

Isabelle/FOL — First-Order Logic

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1 Intuitionistic first-order logic

```
theory IFOL
imports Pure
uses
  ~/src/Provers/splitter.ML
  ~/src/Provers/hypsubst.ML
  ~/src/Tools/IsaPlanner/zipper.ML
  ~/src/Tools/IsaPlanner/isand.ML
  ~/src/Tools/IsaPlanner/rw-tools.ML
  ~/src/Tools/IsaPlanner/rw-inst.ML
  ~/src/Provers/eqsubst.ML
  ~/src/Provers/quantifier1.ML
  ~/src/Provers/project-rule.ML
(fologic.ML)
(hypsubstdata.ML)
(intprover.ML)
begin
```

1.1 Syntax and axiomatic basis

global

classes *term*

defaultsort *term*

typedecl *o*

judgment

Trueprop :: $o \Rightarrow prop$ ((-) 5)

consts

True :: *o*

False :: *o*

op = :: $['a, 'a] \Rightarrow o$ (**infixl** = 50)

Not :: $o \Rightarrow o$ (\sim - [40] 40)

op & :: $[o, o] \Rightarrow o$ (**infixr** & 35)

op | :: $[o, o] \Rightarrow o$ (**infixr** | 30)

op --> :: $[o, o] \Rightarrow o$ (**infixr** --> 25)

op <-> :: $[o, o] \Rightarrow o$ (**infixr** <-> 25)

All :: $('a \Rightarrow o) \Rightarrow o$ (**binder** ALL 10)

Ex :: $('a \Rightarrow o) \Rightarrow o$ (**binder** EX 10)

Ex1 :: $('a \Rightarrow o) \Rightarrow o$ (**binder** EX! 10)

abbreviation

not-equal :: $['a, 'a] \Rightarrow o$ (**infixl** $\sim =$ 50) **where**

$x \sim = y == \sim (x = y)$

notation (*xsymbols*)

not-equal (**infixl** \neq 50)

notation (*HTML output*)

not-equal (**infixl** \neq 50)

notation (*xsymbols*)

Not (\neg - [40] 40) **and**

op & (**infixr** \wedge 35) **and**

op | (**infixr** \vee 30) **and**

All (**binder** \forall 10) **and**

Ex (**binder** \exists 10) **and**

Ex1 (**binder** $\exists!$ 10) **and**

$op \dashrightarrow$ (**infix** \longrightarrow 25) **and**
 $op \longleftrightarrow$ (**infix** \longleftrightarrow 25)

notation (*HTML output*)

Not (\neg - [40] 40) **and**
op & (**infix** \wedge 35) **and**
op | (**infix** \vee 30) **and**
All (**binder** \forall 10) **and**
Ex (**binder** \exists 10) **and**
Ex1 (**binder** $\exists!$ 10)

local

finalconsts

False All Ex
op =
op &
op |
op \dashrightarrow

axioms

refl: $a=a$

conjI: $[[P; Q]] \implies P \& Q$
conjunct1: $P \& Q \implies P$
conjunct2: $P \& Q \implies Q$

disjI1: $P \implies P | Q$
disjI2: $Q \implies P | Q$
disjE: $[[P | Q; P \implies R; Q \implies R]] \implies R$

impI: $(P \implies Q) \implies P \dashrightarrow Q$
mp: $[[P \dashrightarrow Q; P]] \implies Q$

