mult-commute])
  apply clarsimp
done

lemmas uno-simps [THEN le-unat-voi, standard] =
  mod-le-divisor div-le-dividend thd1

lemma word-mod-div-equality:
  (n div b) * b + (n mod b) = (n :: 'a :: len word)
  apply (unfold word-less-nat-alt word-arith-nat-defs)
  apply (cut-tac y=unat b in gt-or-eq-0)
  apply (erule disjE)
  apply (simp add: mod-div-equality uno-simps)
  apply simp

```

done

**lemma** *word-div-mult-le*:  $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$   
**apply** (*unfold word-le-nat-alt word-arith-nat-defs*)  
**apply** (*cut-tac y=unat b in gt-or-eq-0*)  
**apply** (*erule disjE*)  
**apply** (*simp add: div-mult-le uno-simps*)  
**apply** *simp*  
done

**lemma** *word-mod-less-divisor*:  $0 < n \implies m \bmod n < (n :: 'a :: \text{len word})$   
**apply** (*simp only: word-less-nat-alt word-arith-nat-defs*)  
**apply** (*clarsimp simp add: uno-simps*)  
done

**lemma** *word-of-int-power-hom*:  
 $\text{word-of-int } a ^ n = (\text{word-of-int } (a ^ n) :: 'a :: \text{len word})$   
**by** (*induct n*) (*simp-all add: word-of-int-hom-syms power-Suc*)

**lemma** *word-arith-power-alt*:  
 $a ^ n = (\text{word-of-int } (\text{uint } a ^ n) :: 'a :: \text{len word})$   
**by** (*simp add: word-of-int-power-hom [symmetric]*)

**lemma** *of-bl-length-less*:  
 $\text{length } x = k \implies k < \text{len-of TYPE('a)} \implies (\text{of-bl } x :: 'a :: \text{len word}) < 2 ^ k$   
**apply** (*unfold of-bl-no [unfolded word-number-of-def]*  
 $\text{word-less-alt word-number-of-alt}$ )  
**apply** *safe*  
**apply** (*simp (no-asm) add: word-of-int-power-hom word-uint.eq-norm*  
 $\text{del: word-of-int-bin}$ )  
**apply** (*simp add: mod-pos-pos-trivial*)  
**apply** (*subst mod-pos-pos-trivial*)  
**apply** (*rule bl-to-bin-ge0*)  
**apply** (*rule order-less-trans*)  
**apply** (*rule bl-to-bin-lt2p*)  
**apply** *simp*  
**apply** (*rule bl-to-bin-lt2p*)  
done

## 11.9 Cardinality, finiteness of set of words

**lemmas**  $\text{card-lessThan}' = \text{card-lessThan}$  [*unfolded lessThan-def*]

**lemmas**  $\text{card-eq} = \text{word-unat.Abs-inj-on}$  [*THEN card-image,*  
 $\text{unfolded word-unat.image, unfolded unats-def, standard}$ ]

**lemmas**  $\text{card-word} = \text{trans}$  [*OF card-eq card-lessThan', standard*]

**lemma** *finite-word-UNIV*:  $\text{finite } (\text{UNIV} :: 'a :: \text{len word set})$

```

apply (rule contrapos-np)
  prefer 2
  apply (erule card-infinite)
apply (simp add: card-word)
done

lemma card-word-size:
   $\text{card } (UNIV :: 'a :: \text{len word set}) = (2 \wedge \text{size } (x :: 'a \text{ word}))$ 
unfolding word-size by (rule card-word)

end

```

## 12 WordBitwise: Bitwise Operations on Words

**theory** WordBitwise **imports** WordArith **begin**

**lemmas** bin-log-bintrs = bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or

**lemmas** wils1 = bin-log-bintrs [THEN word-ubin.norm-eq-iff [THEN iffD1],  
folded word-ubin.eq-norm, THEN eq-reflection, standard]

**lemmas** word-log-binary-defs =  
word-and-def word-or-def word-xor-def

**lemmas** word-no-log-defs [simp] =  
word-not-def [where a=number-of a,  
unfolded word-no-wi wils1, folded word-no-wi, standard]  
word-log-binary-defs [where a=number-of a and b=number-of b,  
unfolded word-no-wi wils1, folded word-no-wi, standard]

**lemmas** word-wi-log-defs = word-no-log-defs [unfolded word-no-wi]

**lemma** uint-or:  $\text{uint } (x \text{ OR } y) = (\text{uint } x) \text{ OR } (\text{uint } y)$   
**by** (simp add: word-or-def word-no-wi [symmetric] number-of-is-id  
bin-trunc-ao(2) [symmetric])

**lemma** uint-and:  $\text{uint } (x \text{ AND } y) = (\text{uint } x) \text{ AND } (\text{uint } y)$   
**by** (simp add: word-and-def number-of-is-id word-no-wi [symmetric]  
bin-trunc-ao(1) [symmetric])

**lemma** word-ops-nth-size:  
 $n < \text{size } (x :: 'a :: \text{len0 word}) \implies$   
 $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \ \&$

$$(x \text{ AND } y) !! n = (x !! n \& y !! n) \&$$

$$(x \text{ XOR } y) !! n = (x !! n \sim y !! n) \&$$

$$(NOT\ x) !! n = (\sim x !! n)$$

**unfolding** *word-size word-no-wi word-test-bit-def word-log-defs*  
**by** (*clarsimp simp add : word-ubin.eq-norm nth-bintr bin-nth-ops*)

**lemma** *word-ao-nth*:  
**fixes**  $x :: 'a::len0\ word$   
**shows**  $(x \text{ OR } y) !! n = (x !! n \mid y !! n) \&$   
 $(x \text{ AND } y) !! n = (x !! n \& y !! n)$   
**apply** (*cases n < size x*)  
**apply** (*drule-tac y = y in word-ops-nth-size*)  
**apply** *simp*  
**apply** (*simp add : test-bit-bin word-size*)  
**done**

**lemmas** *bwsimps =*  
*word-of-int-homs(2)*  
*word-0-wi-Pls*  
*word-m1-wi-Min*  
*word-wi-log-defs*

**lemma** *word-bw-assocs*:  
**fixes**  $x :: 'a::len0\ word$   
**shows**  
 $(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$   
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$   
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$   
**using** *word-of-int-Ex [where x=x]*  
*word-of-int-Ex [where x=y]*  
*word-of-int-Ex [where x=z]*  
**by** (*auto simp: bwsimps bbw-assocs*)

**lemma** *word-bw-comms*:  
**fixes**  $x :: 'a::len0\ word$   
**shows**  
 $x \text{ AND } y = y \text{ AND } x$   
 $x \text{ OR } y = y \text{ OR } x$   
 $x \text{ XOR } y = y \text{ XOR } x$   
**using** *word-of-int-Ex [where x=x]*  
*word-of-int-Ex [where x=y]*  
**by** (*auto simp: bwsimps bin-ops-comm*)

**lemma** *word-bw-lcs*:  
**fixes**  $x :: 'a::len0\ word$   
**shows**  
 $y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$

```

 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$ 
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$ 
using word-of-int-Ex [where  $x=x$ ]
      word-of-int-Ex [where  $x=y$ ]
      word-of-int-Ex [where  $x=z$ ]
by (auto simp: bwsimps)

```

```

lemma word-log-esimps [simp]:
  fixes  $x :: 'a::len0 \text{ word}$ 
  shows
     $x \text{ AND } 0 = 0$ 
     $x \text{ AND } -1 = x$ 
     $x \text{ OR } 0 = x$ 
     $x \text{ OR } -1 = -1$ 
     $x \text{ XOR } 0 = x$ 
     $x \text{ XOR } -1 = \text{NOT } x$ 
     $0 \text{ AND } x = 0$ 
     $-1 \text{ AND } x = x$ 
     $0 \text{ OR } x = x$ 
     $-1 \text{ OR } x = -1$ 
     $0 \text{ XOR } x = x$ 
     $-1 \text{ XOR } x = \text{NOT } x$ 
  using word-of-int-Ex [where  $x=x$ ]
  by (auto simp: bwsimps)

```

```

lemma word-not-dist:
  fixes  $x :: 'a::len0 \text{ word}$ 
  shows
     $\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$ 
     $\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$ 
  using word-of-int-Ex [where  $x=x$ ]
      word-of-int-Ex [where  $x=y$ ]
  by (auto simp: bwsimps bbw-not-dist)

```

```

lemma word-bw-same:
  fixes  $x :: 'a::len0 \text{ word}$ 
  shows
     $x \text{ AND } x = x$ 
     $x \text{ OR } x = x$ 
     $x \text{ XOR } x = 0$ 
  using word-of-int-Ex [where  $x=x$ ]
  by (auto simp: bwsimps)

```

```

lemma word-ao-absorbs [simp]:
  fixes  $x :: 'a::len0 \text{ word}$ 
  shows
     $x \text{ AND } (y \text{ OR } x) = x$ 
     $x \text{ OR } y \text{ AND } x = x$ 
     $x \text{ AND } (x \text{ OR } y) = x$ 

```

```

y AND x OR x = x
(y OR x) AND x = x
x OR x AND y = x
(x OR y) AND x = x
x AND y OR x = x
using word-of-int-Ex [where x=x]
      word-of-int-Ex [where x=y]
by (auto simp: bwsimps)

```

```

lemma word-not-not [simp]:
  NOT NOT (x::'a::len0 word) = x
  using word-of-int-Ex [where x=x]
  by (auto simp: bwsimps)

```

```

lemma word-ao-dist:
  fixes x :: 'a::len0 word
  shows (x OR y) AND z = x AND z OR y AND z
  using word-of-int-Ex [where x=x]
        word-of-int-Ex [where x=y]
        word-of-int-Ex [where x=z]
  by (auto simp: bwsimps bbw-ao-dist simp del: bin-ops-comm)

```

```

lemma word-oa-dist:
  fixes x :: 'a::len0 word
  shows x AND y OR z = (x OR z) AND (y OR z)
  using word-of-int-Ex [where x=x]
        word-of-int-Ex [where x=y]
        word-of-int-Ex [where x=z]
  by (auto simp: bwsimps bbw-oa-dist simp del: bin-ops-comm)

```

```

lemma word-add-not [simp]:
  fixes x :: 'a::len0 word
  shows x + NOT x = -1
  using word-of-int-Ex [where x=x]
  by (auto simp: bwsimps bin-add-not)

```

```

lemma word-plus-and-or [simp]:
  fixes x :: 'a::len0 word
  shows (x AND y) + (x OR y) = x + y
  using word-of-int-Ex [where x=x]
        word-of-int-Ex [where x=y]
  by (auto simp: bwsimps plus-and-or)

```

```

lemma leoa:
  fixes x :: 'a::len0 word
  shows (w = (x OR y)) ==> (y = (w AND y)) by auto
lemma leao:
  fixes x' :: 'a::len0 word
  shows (w' = (x' AND y')) ==> (x' = (x' OR w')) by auto

```

```

lemmas word-ao-equiv = leao [COMP leoa [COMP iffI]]

lemma le-word-or2: x <= x OR (y::'a::len0 word)
  unfolding word-le-def uint-or
  by (auto intro: le-int-or)

lemmas le-word-or1 = xtr3 [OF word-bw-comms (2) le-word-or2, standard]
lemmas word-and-le1 =
  xtr3 [OF word-ao-absorbs (4) [symmetric] le-word-or2, standard]
lemmas word-and-le2 =
  xtr3 [OF word-ao-absorbs (8) [symmetric] le-word-or2, standard]

lemma bl-word-not: to-bl (NOT w) = map Not (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-not-bin number-of-is-id word-no-wi [symmetric])

lemma bl-word-xor: to-bl (v XOR w) = app2 op ~ = (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs bl-xor-bin
  by (simp add: number-of-is-id word-no-wi [symmetric])

lemma bl-word-or: to-bl (v OR w) = app2 op | (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-or-bin number-of-is-id word-no-wi [symmetric])

lemma bl-word-and: to-bl (v AND w) = app2 op & (to-bl v) (to-bl w)
  unfolding to-bl-def word-log-defs
  by (simp add: bl-and-bin number-of-is-id word-no-wi [symmetric])

lemma word-lsb-alt: lsb (w::'a::len0 word) = test-bit w 0
  by (auto simp: word-test-bit-def word-lsb-def)

lemma word-lsb-1-0: lsb (1::'a::len word) & ~ lsb (0::'b::len0 word)
  unfolding word-lsb-def word-1-no word-0-no by auto

lemma word-lsb-last: lsb (w::'a::len word) = last (to-bl w)
  apply (unfold word-lsb-def uint-bl bin-to-bl-def)
  apply (rule-tac bin=uint w in bin-exhaust)
  apply (cases size w)
  apply auto
  apply (auto simp add: bin-to-bl-aux-alt)
done

lemma word-lsb-int: lsb w = (uint w mod 2 = 1)
  unfolding word-lsb-def bin-last-mod by auto

lemma word-msb-sint: msb w = (sint w < 0)
  unfolding word-msb-def
  by (simp add : sign-Min-lt-0 number-of-is-id)

```

```

lemma word-msb-no':
  w = number-of bin ==> msb (w::'a::len word) = bin-nth bin (size w - 1)
  unfolding word-msb-def word-number-of-def
  by (clarsimp simp add: word-sbin.eq-norm word-size bin-sign-lem)

lemmas word-msb-no = refl [THEN word-msb-no', unfolded word-size]

lemma word-msb-nth': msb (w::'a::len word) = bin-nth (uint w) (size w - 1)
  apply (unfold word-size)
  apply (rule trans [OF - word-msb-no])
  apply (simp add : word-number-of-def)
  done

lemmas word-msb-nth = word-msb-nth' [unfolded word-size]

lemma word-msb-alt: msb (w::'a::len word) = hd (to-bl w)
  apply (unfold word-msb-nth uint-bl)
  apply (subst hd-conv-nth)
  apply (rule length-greater-0-conv [THEN iffD1])
  apply simp
  apply (simp add : nth-bin-to-bl word-size)
  done

lemma word-set-nth:
  set-bit w n (test-bit w n) = (w::'a::len0 word)
  unfolding word-test-bit-def word-set-bit-def by auto

lemma bin-nth-uint':
  bin-nth (uint w) n = (rev (bin-to-bl (size w) (uint w)) ! n & n < size w)
  apply (unfold word-size)
  apply (safe elim!: bin-nth-uint-imp)
  apply (frule bin-nth-uint-imp)
  apply (fast dest!: bin-nth-bl)+
  done

lemmas bin-nth-uint = bin-nth-uint' [unfolded word-size]

lemma test-bit-bl: w !! n = (rev (to-bl w) ! n & n < size w)
  unfolding to-bl-def word-test-bit-def word-size
  by (rule bin-nth-uint)

lemma to-bl-nth: n < size w ==> to-bl w ! n = w !! (size w - Suc n)
  apply (unfold test-bit-bl)
  apply clarsimp
  apply (rule trans)
  apply (rule nth-rev-alt)
  apply (auto simp add: word-size)
  done

```



**lemma** *test-bit-set*:

**fixes**  $w :: 'a::len0$  word  
**shows**  $(set-bit\ w\ n\ x) !!\ n = (n < size\ w \ \&\ x)$   
**unfolding** *word-size word-test-bit-def word-set-bit-def*  
**by** (*clarsimp simp add : word-ubin.eq-norm nth-bintr*)

**lemma** *test-bit-set-gen*:

**fixes**  $w :: 'a::len0$  word  
**shows**  $test-bit\ (set-bit\ w\ n\ x)\ m =$   
 $(if\ m = n\ then\ n < size\ w \ \&\ x\ else\ test-bit\ w\ m)$   
**apply** (*unfold word-size word-test-bit-def word-set-bit-def*)  
**apply** (*clarsimp simp add: word-ubin.eq-norm nth-bintr bin-nth-sc-gen*)  
**apply** (*auto elim!: test-bit-size [unfolded word-size]*  
 $simp\ add: word-test-bit-def [symmetric]$ )  
**done**

**lemma** *of-bl-rep-False*:  $of-bl\ (replicate\ n\ False\ @\ bs) = of-bl\ bs$

**unfolding** *of-bl-def bl-to-bin-rep-F* **by** *auto*

**lemma** *msb-nth'*:

**fixes**  $w :: 'a::len$  word  
**shows**  $msb\ w = w !!\ (size\ w - 1)$   
**unfolding** *word-msb-nth' word-test-bit-def* **by** *simp*

**lemmas**  $msb-nth = msb-nth'\ [unfolded\ word-size]$

**lemmas**  $msb0 = len-gt-0\ [THEN\ diff-Suc-less,\ THEN$

$word-ops-nth-size\ [unfolded\ word-size],\ standard]$

**lemmas**  $msb1 = msb0\ [where\ i = 0]$

**lemmas**  $word-ops-msb = msb1\ [unfolded\ msb-nth\ [symmetric,\ unfolded\ One-nat-def]]$

**lemmas**  $lsb0 = len-gt-0\ [THEN\ word-ops-nth-size\ [unfolded\ word-size],\ standard]$

**lemmas**  $word-ops-lsb = lsb0\ [unfolded\ word-lsb-alt]$

**lemma** *td-ext-nth'*:

$n = size\ (w::'a::len0\ word) ==> ofn = set-bits ==> [w,\ ofn\ g] = l ==>$   
 $td-ext\ test-bit\ ofn\ \{f.\ ALL\ i.\ f\ i \longrightarrow i < n\}\ (\%h\ i.\ h\ i \ \&\ i < n)$

**apply** (*unfold word-size td-ext-def'*)

**apply** *safe*

**apply** (*rule-tac [3] ext*)

**apply** (*rule-tac [4] ext*)

**apply** (*unfold word-size of-nth-def test-bit-bl*)

**apply** *safe*

**defer**

**apply** (*clarsimp simp: word-bl.Abs-inverse*)**+**

**apply** (*rule word-bl.Rep-inverse'*)

**apply** (*rule sym [THEN trans]*)

**apply** (*rule bl-of-nth-nth*)

