

# The Hahn-Banach Theorem for Real Vector Spaces

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## Abstract

The Hahn-Banach Theorem is one of the most fundamental results in functional analysis. We present a fully formal proof of two versions of the theorem, one for general linear spaces and another for normed spaces. This development is based on simply-typed classical set-theory, as provided by Isabelle/HOL.

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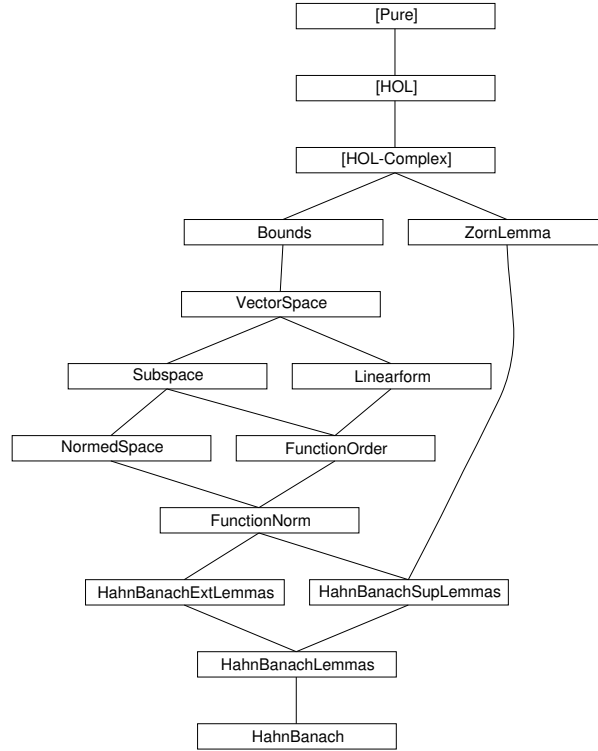
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# 1 Preface

This is a fully formal proof of the Hahn-Banach Theorem. It closely follows the informal presentation given in Heuser's textbook [1, § 36]. Another formal proof of the same theorem has been done in Mizar [3]. A general overview of the relevance and history of the Hahn-Banach Theorem is given by Narici and Beckenstein [2].

The document is structured as follows. The first part contains definitions of basic notions of linear algebra: vector spaces, subspaces, normed spaces, continuous linear-forms, norm of functions and an order on functions by domain extension. The second part contains some lemmas about the supremum (w.r.t. the function order) and extension of non-maximal functions. With these preliminaries, the main proof of the theorem (in its two versions) is conducted in the third part. The dependencies of individual theories are as follows.



## Part I

# Basic Notions

## 2 Bounds

**theory** *Bounds* **imports** *Main Real* **begin**

**locale** *lub* =  
**fixes** *A* **and** *x*  
**assumes** *least* [*intro?*]:  $(\bigwedge a. a \in A \implies a \leq b) \implies x \leq b$   
**and** *upper* [*intro?*]:  $a \in A \implies a \leq x$

**lemmas** [*elim?*] = *lub.least lub.upper*

**constdefs**  
*the-lub* :: '*a*::order set  $\Rightarrow$  '*a*  
*the-lub* *A*  $\equiv$  *The* (*lub* *A*)

**syntax** (*xsymbols*)  
*the-lub* :: '*a*::order set  $\Rightarrow$  '*a* ( $\bigsqcup$  - [90] 90)

**lemma** *the-lub-equality* [*elim?*]:  
**includes** *lub*  
**shows**  $\bigsqcup A = (x::'a::order)$   
**proof** (*unfold the-lub-def*)  
**from** *lub-axioms* **show** *The* (*lub* *A*) = *x*  
**proof**  
**fix** *x'* **assume** *lub'*: *lub* *A* *x'*  
**show**  $x' = x$   
**proof** (*rule order-antisym*)  
**from** *lub'* **show**  $x' \leq x$   
**proof**  
**fix** *a* **assume**  $a \in A$   
**then show**  $a \leq x$  ..  
**qed**  
**show**  $x \leq x'$   
**proof**  
**fix** *a* **assume**  $a \in A$   
**with** *lub'* **show**  $a \leq x'$  ..  
**qed**  
**qed**  
**qed**  
**qed**

**lemma** *the-lubI-ex*:  
**assumes** *ex*:  $\exists x. \text{lub } A \ x$   
**shows** *lub* *A*  $(\bigsqcup A)$   
**proof** –  
**from** *ex* **obtain** *x* **where** *lub* *A* *x* ..  
**also from** *x* **have** [*symmetric*]:  $\bigsqcup A = x$  ..  
**finally show** ?*thesis* .  
**qed**

```

lemma lub-compat: lub A x = isLub UNIV A x
proof –
  have isUb UNIV A = (λx. A *≤ x ∧ x ∈ UNIV)
    by (rule ext) (simp only: isUb-def)
  then show ?thesis
    by (simp only: lub-def isLub-def leastP-def setge-def settle-def) blast
qed

```

```

lemma real-complete:
  fixes A :: real set
  assumes nonempty:  $\exists a. a \in A$ 
    and ex-upper:  $\exists y. \forall a \in A. a \leq y$ 
  shows  $\exists x. \text{lub } A \ x$ 
proof –
  from ex-upper have  $\exists y. \text{isUb } UNIV \ A \ y$ 
    by (unfold isUb-def settle-def) blast
  with nonempty have  $\exists x. \text{isLub } UNIV \ A \ x$ 
    by (rule reals-complete)
  then show ?thesis by (simp only: lub-compat)
qed

end

```

### 3 Vector spaces

**theory** *VectorSpace* **imports** *Real Bounds Zorn* **begin**

#### 3.1 Signature

For the definition of real vector spaces a type  $'a$  of the sort  $\{plus, minus, zero\}$  is considered, on which a real scalar multiplication  $\cdot$  is declared.

```

consts
  prod :: real  $\Rightarrow$   $'a::\{plus, minus, zero\}$   $\Rightarrow$   $'a$     (infixr  $'(*)$  70)

syntax (xsymbols)
  prod :: real  $\Rightarrow$   $'a \Rightarrow 'a$                                 (infixr  $\cdot$  70)
syntax (HTML output)
  prod :: real  $\Rightarrow$   $'a \Rightarrow 'a$                                 (infixr  $\cdot$  70)

```

#### 3.2 Vector space laws

A *vector space* is a non-empty set  $V$  of elements from  $'a$  with the following vector space laws: The set  $V$  is closed under addition and scalar multiplication, addition is associative and commutative;  $-x$  is the inverse of  $x$  w. r. t. addition and  $0$  is the neutral element of addition. Addition and multiplication are distributive; scalar multiplication is associative and the real number  $1$  is the neutral element of scalar multiplication.

```

locale vectorspace = var V +
  assumes non-empty [iff, intro?]:  $V \neq \{\}$ 

```

**and** *add-closed* [*iff*]:  $x \in V \implies y \in V \implies x + y \in V$   
**and** *mult-closed* [*iff*]:  $x \in V \implies a \cdot x \in V$   
**and** *add-assoc*:  $x \in V \implies y \in V \implies z \in V \implies (x + y) + z = x + (y + z)$   
**and** *add-commute*:  $x \in V \implies y \in V \implies x + y = y + x$   
**and** *diff-self* [*simp*]:  $x \in V \implies x - x = 0$   
**and** *add-zero-left* [*simp*]:  $x \in V \implies 0 + x = x$   
**and** *add-mult-distrib1*:  $x \in V \implies y \in V \implies a \cdot (x + y) = a \cdot x + a \cdot y$   
**and** *add-mult-distrib2*:  $x \in V \implies (a + b) \cdot x = a \cdot x + b \cdot x$   
**and** *mult-assoc*:  $x \in V \implies (a * b) \cdot x = a \cdot (b \cdot x)$   
**and** *mult-1* [*simp*]:  $x \in V \implies 1 \cdot x = x$   
**and** *negate-eq1*:  $x \in V \implies -x = (-1) \cdot x$   
**and** *diff-eq1*:  $x \in V \implies y \in V \implies x - y = x + -y$

**lemma** (**in** *vectorspace*) *negate-eq2*:  $x \in V \implies (-1) \cdot x = -x$   
**by** (*rule negate-eq1* [*symmetric*])

**lemma** (**in** *vectorspace*) *negate-eq2a*:  $x \in V \implies -1 \cdot x = -x$   
**by** (*simp add: negate-eq1*)

**lemma** (**in** *vectorspace*) *diff-eq2*:  $x \in V \implies y \in V \implies x + -y = x - y$   
**by** (*rule diff-eq1* [*symmetric*])

**lemma** (**in** *vectorspace*) *diff-closed* [*iff*]:  $x \in V \implies y \in V \implies x - y \in V$   
**by** (*simp add: diff-eq1 negate-eq1*)

**lemma** (**in** *vectorspace*) *neg-closed* [*iff*]:  $x \in V \implies -x \in V$   
**by** (*simp add: negate-eq1*)

**lemma** (**in** *vectorspace*) *add-left-commute*:  
 $x \in V \implies y \in V \implies z \in V \implies x + (y + z) = y + (x + z)$   
**proof** –  
**assume** *xyz*:  $x \in V \ y \in V \ z \in V$   
**hence**  $x + (y + z) = (x + y) + z$   
**by** (*simp only: add-assoc*)  
**also from** *xyz* **have**  $\dots = (y + x) + z$  **by** (*simp only: add-commute*)  
**also from** *xyz* **have**  $\dots = y + (x + z)$  **by** (*simp only: add-assoc*)  
**finally show** *?thesis* .  
**qed**

**theorems** (**in** *vectorspace*) *add-ac* =  
*add-assoc add-commute add-left-commute*

The existence of the zero element of a vector space follows from the non-emptiness of carrier set.

**lemma** (**in** *vectorspace*) *zero* [*iff*]:  $0 \in V$   
**proof** –

**from** *non-empty* **obtain** *x* **where**  $x: x \in V$  **by** *blast*  
**then have**  $0 = x - x$  **by** (*rule diff-self* [*symmetric*])  
**also from** *x* **have**  $\dots \in V$  **by** (*rule diff-closed*)  
**finally show** *?thesis* .  
**qed**

**lemma** (**in** *vectorspace*) *add-zero-right* [*simp*]:  
 $x \in V \implies x + 0 = x$

**proof** –

**assume**  $x: x \in V$

**from this and zero have**  $x + 0 = 0 + x$  **by** (rule add-commute)

**also from x have**  $\dots = x$  **by** (rule add-zero-left)

**finally show** ?thesis .

**qed**

**lemma** (in *vectorspace*) *mult-assoc2*:

$x \in V \implies a \cdot b \cdot x = (a * b) \cdot x$

**by** (simp only: mult-assoc)

**lemma** (in *vectorspace*) *diff-mult-distrib1*:

$x \in V \implies y \in V \implies a \cdot (x - y) = a \cdot x - a \cdot y$

**by** (simp add: diff-eq1 negate-eq1 add-mult-distrib1 mult-assoc2)

**lemma** (in *vectorspace*) *diff-mult-distrib2*:

$x \in V \implies (a - b) \cdot x = a \cdot x - (b \cdot x)$

**proof** –

**assume**  $x: x \in V$

**have**  $(a - b) \cdot x = (a + - b) \cdot x$

**by** (simp add: real-diff-def)

**also have**  $\dots = a \cdot x + (- b) \cdot x$

**by** (rule add-mult-distrib2)

**also from x have**  $\dots = a \cdot x + - (b \cdot x)$

**by** (simp add: negate-eq1 mult-assoc2)

**also from x have**  $\dots = a \cdot x - (b \cdot x)$

**by** (simp add: diff-eq1)

**finally show** ?thesis .

**qed**

**lemmas** (in *vectorspace*) *distrib* =

*add-mult-distrib1 add-mult-distrib2*

*diff-mult-distrib1 diff-mult-distrib2*

Further derived laws:

**lemma** (in *vectorspace*) *mult-zero-left* [simp]:

$x \in V \implies 0 \cdot x = 0$

**proof** –

**assume**  $x: x \in V$

**have**  $0 \cdot x = (1 - 1) \cdot x$  **by** simp

**also have**  $\dots = (1 + - 1) \cdot x$  **by** simp

**also have**  $\dots = 1 \cdot x + (- 1) \cdot x$

**by** (rule add-mult-distrib2)

**also from x have**  $\dots = x + (- 1) \cdot x$  **by** simp

**also from x have**  $\dots = x + - x$  **by** (simp add: negate-eq2a)

**also from x have**  $\dots = x - x$  **by** (simp add: diff-eq2)

**also from x have**  $\dots = 0$  **by** simp

**finally show** ?thesis .

**qed**

**lemma** (in *vectorspace*) *mult-zero-right* [simp]:

$a \cdot 0 = (0::'a)$

**proof** –

**have**  $a \cdot 0 = a \cdot (0 - (0::'a))$  **by** simp

also have  $\dots = a \cdot 0 - a \cdot 0$   
 by (rule diff-mult-distrib1) simp-all  
 also have  $\dots = 0$  by simp  
 finally show ?thesis .  
 qed

lemma (in vectorspace) minus-mult-cancel [simp]:  
 $x \in V \implies (-a) \cdot -x = a \cdot x$   
 by (simp add: negate-eq1 mult-assoc2)

lemma (in vectorspace) add-minus-left-eq-diff:  
 $x \in V \implies y \in V \implies -x + y = y - x$   
 proof -  
 assume xy:  $x \in V \ y \in V$   
 hence  $-x + y = y + -x$  by (simp add: add-commute)  
 also from xy have  $\dots = y - x$  by (simp add: diff-eq1)  
 finally show ?thesis .  
 qed

lemma (in vectorspace) add-minus [simp]:  
 $x \in V \implies x + -x = 0$   
 by (simp add: diff-eq2)

lemma (in vectorspace) add-minus-left [simp]:  
 $x \in V \implies -x + x = 0$   
 by (simp add: diff-eq2 add-commute)

lemma (in vectorspace) minus-minus [simp]:  
 $x \in V \implies -(-x) = x$   
 by (simp add: negate-eq1 mult-assoc2)

lemma (in vectorspace) minus-zero [simp]:  
 $-(0::'a) = 0$   
 by (simp add: negate-eq1)

lemma (in vectorspace) minus-zero-iff [simp]:  
 $x \in V \implies (-x = 0) = (x = 0)$   
 proof  
 assume x:  $x \in V$   
 {  
 from x have  $x = -(-x)$  by (simp add: minus-minus)  
 also assume  $-x = 0$   
 also have  $- \dots = 0$  by (rule minus-zero)  
 finally show  $x = 0$  .  
 }  
 next  
 assume  $x = 0$   
 then show  $-x = 0$  by simp  
 }  
 qed

lemma (in vectorspace) add-minus-cancel [simp]:  
 $x \in V \implies y \in V \implies x + (-x + y) = y$   
 by (simp add: add-assoc [symmetric] del: add-commute)