FalseE: $False \implies P$

allI: $(!!x. P(x)) \implies (ALL x. P(x))$
spec: $(ALL x. P(x)) \implies P(x)$

exI: $P(x) \implies (EX x. P(x))$
exE: $[[EX x. P(x); !!x. P(x) \implies R]] \implies R$

eq-reflection: $(x=y) \implies (x==y)$
iff-reflection: $(P<->Q) \implies (P==Q)$

lemmas *strip* = *impI all*

Thanks to Stephan Merz

theorem *subst*:

assumes *eq*: $a = b$ **and** *p*: $P(a)$

shows $P(b)$

<proof>

defs

True-def: $True == False \dashrightarrow False$

not-def: $\sim P == P \dashrightarrow False$

iff-def: $P <-> Q == (P \dashrightarrow Q) \ \& \ (Q \dashrightarrow P)$

ex1-def: $Ex1(P) == EX\ x.\ P(x) \ \& \ (ALL\ y.\ P(y) \dashrightarrow y=x)$

1.2 Lemmas and proof tools

lemma *TrueI*: *True*

<proof>

lemma *conjE*:

assumes *major*: $P \ \& \ Q$

and *r*: $[| P; Q |] \implies R$

shows R

<proof>

lemma *impE*:

assumes *major*: $P \dashrightarrow Q$

and P

and *r*: $Q \implies R$

shows R

<proof>

lemma *allE*:

assumes *major*: $ALL\ x.\ P(x)$

and $r: P(x) \implies R$
shows R
 $\langle proof \rangle$

lemma *all-dupE*:
assumes *major*: $ALL\ x.\ P(x)$
and $r: [[P(x); ALL\ x.\ P(x)]] \implies R$
shows R
 $\langle proof \rangle$

lemma *notI*: $(P \implies False) \implies \sim P$
 $\langle proof \rangle$

lemma *notE*: $[[\sim P; P]] \implies R$
 $\langle proof \rangle$

lemma *rev-notE*: $[[P; \sim P]] \implies R$
 $\langle proof \rangle$

lemma *not-to-imp*:
assumes $\sim P$
and $r: P \longrightarrow False \implies Q$
shows Q
 $\langle proof \rangle$

lemma *rev-mp*: $[[P; P \longrightarrow Q]] \implies Q$
 $\langle proof \rangle$

lemma *contrapos*:
assumes *major*: $\sim Q$
and *minor*: $P \implies Q$
shows $\sim P$
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *iffI*: $[[P \implies Q; Q \implies P]] \implies P \leftrightarrow Q$
<proof>

lemma *iffE*:
 assumes *major*: $P \leftrightarrow Q$
 and *r*: $P \dashv\vdash Q \implies Q \dashv\vdash P \implies R$
 shows *R*
<proof>

lemma *iffD1*: $[[P \leftrightarrow Q; P]] \implies Q$
<proof>

lemma *iffD2*: $[[P \leftrightarrow Q; Q]] \implies P$
<proof>

lemma *rev-iffD1*: $[[P; P \leftrightarrow Q]] \implies Q$
<proof>

lemma *rev-iffD2*: $[[Q; P \leftrightarrow Q]] \implies P$
<proof>

lemma *iff-refl*: $P \leftrightarrow P$
<proof>

lemma *iff-sym*: $Q \leftrightarrow P \implies P \leftrightarrow Q$
<proof>

lemma *iff-trans*: $[[P \leftrightarrow Q; Q \leftrightarrow R]] \implies P \leftrightarrow R$
<proof>

lemma *exI1*:
 $P(a) \implies (!x. P(x) \implies x=a) \implies EX! x. P(x)$
<proof>

lemma *ex-exI1*:
 $EX x. P(x) \implies (!x y. [[P(x); P(y)]] \implies x=y) \implies EX! x. P(x)$
<proof>

lemma *ex1E*:
 $EX! x. P(x) \implies (!x. [[P(x); ALL y. P(y) \dashv\vdash y=x]] \implies R) \implies R$