```

apply simp
apply (rule bl-of-nth-inj)
apply (clarsimp simp add : test-bit-bl word-size)
done

lemmas td-ext-nth = td-ext-nth' [OF refl refl refl, unfolded word-size]

interpretation test-bit:
  td-ext [op !! :: 'a::len0 word ==> nat ==> bool
    set-bits
    {f. ∀ i. f i ⟶ i < len-of TYPE('a::len0)}
    (λh i. h i ∧ i < len-of TYPE('a::len0))]
  by (rule td-ext-nth)

declare test-bit.Rep' [simp del]
declare test-bit.Rep' [rule del]

lemmas td-nth = test-bit.td-thm

lemma word-set-set-same:
  fixes w :: 'a::len0 word
  shows set-bit (set-bit w n x) n y = set-bit w n y
  by (rule word-eqI) (simp add : test-bit-set-gen word-size)

lemma word-set-set-diff:
  fixes w :: 'a::len0 word
  assumes m ~ = n
  shows set-bit (set-bit w m x) n y = set-bit (set-bit w n y) m x
  by (rule word-eqI) (clarsimp simp add : test-bit-set-gen word-size prems)

lemma test-bit-no':
  fixes w :: 'a::len0 word
  shows w = number-of bin ==> test-bit w n = (n < size w & bin-nth bin n)
  unfolding word-test-bit-def word-number-of-def word-size
  by (simp add : nth-bintr [symmetric] word-ubin.eq-norm)

lemmas test-bit-no =
  refl [THEN test-bit-no', unfolded word-size, THEN eq-reflection, standard]

lemma nth-0: ~ (0::'a::len0 word) !! n
  unfolding test-bit-no word-0-no by auto

lemma nth-sint:
  fixes w :: 'a::len word
  defines l ≡ len-of TYPE ('a)
  shows bin-nth (sint w) n = (if n < l - 1 then w !! n else w !! (l - 1))
  unfolding sint-uint l-def
  by (clarsimp simp add: nth-sbintr word-test-bit-def [symmetric])

```

**lemma** *word-lsb-no*:

*lsb (number-of bin :: 'a :: len word) = (bin-last bin = bit.B1)*

**unfolding** *word-lsb-alt test-bit-no* **by** *auto*

**lemma** *word-set-no*:

*set-bit (number-of bin :: 'a :: len0 word) n b =*

*number-of (bin-sc n (if b then bit.B1 else bit.B0) bin)*

**apply** (*unfold word-set-bit-def word-number-of-def [symmetric]*)

**apply** (*rule word-eqI*)

**apply** (*clarsimp simp: word-size bin-nth-sc-gen number-of-is-id*  
*test-bit-no nth-bintr*)

**done**

**lemmas** *setBit-no = setBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*  
*simplified if-simps, THEN eq-reflection, standard]*

**lemmas** *clearBit-no = clearBit-def [THEN trans [OF meta-eq-to-obj-eq word-set-no],*  
*simplified if-simps, THEN eq-reflection, standard]*

**lemma** *to-bl-n1*:

*to-bl (-1 :: 'a :: len0 word) = replicate (len-of TYPE ('a)) True*

**apply** (*rule word-bl.Abs-inverse'*)

**apply** *simp*

**apply** (*rule word-eqI*)

**apply** (*clarsimp simp add: word-size test-bit-no*)

**apply** (*auto simp add: word-bl.Abs-inverse test-bit-bl word-size*)

**done**

**lemma** *word-msb-n1: msb (-1 :: 'a :: len word)*

**unfolding** *word-msb-alt word-msb-alt to-bl-n1* **by** *simp*

**declare** *word-set-set-same [simp] word-set-nth [simp]*

*test-bit-no [simp] word-set-no [simp] nth-0 [simp]*

*setBit-no [simp] clearBit-no [simp]*

*word-lsb-no [simp] word-msb-no [simp] word-msb-n1 [simp] word-lsb-1-0 [simp]*

**lemma** *word-set-nth-iff*:

*(set-bit w n b = w) = (w !! n = b | n >= size (w :: 'a :: len0 word))*

**apply** (*rule iffI*)

**apply** (*rule disjCI*)

**apply** (*drule word-eqD*)

**apply** (*erule sym [THEN trans]*)

**apply** (*simp add: test-bit-set*)

**apply** (*erule disjE*)

**apply** *clarsimp*

**apply** (*rule word-eqI*)

**apply** (*clarsimp simp add: test-bit-set-gen*)

**apply** (*drule test-bit-size*)

**apply** *force*

**done**

```

lemma test-bit-2p':
  w = word-of-int (2 ^ n) ==>
    w !! m = (m = n & m < size (w :: 'a :: len word))
  unfolding word-test-bit-def word-size
  by (auto simp add: word-ubin.eq-norm nth-bintr nth-2p-bin)

lemmas test-bit-2p = refl [THEN test-bit-2p', unfolded word-size]

lemmas nth-w2p = test-bit-2p [unfolded of-int-number-of-eq
  word-of-int [symmetric] of-int-power]

lemma uint-2p:
  (0::'a::len word) < 2 ^ n ==> uint (2 ^ n::'a::len word) = 2 ^ n
  apply (unfold word-arith-power-alt)
  apply (case-tac len-of TYPE ('a))
  apply clarsimp
  apply (case-tac nat)
  apply clarsimp
  apply (case-tac n)
  apply (clarsimp simp add : word-1-wi [symmetric])
  apply (clarsimp simp add : word-0-wi [symmetric])
  apply (drule word-gt-0 [THEN iffD1])
  apply (safe intro!: word-eqI bin-nth-lem ext)
  apply (auto simp add: test-bit-2p nth-2p-bin word-test-bit-def [symmetric])
  done

lemma word-of-int-2p: (word-of-int (2 ^ n) :: 'a :: len word) = 2 ^ n
  apply (unfold word-arith-power-alt)
  apply (case-tac len-of TYPE ('a))
  apply clarsimp
  apply (case-tac nat)
  apply (rule word-ubin.norm-eq-iff [THEN iffD1])
  apply (rule box-equals)
  apply (rule-tac [2] bintr-ariths (1))+
  apply (clarsimp simp add : number-of-is-id)
  apply simp
  done

lemma bang-is-le: x !! m ==> 2 ^ m <= (x :: 'a :: len word)
  apply (rule xtr3)
  apply (rule-tac [2] y = x in le-word-or2)
  apply (rule word-eqI)
  apply (auto simp add: word-ao-nth nth-w2p word-size)
  done

lemma word-clr-le:
  fixes w :: 'a::len0 word
  shows w >= set-bit w n False

```

```

apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
apply simp
apply (rule order-trans)
  apply (rule bintr-bin-clr-le)
apply simp
done

lemma word-set-ge:
  fixes w :: 'a::len word
  shows w <= set-bit w n True
  apply (unfold word-set-bit-def word-le-def word-ubin.eq-norm)
  apply simp
  apply (rule order-trans [OF - bintr-bin-set-ge])
  apply simp
  done

end

```

## 13 WordShift: Shifting, Rotating, and Splitting Words

**theory** WordShift **imports** WordBitwise **begin**

### 13.1 Bit shifting

```

lemma shiffl1-number [simp] :
  shiffl1 (number-of w) = number-of (w BIT bit.B0)
apply (unfold shiffl1-def word-number-of-def)
apply (simp add: word-ubin.norm-eq-iff [symmetric] word-ubin.eq-norm)
apply (subst refl [THEN bintrunc-BIT-I, symmetric])
apply (subst bintrunc-bintrunc-min)
apply simp
done

```

```

lemma shiffl1-0 [simp] : shiffl1 0 = 0
  unfolding word-0-no shiffl1-number by auto

```

```

lemmas shiffl1-def-u = shiffl1-def [folded word-number-of-def]

```

```

lemma shiffl1-def-s: shiffl1 w = number-of (sint w BIT bit.B0)
  by (rule trans [OF - shiffl1-number]) simp

```

```

lemma shiftr1-0 [simp] : shiftr1 0 = 0
  unfolding shiftr1-def
  by simp (simp add: word-0-wi)

```

```

lemma sshiftr1-0 [simp] : sshiftr1 0 = 0
  apply (unfold sshiftr1-def)
  apply simp
  apply (simp add : word-0-wi)
  done

lemma sshiftr1-n1 [simp] : sshiftr1 -1 = -1
  unfolding sshiftr1-def by auto

lemma shiftr1-0 [simp] : (0::'a::len0 word) << n = 0
  unfolding shiftr1-def by (induct n) auto

lemma shiftr1-0 [simp] : (0::'a::len0 word) >> n = 0
  unfolding shiftr1-def by (induct n) auto

lemma sshiftr1-0 [simp] : 0 >>> n = 0
  unfolding sshiftr1-def by (induct n) auto

lemma sshiftr1-n1 [simp] : -1 >>> n = -1
  unfolding sshiftr1-def by (induct n) auto

lemma nth-shiftr1: shiftr1 w !! n = (n < size w & n > 0 & w !! (n - 1))
  apply (unfold shiftr1-def word-test-bit-def)
  apply (simp add: nth-bintr word-ubin.eq-norm word-size)
  apply (cases n)
  apply auto
  done

lemma nth-shiftr1' [rule-format]:
  ALL n. ((w::'a::len0 word) << m) !! n = (n < size w & n >= m & w !! (n - m))
  apply (unfold shiftr1-def)
  apply (induct m)
  apply (force elim!: test-bit-size)
  apply (clarsimp simp add : nth-shiftr1 word-size)
  apply arith
  done

lemmas nth-shiftr1 = nth-shiftr1' [unfolded word-size]

lemma nth-shiftr1: shiftr1 w !! n = w !! Suc n
  apply (unfold shiftr1-def word-test-bit-def)
  apply (simp add: nth-bintr word-ubin.eq-norm)
  apply safe
  apply (drule bin-nth.Suc [THEN iffD2, THEN bin-nth-uint-imp])
  apply simp
  done

lemma nth-shiftr:

```

```

 $\wedge n. ((w::'a::len0\ word) >> m) !! n = w !! (n + m)$ 
apply (unfold shiftr-def)
apply (induct m)
  apply (auto simp add : nth-shiftr1)
done

```

```

lemma uint-shiftr1: uint (shiftr1 w) = bin-rest (uint w)
apply (unfold shiftr1-def word-ubin.eq-norm bin-rest-trunc-i)
apply (subst bintr-uint [symmetric, OF order-refl])
apply (simp only : bintrunc-bintrunc-l)
apply simp
done

```

```

lemma nth-sshiftr1:
  sshiftr1 w !! n = (if n = size w - 1 then w !! n else w !! Suc n)
apply (unfold sshiftr1-def word-test-bit-def)
apply (simp add: nth-bintr word-ubin.eq-norm
              bin-nth.Suc [symmetric] word-size
              del: bin-nth.simps)
apply (simp add: nth-bintr uint-sint del : bin-nth.simps)
apply (auto simp add: bin-nth-sint)
done

```

```

lemma nth-sshiftr [rule-format] :
  ALL n. sshiftr w m !! n = (n < size w &
    (if n + m >= size w then w !! (size w - 1) else w !! (n + m)))
apply (unfold sshiftr-def)
apply (induct-tac m)
  apply (simp add: test-bit-bl)
apply (clarsimp simp add: nth-sshiftr1 word-size)
apply safe
  apply arith
  apply arith
apply (erule thin-rl)
apply (case-tac n)
  apply safe
  apply simp
apply simp
apply (erule thin-rl)
apply (case-tac n)
  apply safe
  apply simp
apply simp
apply arith+
done

```

```

lemma shiftr1-div-2: uint (shiftr1 w) = uint w div 2

```

```

apply (unfold shiftr1-def bin-rest-div)
apply (rule word-uint.Abs-inverse)
apply (simp add: uints-num pos-imp-zdiv-nonneg-iff)
apply (rule xtr7)
prefer 2
apply (rule zdiv-le-dividend)
apply auto
done

lemma sshiftr1-div-2: sint (sshiftr1 w) = sint w div 2
apply (unfold sshiftr1-def bin-rest-div [symmetric])
apply (simp add: word-sbin.eq-norm)
apply (rule trans)
defer
apply (subst word-sbin.norm-Rep [symmetric])
apply (rule refl)
apply (subst word-sbin.norm-Rep [symmetric])
apply (unfold One-nat-def)
apply (rule sbintrunc-rest)
done

lemma shiftr-div-2n: uint (shiftr w n) = uint w div 2 ^ n
apply (unfold shiftr-def)
apply (induct n)
apply simp
apply (simp add: shiftr1-div-2 mult-commute
               zdiv-zmult2-eq [symmetric])
done

lemma sshiftr-div-2n: sint (sshiftr w n) = sint w div 2 ^ n
apply (unfold sshiftr-def)
apply (induct n)
apply simp
apply (simp add: sshiftr1-div-2 mult-commute
               zdiv-zmult2-eq [symmetric])
done

```

### 13.1.1 shift functions in terms of lists of bools

```

lemmas bshiftr1-no-bin [simp] =
  bshiftr1-def [where w=number-of w, unfolded to-bl-no-bin, standard]

lemma bshiftr1-bl: to-bl (bshiftr1 b w) = b # butlast (to-bl w)
unfolding bshiftr1-def by (rule word-bl.Abs-inverse) simp

lemma shiftl1-of-bl: shiftl1 (of-bl bl) = of-bl (bl @ [False])
unfolding uint-bl of-bl-no
by (simp add: bl-to-bin-aux-append bl-to-bin-def)

```



```

lemma shifftl1-bl: shifftl1 (w::'a::len0 word) = of-bl (to-bl w @ [False])
proof –
  have shifftl1 w = shifftl1 (of-bl (to-bl w)) by simp
  also have ... = of-bl (to-bl w @ [False]) by (rule shifftl1-of-bl)
  finally show ?thesis .
qed

```

```

lemma bl-shifftl1:
  to-bl (shifftl1 (w :: 'a :: len word)) = tl (to-bl w) @ [False]
  apply (simp add: shifftl1-bl word-rep-drop drop-Suc drop-Cons)
  apply (fast intro!: Suc-leI)
  done

```