**lemma** (in *vectorspace*) *minus-add-cancel* [simp]:

$x \in V \implies y \in V \implies -x + (x + y) = y$   
**by** (simp add: add-assoc [symmetric] del: add-commute)

**lemma** (in *vectorspace*) *minus-add-distrib* [simp]:

$x \in V \implies y \in V \implies -(x + y) = -x + -y$   
**by** (simp add: negate-eq1 add-mult-distrib1)

**lemma** (in *vectorspace*) *diff-zero* [simp]:

$x \in V \implies x - 0 = x$   
**by** (simp add: diff-eq1)

**lemma** (in *vectorspace*) *diff-zero-right* [simp]:

$x \in V \implies 0 - x = -x$   
**by** (simp add: diff-eq1)

**lemma** (in *vectorspace*) *add-left-cancel*:

$x \in V \implies y \in V \implies z \in V \implies (x + y = x + z) = (y = z)$

**proof**

**assume**  $x: x \in V$  **and**  $y: y \in V$  **and**  $z: z \in V$

{

**from**  $y$  **have**  $y = 0 + y$  **by** simp

**also from**  $x$   $y$  **have**  $\dots = (-x + x) + y$  **by** simp

**also from**  $x$   $y$  **have**  $\dots = -x + (x + y)$

**by** (simp add: add-assoc neg-closed)

**also assume**  $x + y = x + z$

**also from**  $x$   $z$  **have**  $-x + (x + z) = -x + x + z$

**by** (simp add: add-assoc [symmetric] neg-closed)

**also from**  $x$   $z$  **have**  $\dots = z$  **by** simp

**finally show**  $y = z$  .

**next**

**assume**  $y = z$

**then show**  $x + y = x + z$  **by** (simp only:)

}

**qed**

**lemma** (in *vectorspace*) *add-right-cancel*:

$x \in V \implies y \in V \implies z \in V \implies (y + x = z + x) = (y = z)$

**by** (simp only: add-commute add-left-cancel)

**lemma** (in *vectorspace*) *add-assoc-cong*:

$x \in V \implies y \in V \implies x' \in V \implies y' \in V \implies z \in V$   
 $\implies x + y = x' + y' \implies x + (y + z) = x' + (y' + z)$   
**by** (simp only: add-assoc [symmetric])

**lemma** (in *vectorspace*) *mult-left-commute*:

$x \in V \implies a \cdot b \cdot x = b \cdot a \cdot x$

**by** (simp add: real-mult-commute mult-assoc2)

**lemma** (in *vectorspace*) *mult-zero-uniq*:

$x \in V \implies x \neq 0 \implies a \cdot x = 0 \implies a = 0$

**proof** (rule classical)

**assume**  $a: a \neq 0$

**assume**  $x: x \in V$   $x \neq 0$  **and**  $ax: a \cdot x = 0$

from  $x$   $a$  have  $x = (\text{inverse } a * a) \cdot x$  **by** *simp*  
 also have  $\dots = \text{inverse } a \cdot (a \cdot x)$  **by** *(rule mult-assoc)*  
 also from  $ax$  have  $\dots = \text{inverse } a \cdot 0$  **by** *simp*  
 also have  $\dots = 0$  **by** *simp*  
 finally have  $x = 0$  .  
 thus  $a = 0$  **by** *contradiction*  
 qed

**lemma** (in *vectorspace*) *mult-left-cancel*:  
 $x \in V \implies y \in V \implies a \neq 0 \implies (a \cdot x = a \cdot y) = (x = y)$   
**proof**  
 assume  $x: x \in V$  and  $y: y \in V$  and  $a: a \neq 0$   
 from  $x$  have  $x = 1 \cdot x$  **by** *simp*  
 also from  $a$  have  $\dots = (\text{inverse } a * a) \cdot x$  **by** *simp*  
 also from  $x$  have  $\dots = \text{inverse } a \cdot (a \cdot x)$   
     **by** *(simp only: mult-assoc)*  
 also assume  $a \cdot x = a \cdot y$   
 also from  $a$   $y$  have  $\text{inverse } a \cdot \dots = y$   
     **by** *(simp add: mult-assoc2)*  
 finally show  $x = y$  .  
**next**  
 assume  $x = y$   
 then show  $a \cdot x = a \cdot y$  **by** *(simp only:)*  
 qed

**lemma** (in *vectorspace*) *mult-right-cancel*:  
 $x \in V \implies x \neq 0 \implies (a \cdot x = b \cdot x) = (a = b)$   
**proof**  
 assume  $x: x \in V$  and *neg*:  $x \neq 0$   
 {  
     from  $x$  have  $(a - b) \cdot x = a \cdot x - b \cdot x$   
         **by** *(simp add: diff-mult-distrib2)*  
     also assume  $a \cdot x = b \cdot x$   
     with  $x$  have  $a \cdot x - b \cdot x = 0$  **by** *simp*  
     finally have  $(a - b) \cdot x = 0$  .  
     with  $x$  *neg* have  $a - b = 0$  **by** *(rule mult-zero-uniq)*  
     thus  $a = b$  **by** *simp*  
 }  
**next**  
 assume  $a = b$   
 then show  $a \cdot x = b \cdot x$  **by** *(simp only:)*  
 }  
 qed

**lemma** (in *vectorspace*) *eq-diff-eq*:  
 $x \in V \implies y \in V \implies z \in V \implies (x = z - y) = (x + y = z)$   
**proof**  
 assume  $x: x \in V$  and  $y: y \in V$  and  $z: z \in V$   
 {  
     assume  $x = z - y$   
     hence  $x + y = z - y + y$  **by** *simp*  
     also from  $y$   $z$  have  $\dots = z + - y + y$   
         **by** *(simp add: diff-eq1)*  
     also have  $\dots = z + (- y + y)$   
         **by** *(rule add-assoc) (simp-all add: y z)*

```

    also from  $y\ z$  have  $\dots = z + 0$ 
      by (simp only: add-minus-left)
    also from  $z$  have  $\dots = z$ 
      by (simp only: add-zero-right)
    finally show  $x + y = z$  .
  next
    assume  $x + y = z$ 
    hence  $z - y = (x + y) - y$  by simp
    also from  $x\ y$  have  $\dots = x + y + -\ y$ 
      by (simp add: diff-eq1)
    also have  $\dots = x + (y + -\ y)$ 
      by (rule add-assoc) (simp-all add: x y)
    also from  $x\ y$  have  $\dots = x$  by simp
    finally show  $x = z - y$  ..
  }
qed

```

lemma (in vectorspace) add-minus-eq-minus:  
 $x \in V \implies y \in V \implies x + y = 0 \implies x = -\ y$

```

proof -
  assume  $x: x \in V$  and  $y: y \in V$ 
  from  $x\ y$  have  $x = (-\ y + y) + x$  by simp
  also from  $x\ y$  have  $\dots = -\ y + (x + y)$  by (simp add: add-ac)
  also assume  $x + y = 0$ 
  also from  $y$  have  $- y + 0 = -\ y$  by simp
  finally show  $x = -\ y$  .
qed

```

lemma (in vectorspace) add-minus-eq:  
 $x \in V \implies y \in V \implies x - y = 0 \implies x = y$

```

proof -
  assume  $x: x \in V$  and  $y: y \in V$ 
  assume  $x - y = 0$ 
  with  $x\ y$  have eq:  $x + -\ y = 0$  by (simp add: diff-eq1)
  with - - have  $x = -\ (-\ y)$ 
    by (rule add-minus-eq-minus) (simp-all add: x y)
  with  $x\ y$  show  $x = y$  by simp
qed

```

lemma (in vectorspace) add-diff-swap:  
 $a \in V \implies b \in V \implies c \in V \implies d \in V \implies a + b = c + d$   
 $\implies a - c = d - b$

```

proof -
  assume  $vs: a \in V\ b \in V\ c \in V\ d \in V$ 
  and eq:  $a + b = c + d$ 
  then have  $- c + (a + b) = - c + (c + d)$ 
    by (simp add: add-left-cancel)
  also have  $\dots = d$  by (rule minus-add-cancel)
  finally have eq:  $- c + (a + b) = d$  .
  from  $vs$  have  $a - c = (- c + (a + b)) + -\ b$ 
    by (simp add: add-ac diff-eq1)
  also from  $vs\ eq$  have  $\dots = d + -\ b$ 
    by (simp add: add-right-cancel)
  also from  $vs$  have  $\dots = d - b$  by (simp add: diff-eq2)

```

finally show  $a - c = d - b$  .  
qed

lemma (in *vectorspace*) *vs-add-cancel-21*:

$x \in V \implies y \in V \implies z \in V \implies u \in V$   
 $\implies (x + (y + z) = y + u) = (x + z = u)$

proof

assume *vs*:  $x \in V \ y \in V \ z \in V \ u \in V$

{  
 from *vs* have  $x + z = -y + y + (x + z)$  by *simp*  
 also have  $\dots = -y + (y + (x + z))$   
 by (*rule add-assoc*) (*simp-all add: vs*)  
 also from *vs* have  $y + (x + z) = x + (y + z)$   
 by (*simp add: add-ac*)  
 also assume  $x + (y + z) = y + u$   
 also from *vs* have  $-y + (y + u) = u$  by *simp*  
 finally show  $x + z = u$  .

next

assume  $x + z = u$   
 with *vs* show  $x + (y + z) = y + u$   
 by (*simp only: add-left-commute [of x]*)

}

qed

lemma (in *vectorspace*) *add-cancel-end*:

$x \in V \implies y \in V \implies z \in V \implies (x + (y + z) = y) = (x = -z)$

proof

assume *vs*:  $x \in V \ y \in V \ z \in V$

{  
 assume  $x + (y + z) = y$   
 with *vs* have  $(x + z) + y = 0 + y$   
 by (*simp add: add-ac*)  
 with *vs* have  $x + z = 0$   
 by (*simp only: add-right-cancel add-closed zero*)  
 with *vs* show  $x = -z$  by (*simp add: add-minus-eq-minus*)

next

assume *eq*:  $x = -z$   
 hence  $x + (y + z) = -z + (y + z)$  by *simp*  
 also have  $\dots = y + (-z + z)$   
 by (*rule add-left-commute*) (*simp-all add: vs*)  
 also from *vs* have  $\dots = y$  by *simp*  
 finally show  $x + (y + z) = y$  .

}

qed

end

## 4 Subspaces

theory *Subspace* imports *VectorSpace* begin

## 4.1 Definition

A non-empty subset  $U$  of a vector space  $V$  is a *subspace* of  $V$ , iff  $U$  is closed under addition and scalar multiplication.

```

locale subspace = var U + var V +
  assumes non-empty [iff, intro]:  $U \neq \{\}$ 
  and subset [iff]:  $U \subseteq V$ 
  and add-closed [iff]:  $x \in U \implies y \in U \implies x + y \in U$ 
  and mult-closed [iff]:  $x \in U \implies a \cdot x \in U$ 

```

```

declare vectorspace.intro [intro?] subspace.intro [intro?]

```

```

syntax (symbols)
  subspace :: 'a set  $\Rightarrow$  'a set  $\Rightarrow$  bool    (infix  $\leq$  50)

```

```

lemma subspace-subset [elim]:  $U \leq V \implies U \subseteq V$ 
by (rule subspace.subset)

```

```

lemma (in subspace) subsetD [iff]:  $x \in U \implies x \in V$ 
using subset by blast

```

```

lemma subspaceD [elim]:  $U \leq V \implies x \in U \implies x \in V$ 
by (rule subspace.subsetD)

```

```

lemma rev-subspaceD [elim?]:  $x \in U \implies U \leq V \implies x \in V$ 
by (rule subspace.subsetD)

```

```

lemma (in subspace) diff-closed [iff]:
  includes vectorspace
  shows  $x \in U \implies y \in U \implies x - y \in U$ 
by (simp add: diff-eq1 negate-eq1)

```

Similar as for linear spaces, the existence of the zero element in every subspace follows from the non-emptiness of the carrier set and by vector space laws.

```

lemma (in subspace) zero [intro]:
  includes vectorspace
  shows  $0 \in U$ 
proof -
  have  $U \neq \{\}$  by (rule U-V.non-empty)
  then obtain  $x$  where  $x: x \in U$  by blast
  hence  $x \in V$  .. hence  $0 = x - x$  by simp
  also have  $\dots \in U$  by (rule U-V.diff-closed)
  finally show ?thesis .
qed

```

```

lemma (in subspace) neg-closed [iff]:
  includes vectorspace
  shows  $x \in U \implies -x \in U$ 
by (simp add: negate-eq1)

```

Further derived laws: every subspace is a vector space.

```

lemma (in subspace) vectorspace [iff]:

```

```

includes vectorspace
shows vectorspace U
proof
  show  $U \neq \{\}$  ..
  fix  $x\ y\ z$  assume  $x: x \in U$  and  $y: y \in U$  and  $z: z \in U$ 
  fix  $a\ b :: \text{real}$ 
  from  $x\ y$  show  $x + y \in U$  by simp
  from  $x$  show  $a \cdot x \in U$  by simp
  from  $x\ y\ z$  show  $(x + y) + z = x + (y + z)$  by (simp add: add-ac)
  from  $x\ y$  show  $x + y = y + x$  by (simp add: add-ac)
  from  $x$  show  $x - x = 0$  by simp
  from  $x$  show  $0 + x = x$  by simp
  from  $x\ y$  show  $a \cdot (x + y) = a \cdot x + a \cdot y$  by (simp add: distrib)
  from  $x$  show  $(a + b) \cdot x = a \cdot x + b \cdot x$  by (simp add: distrib)
  from  $x$  show  $(a * b) \cdot x = a \cdot b \cdot x$  by (simp add: mult-assoc)
  from  $x$  show  $1 \cdot x = x$  by simp
  from  $x$  show  $-x = -1 \cdot x$  by (simp add: negate-eq1)
  from  $x\ y$  show  $x - y = x + -y$  by (simp add: diff-eq1)
qed