$\langle proof \rangle$

$\langle ML \rangle$

lemma *conj-cong*:

assumes $P \leftrightarrow P'$
and $P' \implies Q \leftrightarrow Q'$
shows $(P \& Q) \leftrightarrow (P' \& Q')$
 $\langle proof \rangle$

lemma *conj-cong2*:

assumes $P \leftrightarrow P'$
and $P' \implies Q \leftrightarrow Q'$
shows $(Q \& P) \leftrightarrow (Q' \& P')$
 $\langle proof \rangle$

lemma *disj-cong*:

assumes $P \leftrightarrow P'$ **and** $Q \leftrightarrow Q'$
shows $(P | Q) \leftrightarrow (P' | Q')$
 $\langle proof \rangle$

lemma *imp-cong*:

assumes $P \leftrightarrow P'$
and $P' \implies Q \leftrightarrow Q'$
shows $(P \dashrightarrow Q) \leftrightarrow (P' \dashrightarrow Q')$
 $\langle proof \rangle$

lemma *iff-cong*: $[[P \leftrightarrow P'; Q \leftrightarrow Q']] \implies (P \leftrightarrow Q) \leftrightarrow (P' \leftrightarrow Q')$
 $\langle proof \rangle$

lemma *not-cong*: $P \leftrightarrow P' \implies \sim P \leftrightarrow \sim P'$
 $\langle proof \rangle$

lemma *all-cong*:

assumes $\forall x. P(x) \leftrightarrow Q(x)$
shows $(\forall x. P(x)) \leftrightarrow (\forall x. Q(x))$
 $\langle proof \rangle$

lemma *ex-cong*:

assumes $\forall x. P(x) \leftrightarrow Q(x)$
shows $(\exists x. P(x)) \leftrightarrow (\exists x. Q(x))$
 $\langle proof \rangle$

lemma *ex1-cong*:

assumes $\exists!x. P(x) \leftrightarrow Q(x)$
shows $(\exists!x. P(x)) \leftrightarrow (\exists!x. Q(x))$
<proof>

lemma *sym*: $a=b \implies b=a$
<proof>

lemma *trans*: $[| a=b; b=c |] \implies a=c$
<proof>

lemma *not-sym*: $b \sim a \implies a \sim b$
<proof>

lemma *def-imp-iff*: $(A == B) \implies A \leftrightarrow B$
<proof>

lemma *meta-eq-to-obj-eq*: $(A == B) \implies A = B$
<proof>

lemma *meta-eq-to-iff*: $x==y \implies x \leftrightarrow y$
<proof>

lemma *ssubst*: $[| b = a; P(a) |] \implies P(b)$
<proof>

lemma *ex1-equalsE*:
 $[| \exists!x. P(x); P(a); P(b) |] \implies a=b$
<proof>

lemma *subst-context*: $[| a=b |] \implies t(a)=t(b)$
<proof>

lemma *subst-context2*: $[| a=b; c=d |] \implies t(a,c)=t(b,d)$
<proof>

lemma *subst-context3*: $[| a=b; c=d; e=f |] \implies t(a,c,e)=t(b,d,f)$
<proof>

lemma *box-equals*: $[| a=b; a=c; b=d |] \implies c=d$

<proof>

lemma *simp-equals*: $[[a=c; b=d; c=d]] ==> a=b$
<proof>

lemma *pred1-cong*: $a=a' ==> P(a) <-> P(a')$
<proof>

lemma *pred2-cong*: $[[a=a'; b=b']] ==> P(a,b) <-> P(a',b')$
<proof>

lemma *pred3-cong*: $[[a=a'; b=b'; c=c']] ==> P(a,b,c) <-> P(a',b',c')$
<proof>

<ML>

lemma *eq-cong*: $[[a = a'; b = b']] ==> a = b <-> a' = b'$
<proof>

lemma *conj-impE*:
 assumes *major*: $(P \& Q) \longrightarrow S$
 and *r*: $P \longrightarrow (Q \longrightarrow S) ==> R$
 shows *R*
<proof>

lemma *disj-impE*:
 assumes *major*: $(P | Q) \longrightarrow S$
 and *r*: $[[P \longrightarrow S; Q \longrightarrow S]] ==> R$
 shows *R*
<proof>

lemma *imp-impE*:
 assumes *major*: $(P \longrightarrow Q) \longrightarrow S$
 and *r1*: $[[P; Q \longrightarrow S]] ==> Q$
 and *r2*: $S ==> R$
 shows *R*
<proof>

lemma not-impE:

$\sim P \dashrightarrow S \implies (P \implies \text{False}) \implies (S \implies R) \implies R$
<proof>

lemma iff-impE:

assumes major: $(P \leftrightarrow Q) \dashrightarrow S$
and r1: $[[P; Q \dashrightarrow S]] \implies Q$
and r2: $[[Q; P \dashrightarrow S]] \implies P$
and r3: $S \implies R$
shows R
<proof>

lemma all-impE:

assumes major: $(\text{ALL } x. P(x)) \dashrightarrow S$
and r1: $!!x. P(x)$
and r2: $S \implies R$
shows R
<proof>

lemma ex-impE:

assumes major: $(\text{EX } x. P(x)) \dashrightarrow S$
and r: $P(x) \dashrightarrow S \implies R$
shows R
<proof>

lemma disj-imp-disj:

$P \mid Q \implies (P \implies R) \implies (Q \implies S) \implies R \mid S$
<proof>

<ML>

lemma thin-refl: $!!X. [[x=x; \text{PROP } W]] \implies \text{PROP } W$ *<proof>*

<ML>

1.3 Intuitionistic Reasoning

lemma impE':

assumes $1: P \dashrightarrow Q$
and $2: Q \implies R$
and $3: P \dashrightarrow Q \implies P$
shows R
<proof>

lemma *allE'*:
assumes 1: $\text{ALL } x. P(x)$
and 2: $P(x) \implies \text{ALL } x. P(x) \implies Q$
shows Q
 $\langle \text{proof} \rangle$

lemma *notE'*:
assumes 1: $\sim P$
and 2: $\sim P \implies P$
shows R
 $\langle \text{proof} \rangle$

lemmas [*Pure.elim!*] = *disjE iffE FalseE conjE exE*
and [*Pure.intro!*] = *iffI conjI impI TrueI notI allI refl*
and [*Pure.elim 2*] = *allE notE' impE'*
and [*Pure.intro*] = *exI disjI2 disjI1*

$\langle \text{ML} \rangle$

lemma *iff-not-sym*: $\sim (Q \longleftrightarrow P) \implies \sim (P \longleftrightarrow Q)$
 $\langle \text{proof} \rangle$

lemmas [*sym*] = *sym iff-sym not-sym iff-not-sym*
and [*Pure.elim?*] = *iffD1 iffD2 impE*

lemma *eq-commute*: $a=b \longleftrightarrow b=a$
 $\langle \text{proof} \rangle$

1.4 Atomizing meta-level rules

lemma *atomize-all* [*atomize*]: $(!!x. P(x)) \implies \text{Trueprop } (\text{ALL } x. P(x))$
 $\langle \text{proof} \rangle$

lemma *atomize-imp* [*atomize*]: $(A \implies B) \implies \text{Trueprop } (A \longrightarrow B)$
 $\langle \text{proof} \rangle$

lemma *atomize-eq* [*atomize*]: $(x == y) \implies \text{Trueprop } (x = y)$
 $\langle \text{proof} \rangle$

lemma *atomize-iff* [*atomize*]: $(A == B) \implies \text{Trueprop } (A \longleftrightarrow B)$
 $\langle \text{proof} \rangle$

lemma *atomize-conj* [*atomize*]:
includes *meta-conjunction-syntax*
shows $(A \ \&\& \ B) \implies \text{Trueprop } (A \ \& \ B)$
 $\langle \text{proof} \rangle$

lemmas [*symmetric, rulify*] = *atomize-all atomize-imp*
and [*symmetric, defn*] = *atomize-all atomize-imp atomize-eq atomize-iff*

1.5 Calculational rules

lemma *forw-subst*: $a = b \implies P(b) \implies P(a)$
 ⟨*proof*⟩

lemma *back-subst*: $P(a) \implies a = b \implies P(b)$
 ⟨*proof*⟩

Note that this list of rules is in reverse order of priorities.