```

lemma shiftr1-bl: shiftr1 w = of-bl (butlast (to-bl w))
  apply (unfold shiftr1-def uint-bl of-bl-def)
  apply (simp add: butlast-rest-bin word-size)
  apply (simp add: bin-rest-trunc [symmetric, unfolded One-nat-def])
  done

```

```

lemma bl-shiftr1:
  to-bl (shiftr1 (w :: 'a :: len word)) = False # butlast (to-bl w)
  unfolding shiftr1-bl
  by (simp add : word-rep-drop len-gt-0 [THEN Suc-leI])

```

```

lemma shifftl1-rev:
  shifftl1 (w :: 'a :: len word) = word-reverse (shiftr1 (word-reverse w))
  apply (unfold word-reverse-def)
  apply (rule word-bl.Rep-inverse' [symmetric])
  apply (simp add: bl-shifftl1 bl-shiftr1 word-bl.Abs-inverse)
  apply (cases to-bl w)
  apply auto
  done

```

```

lemma shiftr1-rev:
  shiftr1 (w :: 'a :: len word) n = word-reverse (shiftr (word-reverse w) n)
  apply (unfold shiftr1-def shiftr-def)
  apply (induct n)
  apply (auto simp add : shifftl1-rev)
  done

```

```

lemmas rev-shifftl =
  shifftl-rev [where w = word-reverse w, simplified, standard]

```

```

lemmas shiftr-rev = rev-shifftl [THEN word-rev-gal', standard]
lemmas rev-shiftr = shiftr-rev [THEN word-rev-gal', standard]

```

```

lemma bl-sshiftr1:

```

```

to-bl (sshiftr1 (w :: 'a :: len word)) = hd (to-bl w) # butlast (to-bl w)
apply (unfold sshiftr1-def uint-bl word-size)
apply (simp add: butlast-rest-bin word-ubin.eq-norm)
apply (simp add: sint-uint)
apply (rule nth-equalityI)
  apply clarsimp
  apply clarsimp
  apply (case-tac i)
    apply (simp-all add: hd-conv-nth length-0-conv [symmetric]
      nth-bin-to-bl bin-nth.Suc [symmetric]
      nth-sbintr
      del: bin-nth.Suc)
    apply force
  apply (rule impI)
  apply (rule-tac f = bin-nth (uint w) in arg-cong)
  apply simp
done

```

**lemma** *drop-shiftr*:

```

drop n (to-bl ((w :: 'a :: len word) >> n)) = take (size w - n) (to-bl w)
apply (unfold shiftr-def)
apply (induct n)
  prefer 2
  apply (simp add: drop-Suc bl-shiftr1 butlast-drop [symmetric])
  apply (rule butlast-take [THEN trans])
  apply (auto simp: word-size)
done

```

**lemma** *drop-sshiftr*:

```

drop n (to-bl ((w :: 'a :: len word) >>> n)) = take (size w - n) (to-bl w)
apply (unfold sshiftr-def)
apply (induct n)
  prefer 2
  apply (simp add: drop-Suc bl-sshiftr1 butlast-drop [symmetric])
  apply (rule butlast-take [THEN trans])
  apply (auto simp: word-size)
done

```

**lemma** *take-shiftr* [rule-format] :

```

n <= size (w :: 'a :: len word) --> take n (to-bl (w >> n)) =
  replicate n False
apply (unfold shiftr-def)
apply (induct n)
  prefer 2
  apply (simp add: bl-shiftr1)
  apply (rule impI)
  apply (rule take-butlast [THEN trans])
  apply (auto simp: word-size)
done

```

```

lemma take-sshiftr' [rule-format] :
  n <= size (w :: 'a :: len word) --> hd (to-bl (w >>> n)) = hd (to-bl w) &
  take n (to-bl (w >>> n)) = replicate n (hd (to-bl w))
apply (unfold sshiftr-def)
apply (induct n)
prefer 2
apply (simp add: bl-sshiftr1)
apply (rule impI)
apply (rule take-butlast [THEN trans])
apply (auto simp: word-size)
done

lemmas hd-sshiftr = take-sshiftr' [THEN conjunct1, standard]
lemmas take-sshiftr = take-sshiftr' [THEN conjunct2, standard]

lemma atd-lem: take n xs = t ==> drop n xs = d ==> xs = t @ d
by (auto intro: append-take-drop-id [symmetric])

lemmas bl-shiftr = atd-lem [OF take-shiftr drop-shiftr]
lemmas bl-sshiftr = atd-lem [OF take-sshiftr drop-sshiftr]

lemma shiftl-of-bl: of-bl bl << n = of-bl (bl @ replicate n False)
unfolding shiftl-def
by (induct n) (auto simp: shiftl1-of-bl replicate-app-Cons-same)

lemma shiftl-bl:
  (w::'a::len0 word) << (n::nat) = of-bl (to-bl w @ replicate n False)
proof –
  have w << n = of-bl (to-bl w) << n by simp
  also have ... = of-bl (to-bl w @ replicate n False) by (rule shiftl-of-bl)
  finally show ?thesis .
qed

lemmas shiftl-number [simp] = shiftl-def [where w=number-of w, standard]

lemma bl-shiftl:
  to-bl (w << n) = drop n (to-bl w) @ replicate (min (size w) n) False
by (simp add: shiftl-bl word-rep-drop word-size min-def)

lemma shiftl-zero-size:
  fixes x :: 'a::len0 word
  shows size x <= n ==> x << n = 0
  apply (unfold word-size)
  apply (rule word-eqI)
  apply (clarsimp simp add: shiftl-bl word-size test-bit-of-bl nth-append)
done

```

```

lemma shiffl1-2t: shiffl1 (w :: 'a :: len word) = 2 * w
  apply (simp add: shiffl1-def-u)
  apply (simp only: double-number-of-BIT [symmetric])
  apply simp
  done

```

```

lemma shiffl1-p: shiffl1 (w :: 'a :: len word) = w + w
  apply (simp add: shiffl1-def-u)
  apply (simp only: double-number-of-BIT [symmetric])
  apply simp
  done

```

```

lemma shiffl-t2n: shiffl (w :: 'a :: len word) n = 2 ^ n * w
  unfolding shiffl-def
  by (induct n) (auto simp: shiffl1-2t power-Suc)

```

```

lemma shiftr1-bintr [simp]:
  (shiftr1 (number-of w) :: 'a :: len0 word) =
    number-of (bin-rest (bintrunc (len-of TYPE ('a)) w))
  unfolding shiftr1-def word-number-of-def
  by (simp add : word-ubin.eq-norm)

```

```

lemma sshiftr1-sbintr [simp] :
  (sshiftr1 (number-of w) :: 'a :: len word) =
    number-of (bin-rest (sbintrunc (len-of TYPE ('a) - 1) w))
  unfolding sshiftr1-def word-number-of-def
  by (simp add : word-sbin.eq-norm)

```

```

lemma shiftr-no':
  w = number-of bin ==>
    (w::'a::len0 word) >> n = number-of ((bin-rest ^ n) (bintrunc (size w) bin))
  apply clarsimp
  apply (rule word-eqI)
  apply (auto simp: nth-shiftr nth-rest-power-bin nth-bintr word-size)
  done

```

```

lemma sshiftr-no':
  w = number-of bin ==> w >>> n = number-of ((bin-rest ^ n)
    (sbintrunc (size w - 1) bin))
  apply clarsimp
  apply (rule word-eqI)
  apply (auto simp: nth-sshiftr nth-rest-power-bin nth-sbintr word-size)
  apply (subgoal-tac na + n = len-of TYPE('a) - Suc 0, simp, simp) +
  done

```

```

lemmas sshiftr-no [simp] =
  sshiftr-no' [where w = number-of w, OF refl, unfolded word-size, standard]

```

**lemmas** *shiftr-no* [*simp*] =  
*shiftr-no'* [**where** *w* = *number-of w*, *OF refl*, *unfolded word-size*, *standard*]

**lemma** *shiftr1-bl-of'*:  
 $us = \text{shiftr1 } (of\text{-bl } bl) \implies \text{length } bl \leq \text{size } us \implies$   
 $us = of\text{-bl } (\text{butlast } bl)$   
**by** (*clarsimp simp: shiftr1-def of-bl-def word-size butlast-rest-bl2bin*  
*word-ubin.eq-norm trunc-bl2bin*)

**lemmas** *shiftr1-bl-of* = *refl* [*THEN shiftr1-bl-of'*, *unfolded word-size*]

**lemma** *shiftr-bl-of'* [*rule-format*]:  
 $us = of\text{-bl } bl \gg n \implies \text{length } bl \leq \text{size } us \implies$   
 $us = of\text{-bl } (\text{take } (\text{length } bl - n) \text{ } bl)$   
**apply** (*unfold shiftr-def*)  
**apply** *hypsubst*  
**apply** (*unfold word-size*)  
**apply** (*induct n*)  
**apply** *clarsimp*  
**apply** *clarsimp*  
**apply** (*subst shiftr1-bl-of*)  
**apply** *simp*  
**apply** (*simp add: butlast-take*)  
**done**

**lemmas** *shiftr-bl-of* = *refl* [*THEN shiftr-bl-of'*, *unfolded word-size*]

**lemmas** *shiftr-bl* = *word-bl.Rep'* [*THEN eq-imp-le*, *THEN shiftr-bl-of*,  
*simplified word-size*, *simplified*, *THEN eq-reflection*, *standard*]

**lemma** *msb-shift'*:  $\text{msb } (w::'a::\text{len word}) \longleftrightarrow (w \gg (\text{size } w - 1)) \sim = 0$   
**apply** (*unfold shiftr-bl word-msb-alt*)  
**apply** (*simp add: word-size Suc-le-eq take-Suc*)  
**apply** (*cases hd (to-bl w)*)  
**apply** (*auto simp: word-1-bl word-0-bl*  
*of-bl-rep-False [where n=1 and bs=[], simplified]*)  
**done**

**lemmas** *msb-shift* = *msb-shift'* [*unfolded word-size*]

**lemma** *align-lem-or* [*rule-format*] :  
 $\text{ALL } x \text{ } m. \text{length } x = n + m \implies \text{length } y = n + m \implies$   
 $\text{drop } m \text{ } x = \text{replicate } n \text{ False} \implies \text{take } m \text{ } y = \text{replicate } m \text{ False} \implies$   
 $\text{app2 } op \mid x \text{ } y = \text{take } m \text{ } x @ \text{drop } m \text{ } y$   
**apply** (*induct-tac y*)  
**apply** *force*  
**apply** *clarsimp*  
**apply** (*case-tac x, force*)  
**apply** (*case-tac m, auto*)

```

apply (drule sym)
apply auto
apply (induct-tac list, auto)
done

```

```

lemma align-lem-and [rule-format] :
  ALL x m. length x = n + m --> length y = n + m -->
    drop m x = replicate n False --> take m y = replicate m False -->
    app2 op & x y = replicate (n + m) False
apply (induct-tac y)
apply force
apply clarsimp
apply (case-tac x, force)
apply (case-tac m, auto)
apply (drule sym)
apply auto
apply (induct-tac list, auto)
done

```

```

lemma aligned-bl-add-size':
  size x - n = m ==> n <= size x ==> drop m (to-bl x) = replicate n False
==>
  take m (to-bl y) = replicate m False ==>
  to-bl (x + y) = take m (to-bl x) @ drop m (to-bl y)
apply (subgoal-tac x AND y = 0)
prefer 2
apply (rule word-bl.Rep-eqD)
apply (simp add: bl-word-and to-bl-0)
apply (rule align-lem-and [THEN trans])
  apply (simp-all add: word-size)[5]
apply (rule-tac f = %n. replicate n False in arg-cong)
apply simp
apply (subst word-plus-and-or [symmetric])
apply (simp add : bl-word-or)
apply (rule align-lem-or)
  apply (simp-all add: word-size)
done

```

```

lemmas aligned-bl-add-size = refl [THEN aligned-bl-add-size']

```

### 13.1.2 Mask

```

lemma nth-mask': m = mask n ==> test-bit m i = (i < n & i < size m)
apply (unfold mask-def test-bit-bl)
apply (simp only: word-1-bl [symmetric] shiftl-of-bl)
apply (clarsimp simp add: word-size)
apply (simp only: of-bl-no mask-lem number-of-succ add-diff-cancel2)
apply (fold of-bl-no)
apply (simp add: word-1-bl)

```

```

apply (rule test-bit-of-bl [THEN trans, unfolded test-bit-bl word-size])
apply auto
done

lemmas nth-mask [simp] = refl [THEN nth-mask']

lemma mask-bl: mask n = of-bl (replicate n True)
  by (auto simp add : test-bit-of-bl word-size intro: word-eqI)

lemma mask-bin: mask n = number-of (bintrunc n Numeral.Min)
  by (auto simp add: nth-bintr word-size intro: word-eqI)

lemma and-mask-bintr: w AND mask n = number-of (bintrunc n (uint w))
  apply (rule word-eqI)
  apply (simp add: nth-bintr word-size word-ops-nth-size)
  apply (auto simp add: test-bit-bin)
  done

lemma and-mask-no: number-of i AND mask n = number-of (bintrunc n i)
  by (auto simp add : nth-bintr word-size word-ops-nth-size intro: word-eqI)

lemmas and-mask-wi = and-mask-no [unfolded word-number-of-def]

lemma bl-and-mask:
  to-bl (w AND mask n :: 'a :: len word) =
    replicate (len-of TYPE('a) - n) False @
    drop (len-of TYPE('a) - n) (to-bl w)
  apply (rule nth-equalityI)
  apply simp
  apply (clarsimp simp add: to-bl-nth word-size)
  apply (simp add: word-size word-ops-nth-size)
  apply (auto simp add: word-size test-bit-bl nth-append nth-rev)
  done

lemmas and-mask-mod-2p =
  and-mask-bintr [unfolded word-number-of-alt no-bintr-alt]

lemma and-mask-lt-2p: uint (w AND mask n) < 2 ^ n
  apply (simp add : and-mask-bintr no-bintr-alt)
  apply (rule xtr8)
  prefer 2
  apply (rule pos-mod-bound)
  apply auto
  done

lemmas eq-mod-iff = trans [symmetric, OF int-mod-lem eq-sym-conv]

lemma mask-eq-iff: (w AND mask n) = w <-> uint w < 2 ^ n
  apply (simp add: and-mask-bintr word-number-of-def)

```