```

The subspace relation is reflexive.

```

lemma (in vectorspace) subspace-refl [intro]:  $V \trianglelefteq V$ 
proof
  show  $V \neq \{\}$  ..
  show  $V \subseteq V$  ..
  fix  $x\ y$  assume  $x: x \in V$  and  $y: y \in V$ 
  fix  $a :: \text{real}$ 
  from  $x\ y$  show  $x + y \in V$  by simp
  from  $x$  show  $a \cdot x \in V$  by simp
qed

```

The subspace relation is transitive.

```

lemma (in vectorspace) subspace-trans [trans]:
   $U \trianglelefteq V \implies V \trianglelefteq W \implies U \trianglelefteq W$ 
proof
  assume  $uv: U \trianglelefteq V$  and  $vw: V \trianglelefteq W$ 
  from  $uv$  show  $U \neq \{\}$  by (rule subspace.non-empty)
  show  $U \subseteq W$ 
  proof -
    from  $uv$  have  $U \subseteq V$  by (rule subspace.subset)
    also from  $vw$  have  $V \subseteq W$  by (rule subspace.subset)
    finally show ?thesis .
  qed
  fix  $x\ y$  assume  $x: x \in U$  and  $y: y \in U$ 
  from  $uv$  and  $x\ y$  show  $x + y \in U$  by (rule subspace.add-closed)
  from  $uv$  and  $x$  show  $\bigwedge a. a \cdot x \in U$  by (rule subspace.mult-closed)
qed

```

## 4.2 Linear closure

The *linear closure* of a vector  $x$  is the set of all scalar multiples of  $x$ .

```

constdefs
  lin :: ('a::{minus, plus, zero})  $\Rightarrow$  'a set

```

$\text{lin } x \equiv \{a \cdot x \mid a. \text{ True}\}$

**lemma** *linI* [*intro*]:  $y = a \cdot x \implies y \in \text{lin } x$   
**by** (*unfold lin-def*) *blast*

**lemma** *linI'* [*iff*]:  $a \cdot x \in \text{lin } x$   
**by** (*unfold lin-def*) *blast*

**lemma** *linE* [*elim*]:  
 $x \in \text{lin } v \implies (\bigwedge a::\text{real}. x = a \cdot v \implies C) \implies C$   
**by** (*unfold lin-def*) *blast*

Every vector is contained in its linear closure.

**lemma** (*in vectorspace*) *x-lin-x* [*iff*]:  $x \in V \implies x \in \text{lin } x$   
**proof** –  
**assume**  $x \in V$   
**hence**  $x = 1 \cdot x$  **by** *simp*  
**also have**  $\dots \in \text{lin } x$  ..  
**finally show** *?thesis* .  
**qed**

**lemma** (*in vectorspace*) *0-lin-x* [*iff*]:  $x \in V \implies 0 \in \text{lin } x$   
**proof**  
**assume**  $x \in V$   
**thus**  $0 = 0 \cdot x$  **by** *simp*  
**qed**

Any linear closure is a subspace.

**lemma** (*in vectorspace*) *lin-subspace* [*intro*]:  
 $x \in V \implies \text{lin } x \trianglelefteq V$   
**proof**  
**assume**  $x: x \in V$   
**thus**  $\text{lin } x \neq \{\}$  **by** (*auto simp add: x-lin-x*)  
**show**  $\text{lin } x \subseteq V$   
**proof**  
**fix**  $x'$  **assume**  $x' \in \text{lin } x$   
**then obtain**  $a$  **where**  $x' = a \cdot x$  ..  
**with**  $x$  **show**  $x' \in V$  **by** *simp*  
**qed**  
**fix**  $x' x''$  **assume**  $x': x' \in \text{lin } x$  **and**  $x'': x'' \in \text{lin } x$   
**show**  $x' + x'' \in \text{lin } x$   
**proof** –  
**from**  $x'$  **obtain**  $a'$  **where**  $x' = a' \cdot x$  ..  
**moreover from**  $x''$  **obtain**  $a''$  **where**  $x'' = a'' \cdot x$  ..  
**ultimately have**  $x' + x'' = (a' + a'') \cdot x$   
**using**  $x$  **by** (*simp add: distrib*)  
**also have**  $\dots \in \text{lin } x$  ..  
**finally show** *?thesis* .  
**qed**  
**fix**  $a :: \text{real}$   
**show**  $a \cdot x' \in \text{lin } x$   
**proof** –  
**from**  $x'$  **obtain**  $a'$  **where**  $x' = a' \cdot x$  ..  
**with**  $x$  **have**  $a \cdot x' = (a * a') \cdot x$  **by** (*simp add: mult-assoc*)

```

    also have ... ∈ lin x ..
    finally show ?thesis .
qed
qed

```

Any linear closure is a vector space.

```

lemma (in vectorspace) lin-vectorspace [intro]:
  x ∈ V ⇒ vectorspace (lin x)
by (rule subspace.vectorspace) (rule lin-subspace)

```

### 4.3 Sum of two vectorspaces

The *sum* of two vectorspaces  $U$  and  $V$  is the set of all sums of elements from  $U$  and  $V$ .

```

instance set :: (plus) plus ..

```

```

defs (overloaded)
  sum-def: U + V ≡ {u + v | u v. u ∈ U ∧ v ∈ V}

```

```

lemma sumE [elim]:
  x ∈ U + V ⇒ (⋀u v. x = u + v ⇒ u ∈ U ⇒ v ∈ V ⇒ C) ⇒ C
by (unfold sum-def) blast

```

```

lemma sumI [intro]:
  u ∈ U ⇒ v ∈ V ⇒ x = u + v ⇒ x ∈ U + V
by (unfold sum-def) blast

```

```

lemma sumI' [intro]:
  u ∈ U ⇒ v ∈ V ⇒ u + v ∈ U + V
by (unfold sum-def) blast

```

$U$  is a subspace of  $U + V$ .

```

lemma subspace-sum1 [iff]:
  includes vectorspace U + vectorspace V
shows U ≤ U + V

```

```

proof
  show U ≠ {} ..
  show U ⊆ U + V
proof
  fix x assume x: x ∈ U
  moreover have 0 ∈ V ..
  ultimately have x + 0 ∈ U + V ..
  with x show x ∈ U + V by simp
qed
fix x y assume x: x ∈ U and y ∈ U
thus x + y ∈ U by simp
from x show ⋀a. a · x ∈ U by simp
qed

```

The sum of two subspaces is again a subspace.

```

lemma sum-subspace [intro?]:
  includes subspace U E + vectorspace E + subspace V E

```



```

shows  $U + V \trianglelefteq E$ 
proof
  have  $0 \in U + V$ 
  proof
    show  $0 \in U$  ..
    show  $0 \in V$  ..
    show  $(0::'a) = 0 + 0$  by simp
  qed
thus  $U + V \neq \{\}$  by blast
show  $U + V \subseteq E$ 
proof
  fix  $x$  assume  $x \in U + V$ 
  then obtain  $u\ v$  where  $x = u + v$  and
     $u: u \in U$  and  $v: v \in V$  ..
  have  $U \trianglelefteq E$  . with  $u$  have  $u \in E$  ..
  moreover have  $V \trianglelefteq E$  . with  $v$  have  $v \in E$  ..
  ultimately show  $x \in E$  using  $x$  by simp
qed
fix  $x\ y$  assume  $x: x \in U + V$  and  $y: y \in U + V$ 
show  $x + y \in U + V$ 
proof -
  from  $x$  obtain  $ux\ vx$  where  $x = ux + vx$  and  $ux \in U$  and  $vx \in V$  ..
  moreover
  from  $y$  obtain  $uy\ vy$  where  $y = uy + vy$  and  $uy \in U$  and  $vy \in V$  ..
  ultimately
  have  $ux + uy \in U$ 
  and  $vx + vy \in V$ 
  and  $x + y = (ux + uy) + (vx + vy)$ 
  using  $x\ y$  by (simp-all add: add-ac)
  thus ?thesis ..
qed
fix  $a$  show  $a \cdot x \in U + V$ 
proof -
  from  $x$  obtain  $u\ v$  where  $x = u + v$  and  $u \in U$  and  $v \in V$  ..
  hence  $a \cdot u \in U$  and  $a \cdot v \in V$ 
  and  $a \cdot x = (a \cdot u) + (a \cdot v)$  by (simp-all add: distrib)
  thus ?thesis ..
qed
qed

```

The sum of two subspaces is a vectorspace.

**lemma** *sum-vs* [*intro?*]:

$U \trianglelefteq E \implies V \trianglelefteq E \implies \text{vectorspace } E \implies \text{vectorspace } (U + V)$   
**by** (*rule subspace.vectorspace*) (*rule sum-subspace*)

## 4.4 Direct sums

The sum of  $U$  and  $V$  is called *direct*, iff the zero element is the only common element of  $U$  and  $V$ . For every element  $x$  of the direct sum of  $U$  and  $V$  the decomposition in  $x = u + v$  with  $u \in U$  and  $v \in V$  is unique.

**lemma** *decomp*:

**includes** *vectorspace*  $E + \text{subspace } U\ E + \text{subspace } V\ E$   
**assumes** *direct*:  $U \cap V = \{0\}$

```

    and u1: u1 ∈ U and u2: u2 ∈ U
    and v1: v1 ∈ V and v2: v2 ∈ V
    and sum: u1 + v1 = u2 + v2
  shows u1 = u2 ∧ v1 = v2
proof
  have U: vectorspace U by (rule subspace.vectorspace)
  have V: vectorspace V by (rule subspace.vectorspace)
  from u1 u2 v1 v2 and sum have eq: u1 - u2 = v2 - v1
    by (simp add: add-diff-swap)
  from u1 u2 have u: u1 - u2 ∈ U
    by (rule vectorspace.diff-closed [OF U])
  with eq have v': v2 - v1 ∈ U by (simp only:)
  from v2 v1 have v: v2 - v1 ∈ V
    by (rule vectorspace.diff-closed [OF V])
  with eq have u': u1 - u2 ∈ V by (simp only:)

  show u1 = u2
proof (rule add-minus-eq)
  show u1 ∈ E ..
  show u2 ∈ E ..
  from u u' and direct show u1 - u2 = 0 by force
qed
show v1 = v2
proof (rule add-minus-eq [symmetric])
  show v1 ∈ E ..
  show v2 ∈ E ..
  from v v' and direct show v2 - v1 = 0 by force
qed
qed

```

An application of the previous lemma will be used in the proof of the Hahn-Banach Theorem (see page 40): for any element  $y + a \cdot x_0$  of the direct sum of a vectorspace  $H$  and the linear closure of  $x_0$  the components  $y \in H$  and  $a$  are uniquely determined.

```

lemma decomp-H':
  includes vectorspace E + subspace H E
  assumes y1: y1 ∈ H and y2: y2 ∈ H
    and x': x' ∉ H x' ∈ E x' ≠ 0
    and eq: y1 + a1 · x' = y2 + a2 · x'
  shows y1 = y2 ∧ a1 = a2
proof
  have c: y1 = y2 ∧ a1 · x' = a2 · x'
proof (rule decomp)
  show a1 · x' ∈ lin x' ..
  show a2 · x' ∈ lin x' ..
  show H ∩ lin x' = {0}
proof
  show H ∩ lin x' ⊆ {0}
proof
  fix x assume x: x ∈ H ∩ lin x'
  then obtain a where xx': x = a · x'
  by blast
  have x = 0
proof cases

```

```

    assume  $a = 0$ 
    with  $xx'$  and  $x'$  show ?thesis by simp
  next
    assume  $a: a \neq 0$ 
    from  $x$  have  $x \in H$  ..
    with  $xx'$  have  $\text{inverse } a \cdot a \cdot x' \in H$  by simp
    with  $a$  and  $x'$  have  $x' \in H$  by (simp add: mult-assoc2)
    thus ?thesis by contradiction
  qed
  thus  $x \in \{0\}$  ..
qed
show  $\{0\} \subseteq H \cap \text{lin } x'$ 
proof -
  have  $0 \in H$  ..
  moreover have  $0 \in \text{lin } x'$  ..
  ultimately show ?thesis by blast
qed
qed
show  $\text{lin } x' \trianglelefteq E$  ..
qed
thus  $y1 = y2$  ..
from  $c$  have  $a1 \cdot x' = a2 \cdot x'$  ..
with  $x'$  show  $a1 = a2$  by (simp add: mult-right-cancel)
qed

```

Since for any element  $y + a \cdot x'$  of the direct sum of a vectorspace  $H$  and the linear closure of  $x'$  the components  $y \in H$  and  $a$  are unique, it follows from  $y \in H$  that  $a = 0$ .

**lemma** *decomp- $H'$ - $H$ :*  
**includes** *vectorspace*  $E + \text{subspace } H \ E$   
**assumes**  $t: t \in H$   
**and**  $x': x' \notin H \ x' \in E \ x' \neq 0$   
**shows**  $(\text{SOME } (y, a). t = y + a \cdot x' \wedge y \in H) = (t, 0)$   
**proof** (*rule, simp-all only: split-paired-all split-conv*)  
**from**  $t \ x'$  **show**  $t = t + 0 \cdot x' \wedge t \in H$  **by** *simp*  
**fix**  $y$  **and**  $a$  **assume**  $ya: t = y + a \cdot x' \wedge y \in H$   
**have**  $y = t \wedge a = 0$   
**proof** (*rule decomp- $H'$* )  
**from**  $ya \ x'$  **show**  $y + a \cdot x' = t + 0 \cdot x'$  **by** *simp*  
**from**  $ya$  **show**  $y \in H$  ..  
**qed**  
**with**  $t \ x'$  **show**  $(y, a) = (y + a \cdot x', 0)$  **by** *simp*  
**qed**

The components  $y \in H$  and  $a$  in  $y + a \cdot x'$  are unique, so the function  $h'$  defined by  $h'(y + a \cdot x') = h y + a \cdot \xi$  is definite.