lemmas *basic-trans-rules* [*trans*] =
forw-subst
back-subst
rev-mp
mp
trans

1.6 “Let” declarations

nonterminals *letbinds letbind*

constdefs

Let :: [*a*::{}], *'a* => *'b*] => (*'b*::{})
Let(*s*, *f*) == *f*(*s*)

syntax

-bind :: [*pttrn*, *'a*] => *letbind* ((*?*- / -) 10)
 :: *letbind* => *letbinds* (-)
-binds :: [*letbind*, *letbinds*] => *letbinds* (-; / -)
-Let :: [*letbinds*, *'a*] => *'a* ((*let* (-) / *in* (-)) 10)

translations

-Let(*-binds*(*b*, *bs*), *e*) == *-Let*(*b*, *-Let*(*bs*, *e*))
let *x* = *a* *in* *e* == *Let*(*a*, %*x*. *e*)

lemma *LetI*:

assumes !!*x*. $x=t \implies P(u(x))$
shows $P(\text{let } x=t \text{ in } u(x))$
 ⟨*proof*⟩

1.7 ML bindings

⟨*ML*⟩

end

2 Classical first-order logic

```
theory FOL
imports IFOL
uses
  ~/src/Provers/classical.ML
  ~/src/Provers/blast.ML
  ~/src/Provers/clasimp.ML
  ~/src/Tools/induct.ML
  (cladata.ML)
  (blastdata.ML)
  (simpdata.ML)
begin
```

2.1 The classical axiom

```
axioms
  classical: ( $\sim P \implies P$ )  $\implies P$ 
```

2.2 Lemmas and proof tools

```
lemma ccontr: ( $\neg P \implies False$ )  $\implies P$ 
  <proof>
```

```
lemma disjCI: ( $\sim Q \implies P$ )  $\implies P \mid Q$ 
  <proof>
```

```
lemma ex-classical:
  assumes r:  $\sim(EX\ x.\ P(x)) \implies P(a)$ 
  shows  $EX\ x.\ P(x)$ 
  <proof>
```

```
lemma exCI:
  assumes r:  $ALL\ x.\ \sim P(x) \implies P(a)$ 
  shows  $EX\ x.\ P(x)$ 
  <proof>
```

```
lemma excluded-middle:  $\sim P \mid P$ 
  <proof>
```

```
<ML>
```

```
lemma case-split-thm:
  assumes r1:  $P \implies Q$ 
  and r2:  $\sim P \implies Q$ 
```

shows Q
<proof>

lemmas *case-split* = *case-split-thm* [*case-names True False*]

<ML>

lemma *impCE*:
 assumes *major*: $P \dashrightarrow Q$
 and *r1*: $\sim P \implies R$
 and *r2*: $Q \implies R$
 shows R
 <proof>

lemma *impCE'*:
 assumes *major*: $P \dashrightarrow Q$
 and *r1*: $Q \implies R$
 and *r2*: $\sim P \implies R$
 shows R
 <proof>

lemma *notnotD*: $\sim\sim P \implies P$
<proof>

lemma *contrapos2*: $[[Q; \sim P \implies \sim Q]] \implies P$
<proof>

lemma *iffCE*:
 assumes *major*: $P \leftrightarrow Q$
 and *r1*: $[[P; Q]] \implies R$
 and *r2*: $[[\sim P; \sim Q]] \implies R$
 shows R
 <proof>

lemma *alt-ex1E*:
 assumes *major*: $EX! x. P(x)$

and $r: !!x. [| P(x); ALL y y'. P(y) \& P(y') \dashrightarrow y=y' |] \implies R$
shows R
 $\langle proof \rangle$