```

apply (simp add: word-ubin.inverse-norm)
apply (simp add: eq-mod-iff bintrunc-mod2p min-def)
apply (fast intro!: lt2p-lem)
done

lemma and-mask-dvd:  $2^n \text{ dvd } \text{uint } w = (w \text{ AND } \text{mask } n = 0)$ 
apply (simp add: zdvd-iff-zmod-eq-0 and-mask-mod-2p)
apply (simp add: word-uint.norm-eq-iff [symmetric] word-of-int-homs)
apply (subst word-uint.norm-Rep [symmetric])
apply (simp only: bintrunc-bintrunc-min bintrunc-mod2p [symmetric] min-def)
apply auto
done

lemma and-mask-dvd-nat:  $2^n \text{ dvd } \text{unat } w = (w \text{ AND } \text{mask } n = 0)$ 
apply (unfold unat-def)
apply (rule trans [OF - and-mask-dvd])
apply (unfold dvd-def)
apply auto
apply (drule uint-ge-0 [THEN nat-int.Abs-inverse' [simplified], symmetric])
apply (simp add : int-mult int-power)
apply (simp add : nat-mult-distrib nat-power-eq)
done

lemma word-2p-lem:
   $n < \text{size } w \implies w < 2^n = (\text{uint } (w :: 'a :: \text{len word}) < 2^n)$ 
apply (unfold word-size word-less-alt word-number-of-alt)
apply (clarsimp simp add: word-of-int-power-hom word-uint.eq-norm
    int-mod-eq'
    simp del: word-of-int-bin)
done

lemma less-mask-eq:  $x < 2^n \implies x \text{ AND } \text{mask } n = (x :: 'a :: \text{len word})$ 
apply (unfold word-less-alt word-number-of-alt)
apply (clarsimp simp add: and-mask-mod-2p word-of-int-power-hom
    word-uint.eq-norm
    simp del: word-of-int-bin)
apply (drule xtr8 [rotated])
apply (rule int-mod-le)
apply (auto simp add : mod-pos-pos-trivial)
done

lemmas mask-eq-iff-w2p =
  trans [OF mask-eq-iff word-2p-lem [symmetric], standard]

lemmas and-mask-less' =
  iffD2 [OF word-2p-lem and-mask-lt-2p, simplified word-size, standard]

lemma and-mask-less-size:  $n < \text{size } x \implies x \text{ AND } \text{mask } n < 2^n$ 
unfolding word-size by (erule and-mask-less')

```



**lemma** *word-mod-2p-is-mask'*:

$c = 2 \wedge n \implies c > 0 \implies x \bmod c = (x :: 'a :: \text{len word}) \text{ AND } \text{mask } n$   
**by** (*clarsimp simp add: word-mod-def uint-2p and-mask-mod-2p*)

**lemmas** *word-mod-2p-is-mask* = *refl [THEN word-mod-2p-is-mask']*

**lemma** *mask-egs*:

$(a \text{ AND } \text{mask } n) + b \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$   
 $a + (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$   
 $(a \text{ AND } \text{mask } n) - b \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$   
 $a - (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$   
 $a * (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a * b \text{ AND } \text{mask } n$   
 $(b \text{ AND } \text{mask } n) * a \text{ AND } \text{mask } n = b * a \text{ AND } \text{mask } n$   
 $(a \text{ AND } \text{mask } n) + (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a + b \text{ AND } \text{mask } n$   
 $(a \text{ AND } \text{mask } n) - (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a - b \text{ AND } \text{mask } n$   
 $(a \text{ AND } \text{mask } n) * (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a * b \text{ AND } \text{mask } n$   
 $-(a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = -a \text{ AND } \text{mask } n$   
 $\text{word-succ } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-succ } a \text{ AND } \text{mask } n$   
 $\text{word-pred } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-pred } a \text{ AND } \text{mask } n$   
**using** *word-of-int-Ex [where x=a] word-of-int-Ex [where x=b]*  
**by** (*auto simp: and-mask-wi bintr-ariths bintr-arith1s new-word-of-int-homs*)

**lemma** *mask-power-eq*:

$(x \text{ AND } \text{mask } n) \wedge^k \text{ AND } \text{mask } n = x \wedge^k \text{ AND } \text{mask } n$   
**using** *word-of-int-Ex [where x=x]*  
**by** (*clarsimp simp: and-mask-wi word-of-int-power-hom bintr-ariths*)

### 13.1.3 Recast

**lemmas** *revcast-def' = revcast-def [simplified]*

**lemmas** *revcast-def'' = revcast-def' [simplified word-size]*

**lemmas** *revcast-no-def [simp] =*  
*revcast-def' [where w=number-of w, unfolded word-size, standard]*

**lemma** *to-bl-revcast*:

$\text{to-bl } (\text{revcast } w :: 'a :: \text{len0 word}) =$   
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$   
**apply** (*unfold revcast-def' word-size*)  
**apply** (*rule word-bl.Abs-inverse*)  
**apply** *simp*  
**done**

**lemma** *revcast-rev-ucast'*:

$cs = [rc, uc] \implies rc = \text{revcast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$   
 $rc = \text{word-reverse } uc$   
**apply** (*unfold ucast-def revcast-def' Let-def word-reverse-def*)  
**apply** (*clarsimp simp add : to-bl-of-bin takefill-bintrunc*)  
**apply** (*simp add : word-bl.Abs-inverse word-size*)

done

lemmas *revcast-rev-ucast* = *revcast-rev-ucast'* [*OF refl refl refl*]

lemmas *revcast-ucast* = *revcast-rev-ucast*  
[where *w* = *word-reverse w*, *simplified word-rev-rev*, *standard*]

lemmas *ucast-revcast* = *revcast-rev-ucast* [*THEN word-rev-gal'*, *standard*]

lemmas *ucast-rev-revcast* = *revcast-ucast* [*THEN word-rev-gal'*, *standard*]

— linking revcast and cast via shift

lemmas *wsst-TYs* = *source-size target-size word-size*

lemma *revcast-down-uu'*:  
 $rc = revcast ==> source-size\ rc = target-size\ rc + n ==>$   
 $rc\ (w :: 'a :: len\ word) = ucast\ (w >> n)$   
 apply (*simp add: revcast-def'*)  
 apply (*rule word-bl.Rep-inverse'*)  
 apply (*rule trans, rule ucast-down-drop*)  
 prefer 2  
 apply (*rule trans, rule drop-shiftr*)  
 apply (*auto simp: takefill-alt wsst-TYs*)  
 done

lemma *revcast-down-us'*:  
 $rc = revcast ==> source-size\ rc = target-size\ rc + n ==>$   
 $rc\ (w :: 'a :: len\ word) = ucast\ (w >>> n)$   
 apply (*simp add: revcast-def'*)  
 apply (*rule word-bl.Rep-inverse'*)  
 apply (*rule trans, rule ucast-down-drop*)  
 prefer 2  
 apply (*rule trans, rule drop-sshiftr*)  
 apply (*auto simp: takefill-alt wsst-TYs*)  
 done

lemma *revcast-down-su'*:  
 $rc = revcast ==> source-size\ rc = target-size\ rc + n ==>$   
 $rc\ (w :: 'a :: len\ word) = scast\ (w >> n)$   
 apply (*simp add: revcast-def'*)  
 apply (*rule word-bl.Rep-inverse'*)  
 apply (*rule trans, rule scast-down-drop*)  
 prefer 2  
 apply (*rule trans, rule drop-shiftr*)  
 apply (*auto simp: takefill-alt wsst-TYs*)  
 done

lemma *revcast-down-ss'*:

```

rc = revcast ==> source-size rc = target-size rc + n ==>
  rc (w :: 'a :: len word) = scast (w >>> n)
apply (simp add: revcast-def')
apply (rule word-bl.Rep-inverse')
apply (rule trans, rule scast-down-drop)
prefer 2
apply (rule trans, rule drop-sshiftr)
apply (auto simp: takefill-alt wsst-TYs)
done

```

```

lemmas revcast-down-uu = refl [THEN revcast-down-uu']
lemmas revcast-down-us = refl [THEN revcast-down-us']
lemmas revcast-down-su = refl [THEN revcast-down-su']
lemmas revcast-down-ss = refl [THEN revcast-down-ss']

```

```

lemma cast-down-rev:
  uc = ucast ==> source-size uc = target-size uc + n ==>
    uc w = revcast ((w :: 'a :: len word) << n)
apply (unfold shiftl-rev)
apply clarify
apply (simp add: revcast-rev-ucast)
apply (rule word-rev-gal')
apply (rule trans [OF - revcast-rev-ucast])
apply (rule revcast-down-uu [symmetric])
apply (auto simp add: wsst-TYs)
done

```

```

lemma revcast-up':
  rc = revcast ==> source-size rc + n = target-size rc ==>
    rc w = (ucast w :: 'a :: len word) << n
apply (simp add: revcast-def')
apply (rule word-bl.Rep-inverse')
apply (simp add: takefill-alt)
apply (rule bl-shiftl [THEN trans])
apply (subst ucast-up-app)
apply (auto simp add: wsst-TYs)
apply (drule sym)
apply (simp add: min-def)
done

```

```

lemmas revcast-up = refl [THEN revcast-up']

```

```

lemmas rc1 = revcast-up [THEN
  revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]
lemmas rc2 = revcast-down-uu [THEN
  revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

```

```

lemmas ucast-up =
  rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]

```

**lemmas** *ucast-down* =  
*rc2* [*simplified rev-shiftr revcast-ucast* [*symmetric*]]

### 13.1.4 Slices

**lemmas** *slice1-no-bin* [*simp*] =  
*slice1-def* [**where** *w=number-of w, unfolded to-bl-no-bin, standard*]

**lemmas** *slice-no-bin* [*simp*] =  
*trans* [*OF slice-def* [*THEN meta-eq-to-obj-eq*]  
*slice1-no-bin* [*THEN meta-eq-to-obj-eq*],  
*unfolded word-size, standard*]

**lemma** *slice1-0* [*simp*] : *slice1 n 0 = 0*  
**unfolding** *slice1-def* **by** (*simp add : to-bl-0*)

**lemma** *slice-0* [*simp*] : *slice n 0 = 0*  
**unfolding** *slice-def* **by** *auto*

**lemma** *slice-take'*: *slice n w = of-bl (take (size w - n) (to-bl w))*  
**unfolding** *slice-def' slice1-def*  
**by** (*simp add : takefill-alt word-size*)

**lemmas** *slice-take* = *slice-take'* [*unfolded word-size*]

— shiftr to a word of the same size is just slice, slice is just shiftr then ucast

**lemmas** *shiftr-slice* = *trans*  
*[OF shiftr-bl* [*THEN meta-eq-to-obj-eq*] *slice-take* [*symmetric*], *standard*]

**lemma** *slice-shiftr*: *slice n w = ucast (w >> n)*  
**apply** (*unfold slice-take shiftr-bl*)  
**apply** (*rule ucast-of-bl-up* [*symmetric*])  
**apply** (*simp add: word-size*)  
**done**

**lemma** *nth-slice*:  
(*slice n w :: 'a :: len0 word*) !! *m* =  
(*w* !! (*m* + *n*) & *m* < *len-of TYPE ('a)*)  
**unfolding** *slice-shiftr*  
**by** (*simp add : nth-ucast nth-shiftr*)

**lemma** *slice1-down-alt'*:  
*sl = slice1 n w ==> fs = size sl ==> fs + k = n ==>*  
*to-bl sl = takefill False fs (drop k (to-bl w))*  
**unfolding** *slice1-def word-size of-bl-def uint-bl*  
**by** (*clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop drop-takefill*)

**lemma** *slice1-up-alt'*:  
*sl = slice1 n w ==> fs = size sl ==> fs = n + k ==>*

```

    to-bl sl = takefill False fs (replicate k False @ (to-bl w))
  apply (unfold slice1-def word-size of-bl-def uint-bl)
  apply (clarsimp simp: word-ubin.eq-norm bl-bin-bl-rep-drop
           takefill-append [symmetric])
  apply (rule-tac f = %k. takefill False (len-of TYPE('a))
         (replicate k False @ bin-to-bl (len-of TYPE('b)) (uint w)) in arg-cong)
  apply arith
done

```

```

lemmas sd1 = slice1-down-alt' [OF refl refl, unfolded word-size]
lemmas su1 = slice1-up-alt' [OF refl refl, unfolded word-size]
lemmas slice1-down-alt = le-add-diff-inverse [THEN sd1]
lemmas slice1-up-alt =
  le-add-diff-inverse [symmetric, THEN su1]
  le-add-diff-inverse2 [symmetric, THEN su1]

```

```

lemma ucast-slice1: ucast w = slice1 (size w) w
  unfolding slice1-def ucast-bl
  by (simp add : takefill-same' word-size)

```

```

lemma ucast-slice: ucast w = slice 0 w
  unfolding slice-def by (simp add : ucast-slice1)

```

```

lemmas slice-id = trans [OF ucast-slice [symmetric] ucast-id]

```

```

lemma revcast-slice1':
  rc = revcast w ==> slice1 (size rc) w = rc
  unfolding slice1-def revcast-def' by (simp add : word-size)

```

```

lemmas revcast-slice1 = refl [THEN revcast-slice1']

```

```

lemma slice1-tf-tf':
  to-bl (slice1 n w :: 'a :: len0 word) =
    rev (takefill False (len-of TYPE('a)) (rev (takefill False n (to-bl w))))
  unfolding slice1-def by (rule word-rev-tf)

```

```

lemmas slice1-tf-tf = slice1-tf-tf'
  [THEN word-bl.Rep-inverse', symmetric, standard]

```

```

lemma rev-slice1:
  n + k = len-of TYPE('a) + len-of TYPE('b) ==>
  slice1 n (word-reverse w :: 'b :: len0 word) =
  word-reverse (slice1 k w :: 'a :: len0 word)
  apply (unfold word-reverse-def slice1-tf-tf)
  apply (rule word-bl.Rep-inverse')
  apply (rule rev-swap [THEN iffD1])
  apply (rule trans [symmetric])
  apply (rule tf-rev)
  apply (simp add: word-bl.Abs-inverse)

```

```

apply (simp add: word-bl.Abs-inverse)
done

```

**lemma** *rev-slice'*:

```

  res = slice n (word-reverse w) ==> n + k + size res = size w ==>
    res = word-reverse (slice k w)
apply (unfold slice-def word-size)
apply clarify
apply (rule rev-slice1)
apply arith
done

```

**lemmas** *rev-slice* = refl [THEN *rev-slice'*, unfolded word-size]

**lemmas** *sym-notr* =

```

  not-iff [THEN iffD2, THEN not-sym, THEN not-iff [THEN iffD1]]

```

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

**lemma** *soft-test*:

```

  (sint (x :: 'a :: len word) + sint y = sint (x + y)) =
    (((x+y) XOR x) AND ((x+y) XOR y)) >> (size x - 1) = 0)
apply (unfold word-size)
apply (cases len-of TYPE('a), simp)
apply (subst msb-shift [THEN sym-notr])
apply (simp add: word-ops-msb)
apply (simp add: word-msb-sint)
apply safe
  apply simp-all
apply (unfold sint-word-ariths)
apply (unfold word-sbin.set-iff-norm [symmetric] sints-num)
apply safe
  apply (insert sint-range' [where x=x])
  apply (insert sint-range' [where x=y])
  defer
  apply (simp (no-asm), arith)
  apply (simp (no-asm), arith)
  defer
  defer
  apply (simp (no-asm), arith)
  apply (simp (no-asm), arith)
apply (rule notI [THEN notnotD],
  drule leI not-leE,
  drule sbintrunc-inc sbintrunc-dec,
  simp)+
done

```

### 13.2 Split and cat

**lemmas** *word-split-bin'* = *word-split-def* [*THEN meta-eq-to-obj-eq, standard*]

**lemmas** *word-cat-bin'* = *word-cat-def* [*THEN meta-eq-to-obj-eq, standard*]

**lemma** *word-rsplit-no*:

(*word-rsplit* (*number-of* *bin* :: '*b* :: len0 *word*) :: '*a* *word list*) =  
 map *number-of* (*bin-rsplit* (*len-of TYPE*('a :: len))  
 (*len-of TYPE*('b), *bintrunc* (*len-of TYPE*('b)) *bin*))

**apply** (*unfold word-rsplit-def word-no-wi*)

**apply** (*simp add: word-ubin.eq-norm*)

**done**

**lemmas** *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*

[*unfolded bin-rsplittl-def bin-rsplit-l [symmetric]*]

**lemma** *test-bit-cat*:

*wc* = *word-cat* *a b* ==> *wc* !! *n* = (*n* < *size wc* &  
 (*if n* < *size b* then *b* !! *n* else *a* !! (*n* - *size b*)))

**apply** (*unfold word-cat-bin' test-bit-bin*)

**apply** (*auto simp add : word-ubin.eq-norm nth-bintr bin-nth-cat word-size*)

**apply** (*erule bin-nth-uint-imp*)

**done**

**lemma** *word-cat-bl*: *word-cat* *a b* = *of-bl* (*to-bl* *a* @ *to-bl* *b*)

**apply** (*unfold of-bl-def to-bl-def word-cat-bin'*)

**apply** (*simp add: bl-to-bin-app-cat*)

**done**

**lemma** *of-bl-append*:

(*of-bl* (*xs* @ *ys*) :: '*a* :: len *word*) = *of-bl* *xs* \*  $2^{length\ ys}$  + *of-bl* *ys*

**apply** (*unfold of-bl-def*)

**apply** (*simp add: bl-to-bin-app-cat bin-cat-num*)

**apply** (*simp add: word-of-int-power-hom [symmetric] new-word-of-int-hom-syms*)

**done**

**lemma** *of-bl-False* [*simp*]:

*of-bl* (*False*#*xs*) = *of-bl* *xs*

**by** (*rule word-eqI*)

(*auto simp add: test-bit-of-bl nth-append*)

**lemma** *of-bl-True*:

(*of-bl* (*True*#*xs*)::'*a*::len *word*) =  $2^{length\ xs}$  + *of-bl* *xs*

**by** (*subst of-bl-append [where xs=[True], simplified]*)

(*simp add: word-1-bl*)

**lemma** *of-bl-Cons*:

*of-bl* (*x*#*xs*) = *of-bool* *x* \*  $2^{length\ xs}$  + *of-bl* *xs*

**by** (*cases x*) (*simp-all add: of-bl-True*)