**lemma**  *$h'$ -definite:*  
**includes** *var*  $H$   
**assumes**  $h'$ -def:  
 $h' \equiv (\lambda x. \text{let } (y, a) = \text{SOME } (y, a). (x = y + a \cdot x' \wedge y \in H) \\ \text{in } (h y) + a * \xi)$   
**and**  $x: x = y + a \cdot x'$   
**includes** *vectorspace*  $E + \text{subspace } H \ E$

```

assumes  $y: y \in H$ 
and  $x': x' \notin H \ x' \in E \ x' \neq 0$ 
shows  $h' x = h y + a * x$ 
proof -
  from  $x y x'$  have  $x \in H + \text{lin } x'$  by auto
  have  $\exists! p. (\lambda(y, a). x = y + a \cdot x' \wedge y \in H) p$  (is  $\exists! p. ?P p$ )
  proof
    from  $x y$  show  $\exists p. ?P p$  by blast
    fix  $p q$  assume  $p: ?P p$  and  $q: ?P q$ 
    show  $p = q$ 
    proof -
      from  $p$  have  $xp: x = \text{fst } p + \text{snd } p \cdot x' \wedge \text{fst } p \in H$ 
      by (cases  $p$ ) simp
      from  $q$  have  $xq: x = \text{fst } q + \text{snd } q \cdot x' \wedge \text{fst } q \in H$ 
      by (cases  $q$ ) simp
      have  $\text{fst } p = \text{fst } q \wedge \text{snd } p = \text{snd } q$ 
      proof (rule decomp-H')
        from  $xp$  show  $\text{fst } p \in H$  ..
        from  $xq$  show  $\text{fst } q \in H$  ..
        from  $xp$  and  $xq$  show  $\text{fst } p + \text{snd } p \cdot x' = \text{fst } q + \text{snd } q \cdot x'$ 
        by simp
      apply-end assumption+
    qed
    thus  $?thesis$  by (cases  $p$ , cases  $q$ ) simp
  qed
qed
end

```

## 5 Normed vector spaces

theory *NormedSpace* imports *Subspace* begin

### 5.1 Quasinorms

A *seminorm*  $\|\cdot\|$  is a function on a real vector space into the reals that has the following properties: it is positive definite, absolute homogenous and subadditive.

```

locale norm-syntax =
  fixes norm :: 'a  $\Rightarrow$  real  ( $\|\cdot\|$ )

locale seminorm = var  $V$  + norm-syntax +
  assumes ge-zero [iff?]:  $x \in V \Longrightarrow 0 \leq \|x\|$ 
  and abs-homogenous [iff?]:  $x \in V \Longrightarrow \|a \cdot x\| = |a| * \|x\|$ 
  and subadditive [iff?]:  $x \in V \Longrightarrow y \in V \Longrightarrow \|x + y\| \leq \|x\| + \|y\|$ 

declare seminorm.intro [intro?]

```

```

lemma (in seminorm) diff-subadditive:
  includes vectorspace
  shows  $x \in V \implies y \in V \implies \|x - y\| \leq \|x\| + \|y\|$ 
proof -
  assume  $x: x \in V$  and  $y: y \in V$ 
  hence  $x - y = x + -1 \cdot y$ 
    by (simp add: diff-eq2 negate-eq2a)
  also from  $x\ y$  have  $\|\dots\| \leq \|x\| + \|-1 \cdot y\|$ 
    by (simp add: subadditive)
  also from  $y$  have  $\|-1 \cdot y\| = |-1| * \|y\|$ 
    by (rule abs-homogenous)
  also have  $\dots = \|y\|$  by simp
  finally show ?thesis .
qed

```

```

lemma (in seminorm) minus:
  includes vectorspace
  shows  $x \in V \implies \|-x\| = \|x\|$ 
proof -
  assume  $x: x \in V$ 
  hence  $-x = -1 \cdot x$  by (simp only: negate-eq1)
  also from  $x$  have  $\|\dots\| = |-1| * \|x\|$ 
    by (rule abs-homogenous)
  also have  $\dots = \|x\|$  by simp
  finally show ?thesis .
qed

```

## 5.2 Norms

A norm  $\|\cdot\|$  is a seminorm that maps only the  $0$  vector to  $0$ .

```

locale norm = seminorm +
  assumes zero-iff [iff]:  $x \in V \implies (\|x\| = 0) = (x = 0)$ 

```

## 5.3 Normed vector spaces

A vector space together with a norm is called a *normed space*.

```

locale normed-vectorspace = vectorspace + norm

```

```

declare normed-vectorspace.intro [intro?]

```

```

lemma (in normed-vectorspace) gt-zero [intro?]:
   $x \in V \implies x \neq 0 \implies 0 < \|x\|$ 
proof -
  assume  $x: x \in V$  and  $neg: x \neq 0$ 
  from  $x$  have  $0 \leq \|x\|$  ..
  also have [symmetric]:  $\dots \neq 0$ 
  proof
    assume  $\|x\| = 0$ 
    with  $x$  have  $x = 0$  by simp
    with  $neg$  show False by contradiction
  qed
  finally show ?thesis .
qed

```

Any subspace of a normed vector space is again a normed vectorspace.

```

lemma subspace-normed-vs [intro?]:
  includes subspace  $F\ E + \text{normed-vectorspace}\ E$ 
  shows normed-vectorspace  $F\ \text{norm}$ 
proof
  show vectorspace  $F$  by (rule vectorspace)
  have seminorm  $E\ \text{norm}$  . with subset show seminorm  $F\ \text{norm}$ 
    by (simp add: seminorm-def)
  have norm-axioms  $E\ \text{norm}$  . with subset show norm-axioms  $F\ \text{norm}$ 
    by (simp add: norm-axioms-def)
qed

end

```

## 6 Linearforms

**theory** *Linearform* **imports** *VectorSpace* **begin**

A *linear form* is a function on a vector space into the reals that is additive and multiplicative.

```

locale linearform = var  $V + \text{var}\ f +$ 
  assumes add [iff]:  $x \in V \implies y \in V \implies f\ (x + y) = f\ x + f\ y$ 
  and mult [iff]:  $x \in V \implies f\ (a \cdot x) = a * f\ x$ 

```

**declare** *linearform.intro* [intro?]

```

lemma (in linearform) neg [iff]:
  includes vectorspace
  shows  $x \in V \implies f\ (-\ x) = -\ f\ x$ 
proof -
  assume  $x: x \in V$ 
  hence  $f\ (-\ x) = f\ ((- \ 1) \cdot x)$  by (simp add: negate-eq1)
  also from  $x$  have  $\dots = (- \ 1) * (f\ x)$  by (rule mult)
  also from  $x$  have  $\dots = -\ (f\ x)$  by simp
  finally show ?thesis .
qed

```

```

lemma (in linearform) diff [iff]:
  includes vectorspace
  shows  $x \in V \implies y \in V \implies f\ (x - y) = f\ x - f\ y$ 
proof -
  assume  $x: x \in V$  and  $y: y \in V$ 
  hence  $x - y = x + -\ y$  by (rule diff-eq1)
  also have  $f\ \dots = f\ x + f\ (-\ y)$  by (rule add) (simp-all add:  $x\ y$ )
  also from  $- \ y$  have  $f\ (-\ y) = -\ f\ y$  by (rule neg)
  finally show ?thesis by simp
qed

```

Every linear form yields 0 for the 0 vector.

```

lemma (in linearform) zero [iff]:
  includes vectorspace

```

```

  shows  $f\ 0 = 0$ 
proof -
  have  $f\ 0 = f\ (0 - 0)$  by simp
  also have  $\dots = f\ 0 - f\ 0$  by (rule diff) simp-all
  also have  $\dots = 0$  by simp
  finally show ?thesis .
qed

end

```

## 7 An order on functions

**theory** *FunctionOrder* **imports** *Subspace Linearform* **begin**

### 7.1 The graph of a function

We define the *graph* of a (real) function  $f$  with domain  $F$  as the set

$$\{(x, f\ x). x \in F\}$$

So we are modeling partial functions by specifying the domain and the mapping function. We use the term “function” also for its graph.

**types**  $'a\ graph = ('a \times real)\ set$

**constdefs**

```

graph :: 'a set  $\Rightarrow$  ('a  $\Rightarrow$  real)  $\Rightarrow$  'a graph
graph F f  $\equiv$   $\{(x, f\ x) \mid x. x \in F\}$ 

```

**lemma** *graphI* [*intro*]:  $x \in F \Longrightarrow (x, f\ x) \in graph\ F\ f$   
**by** (*unfold graph-def*) *blast*

**lemma** *graphI2* [*intro?*]:  $x \in F \Longrightarrow \exists t \in graph\ F\ f. t = (x, f\ x)$   
**by** (*unfold graph-def*) *blast*

**lemma** *graphE* [*elim?*]:  
 $(x, y) \in graph\ F\ f \Longrightarrow (x \in F \Longrightarrow y = f\ x \Longrightarrow C) \Longrightarrow C$   
**by** (*unfold graph-def*) *blast*

### 7.2 Functions ordered by domain extension

A function  $h'$  is an extension of  $h$ , iff the graph of  $h$  is a subset of the graph of  $h'$ .

**lemma** *graph-extI*:  
 $(\bigwedge x. x \in H \Longrightarrow h\ x = h'\ x) \Longrightarrow H \subseteq H'$   
 $\Longrightarrow graph\ H\ h \subseteq graph\ H'\ h'$   
**by** (*unfold graph-def*) *blast*

**lemma** *graph-extD1* [*dest?*]:  
 $graph\ H\ h \subseteq graph\ H'\ h' \Longrightarrow x \in H \Longrightarrow h\ x = h'\ x$   
**by** (*unfold graph-def*) *blast*

```

lemma graph-extD2 [dest?]:
  graph H h  $\subseteq$  graph H' h'  $\implies$  H  $\subseteq$  H'
  by (unfold graph-def) blast

```

### 7.3 Domain and function of a graph

The inverse functions to *graph* are *domain* and *funct*.

```

constdefs
  domain :: 'a graph  $\Rightarrow$  'a set
  domain g  $\equiv$  {x.  $\exists$  y. (x, y)  $\in$  g}

  funct :: 'a graph  $\Rightarrow$  ('a  $\Rightarrow$  real)
  funct g  $\equiv$   $\lambda$ x. (SOME y. (x, y)  $\in$  g)

```

The following lemma states that *g* is the graph of a function if the relation induced by *g* is unique.

```

lemma graph-domain-funct:
  assumes uniq:  $\bigwedge$ x y z. (x, y)  $\in$  g  $\implies$  (x, z)  $\in$  g  $\implies$  z = y
  shows graph (domain g) (funct g) = g
proof (unfold domain-def funct-def graph-def, auto)
  fix a b assume (a, b)  $\in$  g
  show (a, SOME y. (a, y)  $\in$  g)  $\in$  g by (rule someI2)
  show  $\exists$  y. (a, y)  $\in$  g ..
  show b = (SOME y. (a, y)  $\in$  g)
  proof (rule some-equality [symmetric])
    fix y assume (a, y)  $\in$  g
    show y = b by (rule uniq)
  qed
qed

```

### 7.4 Norm-preserving extensions of a function

Given a linear form *f* on the space *F* and a seminorm *p* on *E*. The set of all linear extensions of *f*, to superspaces *H* of *F*, which are bounded by *p*, is defined as follows.

```

constdefs
  norm-pres-extensions ::
    'a::{plus, minus, zero} set  $\Rightarrow$  ('a  $\Rightarrow$  real)  $\Rightarrow$  'a set  $\Rightarrow$  ('a  $\Rightarrow$  real)
     $\Rightarrow$  'a graph set
  norm-pres-extensions E p F f
     $\equiv$  {g.  $\exists$  H h. g = graph H h
       $\wedge$  linearform H h
       $\wedge$  H  $\trianglelefteq$  E
       $\wedge$  F  $\trianglelefteq$  H
       $\wedge$  graph F f  $\subseteq$  graph H h
       $\wedge$  ( $\forall$  x  $\in$  H. h x  $\leq$  p x)}

```

```

lemma norm-pres-extensionE [elim]:
  g  $\in$  norm-pres-extensions E p F f
   $\implies$  ( $\bigwedge$  H h. g = graph H h  $\implies$  linearform H h
     $\implies$  H  $\trianglelefteq$  E  $\implies$  F  $\trianglelefteq$  H  $\implies$  graph F f  $\subseteq$  graph H h
     $\implies$   $\forall$  x  $\in$  H. h x  $\leq$  p x  $\implies$  C)  $\implies$  C

```



**by** (*unfold norm-pres-extensions-def*) *blast*

**lemma** *norm-pres-extensionI2* [*intro*]:

*linearform H h*  $\implies H \trianglelefteq E \implies F \trianglelefteq H$   
 $\implies \text{graph } F f \subseteq \text{graph } H h \implies \forall x \in H. h\ x \leq p\ x$   
 $\implies \text{graph } H h \in \text{norm-pres-extensions } E\ p\ F\ f$   
**by** (*unfold norm-pres-extensions-def*) *blast*

**lemma** *norm-pres-extensionI*:

$\exists H\ h. g = \text{graph } H\ h$   
 $\wedge \text{linearform } H\ h$   
 $\wedge H \trianglelefteq E$   
 $\wedge F \trianglelefteq H$   
 $\wedge \text{graph } F\ f \subseteq \text{graph } H\ h$   
 $\wedge (\forall x \in H. h\ x \leq p\ x) \implies g \in \text{norm-pres-extensions } E\ p\ F\ f$   
**by** (*unfold norm-pres-extensions-def*) *blast*

**end**

## 8 The norm of a function

**theory** *FunctionNorm* **imports** *NormedSpace FunctionOrder* **begin**

### 8.1 Continuous linear forms

A linear form  $f$  on a normed vector space  $(V, \|\cdot\|)$  is *continuous*, iff it is bounded, i.e.

$$\exists c \in R. \forall x \in V. |f\ x| \leq c \cdot \|x\|$$

In our application no other functions than linear forms are considered, so we can define continuous linear forms as bounded linear forms:

**locale** *continuous* = *var*  $V + \text{norm-syntax} + \text{linearform} +$   
**assumes** *bounded*:  $\exists c. \forall x \in V. |f\ x| \leq c * \|x\|$

**declare** *continuous.intro* [*intro?*] *continuous-axioms.intro* [*intro?*]

**lemma** *continuousI* [*intro*]:

**includes** *norm-syntax* + *linearform*  
**assumes**  $r$ :  $\bigwedge x. x \in V \implies |f\ x| \leq c * \|x\|$   
**shows** *continuous V norm f*

**proof**

**show** *linearform V f* .  
**from**  $r$  **have**  $\exists c. \forall x \in V. |f\ x| \leq c * \|x\|$  **by** *blast*  
**then show** *continuous-axioms V norm f ..*

**qed**

### 8.2 The norm of a linear form

The least real number  $c$  for which holds

$$\forall x \in V. |f\ x| \leq c \cdot \|x\|$$

is called the *norm* of  $f$ .