$\langle ML \rangle$

lemma *ex1-functional*: $[| EX! z. P(a,z); P(a,b); P(a,c) |] \implies b = c$
 $\langle proof \rangle$

lemma *True-implies-equals*: $(True \implies PROP P) == PROP P$
 $\langle proof \rangle$

lemma *uncurry*: $P \dashrightarrow Q \dashrightarrow R \implies P \& Q \dashrightarrow R$
 $\langle proof \rangle$

lemma *iff-allI*: $(!!x. P(x) \leftrightarrow Q(x)) \implies (ALL x. P(x)) \leftrightarrow (ALL x. Q(x))$
 $\langle proof \rangle$

lemma *iff-exI*: $(!!x. P(x) \leftrightarrow Q(x)) \implies (EX x. P(x)) \leftrightarrow (EX x. Q(x))$
 $\langle proof \rangle$

lemma *all-comm*: $(ALL x y. P(x,y)) \leftrightarrow (ALL y x. P(x,y))$ $\langle proof \rangle$

lemma *ex-comm*: $(EX x y. P(x,y)) \leftrightarrow (EX y x. P(x,y))$ $\langle proof \rangle$

$\langle ML \rangle$

2.3 Other simple lemmas

lemma [*simp*]: $((P \dashrightarrow R) \leftrightarrow (Q \dashrightarrow R)) \leftrightarrow ((P \leftrightarrow Q) | R)$
 $\langle proof \rangle$

lemma [*simp*]: $((P \dashrightarrow Q) \leftrightarrow (P \dashrightarrow R)) \leftrightarrow (P \dashrightarrow (Q \leftrightarrow R))$
 $\langle proof \rangle$

lemma *not-disj-iff-imp*: $\sim P | Q \leftrightarrow (P \dashrightarrow Q)$
 $\langle proof \rangle$

lemma *conj-mono*: $[| P1 \dashrightarrow Q1; P2 \dashrightarrow Q2 |] \implies (P1 \& P2) \dashrightarrow (Q1 \& Q2)$
 $\langle proof \rangle$

lemma *disj-mono*: $[| P1 \dashrightarrow Q1; P2 \dashrightarrow Q2 |] \implies (P1 | P2) \dashrightarrow (Q1 | Q2)$
 $\langle proof \rangle$

lemma *imp-mono*: $[| Q1 \dashrightarrow P1; P2 \dashrightarrow Q2 |] \implies (P1 \dashrightarrow P2) \dashrightarrow (Q1 \dashrightarrow Q2)$

<proof>

lemma *imp-refl*: $P \dashv\vdash P$

<proof>

lemma *ex-mono*: $(\exists x. P(x) \dashv\vdash Q(x)) \implies (EX x. P(x) \dashv\vdash EX x. Q(x))$

<proof>

lemma *all-mono*: $(\forall x. P(x) \dashv\vdash Q(x)) \implies (ALL x. P(x) \dashv\vdash ALL x. Q(x))$

<proof>

2.4 Proof by cases and induction

Proper handling of non-atomic rule statements.

constdefs

induct-forall **where** $induct-forall(P) == \forall x. P(x)$

induct-implies **where** $induct-implies(A, B) == A \longrightarrow B$

induct-equal **where** $induct-equal(x, y) == x = y$

induct-conj **where** $induct-conj(A, B) == A \wedge B$

lemma *induct-forall-eq*: $(\exists x. P(x)) == Trueprop(induct-forall(\lambda x. P(x)))$

<proof>

lemma *induct-implies-eq*: $(A \implies B) == Trueprop(induct-implies(A, B))$

<proof>

lemma *induct-equal-eq*: $(x == y) == Trueprop(induct-equal(x, y))$

<proof>

lemma *induct-conj-eq*:

includes *meta-conjunction-syntax*

shows $(A \ \&\& \ B) == Trueprop(induct-conj(A, B))$

<proof>

lemmas *induct-atomize* = *induct-forall-eq induct-implies-eq induct-equal-eq induct-conj-eq*

lemmas *induct-rulify* [*symmetric, standard*] = *induct-atomize*

lemmas *induct-rulify-fallback* =

induct-forall-def induct-implies-def induct-equal-def induct-conj-def

hide *const induct-forall induct-implies induct-equal induct-conj*

Method setup.

<ML>

declare *case-split* [*cases type: o*]

end