```

lemma split-uint-lem: bin-split n (uint (w :: 'a :: len0 word)) = (a, b) ==>
  a = bintrunc (len-of TYPE('a) - n) a & b = bintrunc (len-of TYPE('a)) b
apply (frule word-ubin.norm-Rep [THEN ssubst])
apply (drule bin-split-trunc1)
apply (drule sym [THEN trans])
apply assumption
apply safe
done

```

```

lemma word-split-bl':
  std = size c - size b ==> (word-split c = (a, b)) ==>
    (a = of-bl (take std (to-bl c)) & b = of-bl (drop std (to-bl c)))
apply (unfold word-split-bin')
apply safe
defer
apply (clarsimp split: prod.splits)
apply (drule word-ubin.norm-Rep [THEN ssubst])
apply (drule split-bintrunc)
apply (simp add : of-bl-def bl2bin-drop word-size
  word-ubin.norm-eq-iff [symmetric] min-def del : word-ubin.norm-Rep)
apply (clarsimp split: prod.splits)
apply (frule split-uint-lem [THEN conjunct1])
apply (unfold word-size)
apply (cases len-of TYPE('a) >= len-of TYPE('b))
defer
apply (simp add: word-0-bl word-0-wi-Pls)
apply (simp add : of-bl-def to-bl-def)
apply (subst bin-split-take1 [symmetric])
prefer 2
apply assumption
apply simp
apply (erule thin-rl)
apply (erule arg-cong [THEN trans])
apply (simp add : word-ubin.norm-eq-iff [symmetric])
done

```

```

lemma word-split-bl: std = size c - size b ==>
  (a = of-bl (take std (to-bl c)) & b = of-bl (drop std (to-bl c))) <->
  word-split c = (a, b)
apply (rule iffI)
defer
apply (erule (1) word-split-bl')
apply (case-tac word-split c)
apply (auto simp add : word-size)
apply (frule word-split-bl' [rotated])
apply (auto simp add : word-size)
done

```

```

lemma word-split-bl-eq:

```



```

(word-split (c::'a::len word) :: ('c :: len0 word * 'd :: len0 word)) =
  (of-bl (take (len-of TYPE('a::len) - len-of TYPE('d::len0)) (to-bl c)),
   of-bl (drop (len-of TYPE('a) - len-of TYPE('d)) (to-bl c)))
apply (rule word-split-bl [THEN iffD1])
apply (unfold word-size)
apply (rule refl conjI)+
done

```

— keep quantifiers for use in simplification

```

lemma test-bit-split':
  word-split c = (a, b) --> (ALL n m. b !! n = (n < size b & c !! n) &
    a !! m = (m < size a & c !! (m + size b)))
apply (unfold word-split-bin' test-bit-bin)
apply (clarify)
apply (clarsimp simp: word-ubin.eq-norm nth-bintr word-size split: prod.splits)
apply (drule bin-nth-split)
apply safe
  apply (simp-all add: add-commute)
apply (erule bin-nth-uint-imp)+
done

```

```

lemmas test-bit-split =
  test-bit-split' [THEN mp, simplified all-simps, standard]

```

```

lemma test-bit-split-eq: word-split c = (a, b) <->
  ((ALL n::nat. b !! n = (n < size b & c !! n)) &
   (ALL m::nat. a !! m = (m < size a & c !! (m + size b))))
apply (rule-tac iffI)
apply (rule-tac conjI)
  apply (erule test-bit-split [THEN conjunct1])
  apply (erule test-bit-split [THEN conjunct2])
apply (case-tac word-split c)
apply (frule test-bit-split)
apply (erule trans)
apply (fastsimp intro ! : word-eqI simp add : word-size)
done

```

— this odd result is analogous to ucast\_id, result to the length given by the result type

```

lemma word-cat-id: word-cat a b = b
  unfolding word-cat-bin' by (simp add: word-ubin.inverse-norm)

```

— limited hom result

```

lemma word-cat-hom:
  len-of TYPE('a::len0) <= len-of TYPE('b::len0) + len-of TYPE('c::len0)
  ==>
  (word-cat (word-of-int w :: 'b word) (b :: 'c word) :: 'a word) =
  word-of-int (bin-cat w (size b) (uint b))

```

```

apply (unfold word-cat-def word-size)
apply (clarsimp simp add : word-ubin.norm-eq-iff [symmetric]
      word-ubin.eq-norm bintr-cat min-def)
apply arith
done

```

```

lemma word-cat-split-alt:
  size w <= size u + size v ==> word-split w = (u, v) ==> word-cat u v = w
apply (rule word-eqI)
apply (drule test-bit-split)
apply (clarsimp simp add : test-bit-cat word-size)
apply safe
apply arith
done

```

```

lemmas word-cat-split-size =
  sym [THEN [2] word-cat-split-alt [symmetric], standard]

```

### 13.2.1 Split and slice

```

lemma split-slices:
  word-split w = (u, v) ==> u = slice (size v) w & v = slice 0 w
apply (drule test-bit-split)
apply (rule conjI)
apply (rule word-eqI, clarsimp simp: nth-slice word-size)+
done

```

```

lemma slice-cat1':
  wc = word-cat a b ==> size wc >= size a + size b ==> slice (size b) wc = a
apply safe
apply (rule word-eqI)
apply (simp add: nth-slice test-bit-cat word-size)
done

```

```

lemmas slice-cat1 = refl [THEN slice-cat1']
lemmas slice-cat2 = trans [OF slice-id word-cat-id]

```

```

lemma cat-slices:
  a = slice n c ==> b = slice 0 c ==> n = size b ==>
    size a + size b >= size c ==> word-cat a b = c
apply safe
apply (rule word-eqI)
apply (simp add: nth-slice test-bit-cat word-size)
apply safe
apply arith
done

```

```

lemma word-split-cat-alt:
  w = word-cat u v ==> size u + size v <= size w ==> word-split w = (u, v)

```

```

apply (case-tac word-split ?w)
apply (rule trans, assumption)
apply (drule test-bit-split)
apply safe
apply (rule word-eqI, clarsimp simp: test-bit-cat word-size)+
done

```

```

lemmas word-cat-bl-no-bin [simp] =
  word-cat-bl [where a=number-of a
    and b=number-of b,
    unfolded to-bl-no-bin, standard]

```

```

lemmas word-split-bl-no-bin [simp] =
  word-split-bl-eq [where c=number-of c, unfolded to-bl-no-bin, standard]

```

— this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

```

lemma word-rsplit-same: word-rsplit w = [w]
  unfolding word-rsplit-def by (simp add : bin-rsplit-all)

```

```

lemma word-rsplit-empty-iff-size:
  (word-rsplit w = []) = (size w = 0)
  unfolding word-rsplit-def bin-rsplit-def word-size
  by (simp add: bin-rsplit-aux-simp-alt Let-def split: split-split)

```

```

lemma test-bit-rsplit:
  sw = word-rsplit w ==> m < size (hd sw :: 'a :: len word) ==>
    k < length sw ==> (rev sw ! k) !! m = (w !! (k * size (hd sw) + m))
  apply (unfold word-rsplit-def word-test-bit-def)
  apply (rule trans)
  apply (rule-tac f = %x. bin-nth x m in arg-cong)
  apply (rule nth-map [symmetric])
  apply simp
  apply (rule bin-nth-rsplit)
  apply simp-all
  apply (simp add : word-size rev-map map-compose [symmetric])
  apply (rule trans)
  defer
  apply (rule map-ident [THEN fun-cong])
  apply (rule refl [THEN map-cong])
  apply (simp add : word-ubin.eq-norm)
  apply (erule bin-rsplit-size-sign [OF len-gt-0 refl])
  done

```

```

lemma word-rcat-bl: word-rcat wl == of-bl (concat (map to-bl wl))
  unfolding word-rcat-def to-bl-def' of-bl-def
  by (clarsimp simp add : bin-rcat-bl map-compose)

```

```

lemma size-rcat-lem':
  size (concat (map to-bl wl)) = length wl * size (hd wl)
  unfolding word-size by (induct wl) auto

lemmas size-rcat-lem = size-rcat-lem' [unfolded word-size]

lemmas td-gal-lt-len = len-gt-0 [THEN td-gal-lt, standard]

lemma nth-rcat-lem' [rule-format] :
  sw = size (hd wl :: 'a :: len word) ==> (ALL n. n < size wl * sw -->
    rev (concat (map to-bl wl)) ! n =
    rev (to-bl (rev wl ! (n div sw))) ! (n mod sw))
  apply (unfold word-size)
  apply (induct wl)
  apply clarsimp
  apply (clarsimp simp add : nth-append size-rcat-lem)
  apply (simp (no-asm-use) only: mult-Suc [symmetric]
    td-gal-lt-len less-Suc-eq-le mod-div-equality')
  apply clarsimp
  done

lemmas nth-rcat-lem = refl [THEN nth-rcat-lem', unfolded word-size]

lemma test-bit-rcat:
  sw = size (hd wl :: 'a :: len word) ==> rc = word-rcat wl ==> rc !! n =
    (n < size rc & n div sw < size wl & (rev wl) ! (n div sw) !! (n mod sw))
  apply (unfold word-rcat-bl word-size)
  apply (clarsimp simp add :
    test-bit-of-bl size-rcat-lem word-size td-gal-lt-len)
  apply safe
  apply (auto simp add :
    test-bit-bl word-size td-gal-lt-len [THEN iffD2, THEN nth-rcat-lem])
  done

lemma foldl-eq-foldr [rule-format] :
  ALL x. foldl op + x xs = foldr op + (x # xs) (0 :: 'a :: comm-monoid-add)
  by (induct xs) (auto simp add : add-assoc)

lemmas test-bit-cong = arg-cong [where f = test-bit, THEN fun-cong]

lemmas test-bit-rsplit-alt =
  trans [OF nth-rev-alt [THEN test-bit-cong]
    test-bit-rsplit [OF refl asm-rl diff-Suc-less]]

— lazy way of expressing that u and v, and su and sv, have same types
lemma word-rsplit-len-indep':
  [u,v] = p ==> [su,sv] = q ==> word-rsplit u = su ==>
    word-rsplit v = sv ==> length su = length sv

```

```

apply (unfold word-rsplit-def)
apply (auto simp add : bin-rsplit-len-indep)
done

```

```

lemmas word-rsplit-len-indep = word-rsplit-len-indep' [OF refl refl refl refl]

```

```

lemma length-word-rsplit-size:
  n = len-of TYPE ('a :: len) ==>
    (length (word-rsplit w :: 'a word list) <= m) = (size w <= m * n)
apply (unfold word-rsplit-def word-size)
apply (clarsimp simp add : bin-rsplit-len-le)
done

```

```

lemmas length-word-rsplit-lt-size =
  length-word-rsplit-size [unfolded Not-eq-iff linorder-not-less [symmetric]]

```

```

lemma length-word-rsplit-exp-size:
  n = len-of TYPE ('a :: len) ==>
    length (word-rsplit w :: 'a word list) = (size w + n - 1) div n
unfolding word-rsplit-def by (clarsimp simp add : word-size bin-rsplit-len)

```

```

lemma length-word-rsplit-even-size:
  n = len-of TYPE ('a :: len) ==> size w = m * n ==>
    length (word-rsplit w :: 'a word list) = m
by (clarsimp simp add : length-word-rsplit-exp-size given-quot-alt)

```

```

lemmas length-word-rsplit-exp-size' = refl [THEN length-word-rsplit-exp-size]

```

```

lemmas tdle = iffD2 [OF split-div-lemma refl, THEN conjunct1]
lemmas dtle = xtr4 [OF tdle mult-commute]

```

```

lemma word-rcat-rsplit: word-rcat (word-rsplit w) = w
apply (rule word-eqI)
apply (clarsimp simp add : test-bit-rcat word-size)
apply (subst refl [THEN test-bit-rsplit])
apply (simp-all add: word-size
  refl [THEN length-word-rsplit-size [simplified le-def, simplified]])
apply safe
apply (erule xtr7, rule len-gt-0 [THEN dtle]) +
done

```

```

lemma size-word-rsplit-rcat-size':
  word-rcat (ws :: 'a :: len word list) = frcw ==>
    size frcw = length ws * len-of TYPE ('a) ==>
    size (hd [word-rsplit frcw, ws]) = size ws
apply (clarsimp simp add : word-size length-word-rsplit-exp-size')
apply (fast intro: given-quot-alt)
done

```

```

lemmas size-word-rsplit-rcat-size =
  size-word-rsplit-rcat-size' [simplified]

```

```

lemma msreus:
  fixes n::nat
  shows  $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$ 
  and  $(k * n + m) \text{ mod } n = m \text{ mod } n$ 
  by (auto simp: add-commute)

```

```

lemma word-rsplit-rcat-size':
  word-rcat (ws :: 'a :: len word list) = frcw ==>
    size frcw = length ws * len-of TYPE ('a) ==> word-rsplit frcw = ws
  apply (frule size-word-rsplit-rcat-size, assumption)
  apply (clarsimp simp add: word-size)
  apply (rule nth-equalityI, assumption)
  applyclarsimp
  apply (rule word-eqI)
  apply (rule trans)
  apply (rule test-bit-rsplit-alt)
  apply (clarsimp simp: word-size)+
  apply (rule trans)
  apply (rule test-bit-rcat [OF refl refl])
  apply (simp add: word-size msreus)
  apply (subst nth-rev)
  applyarith
  apply (simp add: le0 [THEN [2] xtr7, THEN diff-Suc-less])
  apply safe
  apply (simp add: diff-mult-distrib)
  apply (rule mpl-lem)
  apply (cases size ws)
  apply simp-all
done

```

```

lemmas word-rsplit-rcat-size = refl [THEN word-rsplit-rcat-size']

```

### 13.3 Rotation

```

lemmas rotater-0' [simp] = rotater-def [where n = 0, simplified]

```

```

lemmas word-rot-defs = word-roti-def word-rotr-def word-rotl-def

```

```

lemma rotate-eq-mod:
   $m \text{ mod } \text{length } xs = n \text{ mod } \text{length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$ 
  apply (rule box-equals)
  defer
  apply (rule rotate-conv-mod [symmetric])+
  apply simp
done

```

```

lemmas rotate-eqs [standard] =
  trans [OF rotate0 [THEN fun-cong] id-apply]
  rotate-rotate [symmetric]
  rotate-id
  rotate-conv-mod
  rotate-eq-mod

```

### 13.3.1 Rotation of list to right

```

lemma rotate1-rl': rotater1 (l @ [a]) = a # l
  unfolding rotater1-def by (cases l) auto

```

```

lemma rotate1-rl [simp] : rotater1 (rotate1 l) = l
  apply (unfold rotater1-def)
  apply (cases l)
  apply (case-tac [2] list)
  apply auto
  done

```

```

lemma rotate1-lr [simp] : rotate1 (rotater1 l) = l
  unfolding rotater1-def by (cases l) auto

```

```

lemma rotater1-rev': rotater1 (rev xs) = rev (rotate1 xs)
  apply (cases xs)
  apply (simp add : rotater1-def)
  apply (simp add : rotate1-rl')
  done

```