For non-trivial vector spaces  $V \neq \{0\}$  the norm can be defined as

$$\|f\| = \sup_{x \neq 0} |f x| / \|x\|$$

For the case  $V = \{0\}$  the supremum would be taken from an empty set. Since  $\mathbb{R}$  is unbounded, there would be no supremum. To avoid this situation it must be guaranteed that there is an element in this set. This element must be  $\{ \} \geq 0$  so that *fn-norm* has the norm properties. Furthermore it does not have to change the norm in all other cases, so it must be 0, as all other elements are  $\{ \} \geq 0$ .

Thus we define the set  $B$  where the supremum is taken from as follows:

$$\{0\} \cup \{|f x| / \|x\| \mid x \neq 0 \wedge x \in V\}$$

*fn-norm* is equal to the supremum of  $B$ , if the supremum exists (otherwise it is undefined).

```

locale fn-norm = norm-syntax +
  fixes  $B$  defines  $B \ V \ f \equiv \{0\} \cup \{|f x| / \|x\| \mid x. x \neq 0 \wedge x \in V\}$ 
  fixes fn-norm ( $\| \cdot \|$ ) --  $[0, 1000]$  999)
  defines  $\|f\| \equiv \bigsqcup (B \ V \ f)$ 

```

```

lemma (in fn-norm) B-not-empty [intro]:  $0 \in B \ V \ f$ 
  by (simp add: B-def)

```

The following lemma states that every continuous linear form on a normed space  $(V, \|\cdot\|)$  has a function norm.

```

lemma (in normed-vectorspace) fn-norm-works:
  includes fn-norm + continuous
  shows  $\text{lub } (B \ V \ f) (\|f\| \cdot V)$ 
proof –

```

The existence of the supremum is shown using the completeness of the reals. Completeness means, that every non-empty bounded set of reals has a supremum.

```

  have  $\exists a. \text{lub } (B \ V \ f) \ a$ 
  proof (rule real-complete)

```

First we have to show that  $B$  is non-empty:

```

    have  $0 \in B \ V \ f$  ..
    thus  $\exists x. x \in B \ V \ f$  ..

```

Then we have to show that  $B$  is bounded:

```

    show  $\exists c. \forall y \in B \ V \ f. y \leq c$ 
    proof –

```

We know that  $f$  is bounded by some value  $c$ .

```

      from bounded obtain  $c$  where  $c: \forall x \in V. |f x| \leq c * \|x\|$  ..

```

To prove the thesis, we have to show that there is some  $b$ , such that  $y \leq b$  for all  $y \in B$ . Due to the definition of  $B$  there are two cases.

```

      def  $b \equiv \max c \ 0$ 

```

```

have  $\forall y \in B \ V f. y \leq b$ 
proof
  fix y assume y:  $y \in B \ V f$ 
  show  $y \leq b$ 
  proof cases
    assume  $y = 0$ 
    thus ?thesis by (unfold b-def) arith
  next

```

The second case is  $y = |f x| / \|x\|$  for some  $x \in V$  with  $x \neq 0$ .

```

  assume  $y \neq 0$ 
  with y obtain x where y-rep:  $y = |f x| * \text{inverse } \|x\|$ 
    and  $x: x \in V$  and neg:  $x \neq 0$ 
  by (auto simp add: B-def real-divide-def)
  from x neg have gt:  $0 < \|x\|$  ..

```

The thesis follows by a short calculation using the fact that  $f$  is bounded.

```

  note y-rep
  also have  $|f x| * \text{inverse } \|x\| \leq (c * \|x\|) * \text{inverse } \|x\|$ 
  proof (rule mult-right-mono)
    from c show  $|f x| \leq c * \|x\|$  ..
    from gt have  $0 < \text{inverse } \|x\|$ 
      by (rule positive-imp-inverse-positive)
    thus  $0 \leq \text{inverse } \|x\|$  by (rule order-less-imp-le)
  qed
  also have  $\dots = c * (\|x\| * \text{inverse } \|x\|)$ 
    by (rule real-mult-assoc)
  also
  from gt have  $\|x\| \neq 0$  by simp
  hence  $\|x\| * \text{inverse } \|x\| = 1$  by simp
  also have  $c * 1 \leq b$  by (simp add: b-def le-maxI1)
  finally show  $y \leq b$  .
qed
qed
thus ?thesis ..
qed
qed
then show ?thesis by (unfold fn-norm-def) (rule the-lubI-ex)
qed

```

**lemma** (in *normed-vectorspace*) *fn-norm-ub* [iff?]:

includes *fn-norm* + *continuous*

assumes  $b: b \in B \ V f$

shows  $b \leq \|f\| - V$

**proof** –

have  $\text{lub } (B \ V f) (\|f\| - V)$

by (unfold B-def fn-norm-def) (rule fn-norm-works [OF continuous.intro])

from this and b show ?thesis ..

**qed**

**lemma** (in *normed-vectorspace*) *fn-norm-leastB*:

includes *fn-norm* + *continuous*

assumes  $b: \bigwedge b. b \in B \ V f \implies b \leq y$

shows  $\|f\| - V \leq y$

**proof** –  
 have  $\text{lub } (B \ V \ f) \ (\|f\| - V)$   
 by  $(\text{unfold } B\text{-def } fn\text{-norm-def}) \ (\text{rule } fn\text{-norm-works } [OF \ continuous.intro])$   
 from *this* and *b* show *?thesis* ..  
**qed**

The norm of a continuous function is always  $\geq 0$ .

**lemma** (in *normed-vectorspace*) *fn-norm-ge-zero* [*iff*]:

**includes** *fn-norm* + *continuous*

**shows**  $0 \leq \|f\| - V$

**proof** –

The function norm is defined as the supremum of  $B$ . So it is  $\geq 0$  if all elements in  $B$  are  $\geq 0$ , provided the supremum exists and  $B$  is not empty.

have  $\text{lub } (B \ V \ f) \ (\|f\| - V)$   
 by  $(\text{unfold } B\text{-def } fn\text{-norm-def}) \ (\text{rule } fn\text{-norm-works } [OF \ continuous.intro])$   
 moreover have  $0 \in B \ V \ f$  ..  
 ultimately show *?thesis* ..  
**qed**

The fundamental property of function norms is:

$$|f \ x| \leq \|f\| \cdot \|x\|$$

**lemma** (in *normed-vectorspace*) *fn-norm-le-cong*:

**includes** *fn-norm* + *continuous* + *linearform*

**assumes**  $x: x \in V$

**shows**  $|f \ x| \leq \|f\| - V * \|x\|$

**proof** *cases*

**assume**  $x = 0$

**then have**  $|f \ x| = |f \ 0|$  **by** *simp*

**also have**  $f \ 0 = 0$  ..

**also have**  $|\dots| = 0$  **by** *simp*

**also have**  $a: 0 \leq \|f\| - V$

by  $(\text{unfold } B\text{-def } fn\text{-norm-def})$   
 $(\text{rule } fn\text{-norm-ge-zero } [OF \ continuous.intro])$

**have**  $0 \leq \text{norm } x$  ..

**with**  $a$  **have**  $0 \leq \|f\| - V * \|x\|$  **by**  $(\text{simp add: zero-le-mult-iff})$

**finally show**  $|f \ x| \leq \|f\| - V * \|x\|$  .

**next**

**assume**  $x \neq 0$

**with**  $x$  **have**  $\text{neg: } \|x\| \neq 0$  **by** *simp*

**then have**  $|f \ x| = (|f \ x| * \text{inverse } \|x\|) * \|x\|$  **by** *simp*

**also have**  $\dots \leq \|f\| - V * \|x\|$

**proof**  $(\text{rule mult-right-mono})$

**from**  $x$  **show**  $0 \leq \|x\|$  ..

**from**  $x$  **and**  $\text{neg}$  **have**  $|f \ x| * \text{inverse } \|x\| \in B \ V \ f$

by  $(\text{auto simp add: } B\text{-def real-divide-def})$

**then show**  $|f \ x| * \text{inverse } \|x\| \leq \|f\| - V$

by  $(\text{unfold } B\text{-def } fn\text{-norm-def}) \ (\text{rule } fn\text{-norm-ub } [OF \ continuous.intro])$

**qed**

**finally show** *?thesis* .

**qed**

The function norm is the least positive real number for which the following inequation holds:

$$|f\ x| \leq c \cdot \|x\|$$

```

lemma (in normed-vectorspace) fn-norm-least [intro?]:
  includes fn-norm + continuous
  assumes ineq:  $\forall x \in V. |f\ x| \leq c * \|x\|$  and ge:  $0 \leq c$ 
  shows  $\|f\| - V \leq c$ 
proof (rule fn-norm-leastB [folded B-def fn-norm-def])
  fix b assume b:  $b \in B\ V\ f$ 
  show  $b \leq c$ 
proof cases
  assume  $b = 0$ 
  with ge show ?thesis by simp
next
  assume  $b \neq 0$ 
  with b obtain x where b-rep:  $b = |f\ x| * \text{inverse } \|x\|$ 
    and x-neq:  $x \neq 0$  and x:  $x \in V$ 
  by (auto simp add: B-def real-divide-def)
  note b-rep
  also have  $|f\ x| * \text{inverse } \|x\| \leq (c * \|x\|) * \text{inverse } \|x\|$ 
  proof (rule mult-right-mono)
    have  $0 < \|x\|$  ..
    then show  $0 \leq \text{inverse } \|x\|$  by simp
    from ineq and x show  $|f\ x| \leq c * \|x\|$  ..
  qed
  also have  $\dots = c$ 
  proof -
    from x-neq and x have  $\|x\| \neq 0$  by simp
    then show ?thesis by simp
  qed
  finally show ?thesis .
qed
qed (simp-all! add: continuous-def)

end

```

## 9 Zorn's Lemma

**theory** *ZornLemma* **imports** *Zorn* **begin**

Zorn's Lemmas states: if every linear ordered subset of an ordered set  $S$  has an upper bound in  $S$ , then there exists a maximal element in  $S$ . In our application,  $S$  is a set of sets ordered by set inclusion. Since the union of a chain of sets is an upper bound for all elements of the chain, the conditions of Zorn's lemma can be modified: if  $S$  is non-empty, it suffices to show that for every non-empty chain  $c$  in  $S$  the union of  $c$  also lies in  $S$ .

**theorem** *Zorn's-Lemma*:

```

assumes r:  $\bigwedge c. c \in \text{chain } S \implies \exists x. x \in c \implies \bigcup c \in S$ 
and aS:  $a \in S$ 
shows  $\exists y \in S. \forall z \in S. y \subseteq z \implies y = z$ 

```

```

proof (rule Zorn-Lemma2)1
  show  $\forall c \in \text{chain } S. \exists y \in S. \forall z \in c. z \subseteq y$ 
proof
  fix  $c$  assume  $c \in \text{chain } S$ 
  show  $\exists y \in S. \forall z \in c. z \subseteq y$ 
proof cases

```

If  $c$  is an empty chain, then every element in  $S$  is an upper bound of  $c$ .

```

  assume  $c = \{\}$ 
  with  $aS$  show ?thesis by fast

```

If  $c$  is non-empty, then  $\bigcup c$  is an upper bound of  $c$ , lying in  $S$ .

```

  next
  assume  $c: c \neq \{\}$ 
  show ?thesis
proof
  show  $\forall z \in c. z \subseteq \bigcup c$  by fast
  show  $\bigcup c \in S$ 
  proof (rule  $r$ )
    from  $c$  show  $\exists x. x \in c$  by fast
  qed
qed
qed
qed
qed
end

```

---

<sup>1</sup>See <http://isabelle.in.tum.de/library/HOL/HOL-Complex/Zorn.html>

## Part II

# Lemmas for the Proof

## 10 The supremum w.r.t. the function order

**theory** *HahnBanachSupLemmas* **imports** *FunctionNorm ZornLemma* **begin**

This section contains some lemmas that will be used in the proof of the Hahn-Banach Theorem. In this section the following context is presumed. Let  $E$  be a real vector space with a seminorm  $p$  on  $E$ .  $F$  is a subspace of  $E$  and  $f$  a linear form on  $F$ . We consider a chain  $c$  of norm-preserving extensions of  $f$ , such that  $\bigcup c = \text{graph } H \ h$ . We will show some properties about the limit function  $h$ , i.e. the supremum of the chain  $c$ .

Let  $c$  be a chain of norm-preserving extensions of the function  $f$  and let  $\text{graph } H \ h$  be the supremum of  $c$ . Every element in  $H$  is member of one of the elements of the chain.

**lemmas**  $[\text{dest?}] = \text{chainD}$

**lemmas**  $\text{chainE2} [\text{elim?}] = \text{chainD2} [\text{elim-format}, \text{standard}]$

**lemma** *some- $H' h'$* :

**assumes**  $M: M = \text{norm-pres-extensions } E \ p \ F \ f$

**and**  $cM: c \in \text{chain } M$

**and**  $u: \text{graph } H \ h = \bigcup c$

**and**  $x: x \in H$

**shows**  $\exists H' \ h'. \text{graph } H' \ h' \in c$

$\wedge (x, h \ x) \in \text{graph } H' \ h'$

$\wedge \text{linearform } H' \ h' \wedge H' \leq E$

$\wedge F \leq H' \wedge \text{graph } F \ f \subseteq \text{graph } H' \ h'$

$\wedge (\forall x \in H'. h' \ x \leq p \ x)$

**proof** –

**from**  $x$  **have**  $(x, h \ x) \in \text{graph } H \ h \ ..$

**also from**  $u$  **have**  $\dots = \bigcup c \ .$

**finally obtain**  $g$  **where**  $gc: g \in c$  **and**  $gh: (x, h \ x) \in g$  **by** *blast*

**from**  $cM$  **have**  $c \subseteq M \ ..$

**with**  $gc$  **have**  $g \in M \ ..$

**also from**  $M$  **have**  $\dots = \text{norm-pres-extensions } E \ p \ F \ f \ .$

**finally obtain**  $H'$  **and**  $h'$  **where**  $g: g = \text{graph } H' \ h'$

**and**  $*$  :  $\text{linearform } H' \ h' \ H' \leq E \ F \leq H'$

$\text{graph } F \ f \subseteq \text{graph } H' \ h' \ \forall x \in H'. h' \ x \leq p \ x \ ..$

**from**  $gc$  **and**  $g$  **have**  $\text{graph } H' \ h' \in c$  **by** (*simp only*.)

**moreover from**  $gh$  **and**  $g$  **have**  $(x, h \ x) \in \text{graph } H' \ h'$  **by** (*simp only*.)

**ultimately show** *?thesis* **using**  $*$  **by** *blast*

**qed**

Let  $c$  be a chain of norm-preserving extensions of the function  $f$  and let  $\text{graph } H \ h$  be the supremum of  $c$ . Every element in the domain  $H$  of the supremum function is member of the domain  $H'$  of some function  $h'$ , such that  $h$  extends  $h'$ .