```

lemma rotater-rev': rotater n (rev xs) = rev (rotate n xs)
  unfolding rotater-def by (induct n) (auto intro: rotater1-rev')

```

```

lemmas rotater-rev = rotater-rev' [where xs = rev ys, simplified, standard]

```

```

lemma rotater-drop-take:
  rotater n xs =
    drop (length xs - n mod length xs) xs @
    take (length xs - n mod length xs) xs
  by (clarsimp simp add : rotater-rev rotate-drop-take rev-take rev-drop)

```

```

lemma rotater-Suc [simp] :
  rotater (Suc n) xs = rotater1 (rotater n xs)
  unfolding rotater-def by auto

```

```

lemma rotate-inv-plus [rule-format] :
  ALL k. k = m + n --> rotater k (rotate n xs) = rotater m xs &
  rotate k (rotater n xs) = rotate m xs &
  rotater n (rotate k xs) = rotate m xs &
  rotate n (rotater k xs) = rotater m xs

```

**unfolding** *rotater-def rotate-def*  
**by** (*induct n*) (*auto intro: funpow-swap1 [THEN trans]*)

**lemmas** *rotate-inv-rel = le-add-diff-inverse2 [symmetric, THEN rotate-inv-plus]*

**lemmas** *rotate-inv-eq = order-refl [THEN rotate-inv-rel, simplified]*

**lemmas** *rotate-lr [simp] = rotate-inv-eq [THEN conjunct1, standard]*  
**lemmas** *rotate-rl [simp] =*  
*rotate-inv-eq [THEN conjunct2, THEN conjunct1, standard]*

**lemma** *rotate-gal: (rotater n xs = ys) = (rotate n ys = xs)*  
**by** *auto*

**lemma** *rotate-gal': (ys = rotater n xs) = (xs = rotate n ys)*  
**by** *auto*

**lemma** *length-rotater [simp]:*  
*length (rotater n xs) = length xs*  
**by** (*simp add : rotater-rev*)

**lemmas** *rrs0 = rotate-eqs [THEN restrict-to-left,*  
*simplified rotate-gal [symmetric] rotate-gal' [symmetric], standard]*  
**lemmas** *rrs1 = rrs0 [THEN refl [THEN rev-iffD1]]*  
**lemmas** *rotater-eqs = rrs1 [simplified length-rotater, standard]*  
**lemmas** *rotater-0 = rotater-eqs (1)*  
**lemmas** *rotater-add = rotater-eqs (2)*

### 13.3.2 map, app2, commuting with rotate(r)

**lemma** *last-map: xs ~ = [] ==> last (map f xs) = f (last xs)*  
**by** (*induct xs*) *auto*

**lemma** *butlast-map:*  
*xs ~ = [] ==> butlast (map f xs) = map f (butlast xs)*  
**by** (*induct xs*) *auto*

**lemma** *rotater1-map: rotater1 (map f xs) = map f (rotater1 xs)*  
**unfolding** *rotater1-def*  
**by** (*cases xs*) (*auto simp add: last-map butlast-map*)

**lemma** *rotater-map:*  
*rotater n (map f xs) = map f (rotater n xs)*  
**unfolding** *rotater-def*  
**by** (*induct n*) (*auto simp add : rotater1-map*)

**lemma** *but-last-zip [rule-format] :*  
*ALL ys. length xs = length ys --> xs ~ = [] -->*  
*last (zip xs ys) = (last xs, last ys) &*



```

    butlast (zip xs ys) = zip (butlast xs) (butlast ys)
  apply (induct xs)
  apply auto
  apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
done

```

```

lemma but-last-app2 [rule-format] :
  ALL ys. length xs = length ys ==> xs ~ = [] ==>
    last (app2 f xs ys) = f (last xs) (last ys) &
    butlast (app2 f xs ys) = app2 f (butlast xs) (butlast ys)
  apply (induct xs)
  apply auto
  apply (unfold app2-def)
  apply ((case-tac ys, auto simp: neq-Nil-conv)[1])+
done

```

```

lemma rotater1-zip:
  length xs = length ys ==>
    rotater1 (zip xs ys) = zip (rotater1 xs) (rotater1 ys)
  apply (unfold rotater1-def)
  apply (cases xs)
  apply auto
  apply ((case-tac ys, auto simp: neq-Nil-conv but-last-zip)[1])+
done

```

```

lemma rotater1-app2:
  length xs = length ys ==>
    rotater1 (app2 f xs ys) = app2 f (rotater1 xs) (rotater1 ys)
  unfolding app2-def by (simp add: rotater1-map rotater1-zip)

```

```

lemmas lrth =
  box-equals [OF asm-rl length-rotater [symmetric]
    length-rotater [symmetric],
    THEN rotater1-app2]

```

```

lemma rotater-app2:
  length xs = length ys ==>
    rotater n (app2 f xs ys) = app2 f (rotater n xs) (rotater n ys)
  by (induct n) (auto intro!: lrth)

```

```

lemma rotate1-app2:
  length xs = length ys ==>
    rotate1 (app2 f xs ys) = app2 f (rotate1 xs) (rotate1 ys)
  apply (unfold app2-def)
  apply (cases xs)
  apply (cases ys, auto simp add : rotate1-def)+
done

```

```

lemmas lth = box-equals [OF asm-rl length-rotate [symmetric]

```

*length-rotate [symmetric], THEN rotate1-app2]*

**lemma** *rotate-app2*:

*length xs = length ys ==>*

*rotate n (app2 f xs ys) = app2 f (rotate n xs) (rotate n ys)*

**by** (*induct n*) (*auto intro! lth*)

— corresponding equalities for word rotation

**lemma** *to-bl-rotl*:

*to-bl (word-rotl n w) = rotate n (to-bl w)*

**by** (*simp add: word-bl.Abs-inverse' word-rotl-def*)

**lemmas** *blrs0 = rotate-egs [THEN to-bl-rotl [THEN trans]]*

**lemmas** *word-rotl-egs =*

*blrs0 [simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotl [symmetric]]*

**lemma** *to-bl-rotr*:

*to-bl (word-rotr n w) = rotater n (to-bl w)*

**by** (*simp add: word-bl.Abs-inverse' word-rotr-def*)

**lemmas** *brrs0 = rotater-egs [THEN to-bl-rotr [THEN trans]]*

**lemmas** *word-rotr-egs =*

*brrs0 [simplified word-bl.Rep' word-bl.Rep-inject to-bl-rotr [symmetric]]*

**declare** *word-rotr-egs (1) [simp]*

**declare** *word-rotl-egs (1) [simp]*

**lemma**

*word-rot-rl [simp]:*

*word-rotl k (word-rotr k v) = v and*

*word-rot-lr [simp]:*

*word-rotr k (word-rotl k v) = v*

**by** (*auto simp add: to-bl-rotr to-bl-rotl word-bl.Rep-inject [symmetric]*)

**lemma**

*word-rot-gal:*

*(word-rotr n v = w) = (word-rotl n w = v) and*

*word-rot-gal':*

*(w = word-rotr n v) = (v = word-rotl n w)*

**by** (*auto simp: to-bl-rotr to-bl-rotl word-bl.Rep-inject [symmetric]*  
*dest: sym*)

**lemma** *word-rotr-rev:*

*word-rotr n w = word-reverse (word-rotl n (word-reverse w))*

**by** (*simp add: word-bl.Rep-inject [symmetric] to-bl-word-rev*)

*to-bl-rotr to-bl-rotl rotater-rev*)

**lemma** *word-roti-0* [*simp*]: *word-roti 0 w = w*  
**by** (*unfold word-rot-defs*) *auto*

**lemmas** *abl-cong* = *arg-cong* [**where** *f* = *of-bl*]

**lemma** *word-roti-add*:

*word-roti (m + n) w = word-roti m (word-roti n w)*

**proof** –

**have** *rotater-eq-lem*:

$\bigwedge m\ n\ xs.\ m = n \implies \text{rotater } m\ xs = \text{rotater } n\ xs$   
**by** *auto*

**have** *rotate-eq-lem*:

$\bigwedge m\ n\ xs.\ m = n \implies \text{rotate } m\ xs = \text{rotate } n\ xs$   
**by** *auto*

**note** *rpts* [*symmetric, standard*] =

*rotate-inv-plus* [*THEN conjunct1*]

*rotate-inv-plus* [*THEN conjunct2, THEN conjunct1*]

*rotate-inv-plus* [*THEN conjunct2, THEN conjunct2, THEN conjunct1*]

*rotate-inv-plus* [*THEN conjunct2, THEN conjunct2, THEN conjunct2*]

**note** *rrp* = *trans* [*symmetric, OF rotate-rotate rotate-eq-lem*]

**note** *rrrp* = *trans* [*symmetric, OF rotater-add* [*symmetric*] *rotater-eq-lem*]

**show** *?thesis*

**apply** (*unfold word-rot-defs*)

**apply** (*simp only: split: split-if*)

**apply** (*safe intro!: abl-cong*)

**apply** (*simp-all only: to-bl-rotl* [*THEN word-bl.Rep-inverse*]

*to-bl-rotl*

*to-bl-rotr* [*THEN word-bl.Rep-inverse*]

*to-bl-rotr*)

**apply** (*rule rrp rrrp rpts,*

*simp add: nat-add-distrib* [*symmetric*]

*nat-diff-distrib* [*symmetric*])+

**done**

**qed**

**lemma** *word-roti-conv-mod*!: *word-roti n w = word-roti (n mod int (size w)) w*

**apply** (*unfold word-rot-defs*)

**apply** (*cut-tac y=size w in gt-or-eq-0*)

**apply** (*erule disjE*)

**apply** *simp-all*

**apply** (*safe intro!: abl-cong*)

**apply** (*rule rotater-eqs*)

**apply** (*simp add: word-size nat-mod-distrib*)

```

apply (simp add: rotater-add [symmetric] rotate-gal [symmetric])
apply (rule rotater-eqs)
apply (simp add: word-size nat-mod-distrib)
apply (rule int-eq-0-conv [THEN iffD1])
apply (simp only: zmod-int zadd-int [symmetric])
apply (simp add: rdmodes)
done

```

**lemmas** *word-roti-conv-mod* = *word-roti-conv-mod'* [unfolded word-size]

### 13.3.3 Word rotation commutes with bit-wise operations

**locale** *word-rotate*

**context** *word-rotate*

**begin**

**lemmas** *word-rot-defs'* = *to-bl-rotl to-bl-rotr*

**lemmas** *blwl-syms* [symmetric] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

**lemmas** *lbl-lbl* = *trans* [OF *word-bl.Rep'* *word-bl.Rep'* [symmetric]]

**lemmas** *ths-app2* [OF *lbl-lbl*] = *rotate-app2 rotater-app2*

**lemmas** *ths-map* [where *xs* = *to-bl v*] = *rotate-map rotater-map*

**lemmas** *th1s* [simplified *word-rot-defs'* [symmetric]] = *ths-app2 ths-map*

**lemma** *word-rot-logs*:

```

word-rotl n (NOT v) = NOT word-rotl n v
word-rotr n (NOT v) = NOT word-rotr n v
word-rotl n (x AND y) = word-rotl n x AND word-rotl n y
word-rotr n (x AND y) = word-rotr n x AND word-rotr n y
word-rotl n (x OR y) = word-rotl n x OR word-rotl n y
word-rotr n (x OR y) = word-rotr n x OR word-rotr n y
word-rotl n (x XOR y) = word-rotl n x XOR word-rotl n y
word-rotr n (x XOR y) = word-rotr n x XOR word-rotr n y
by (rule word-bl.Rep-eqD,
      rule word-rot-defs' [THEN trans],
      simp only: blwl-syms [symmetric],
      rule th1s [THEN trans],
      rule refl) +

```

**end**

**lemmas** *word-rot-logs* = *word-rotate.word-rot-logs*

**lemmas** *bl-word-rotl-dt* = *trans* [OF *to-bl-rotl rotate-drop-take*,  
simplified *word-bl.Rep'*, *standard*]

```

lemmas bl-word-rotr-dt = trans [OF to-bl-rotr rotater-drop-take,
  simplified word-bl.Rep', standard]

lemma bl-word-roti-dt':
  n = nat ((- i) mod int (size (w :: 'a :: len word))) ==>
    to-bl (word-roti i w) = drop n (to-bl w) @ take n (to-bl w)
apply (unfold word-roti-def)
apply (simp add: bl-word-rotl-dt bl-word-rotr-dt word-size)
apply safe
apply (simp add: zmod-zminus1-eq-if)
apply safe
apply (simp add: nat-mult-distrib)
apply (simp add: nat-diff-distrib [OF pos-mod-sign pos-mod-conj
  [THEN conjunct2, THEN order-less-imp-le]]
  nat-mod-distrib)
apply (simp add: nat-mod-distrib)
done

lemmas bl-word-roti-dt = bl-word-roti-dt' [unfolded word-size]

lemmas word-rotl-dt = bl-word-rotl-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-rotr-dt = bl-word-rotr-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]
lemmas word-roti-dt = bl-word-roti-dt
  [THEN word-bl.Rep-inverse' [symmetric], standard]

lemma word-rotx-0 [simp] : word-rotr i 0 = 0 & word-rotl i 0 = 0
  by (simp add : word-rotr-dt word-rotl-dt to-bl-0 replicate-add [symmetric])

lemma word-roti-0' [simp] : word-roti n 0 = 0
  unfolding word-roti-def by auto

lemmas word-rotr-dt-no-bin' [simp] =
  word-rotr-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

lemmas word-rotl-dt-no-bin' [simp] =
  word-rotl-dt [where w=number-of w, unfolded to-bl-no-bin, standard]

declare word-roti-def [simp]

end

```

## 14 Boolean-Algebra: Boolean Algebras

**theory** *Boolean-Algebra*

```

imports Main
begin

locale boolean =
  fixes conj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\sqcap$  70)
  fixes disj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\sqcup$  65)
  fixes compl :: 'a  $\Rightarrow$  'a ( $\sim$  - [81] 80)
  fixes zero :: 'a (0)
  fixes one  :: 'a (1)
  assumes conj-assoc:  $(x \sqcap y) \sqcap z = x \sqcap (y \sqcap z)$ 
  assumes disj-assoc:  $(x \sqcup y) \sqcup z = x \sqcup (y \sqcup z)$ 
  assumes conj-commute:  $x \sqcap y = y \sqcap x$ 
  assumes disj-commute:  $x \sqcup y = y \sqcup x$ 
  assumes conj-disj-distrib:  $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$ 
  assumes disj-conj-distrib:  $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$ 
  assumes conj-one-right [simp]:  $x \sqcap \mathbf{1} = x$ 
  assumes disj-zero-right [simp]:  $x \sqcup \mathbf{0} = x$ 
  assumes conj-cancel-right [simp]:  $x \sqcap \sim x = \mathbf{0}$ 
  assumes disj-cancel-right [simp]:  $x \sqcup \sim x = \mathbf{1}$ 
begin

lemmas disj-ac =
  disj-assoc disj-commute
  mk-left-commute [where 'a = 'a, of disj, OF disj-assoc disj-commute]

lemmas conj-ac =
  conj-assoc conj-commute
  mk-left-commute [where 'a = 'a, of conj, OF conj-assoc conj-commute]

lemma dual: boolean disj conj compl one zero
apply (rule boolean.intro)
apply (rule disj-assoc)
apply (rule conj-assoc)
apply (rule disj-commute)
apply (rule conj-commute)
apply (rule disj-conj-distrib)
apply (rule conj-disj-distrib)
apply (rule disj-zero-right)
apply (rule conj-one-right)
apply (rule disj-cancel-right)
apply (rule conj-cancel-right)
done

14.1 Complement

lemma complement-unique:
  assumes 1:  $a \sqcap x = \mathbf{0}$ 
  assumes 2:  $a \sqcup x = \mathbf{1}$ 
  assumes 3:  $a \sqcap y = \mathbf{0}$ 

```

**assumes**  $4: a \sqcup y = 1$   
**shows**  $x = y$   
**proof** –  
**have**  $(a \sqcap x) \sqcup (x \sqcap y) = (a \sqcap y) \sqcup (x \sqcap y)$  **using** 1 3 **by** *simp*  
**hence**  $(x \sqcap a) \sqcup (x \sqcap y) = (y \sqcap a) \sqcup (y \sqcap x)$  **using** *conj-commute* **by** *simp*  
**hence**  $x \sqcap (a \sqcup y) = y \sqcap (a \sqcup x)$  **using** *conj-disj-distrib* **by** *simp*  
**hence**  $x \sqcap 1 = y \sqcap 1$  **using** 2 4 **by** *simp*  
**thus**  $x = y$  **using** *conj-one-right* **by** *simp*  
**qed**

**lemma** *compl-unique*:  $\llbracket x \sqcap y = 0; x \sqcup y = 1 \rrbracket \implies \sim x = y$   
**by** (*rule complement-unique* [*OF conj-cancel-right disj-cancel-right*])

**lemma** *double-compl* [*simp*]:  $\sim(\sim x) = x$   
**proof** (*rule compl-unique*)  
**from** *conj-cancel-right* **show**  $\sim x \sqcap x = 0$  **by** (*simp only: conj-commute*)  
**from** *disj-cancel-right* **show**  $\sim x \sqcup x = 1$  **by** (*simp only: disj-commute*)  
**qed**

**lemma** *compl-eq-compl-iff* [*simp*]:  $(\sim x = \sim y) = (x = y)$   
**by** (*rule inj-eq* [*OF inj-on-inverseI*], *rule double-compl*)

## 14.2 Conjunction

**lemma** *conj-absorb* [*simp*]:  $x \sqcap x = x$   
**proof** –  
**have**  $x \sqcap x = (x \sqcap x) \sqcup 0$  **using** *disj-zero-right* **by** *simp*  
**also have**  $\dots = (x \sqcap x) \sqcup (x \sqcap \sim x)$  **using** *conj-cancel-right* **by** *simp*  
**also have**  $\dots = x \sqcap (x \sqcup \sim x)$  **using** *conj-disj-distrib* **by** (*simp only:*)  
**also have**  $\dots = x \sqcap 1$  **using** *disj-cancel-right* **by** *simp*  
**also have**  $\dots = x$  **using** *conj-one-right* **by** *simp*  
**finally show** *?thesis* .  
**qed**

**lemma** *conj-zero-right* [*simp*]:  $x \sqcap 0 = 0$   
**proof** –  
**have**  $x \sqcap 0 = x \sqcap (x \sqcap \sim x)$  **using** *conj-cancel-right* **by** *simp*  
**also have**  $\dots = (x \sqcap x) \sqcap \sim x$  **using** *conj-assoc* **by** (*simp only:*)  
**also have**  $\dots = x \sqcap \sim x$  **using** *conj-absorb* **by** *simp*  
**also have**  $\dots = 0$  **using** *conj-cancel-right* **by** *simp*  
**finally show** *?thesis* .  
**qed**

**lemma** *compl-one* [*simp*]:  $\sim 1 = 0$   
**by** (*rule compl-unique* [*OF conj-zero-right disj-zero-right*])

**lemma** *conj-zero-left* [*simp*]:  $0 \sqcap x = 0$   
**by** (*subst conj-commute*) (*rule conj-zero-right*)

**lemma** *conj-one-left* [*simp*]:  $\mathbf{1} \sqcap x = x$   
**by** (*subst conj-commute*) (*rule conj-one-right*)

**lemma** *conj-cancel-left* [*simp*]:  $\sim x \sqcap x = \mathbf{0}$   
**by** (*subst conj-commute*) (*rule conj-cancel-right*)

**lemma** *conj-left-absorb* [*simp*]:  $x \sqcap (x \sqcap y) = x \sqcap y$   
**by** (*simp only: conj-assoc [symmetric] conj-absorb*)

**lemma** *conj-disj-distrib2*:  
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$   
**by** (*simp only: conj-commute conj-disj-distrib*)

**lemmas** *conj-disj-distribs* =  
*conj-disj-distrib conj-disj-distrib2*

### 14.3 Disjunction

**lemma** *disj-absorb* [*simp*]:  $x \sqcup x = x$   
**by** (*rule boolean.conj-absorb [OF dual]*)

**lemma** *disj-one-right* [*simp*]:  $x \sqcup \mathbf{1} = \mathbf{1}$   
**by** (*rule boolean.conj-zero-right [OF dual]*)

**lemma** *compl-zero* [*simp*]:  $\sim \mathbf{0} = \mathbf{1}$   
**by** (*rule boolean.compl-one [OF dual]*)

**lemma** *disj-zero-left* [*simp*]:  $\mathbf{0} \sqcup x = x$   
**by** (*rule boolean.conj-one-left [OF dual]*)

**lemma** *disj-one-left* [*simp*]:  $\mathbf{1} \sqcup x = \mathbf{1}$   
**by** (*rule boolean.conj-zero-left [OF dual]*)

**lemma** *disj-cancel-left* [*simp*]:  $\sim x \sqcup x = \mathbf{1}$   
**by** (*rule boolean.conj-cancel-left [OF dual]*)

**lemma** *disj-left-absorb* [*simp*]:  $x \sqcup (x \sqcup y) = x \sqcup y$   
**by** (*rule boolean.conj-left-absorb [OF dual]*)

**lemma** *disj-conj-distrib2*:  
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$   
**by** (*rule boolean.conj-disj-distrib2 [OF dual]*)

**lemmas** *disj-conj-distribs* =  
*disj-conj-distrib disj-conj-distrib2*

### 14.4 De Morgan’s Laws

**lemma** *de-Morgan-conj* [*simp*]:  $\sim (x \sqcap y) = \sim x \sqcup \sim y$   
**proof** (*rule compl-unique*)



```

have  $(x \sqcap y) \sqcap (\sim x \sqcup \sim y) = ((x \sqcap y) \sqcap \sim x) \sqcup ((x \sqcap y) \sqcap \sim y)$ 
  by (rule conj-disj-distrib)
also have  $\dots = (y \sqcap (x \sqcap \sim x)) \sqcup (x \sqcap (y \sqcap \sim y))$ 
  by (simp only: conj-ac)
finally show  $(x \sqcap y) \sqcap (\sim x \sqcup \sim y) = \mathbf{0}$ 
  by (simp only: conj-cancel-right conj-zero-right disj-zero-right)
next
have  $(x \sqcap y) \sqcup (\sim x \sqcup \sim y) = (x \sqcup (\sim x \sqcup \sim y)) \sqcap (y \sqcup (\sim x \sqcup \sim y))$ 
  by (rule disj-conj-distrib2)
also have  $\dots = (\sim y \sqcup (x \sqcup \sim x)) \sqcap (\sim x \sqcup (y \sqcup \sim y))$ 
  by (simp only: disj-ac)
finally show  $(x \sqcap y) \sqcup (\sim x \sqcup \sim y) = \mathbf{1}$ 
  by (simp only: disj-cancel-right disj-one-right conj-one-right)
qed

```

```

lemma de-Morgan-disj [simp]:  $\sim (x \sqcup y) = \sim x \sqcap \sim y$ 
by (rule boolean.de-Morgan-conj [OF dual])

```

**end**

## 14.5 Symmetric Difference

```

locale boolean-xor = boolean +
  fixes xor :: 'a => 'a => 'a (infixr  $\oplus$  65)
  assumes xor-def:  $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$ 
begin

```

```

lemma xor-def2:
   $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$ 
by (simp only: xor-def conj-disj-distribs
      disj-ac conj-ac conj-cancel-right disj-zero-left)

```

```

lemma xor-commute:  $x \oplus y = y \oplus x$ 
by (simp only: xor-def conj-commute disj-commute)

```

```

lemma xor-assoc:  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$ 
proof –
  let  $?t = (x \sqcap y \sqcap z) \sqcup (x \sqcap \sim y \sqcap \sim z) \sqcup$ 
     $(\sim x \sqcap y \sqcap \sim z) \sqcup (\sim x \sqcap \sim y \sqcap z)$ 
  have  $?t \sqcup (z \sqcap x \sqcap \sim x) \sqcup (z \sqcap y \sqcap \sim y) =$ 
     $?t \sqcup (x \sqcap y \sqcap \sim y) \sqcup (x \sqcap z \sqcap \sim z)$ 
  by (simp only: conj-cancel-right conj-zero-right)
  thus  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$ 
  apply (simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl)
  apply (simp only: conj-disj-distribs conj-ac disj-ac)
  done
qed

```

```

lemmas xor-ac =

```

*xor-assoc xor-commute*  
*mk-left-commute* [where 'a = 'a, of xor, OF xor-assoc xor-commute]

**lemma** *xor-zero-right* [simp]:  $x \oplus \mathbf{0} = x$   
**by** (simp only: xor-def compl-zero conj-one-right conj-zero-right disj-zero-right)

**lemma** *xor-zero-left* [simp]:  $\mathbf{0} \oplus x = x$   
**by** (subst xor-commute) (rule xor-zero-right)

**lemma** *xor-one-right* [simp]:  $x \oplus \mathbf{1} = \sim x$   
**by** (simp only: xor-def compl-one conj-zero-right conj-one-right disj-zero-left)

**lemma** *xor-one-left* [simp]:  $\mathbf{1} \oplus x = \sim x$   
**by** (subst xor-commute) (rule xor-one-right)

**lemma** *xor-self* [simp]:  $x \oplus x = \mathbf{0}$   
**by** (simp only: xor-def conj-cancel-right conj-cancel-left disj-zero-right)

**lemma** *xor-left-self* [simp]:  $x \oplus (x \oplus y) = y$   
**by** (simp only: xor-assoc [symmetric] xor-self xor-zero-left)

**lemma** *xor-compl-left*:  $\sim x \oplus y = \sim (x \oplus y)$   
**apply** (simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl)  
**apply** (simp only: conj-disj-distrib)  
**apply** (simp only: conj-cancel-right conj-cancel-left)  
**apply** (simp only: disj-zero-left disj-zero-right)  
**apply** (simp only: disj-ac conj-ac)  
**done**

**lemma** *xor-compl-right*:  $x \oplus \sim y = \sim (x \oplus y)$   
**apply** (simp only: xor-def de-Morgan-disj de-Morgan-conj double-compl)  
**apply** (simp only: conj-disj-distrib)  
**apply** (simp only: conj-cancel-right conj-cancel-left)  
**apply** (simp only: disj-zero-left disj-zero-right)  
**apply** (simp only: disj-ac conj-ac)  
**done**

**lemma** *xor-cancel-right* [simp]:  $x \oplus \sim x = \mathbf{1}$   
**by** (simp only: xor-compl-right xor-self compl-zero)

**lemma** *xor-cancel-left* [simp]:  $\sim x \oplus x = \mathbf{1}$   
**by** (subst xor-commute) (rule xor-cancel-right)

**lemma** *conj-xor-distrib*:  $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

**proof** –

**have**  $(x \sqcap y \sqcap \sim z) \sqcup (x \sqcap \sim y \sqcap z) =$   
 $(y \sqcap x \sqcap \sim z) \sqcup (z \sqcap x \sqcap \sim y) \sqcup (x \sqcap y \sqcap \sim z) \sqcup (x \sqcap \sim y \sqcap z)$   
**by** (simp only: conj-cancel-right conj-zero-right disj-zero-left)  
**thus**  $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$

```

  by (simp (no-asm-use) only:
      xor-def de-Morgan-disj de-Morgan-conj double-compl
      conj-disj-distrib conj-ac disj-ac)
qed

lemma conj-xor-distrib2:
   $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$ 
proof -
  have  $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$ 
  by (rule conj-xor-distrib)
  thus  $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$ 
  by (simp only: conj-commute)
qed

lemmas conj-xor-distrib =
  conj-xor-distrib conj-xor-distrib2

end

end

```

## 15 WordGenLib: Miscellaneous Library for Words

```

theory WordGenLib imports WordShift Boolean-Algebra
begin

```

```

declare of-nat-2p [simp]

```

```

lemma word-int-cases:

```

```

   $\llbracket \bigwedge n. \llbracket (x :: 'a::len\ word) = word-of-int\ n; 0 \leq n; n < 2^{len-of\ TYPE('a)} \rrbracket \implies P \rrbracket$ 
 $\implies P$ 
  by (cases x rule: word-uint.Abs-cases) (simp add: uints-num)

```

```

lemma word-nat-cases [cases type: word]:

```

```

   $\llbracket \bigwedge n. \llbracket (x :: 'a::len\ word) = of-nat\ n; n < 2^{len-of\ TYPE('a)} \rrbracket \implies P \rrbracket$ 
 $\implies P$ 
  by (cases x rule: word-unat.Abs-cases) (simp add: unats-def)

```

```

lemma max-word-eq:

```

```

   $(max-word :: 'a::len\ word) = 2^{len-of\ TYPE('a)} - 1$ 
  by (simp add: max-word-def word-of-int-hom-syms word-of-int-2p)

```

```

lemma max-word-max [simp,intro!]:

```

```

   $n \leq max-word$ 
  by (cases n rule: word-int-cases)
  (simp add: max-word-def word-le-def int-word-uint int-mod-eq')

```

**lemma** *word-of-int-2p-len*:

*word-of-int* ( $2 \wedge \text{len-of TYPE('a)}$ ) = ( $0::'a::\text{len0}$  word)

**by** (*subst word-uint.Abs-norm [symmetric]*)

(*simp add: word-of-int-hom-syms*)

**lemma** *word-pow-0*:

( $2::'a::\text{len}$  word)  $\wedge \text{len-of TYPE('a)}$  = 0

**proof** –

**have** *word-of-int* ( $2 \wedge \text{len-of TYPE('a)}$ ) = ( $0::'a$  word)

**by** (*rule word-of-int-2p-len*)

**thus** ?thesis **by** (*simp add: word-of-int-2p*)

**qed**

**lemma** *max-word-wrap*:  $x + 1 = 0 \implies x = \text{max-word}$

**apply** (*simp add: max-word-eq*)

**apply** *uint-arith*

**apply** *auto*

**apply** (*simp add: word-pow-0*)

**done**

**lemma** *max-word-minus*:

*max-word* = ( $-1::'a::\text{len}$  word)

**proof** –

**have**  $-1 + 1 = (0::'a \text{ word})$  **by** *simp*

**thus** ?thesis **by** (*rule max-word-wrap [symmetric]*)

**qed**

**lemma** *max-word-bl [simp]*:

*to-bl* (*max-word*:: $'a::\text{len}$  word) = *replicate* (*len-of TYPE('a)*) *True*

**by** (*subst max-word-minus to-bl-n1*) **+** *simp*

**lemma** *max-test-bit [simp]*:

(*max-word*:: $'a::\text{len}$  word) !!  $n = (n < \text{len-of TYPE('a)})$

**by** (*auto simp add: test-bit-bl word-size*)

**lemma** *word-and-max [simp]*:

$x \text{ AND } \text{max-word} = x$

**by** (*rule word-eqI*) (*simp add: word-ops-nth-size word-size*)

**lemma** *word-or-max [simp]*:

$x \text{ OR } \text{max-word} = \text{max-word}$

**by** (*rule word-eqI*) (*simp add: word-ops-nth-size word-size*)

**lemma** *word-ao-dist2*:

$x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } (z::'a::\text{len0} \text{ word})$

**by** (*rule word-eqI*) (*auto simp add: word-ops-nth-size word-size*)

**lemma** *word-oa-dist2*:

$x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z::'a::\text{len0} \text{ word}))$