**lemma** *some- $H'h'$* :  
**assumes**  $M$ :  $M = \text{norm-pres-extensions } E \ p \ F \ f$   
**and**  $cM$ :  $c \in \text{chain } M$   
**and**  $u$ :  $\text{graph } H \ h = \bigcup c$   
**and**  $x$ :  $x \in H$   
**shows**  $\exists H' \ h'. x \in H' \wedge \text{graph } H' \ h' \subseteq \text{graph } H \ h$   
 $\wedge \text{linearform } H' \ h' \wedge H' \trianglelefteq E \wedge F \trianglelefteq H'$   
 $\wedge \text{graph } F \ f \subseteq \text{graph } H' \ h' \wedge (\forall x \in H'. h' \ x \leq p \ x)$   
**proof** –  
**from**  $M \ cM \ u \ x$  **obtain**  $H' \ h'$  **where**  
 $x\text{-hx}$ :  $(x, h \ x) \in \text{graph } H' \ h'$   
**and**  $c$ :  $\text{graph } H' \ h' \in c$   
**and**  $*$ :  $\text{linearform } H' \ h' \ H' \trianglelefteq E \ F \trianglelefteq H'$   
 $\text{graph } F \ f \subseteq \text{graph } H' \ h' \ \forall x \in H'. h' \ x \leq p \ x$   
**by** (rule *some- $H'h'$*  *[elim-format]*) **blast**  
**from**  $x\text{-hx}$  **have**  $x \in H' ..$   
**moreover from**  $cM \ u \ c$  **have**  $\text{graph } H' \ h' \subseteq \text{graph } H \ h$   
**by** (simp only: *chain-ball-Union-upper*)  
**ultimately show** *?thesis* **using**  $*$  **by** **blast**  
**qed**

Any two elements  $x$  and  $y$  in the domain  $H$  of the supremum function  $h$  are both in the domain  $H'$  of some function  $h'$ , such that  $h$  extends  $h'$ .

**lemma** *some- $H'h'2$* :  
**assumes**  $M$ :  $M = \text{norm-pres-extensions } E \ p \ F \ f$   
**and**  $cM$ :  $c \in \text{chain } M$   
**and**  $u$ :  $\text{graph } H \ h = \bigcup c$   
**and**  $x$ :  $x \in H$   
**and**  $y$ :  $y \in H$   
**shows**  $\exists H' \ h'. x \in H' \wedge y \in H'$   
 $\wedge \text{graph } H' \ h' \subseteq \text{graph } H \ h$   
 $\wedge \text{linearform } H' \ h' \wedge H' \trianglelefteq E \wedge F \trianglelefteq H'$   
 $\wedge \text{graph } F \ f \subseteq \text{graph } H' \ h' \wedge (\forall x \in H'. h' \ x \leq p \ x)$   
**proof** –

$y$  is in the domain  $H''$  of some function  $h''$ , such that  $h$  extends  $h''$ .

**from**  $M \ cM \ u$  **and**  $y$  **obtain**  $H' \ h'$  **where**  
 $y\text{-hy}$ :  $(y, h \ y) \in \text{graph } H' \ h'$   
**and**  $c'$ :  $\text{graph } H' \ h' \in c$   
**and**  $*$ :  
 $\text{linearform } H' \ h' \ H' \trianglelefteq E \ F \trianglelefteq H'$   
 $\text{graph } F \ f \subseteq \text{graph } H' \ h' \ \forall x \in H'. h' \ x \leq p \ x$   
**by** (rule *some- $H'h'$*  *[elim-format]*) **blast**

$x$  is in the domain  $H'$  of some function  $h'$ , such that  $h$  extends  $h'$ .

**from**  $M \ cM \ u$  **and**  $x$  **obtain**  $H'' \ h''$  **where**  
 $x\text{-hx}$ :  $(x, h \ x) \in \text{graph } H'' \ h''$   
**and**  $c''$ :  $\text{graph } H'' \ h'' \in c$   
**and**  $**$ :  
 $\text{linearform } H'' \ h'' \ H'' \trianglelefteq E \ F \trianglelefteq H''$   
 $\text{graph } F \ f \subseteq \text{graph } H'' \ h'' \ \forall x \in H''. h'' \ x \leq p \ x$   
**by** (rule *some- $H'h'$*  *[elim-format]*) **blast**



Since both  $h'$  and  $h''$  are elements of the chain,  $h''$  is an extension of  $h'$  or vice versa. Thus both  $x$  and  $y$  are contained in the greater one.

```

from  $cM$  have  $\text{graph } H'' h'' \subseteq \text{graph } H' h' \vee \text{graph } H' h' \subseteq \text{graph } H'' h''$ 
  (is  $?case1 \vee ?case2$ ) ..
then show  $?thesis$ 
proof
  assume  $?case1$ 
  have  $(x, h x) \in \text{graph } H'' h''$  .
  also have  $\dots \subseteq \text{graph } H' h'$  .
  finally have  $xh:(x, h x) \in \text{graph } H' h'$  .
  then have  $x \in H'$  ..
  moreover from  $y-hy$  have  $y \in H'$  ..
  moreover from  $cM u$  and  $c'$  have  $\text{graph } H' h' \subseteq \text{graph } H h$ 
    by (simp only: chain-ball-Union-upper)
  ultimately show  $?thesis$  using * by blast
next
  assume  $?case2$ 
  from  $x-hx$  have  $x \in H''$  ..
  moreover {
    from  $y-hy$  have  $(y, h y) \in \text{graph } H' h'$  .
    also have  $\dots \subseteq \text{graph } H'' h''$  .
    finally have  $(y, h y) \in \text{graph } H'' h''$  .
  } then have  $y \in H''$  ..
  moreover from  $cM u$  and  $c''$  have  $\text{graph } H'' h'' \subseteq \text{graph } H h$ 
    by (simp only: chain-ball-Union-upper)
  ultimately show  $?thesis$  using ** by blast
qed
qed

```

The relation induced by the graph of the supremum of a chain  $c$  is definite, i. e.  $t$  is the graph of a function.

**lemma** *sup-definite*:

```

assumes  $M\text{-def}: M \equiv \text{norm-pres-extensions } E p F f$ 
  and  $cM: c \in \text{chain } M$ 
  and  $xy: (x, y) \in \bigcup c$ 
  and  $xz: (x, z) \in \bigcup c$ 
shows  $z = y$ 
proof –
  from  $cM$  have  $c: c \subseteq M$  ..
  from  $xy$  obtain  $G1$  where  $xy': (x, y) \in G1$  and  $G1: G1 \in c$  ..
  from  $xz$  obtain  $G2$  where  $xz': (x, z) \in G2$  and  $G2: G2 \in c$  ..

  from  $G1 c$  have  $G1 \in M$  ..
  then obtain  $H1 h1$  where  $G1\text{-rep}: G1 = \text{graph } H1 h1$ 
    by (unfold M-def) blast

  from  $G2 c$  have  $G2 \in M$  ..
  then obtain  $H2 h2$  where  $G2\text{-rep}: G2 = \text{graph } H2 h2$ 
    by (unfold M-def) blast

```

$G_1$  is contained in  $G_2$  or vice versa, since both  $G_1$  and  $G_2$  are members of  $c$ .

```

from  $cM G1 G2$  have  $G1 \subseteq G2 \vee G2 \subseteq G1$  (is  $?case1 \vee ?case2$ ) ..
then show  $?thesis$ 

```

```

proof
  assume ?case1
  with  $xy'$   $G2\text{-rep}$  have  $(x, y) \in \text{graph } H2 \ h2$  by blast
  hence  $y = h2 \ x \ ..$ 
  also
  from  $xz'$   $G2\text{-rep}$  have  $(x, z) \in \text{graph } H2 \ h2$  by (simp only:)
  hence  $z = h2 \ x \ ..$ 
  finally show ?thesis .
next
  assume ?case2
  with  $xz'$   $G1\text{-rep}$  have  $(x, z) \in \text{graph } H1 \ h1$  by blast
  hence  $z = h1 \ x \ ..$ 
  also
  from  $xy'$   $G1\text{-rep}$  have  $(x, y) \in \text{graph } H1 \ h1$  by (simp only:)
  hence  $y = h1 \ x \ ..$ 
  finally show ?thesis ..
qed
qed

```

The limit function  $h$  is linear. Every element  $x$  in the domain of  $h$  is in the domain of a function  $h'$  in the chain of norm preserving extensions. Furthermore,  $h$  is an extension of  $h'$  so the function values of  $x$  are identical for  $h'$  and  $h$ . Finally, the function  $h'$  is linear by construction of  $M$ .

**lemma** *sup-lf*:

```

assumes  $M$ :  $M = \text{norm-pres-extensions } E \ p \ F \ f$ 
  and  $cM$ :  $c \in \text{chain } M$ 
  and  $u$ :  $\text{graph } H \ h = \bigcup c$ 
shows linearform  $H \ h$ 

```

**proof**

```

fix  $x \ y$  assume  $x$ :  $x \in H$  and  $y$ :  $y \in H$ 
with  $M \ cM \ u$  obtain  $H' \ h'$  where
   $x'$ :  $x \in H'$  and  $y'$ :  $y \in H'$ 
  and  $b$ :  $\text{graph } H' \ h' \subseteq \text{graph } H \ h$ 
  and linearform: linearform  $H' \ h'$ 
  and subspace:  $H' \trianglelefteq E$ 
by (rule some- $H'h'2$  [elim-format]) blast

```

**show**  $h \ (x + y) = h \ x + h \ y$

**proof** –

```

from linearform  $x' \ y'$  have  $h' \ (x + y) = h' \ x + h' \ y$ 
  by (rule linearform.add)
also from  $b \ x'$  have  $h' \ x = h \ x \ ..$ 
also from  $b \ y'$  have  $h' \ y = h \ y \ ..$ 
also from subspace  $x' \ y'$  have  $x + y \in H'$ 
  by (rule subspace.add-closed)
with  $b$  have  $h' \ (x + y) = h \ (x + y) \ ..$ 
finally show ?thesis .

```

**qed**

**next**

```

fix  $x \ a$  assume  $x$ :  $x \in H$ 
with  $M \ cM \ u$  obtain  $H' \ h'$  where
   $x'$ :  $x \in H'$ 
  and  $b$ :  $\text{graph } H' \ h' \subseteq \text{graph } H \ h$ 

```

```

    and linearform: linearform H' h'
    and subspace: H'  $\trianglelefteq$  E
    by (rule some-H'h' [elim-format]) blast

show h (a · x) = a * h x
proof -
  from linearform x' have h' (a · x) = a * h' x
  by (rule linearform.mult)
  also from b x' have h' x = h x ..
  also from subspace x' have a · x  $\in$  H'
  by (rule subspace.mult-closed)
  with b have h' (a · x) = h (a · x) ..
  finally show ?thesis .
qed
qed

```

The limit of a non-empty chain of norm preserving extensions of  $f$  is an extension of  $f$ , since every element of the chain is an extension of  $f$  and the supremum is an extension for every element of the chain.

```

lemma sup-ext:
  assumes graph: graph H h =  $\bigcup c$ 
  and M: M = norm-pres-extensions E p F f
  and cM: c  $\in$  chain M
  and ex:  $\exists x. x \in c$ 
  shows graph F f  $\subseteq$  graph H h
proof -
  from ex obtain x where xc: x  $\in$  c ..
  from cM have c  $\subseteq$  M ..
  with xc have x  $\in$  M ..
  with M have x  $\in$  norm-pres-extensions E p F f
  by (simp only:)
  then obtain G g where x = graph G g and graph F f  $\subseteq$  graph G g ..
  then have graph F f  $\subseteq$  x by (simp only:)
  also from xc have ...  $\subseteq \bigcup c$  by blast
  also from graph have ... = graph H h ..
  finally show ?thesis .
qed

```

The domain  $H$  of the limit function is a superspace of  $F$ , since  $F$  is a subset of  $H$ . The existence of the  $0$  element in  $F$  and the closure properties follow from the fact that  $F$  is a vector space.

```

lemma sup-supF:
  assumes graph: graph H h =  $\bigcup c$ 
  and M: M = norm-pres-extensions E p F f
  and cM: c  $\in$  chain M
  and ex:  $\exists x. x \in c$ 
  and FE: F  $\trianglelefteq$  E
  shows F  $\trianglelefteq$  H
proof
  from FE show F  $\neq \{\}$  by (rule subspace.non-empty)
  from graph M cM ex have graph F f  $\subseteq$  graph H h by (rule sup-ext)
  then show F  $\subseteq$  H ..

```

```

fix x y assume x ∈ F and y ∈ F
with FE show x + y ∈ F by (rule subspace.add-closed)
next
fix x a assume x ∈ F
with FE show a · x ∈ F by (rule subspace.mult-closed)
qed

```

The domain  $H$  of the limit function is a subspace of  $E$ .