```

by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-and-not [simp]:
   $x \text{ AND } NOT\ x = (0 :: 'a :: len0\ word)$ 
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-or-not [simp]:
   $x \text{ OR } NOT\ x = max\text{-}word$ 
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

lemma word-boolean:
  boolean (op AND) (op OR) bitNOT 0 max-word
apply (rule boolean.intro)
  apply (rule word-bw-assocs)
  apply (rule word-bw-assocs)
  apply (rule word-bw-comms)
  apply (rule word-bw-comms)
  apply (rule word-ao-dist2)
  apply (rule word-ao-dist2)
  apply (rule word-and-max)
  apply (rule word-log-esimps)
  apply (rule word-and-not)
  apply (rule word-or-not)
done

interpretation word-bool-alg:
  boolean [op AND op OR bitNOT 0 max-word]
by (rule word-boolean)

lemma word-xor-and-or:
   $x \text{ XOR } y = x \text{ AND } NOT\ y \text{ OR } NOT\ x \text{ AND } (y :: 'a :: len0\ word)$ 
by (rule word-eqI) (auto simp add: word-ops-nth-size word-size)

interpretation word-bool-alg:
  boolean-xor [op AND op OR bitNOT 0 max-word op XOR]
apply (rule boolean-xor.intro)
  apply (rule word-boolean)
  apply (rule boolean-xor-axioms.intro)
  apply (rule word-xor-and-or)
done

lemma shiftr-0 [iff]:
   $(x :: 'a :: len0\ word) >> 0 = x$ 
by (simp add: shiftr-bl)

lemma shiffl-0 [simp]:
   $(x :: 'a :: len\ word) << 0 = x$ 
by (simp add: shiffl-t2n)

```

```

lemma shiftr1-1 [simp]:
  (1::a::len word) << n = 2^n
  by (simp add: shiftr1-2n)

lemma uint-lt-0 [simp]:
  uint x < 0 = False
  by (simp add: linorder-not-less)

lemma shiftr1-1 [simp]:
  shiftr1 (1::a::len word) = 0
  by (simp add: shiftr1-def word-0-alt)

lemma shiftr-1 [simp]:
  (1::a::len word) >> n = (if n = 0 then 1 else 0)
  by (induct n) (auto simp: shiftr-def)

lemma word-less-1 [simp]:
  ((x::a::len word) < 1) = (x = 0)
  by (simp add: word-less-nat-alt unat-0-iff)

lemma to-bl-mask:
  to-bl (mask n :: a::len word) =
    replicate (len-of TYPE(a) - n) False @
    replicate (min (len-of TYPE(a)) n) True
  by (simp add: mask-bl word-rep-drop min-def)

lemma map-replicate-True:
  n = length xs ==>
    map (λ(x,y). x & y) (zip xs (replicate n True)) = xs
  by (induct xs arbitrary: n) auto

lemma map-replicate-False:
  n = length xs ==> map (λ(x,y). x & y)
    (zip xs (replicate n False)) = replicate n False
  by (induct xs arbitrary: n) auto

lemma bl-and-mask:
  fixes w :: a::len word
  fixes n
  defines n' ≡ len-of TYPE(a) - n
  shows to-bl (w AND mask n) = replicate n' False @ drop n' (to-bl w)
proof -
  note [simp] = map-replicate-True map-replicate-False
  have to-bl (w AND mask n) =
    app2 op & (to-bl w) (to-bl (mask n::a::len word))
    by (simp add: bl-word-and)
  also
  have to-bl w = take n' (to-bl w) @ drop n' (to-bl w) by simp
  also

```

```

have app2 op & ... (to-bl (mask n::'a::len word)) =
  replicate n' False @ drop n' (to-bl w)
unfolding to-bl-mask n'-def app2-def
by (subst zip-append) auto
finally
show ?thesis .
qed

```

```

lemma drop-rev-takefill:
  length xs ≤ n ==>
    drop (n - length xs) (rev (takefill False n (rev xs))) = xs
by (simp add: takefill-alt rev-take)

```

```

lemma map-nth-0 [simp]:
  map (op !! (0::'a::len0 word)) xs = replicate (length xs) False
by (induct xs) auto

```

```

lemma uint-plus-if-size:
  uint (x + y) =
    (if uint x + uint y < 2^size x then
      uint x + uint y
    else
      uint x + uint y - 2^size x)
by (simp add: word-arith-alts int-word-uint mod-add-if-z
  word-size)

```

```

lemma unat-plus-if-size:
  unat (x + (y::'a::len word)) =
    (if unat x + unat y < 2^size x then
      unat x + unat y
    else
      unat x + unat y - 2^size x)
apply (subst word-arith-nat-defs)
apply (subst unat-of-nat)
apply (simp add: mod-nat-add word-size)
done

```

```

lemma word-neq-0-conv [simp]:
  fixes w :: 'a :: len word
  shows (w ≠ 0) = (0 < w)
proof -
  have 0 ≤ w by (rule word-zero-le)
  thus ?thesis by (auto simp add: word-less-def)
qed

```

```

lemma max-lt:
  unat (max a b div c) = unat (max a b) div unat (c:: 'a :: len word)
apply (subst word-arith-nat-defs)
apply (subst word-unat.eq-norm)

```

```

apply (subst mod-if)
apply clarsimp
apply (erule notE)
apply (insert div-le-dividend [of unat (max a b) unat c])
apply (erule order-le-less-trans)
apply (insert unat-lt2p [of max a b])
apply simp
done

lemma uint-sub-if-size:
  uint (x - y) =
    (if uint y ≤ uint x then
      uint x - uint y
    else
      uint x - uint y + 2^size x)
by (simp add: word-arith-alts int-word-uint mod-sub-if-z
      word-size)

lemma unat-sub-simple:
  x ≤ y ==> unat (y - x) = unat y - unat x
by (simp add: unat-def uint-sub-if-size word-le-def nat-diff-distrib)

lemmas unat-sub = unat-sub-simple

lemma word-less-sub1:
  fixes x :: 'a :: len word
  shows x ≠ 0 ==> 1 < x = (0 < x - 1)
by (simp add: unat-sub-if-size word-less-nat-alt)

lemma word-le-sub1:
  fixes x :: 'a :: len word
  shows x ≠ 0 ==> 1 ≤ x = (0 ≤ x - 1)
by (simp add: unat-sub-if-size order-le-less word-less-nat-alt)

lemmas word-less-sub1-numberof [simp] =
  word-less-sub1 [of number-of w, standard]
lemmas word-le-sub1-numberof [simp] =
  word-le-sub1 [of number-of w, standard]

lemma word-of-int-minus:
  word-of-int (2^len-of TYPE('a) - i) = (word-of-int (-i)::'a::len word)
proof -
  have x: 2^len-of TYPE('a) - i = -i + 2^len-of TYPE('a) by simp
  show ?thesis
    apply (subst x)
    apply (subst word-uint.Abs-norm [symmetric], subst zmod-zadd-self2)
    apply simp
    done
qed

```



**lemmas** *word-of-int-inj* =  
*word-uint.Abs-inject* [*unfolded uints-num*, *simplified*]

**lemma** *word-le-less-eq*:  
 $(x :: 'z :: \text{len } \text{word}) \leq y = (x = y \vee x < y)$   
**by** (*auto simp add: word-less-def*)

**lemma** *mod-plus-cong*:  
**assumes** 1:  $(b :: \text{int}) = b'$   
**and** 2:  $x \bmod b' = x' \bmod b'$   
**and** 3:  $y \bmod b' = y' \bmod b'$   
**and** 4:  $x' + y' = z'$   
**shows**  $(x + y) \bmod b = z' \bmod b'$   
**proof** –  
**from** 1 2[*symmetric*] 3[*symmetric*] **have**  $(x + y) \bmod b = (x' \bmod b' + y' \bmod b') \bmod b'$   
**by** (*simp add: zmod-zadd1-eq[symmetric]*)  
**also have**  $\dots = (x' + y') \bmod b'$   
**by** (*simp add: zmod-zadd1-eq[symmetric]*)  
**finally show** ?thesis **by** (*simp add: 4*)  
**qed**

**lemma** *mod-minus-cong*:  
**assumes** 1:  $(b :: \text{int}) = b'$   
**and** 2:  $x \bmod b' = x' \bmod b'$   
**and** 3:  $y \bmod b' = y' \bmod b'$   
**and** 4:  $x' - y' = z'$   
**shows**  $(x - y) \bmod b = z' \bmod b'$   
**using** *assms*  
**apply** (*subst zmod-zsub-left-eq*)  
**apply** (*subst zmod-zsub-right-eq*)  
**apply** (*simp add: zmod-zsub-left-eq [symmetric] zmod-zsub-right-eq [symmetric]*)  
**done**

**lemma** *word-induct-less*:  
 $\llbracket P (0 :: 'a :: \text{len } \text{word}); \bigwedge n. \llbracket n < m; P\ n \rrbracket \implies P\ (1 + n) \rrbracket \implies P\ m$   
**apply** (*cases m*)  
**apply** *atomize*  
**apply** (*erule rev-mp*)  
**apply** (*rule-tac x=m in spec*)  
**apply** (*induct-tac n*)  
**apply** *simp*  
**apply** *clarsimp*  
**apply** (*erule impE*)  
**apply** *clarsimp*  
**apply** (*erule-tac x=n in allE*)  
**apply** (*erule impE*)  
**apply** (*simp add: unat-arith-simps*)

```

    apply (clarsimp simp: unat-of-nat)
    apply simp
    apply (erule-tac x=of-nat na in allE)
    apply (erule impE)
    apply (simp add: unat-arith-simps)
    apply (clarsimp simp: unat-of-nat)
    apply simp
  done

lemma word-induct:
   $\llbracket P (0::'a::\text{len word}); \bigwedge n. P n \implies P (1 + n) \rrbracket \implies P m$ 
  by (erule word-induct-less, simp)

lemma word-induct2 [induct type]:
   $\llbracket P 0; \bigwedge n. \llbracket 1 + n \neq 0; P n \rrbracket \implies P (1 + n) \rrbracket \implies P (n::'b::\text{len word})$ 
  apply (rule word-induct, simp)
  apply (case-tac 1+n = 0, auto)
  done

constdefs
  word-rec :: 'a  $\Rightarrow$  ('b::len word  $\Rightarrow$  'a  $\Rightarrow$  'a)  $\Rightarrow$  'b word  $\Rightarrow$  'a
  word-rec forZero forSuc n  $\equiv$  nat-rec forZero (forSuc  $\circ$  of-nat) (unat n)

lemma word-rec-0: word-rec z s 0 = z
  by (simp add: word-rec-def)

lemma word-rec-Suc:
   $1 + n \neq (0::'a::\text{len word}) \implies \text{word-rec } z \ s \ (1 + n) = s \ n \ (\text{word-rec } z \ s \ n)$ 
  apply (simp add: word-rec-def unat-word-ariths)
  apply (subst nat-mod-eq')
  apply (cut-tac x=n in unat-lt2p)
  apply (drule Suc-mono)
  apply (simp add: less-Suc-eq-le)
  apply (simp only: order-less-le, simp)
  apply (erule contrapos-pn, simp)
  apply (drule arg-cong[where f=of-nat])
  apply simp
  apply (subst (asm) word-unat.Rep-Abs-A.Rep-inverse[of n])
  apply simp
  apply simp
  done

lemma word-rec-Pred:
   $n \neq 0 \implies \text{word-rec } z \ s \ n = s \ (n - 1) \ (\text{word-rec } z \ s \ (n - 1))$ 
  apply (rule subst[where t=n and s=1 + (n - 1)])
  apply simp
  apply (subst word-rec-Suc)
  apply simp
  apply simp

```

done

**lemma** *word-rec-in*:

$f \text{ (word-rec } z \text{ (}\lambda\text{-. } f) \text{ } n) = \text{word-rec (} f \text{ } z) \text{ (}\lambda\text{-. } f) \text{ } n$   
**by** (*induct*  $n$ ) (*simp-all* *add*: *word-rec-0* *word-rec-Suc*)

**lemma** *word-rec-in2*:

$f \text{ } n \text{ (word-rec } z \text{ } f \text{ } n) = \text{word-rec (} f \text{ } 0 \text{ } z) \text{ (} f \circ op + 1) \text{ } n$   
**by** (*induct*  $n$ ) (*simp-all* *add*: *word-rec-0* *word-rec-Suc*)

**lemma** *word-rec-twice*:

$m \leq n \implies \text{word-rec } z \text{ } f \text{ } n = \text{word-rec (word-rec } z \text{ } f \text{ (} n - m)) \text{ (} f \circ op + (n - m)) \text{ } m$   
**apply** (*erule* *rev-mp*)  
**apply** (*rule-tac*  $x=z$  **in** *spec*)  
**apply** (*rule-tac*  $x=f$  **in** *spec*)  
**apply** (*induct*  $n$ )  
**apply** (*simp* *add*: *word-rec-0*)  
**apply** *clarsimp*  
**apply** (*rule-tac*  $t=1 + n - m$  **and**  $s=1 + (n - m)$  **in** *subst*)  
**apply** *simp*  
**apply** (*case-tac*  $1 + (n - m) = 0$ )  
**apply** (*simp* *add*: *word-rec-0*)  
**apply** (*rule-tac*  $f = \text{word-rec } ?a \text{ } ?b$  **in** *arg-cong*)  
**apply** (*rule-tac*  $t=m$  **and**  $s=m + (1 + (n - m))$  **in** *subst*)  
**apply** *simp*  
**apply** (*simp* (*no-asm-use*))  
**apply** (*simp* *add*: *word-rec-Suc* *word-rec-in2*)  
**apply** (*erule* *impE*)  
**apply** *uint-arith*  
**apply** (*drule-tac*  $x=x \circ op + 1$  **in** *spec*)  
**apply** (*drule-tac*  $x=x \text{ } 0 \text{ } xa$  **in** *spec*)  
**apply** *simp*  
**apply** (*rule-tac*  $t=\lambda a. x (1 + (n - m + a))$  **and**  $s=\lambda a. x (1 + (n - m) + a)$  **in** *subst*)  
**apply** (*clarsimp* *simp* *add*: *expand-fun-eq*)  
**apply** (*rule-tac*  $t=(1 + (n - m + xb))$  **and**  $s=1 + (n - m) + xb$  **in** *subst*)  
**apply** *simp*  
**apply** (*rule* *refl*)  
**apply** (*rule* *refl*)  
done

**lemma** *word-rec-id*:  $\text{word-rec } z \text{ (}\lambda\text{-. } id) \text{ } n = z$

**by** (*induct*  $n$ ) (*auto* *simp* *add*: *word-rec-0* *word-rec-Suc*)

**lemma** *word-rec-id-eq*:  $\forall m < n. f \text{ } m = id \implies \text{word-rec } z \text{ } f \text{ } n = z$

**apply** (*erule* *rev-mp*)

**apply** (*induct*  $n$ )

**apply** (*auto* *simp* *add*: *word-rec-0* *word-rec-Suc*)

```

apply (drule spec, erule mp)
apply uint-arith
apply (drule-tac x=n in spec, erule impE)
apply uint-arith
apply simp
done

lemma word-rec-max:
   $\forall m \geq n. m \neq -1 \longrightarrow f\ m = id \implies \text{word-rec } z\ f\ -1 = \text{word-rec } z\ f\ n$ 
apply (subst word-rec-twice[where n=-1 and m=-1 - n])
apply simp
apply simp
apply (rule word-rec-id-eq)
apply clarsimp
apply (drule spec, rule mp, erule mp)
apply (rule word-plus-mono-right2[OF - order-less-imp-le])
prefer 2
apply assumption
apply simp
apply (erule contrapos-pn)
apply simp
apply (drule arg-cong[where f= $\lambda x. x - n$ ])
apply simp
done

lemma unatSuc:
   $1 + n \neq (0::'a::\text{len word}) \implies \text{unat } (1 + n) = \text{Suc } (\text{unat } n)$ 
by unat-arith

end

```

## 16 WordMain: Main Word Library

```

theory WordMain imports WordGenLib
begin

lemmas word-no-1 [simp] = word-1-no [symmetric]
lemmas word-no-0 [simp] = word-0-no [symmetric]

declare word-0-bl [simp]
declare bin-to-bl-def [simp]
declare to-bl-0 [simp]
declare of-bl-True [simp]

Examples

types word32 = 32 word
types word8 = 8 word
types byte = word8

```

for more see WordExampes.thy  
**end**

## References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.