**lemma** *sup-subE*:

```

assumes graph: graph H h = ⋃ c
and M: M = norm-pres-extensions E p F f
and cM: c ∈ chain M
and ex: ∃ x. x ∈ c
and FE: F ≤ E
and E: vectorspace E
shows H ≤ E
proof
show H ≠ {}
proof -
from FE E have 0 ∈ F by (rule subspace.zero)
also from graph M cM ex FE have F ≤ H by (rule sup-supF)
then have F ⊆ H ..
finally show ?thesis by blast
qed
show H ⊆ E
proof
fix x assume x ∈ H
with M cM graph
obtain H' h' where x: x ∈ H' and H'E: H' ≤ E
by (rule some-H'h' [elim-format]) blast
from H'E have H' ⊆ E ..
with x show x ∈ E ..
qed
fix x y assume x: x ∈ H and y: y ∈ H
show x + y ∈ H
proof -
from M cM graph x y obtain H' h' where
x': x ∈ H' and y': y ∈ H' and H'E: H' ≤ E
and graphs: graph H' h' ⊆ graph H h
by (rule some-H'h'2 [elim-format]) blast
from H'E x' y' have x + y ∈ H'
by (rule subspace.add-closed)
also from graphs have H' ⊆ H ..
finally show ?thesis .
qed
next
fix x a assume x: x ∈ H
show a · x ∈ H
proof -
from M cM graph x
obtain H' h' where x': x ∈ H' and H'E: H' ≤ E
and graphs: graph H' h' ⊆ graph H h
by (rule some-H'h' [elim-format]) blast
from H'E x' have a · x ∈ H' by (rule subspace.mult-closed)

```

```

    also from graphs have  $H' \subseteq H$  ..
    finally show ?thesis .
  qed
qed

```

The limit function is bounded by the norm  $p$  as well, since all elements in the chain are bounded by  $p$ .

```

lemma sup-norm-pres:
  assumes graph:  $\text{graph } H \ h = \bigcup c$ 
    and  $M: M = \text{norm-pres-extensions } E \ p \ F \ f$ 
    and  $cM: c \in \text{chain } M$ 
  shows  $\forall x \in H. h \ x \leq p \ x$ 
proof
  fix  $x$  assume  $x \in H$ 
  with  $M \ cM$  graph obtain  $H' \ h'$  where  $x': x \in H'$ 
    and graphs:  $\text{graph } H' \ h' \subseteq \text{graph } H \ h$ 
    and  $a: \forall x \in H'. h' \ x \leq p \ x$ 
  by (rule some-H'h' [elim-format]) blast
  from graphs  $x'$  have [symmetric]:  $h' \ x = h \ x$  ..
  also from  $a \ x'$  have  $h' \ x \leq p \ x$  ..
  finally show  $h \ x \leq p \ x$  .
qed

```

The following lemma is a property of linear forms on real vector spaces. It will be used for the lemma *abs-HahnBanach* (see page 47). For real vector spaces the following inequations are equivalent:

$$\forall x \in H. |h \ x| \leq p \ x \quad \text{and} \quad \forall x \in H. h \ x \leq p \ x$$

```

lemma abs-ineq-iff:
  includes subspace  $H \ E + \text{vectorspace } E + \text{seminorm } E \ p + \text{linearform } H \ h$ 
  shows  $(\forall x \in H. |h \ x| \leq p \ x) = (\forall x \in H. h \ x \leq p \ x)$  (is ?L = ?R)
proof
  have  $H: \text{vectorspace } H$  ..
  {
    assume  $l: ?L$ 
    show ?R
    proof
      fix  $x$  assume  $x: x \in H$ 
      have  $h \ x \leq |h \ x|$  by arith
      also from  $l \ x$  have  $\dots \leq p \ x$  ..
      finally show  $h \ x \leq p \ x$  .
    qed
  }
next
  assume  $r: ?R$ 
  show ?L
  proof
    fix  $x$  assume  $x: x \in H$ 
    show  $\bigwedge a \ b :: \text{real}. - \ a \leq b \implies b \leq a \implies |b| \leq a$ 
    by arith
    have linearform  $H \ h$  .
    from this  $H \ x$  have  $- \ h \ x = h \ (- \ x)$  by (rule linearform.neg [symmetric])
    also

```

```

    from  $H$   $x$  have  $-x \in H$  by (rule vectorspace.neg-closed)
    with  $r$  have  $h(-x) \leq p(-x)$  ..
    also have  $\dots = p\ x$ 
    proof (rule seminorm.minus)
      from  $x$  show  $x \in E$  ..
    qed
    finally have  $-h\ x \leq p\ x$  .
    then show  $-p\ x \leq h\ x$  by simp
    from  $r\ x$  show  $h\ x \leq p\ x$  ..
  qed
}
qed
end

```

## 11 Extending non-maximal functions

**theory** *HahnBanachExtLemmas* **imports** *FunctionNorm* **begin**

In this section the following context is presumed. Let  $E$  be a real vector space with a seminorm  $q$  on  $E$ .  $F$  is a subspace of  $E$  and  $f$  a linear function on  $F$ . We consider a subspace  $H$  of  $E$  that is a superspace of  $F$  and a linear form  $h$  on  $H$ .  $H$  is not equal to  $E$  and  $x_0$  is an element in  $E - H$ .  $H$  is extended to the direct sum  $H' = H + \text{lin } x_0$ , so for any  $x \in H'$  the decomposition of  $x = y + a \cdot x_0$  with  $y \in H$  is unique.  $h'$  is defined on  $H'$  by  $h' x = h y + a \cdot \xi$  for a certain  $\xi$ .

Subsequently we show some properties of this extension  $h'$  of  $h$ .

This lemma will be used to show the existence of a linear extension of  $f$  (see page 45). It is a consequence of the completeness of  $\mathbb{R}$ . To show

$$\exists \xi. \forall y \in F. a\ y \leq \xi \wedge \xi \leq b\ y$$

it suffices to show that

$$\forall u \in F. \forall v \in F. a\ u \leq b\ v$$

**lemma** *ex-xi*:

```

includes vectorspace  $F$ 
assumes  $r: \bigwedge u\ v. u \in F \implies v \in F \implies a\ u \leq b\ v$ 
shows  $\exists xi::real. \forall y \in F. a\ y \leq xi \wedge xi \leq b\ y$ 
proof -

```

From the completeness of the reals follows: The set  $S = \{a\ u. u \in F\}$  has a supremum, if it is non-empty and has an upper bound.

```

let  $?S = \{a\ u \mid u. u \in F\}$ 
have  $\exists xi. \text{lub } ?S\ xi$ 
proof (rule real-complete)
  have  $a\ 0 \in ?S$  by blast
  then show  $\exists X. X \in ?S$  ..
  have  $\forall y \in ?S. y \leq b\ 0$ 
proof

```

```

    fix y assume y: y ∈ ?S
    then obtain u where u: u ∈ F and y: y = a u by blast
    from u and zero have a u ≤ b 0 by (rule r)
    with y show y ≤ b 0 by (simp only:)
  qed
  then show ∃ u. ∀ y ∈ ?S. y ≤ u ..
qed
then obtain xi where xi: lub ?S xi ..
{
  fix y assume y ∈ F
  then have a y ∈ ?S by blast
  with xi have a y ≤ xi by (rule lub.upper)
} moreover {
  fix y assume y: y ∈ F
  from xi have xi ≤ b y
  proof (rule lub.least)
    fix au assume au ∈ ?S
    then obtain u where u: u ∈ F and au: au = a u by blast
    from u y have a u ≤ b y by (rule r)
    with au show au ≤ b y by (simp only:)
  qed
} ultimately show ∃ xi. ∀ y ∈ F. a y ≤ xi ∧ xi ≤ b y by blast
qed

```

The function  $h'$  is defined as a  $h' x = h y + a \cdot \xi$  where  $x = y + a \cdot \xi$  is a linear extension of  $h$  to  $H'$ .

```

lemma h'-lf:
  includes var H + var h + var E
  assumes h'-def: h' ≡ λx. let (y, a) =
    SOME (y, a). x = y + a · x0 ∧ y ∈ H in h y + a · xi
    and H'-def: H' ≡ H + lin x0
    and HE: H ≤ E
  includes linearform H h
  assumes x0: x0 ∉ H x0 ∈ E x0 ≠ 0
  includes vectorspace E
  shows linearform H' h'
proof
  have H': vectorspace H'
  proof (unfold H'-def)
    have x0 ∈ E .
    then have lin x0 ≤ E ..
    with HE show vectorspace (H + lin x0) ..
  qed
  {
    fix x1 x2 assume x1: x1 ∈ H' and x2: x2 ∈ H'
    show h' (x1 + x2) = h' x1 + h' x2
    proof -
      from H' x1 x2 have x1 + x2 ∈ H'
      by (rule vectorspace.add-closed)
      with x1 x2 obtain y y1 y2 a a1 a2 where
        x1x2: x1 + x2 = y + a · x0 and y: y ∈ H
        and x1-rep: x1 = y1 + a1 · x0 and y1: y1 ∈ H
        and x2-rep: x2 = y2 + a2 · x0 and y2: y2 ∈ H
    proof -

```

```

    by (unfold H'-def sum-def lin-def) blast

have ya:  $y1 + y2 = y \wedge a1 + a2 = a$  using - HE - y x0
proof (rule decomp-H')
  from HE y1 y2 show  $y1 + y2 \in H$ 
  by (rule subspace.add-closed)
  from x0 and HE y y1 y2
  have  $x0 \in E \ y \in E \ y1 \in E \ y2 \in E$  by auto
  with x1-rep x2-rep have  $(y1 + y2) + (a1 + a2) \cdot x0 = x1 + x2$ 
  by (simp add: add-ac add-mult-distrib2)
  also note x1x2
  finally show  $(y1 + y2) + (a1 + a2) \cdot x0 = y + a \cdot x0$  .
qed

from h'-def x1x2 - HE y x0
have  $h'(x1 + x2) = h y + a * xi$ 
  by (rule h'-definite)
also have  $\dots = h (y1 + y2) + (a1 + a2) * xi$ 
  by (simp only: ya)
also from y1 y2 have  $h (y1 + y2) = h y1 + h y2$ 
  by simp
also have  $\dots + (a1 + a2) * xi = (h y1 + a1 * xi) + (h y2 + a2 * xi)$ 
  by (simp add: left-distrib)
also from h'-def x1-rep - HE y1 x0
have  $h y1 + a1 * xi = h' x1$ 
  by (rule h'-definite [symmetric])
also from h'-def x2-rep - HE y2 x0
have  $h y2 + a2 * xi = h' x2$ 
  by (rule h'-definite [symmetric])
finally show ?thesis .
qed
next
fix x1 c assume x1:  $x1 \in H'$ 
show  $h'(c \cdot x1) = c * (h' x1)$ 
proof -
  from H' x1 have ax1:  $c \cdot x1 \in H'$ 
  by (rule vectorspace.mult-closed)
  with x1 obtain y a y1 a1 where
    cx1-rep:  $c \cdot x1 = y + a \cdot x0$  and y:  $y \in H$ 
    and x1-rep:  $x1 = y1 + a1 \cdot x0$  and y1:  $y1 \in H$ 
  by (unfold H'-def sum-def lin-def) blast

have ya:  $c \cdot y1 = y \wedge c * a1 = a$  using - HE - y x0
proof (rule decomp-H')
  from HE y1 show  $c \cdot y1 \in H$ 
  by (rule subspace.mult-closed)
  from x0 and HE y y1
  have  $x0 \in E \ y \in E \ y1 \in E$  by auto
  with x1-rep have  $c \cdot y1 + (c * a1) \cdot x0 = c \cdot x1$ 
  by (simp add: mult-assoc add-mult-distrib1)
  also note cx1-rep
  finally show  $c \cdot y1 + (c * a1) \cdot x0 = y + a \cdot x0$  .
qed

from h'-def cx1-rep - HE y x0 have  $h'(c \cdot x1) = h y + a * xi$ 

```



```

    by (rule h'-definite)
  also have ... = h (c · y1) + (c * a1) * xi
    by (simp only: ya)
  also from y1 have h (c · y1) = c * h y1
    by simp
  also have ... + (c * a1) * xi = c * (h y1 + a1 * xi)
    by (simp only: right-distrib)
  also from h'-def x1-rep - HE y1 x0 have h y1 + a1 * xi = h' x1
    by (rule h'-definite [symmetric])
  finally show ?thesis .
qed
}
qed

```

The linear extension  $h'$  of  $h$  is bounded by the seminorm  $p$ .

**lemma**  $h'$ -norm-pres:

```

includes var H + var h + var E
assumes h'-def:  $h' \equiv \lambda x. \text{let } (y, a) =$ 
  SOME (y, a).  $x = y + a \cdot x0 \wedge y \in H \text{ in } h y + a * xi$ 
  and H'-def:  $H' \equiv H + \text{lin } x0$ 
  and x0:  $x0 \notin H \ x0 \in E \ x0 \neq 0$ 
includes vectorspace E + subspace H E + seminorm E p + linearform H h
assumes a:  $\forall y \in H. h y \leq p y$ 
  and a':  $\forall y \in H. -p (y + x0) - h y \leq xi \wedge xi \leq p (y + x0) - h y$ 
shows  $\forall x \in H'. h' x \leq p x$ 

```

**proof**

```

fix x assume x':  $x \in H'$ 
show  $h' x \leq p x$ 
proof -
  from a' have a1:  $\forall ya \in H. -p (ya + x0) - h ya \leq xi$ 
    and a2:  $\forall ya \in H. xi \leq p (ya + x0) - h ya$  by auto
  from x' obtain y a where
    x-rep:  $x = y + a \cdot x0$  and y:  $y \in H$ 
    by (unfold H'-def sum-def lin-def) blast
  from y have y':  $y \in E$  ..
  from y have ay:  $\text{inverse } a \cdot y \in H$  by simp

  from h'-def x-rep - - y x0 have  $h' x = h y + a * xi$ 
    by (rule h'-definite)
  also have ...  $\leq p (y + a \cdot x0)$ 
  proof (rule linorder-cases)
    assume z:  $a = 0$ 
    then have  $h y + a * xi = h y$  by simp
    also from a y have ...  $\leq p y$  ..
    also from x0 y' z have  $p y = p (y + a \cdot x0)$  by simp
    finally show ?thesis .
  next

```

In the case  $a < 0$ , we use  $a_1$  with  $ya$  taken as  $y / a$ :

```

assume lz:  $a < 0$  hence nz:  $a \neq 0$  by simp
from a1 ay
have  $-p (\text{inverse } a \cdot y + x0) - h (\text{inverse } a \cdot y) \leq xi$  ..
with lz have  $a * xi \leq$ 

```

```

    a * (- p (inverse a · y + x0) - h (inverse a · y))
  by (simp add: mult-left-mono-neg order-less-imp-le)

  also have ... =
    - a * (p (inverse a · y + x0)) - a * (h (inverse a · y))
  by (simp add: right-diff-distrib)
  also from lz x0 y' have - a * (p (inverse a · y + x0)) =
    p (a · (inverse a · y + x0))
  by (simp add: abs-homogenous)
  also from nz x0 y' have ... = p (y + a · x0)
  by (simp add: add-mult-distrib1 mult-assoc [symmetric])
  also from nz y have a * (h (inverse a · y)) = h y
  by simp
  finally have a * xi ≤ p (y + a · x0) - h y .
  then show ?thesis by simp
next

```

In the case  $a > 0$ , we use  $a_2$  with  $ya$  taken as  $y / a$ :

```

  assume gz: 0 < a hence nz: a ≠ 0 by simp
  from a2 ay
  have xi ≤ p (inverse a · y + x0) - h (inverse a · y) ..
  with gz have a * xi ≤
    a * (p (inverse a · y + x0) - h (inverse a · y))
  by simp
  also have ... = a * p (inverse a · y + x0) - a * h (inverse a · y)
  by (simp add: right-diff-distrib)
  also from gz x0 y'
  have a * p (inverse a · y + x0) = p (a · (inverse a · y + x0))
  by (simp add: abs-homogenous)
  also from nz x0 y' have ... = p (y + a · x0)
  by (simp add: add-mult-distrib1 mult-assoc [symmetric])
  also from nz y have a * h (inverse a · y) = h y
  by simp
  finally have a * xi ≤ p (y + a · x0) - h y .
  then show ?thesis by simp
qed
also from x-rep have ... = p x by (simp only:)
finally show ?thesis .
qed
end

```

## Part III

# The Main Proof

## 12 The Hahn-Banach Theorem

**theory** *HahnBanach* **imports** *HahnBanachLemmas* **begin**

We present the proof of two different versions of the Hahn-Banach Theorem, closely following [1, §36].

### 12.1 The Hahn-Banach Theorem for vector spaces

**Hahn-Banach Theorem.** Let  $F$  be a subspace of a real vector space  $E$ , let  $p$  be a semi-norm on  $E$ , and  $f$  be a linear form defined on  $F$  such that  $f$  is bounded by  $p$ , i.e.  $\forall x \in F. f x \leq p x$ . Then  $f$  can be extended to a linear form  $h$  on  $E$  such that  $h$  is norm-preserving, i.e.  $h$  is also bounded by  $p$ .

**Proof Sketch.**

1. Define  $M$  as the set of norm-preserving extensions of  $f$  to subspaces of  $E$ . The linear forms in  $M$  are ordered by domain extension.
2. We show that every non-empty chain in  $M$  has an upper bound in  $M$ .
3. With Zorn's Lemma we conclude that there is a maximal function  $g$  in  $M$ .
4. The domain  $H$  of  $g$  is the whole space  $E$ , as shown by classical contradiction:
  - Assuming  $g$  is not defined on whole  $E$ , it can still be extended in a norm-preserving way to a super-space  $H'$  of  $H$ .
  - Thus  $g$  can not be maximal. Contradiction!

**theorem** *HahnBanach*:

**includes** *vectorspace*  $E$  + *subspace*  $F$   $E$  + *seminorm*  $E$   $p$  + *linearform*  $F$   $f$

**assumes**  $fp$ :  $\forall x \in F. f x \leq p x$

**shows**  $\exists h. \text{linearform } E h \wedge (\forall x \in F. h x = f x) \wedge (\forall x \in E. h x \leq p x)$

— Let  $E$  be a vector space,  $F$  a subspace of  $E$ ,  $p$  a seminorm on  $E$ ,

— and  $f$  a linear form on  $F$  such that  $f$  is bounded by  $p$ ,

— then  $f$  can be extended to a linear form  $h$  on  $E$  in a norm-preserving way.

**proof** —

**def**  $M \equiv \text{norm-pres-extensions } E p F f$

**hence**  $M$ :  $M = \dots$  **by** (*simp only*:)

**have**  $E$ : *vectorspace*  $E$  .

**have**  $F$ : *vectorspace*  $F$  ..

{

**fix**  $c$  **assume**  $cM$ :  $c \in \text{chain } M$  **and**  $ex$ :  $\exists x. x \in c$

**have**  $\bigcup c \in M$

— Show that every non-empty chain  $c$  of  $M$  has an upper bound in  $M$ :

—  $\bigcup c$  is greater than any element of the chain  $c$ , so it suffices to show  $\bigcup c \in M$ .

```

proof (unfold M-def, rule norm-pres-extensionI)
  let ?H = domain ( $\bigcup c$ )
  let ?h = funct ( $\bigcup c$ )

  have a: graph ?H ?h =  $\bigcup c$ 
  proof (rule graph-domain-funct)
    fix x y z assume (x, y)  $\in \bigcup c$  and (x, z)  $\in \bigcup c$ 
    with M-def cM show z = y by (rule sup-definite)
  qed
  moreover from M cM a have linearform ?H ?h
    by (rule sup-lf)
  moreover from a M cM ex have ?H  $\trianglelefteq E$ 
    by (rule sup-subE)
  moreover from a M cM ex have F  $\trianglelefteq$  ?H
    by (rule sup-supF)
  moreover from a M cM ex have graph F f  $\subseteq$  graph ?H ?h
    by (rule sup-ext)
  moreover from a M cM have  $\forall x \in ?H. ?h\ x \leq p\ x$ 
    by (rule sup-norm-pres)
  ultimately show  $\exists H\ h. \bigcup c = \text{graph } H\ h$ 
     $\wedge$  linearform H h
     $\wedge H \trianglelefteq E$ 
     $\wedge F \trianglelefteq H$ 
     $\wedge \text{graph } F\ f \subseteq \text{graph } H\ h$ 
     $\wedge (\forall x \in H. h\ x \leq p\ x)$  by blast
  qed
}
hence  $\exists g \in M. \forall x \in M. g \subseteq x \longrightarrow g = x$ 
— With Zorn's Lemma we can conclude that there is a maximal element in M.
proof (rule Zorn's-Lemma)
  — We show that M is non-empty:
  show graph F f  $\in M$ 
  proof (unfold M-def, rule norm-pres-extensionI2)
    show linearform F f .
    show F  $\trianglelefteq E$  .
    from F show F  $\trianglelefteq$  F by (rule vectorspace.subspace-refl)
    show graph F f  $\subseteq$  graph F f ..
    show  $\forall x \in F. f\ x \leq p\ x$  .
  qed
qed
then obtain g where gM: g  $\in M$  and  $\forall x \in M. g \subseteq x \longrightarrow g = x$ 
  by blast
from gM [unfolded M-def] obtain H h where
  g-rep: g = graph H h
  and linearform: linearform H h
  and HE: H  $\trianglelefteq E$  and FH: F  $\trianglelefteq$  H
  and graphs: graph F f  $\subseteq$  graph H h
  and hp:  $\forall x \in H. h\ x \leq p\ x$  ..
  — g is a norm-preserving extension of f, in other words:
  — g is the graph of some linear form h defined on a subspace H of E,
  — and h is an extension of f that is again bounded by p.
from HE have H: vectorspace H
  by (rule subspace.vectorspace)

```

```

have HE-eq:  $H = E$ 
  — We show that  $h$  is defined on whole  $E$  by classical contradiction.
proof (rule classical)
  assume neg:  $H \neq E$ 
    — Assume  $h$  is not defined on whole  $E$ . Then show that  $h$  can be extended
    — in a norm-preserving way to a function  $h'$  with the graph  $g'$ .
  have  $\exists g' \in M. g \subseteq g' \wedge g \neq g'$ 
  proof —
    from HE have  $H \subseteq E$  ..
    with neg obtain  $x'$  where  $x'E: x' \in E$  and  $x' \notin H$  by blast
    obtain  $x': x' \neq 0$ 
    proof
      show  $x' \neq 0$ 
    proof
      assume  $x' = 0$ 
      with  $H$  have  $x' \in H$  by (simp only: vectorspace.zero)
      then show False by contradiction
    qed
  qed

def  $H' \equiv H + \text{lin } x'$ 
  — Define  $H'$  as the direct sum of  $H$  and the linear closure of  $x'$ .

have HH':  $H \trianglelefteq H'$ 
proof (unfold H'-def)
  have vectorspace (lin x') ..
  with  $H$  show  $H \trianglelefteq H + \text{lin } x'$  ..
qed

obtain  $xi$  where
   $\forall y \in H. -p(y + x') - h y \leq xi$ 
   $\wedge xi \leq p(y + x') - h y$ 
  — Pick a real number  $\xi$  that fulfills certain inequations; this will
  — be used to establish that  $h'$  is a norm-preserving extension of  $h$ .

proof —
  from  $H$  have  $\exists xi. \forall y \in H. -p(y + x') - h y \leq xi$ 
   $\wedge xi \leq p(y + x') - h y$ 
  proof (rule ex-xi)
    fix  $u v$  assume  $u: u \in H$  and  $v: v \in H$ 
    with HE have  $uE: u \in E$  and  $vE: v \in E$  by auto
    from  $H u v$  linearform have  $h v - h u = h (v - u)$ 
    by (simp add: linearform.diff)
    also from  $hp$  and  $H u v$  have  $\dots \leq p(v - u)$ 
    by (simp only: vectorspace.diff-closed)
    also from  $x'E uE vE$  have  $v - u = x' + -x' + v + -u$ 
    by (simp add: diff-eq1)
    also from  $x'E uE vE$  have  $\dots = v + x' + -(u + x')$ 
    by (simp add: add-ac)
    also from  $x'E uE vE$  have  $\dots = (v + x') - (u + x')$ 
    by (simp add: diff-eq1)
    also from  $x'E uE vE E$  have  $p \dots \leq p(v + x') + p(u + x')$ 
    by (simp add: diff-subadditive)
    finally have  $h v - h u \leq p(v + x') + p(u + x')$  .
    then show  $-p(u + x') - h u \leq p(v + x') - h v$  by simp

```

```

qed
then show ?thesis ..
qed

def h'  $\equiv$   $\lambda x. \text{let } (y, a) =$ 
  SOME  $(y, a). x = y + a \cdot x' \wedge y \in H \text{ in } h y + a * xi$ 
  — Define the extension  $h'$  of  $h$  to  $H'$  using  $\xi$ .

have  $g \subseteq \text{graph } H' h' \wedge g \neq \text{graph } H' h'$ 
  —  $h'$  is an extension of  $h \dots$ 

proof
show  $g \subseteq \text{graph } H' h'$ 
proof —
  have  $\text{graph } H h \subseteq \text{graph } H' h'$ 
  proof (rule graph-extI)
    fix t assume  $t: t \in H$ 
    have  $(\text{SOME } (y, a). t = y + a \cdot x' \wedge y \in H) = (t, 0)$ 
    by (rule decomp-H'-H)
    with h'-def show  $h t = h' t$  by (simp add: Let-def)
  next
    from HH' show  $H \subseteq H' ..$ 
  qed
  with g-rep show ?thesis by (simp only:)
qed

show  $g \neq \text{graph } H' h'$ 
proof —
  have  $\text{graph } H h \neq \text{graph } H' h'$ 
  proof
    assume eq:  $\text{graph } H h = \text{graph } H' h'$ 
    have  $x' \in H'$ 
    proof (unfold H'-def, rule)
      from H show  $0 \in H$  by (rule vectorspace.zero)
      from x'E show  $x' \in \text{lin } x'$  by (rule x-lin-x)
      from x'E show  $x' = 0 + x'$  by simp
    qed
    hence  $(x', h' x') \in \text{graph } H' h' ..$ 
    with eq have  $(x', h' x') \in \text{graph } H h$  by (simp only:)
    hence  $x' \in H ..$ 
    thus False by contradiction
  qed
  with g-rep show ?thesis by simp
qed

moreover have  $\text{graph } H' h' \in M$ 
  — and  $h'$  is norm-preserving.

proof (unfold M-def)
show  $\text{graph } H' h' \in \text{norm-pres-extensions } E p F f$ 
proof (rule norm-pres-extensionI2)
  show linearform  $H' h'$  by (rule h'-lf)
  show  $H' \trianglelefteq E$ 
  proof (unfold H'-def, rule)
    show  $H \trianglelefteq E .$ 
  qed

```

```

    show vectorspace E .
    from  $x' \in E$  show  $\text{lin } x' \subseteq E$  ..
  qed
  have  $F \subseteq H$  .
  from  $H$  this  $HH'$  show  $FH'$ :  $F \subseteq H'$ 
    by (rule vectorspace.subspace-trans)
  show  $\text{graph } F \subseteq \text{graph } H' h'$ 
  proof (rule graph-extI)
    fix  $x$  assume  $x: x \in F$ 
    with graphs have  $fx = hx$  ..
    also have  $\dots = hx + 0 * xi$  by simp
    also have  $\dots = (\text{let } (y, a) = (x, 0) \text{ in } hy + a * xi)$ 
      by (simp add: Let-def)
    also have  $(x, 0) =$ 
      ( $\text{SOME } (y, a). x = y + a \cdot x' \wedge y \in H$ )
    proof (rule decomp-H'-H [symmetric])
      from  $FH x$  show  $x \in H$  ..
      from  $x'$  show  $x' \neq 0$  .
    qed
    also have
      ( $\text{let } (y, a) = (\text{SOME } (y, a). x = y + a \cdot x' \wedge y \in H)$ 
        in  $hy + a * xi$ ) =  $h'x$  by (simp only: h'-def)
    finally show  $fx = h'x$  .
  next
    from  $FH'$  show  $F \subseteq H'$  ..
  qed
  show  $\forall x \in H'. h'x \leq p x$  by (rule h'-norm-pres)
  qed
  qed
  ultimately show ?thesis ..
  qed
  hence  $\neg (\forall x \in M. g \subseteq x \longrightarrow g = x)$  by simp
  — So the graph  $g$  of  $h$  cannot be maximal. Contradiction!

  then show  $H = E$  by contradiction
  qed

  from HE-eq and linearform have linearform  $E h$ 
    by (simp only:)
  moreover have  $\forall x \in F. hx = fx$ 
  proof
    fix  $x$  assume  $x \in F$ 
    with graphs have  $fx = hx$  ..
    then show  $hx = fx$  ..
  qed
  moreover from HE-eq and hp have  $\forall x \in E. hx \leq p x$ 
    by (simp only:)
  ultimately show ?thesis by blast
  qed

```

## 12.2 Alternative formulation

The following alternative formulation of the Hahn-Banach Theorem uses the fact that for a real linear form  $f$  and a seminorm  $p$  the following inequations

are equivalent:<sup>2</sup>

$$\forall x \in H. |h x| \leq p x \quad \text{and} \quad \forall x \in H. h x \leq p x$$

**theorem** *abs-HahnBanach*:

**includes** *vectorspace*  $E$  + *subspace*  $F E$  + *linearform*  $F f$  + *seminorm*  $E p$

**assumes**  $fp: \forall x \in F. |f x| \leq p x$

**shows**  $\exists g. \text{linearform } E g$

$\wedge (\forall x \in F. g x = f x)$

$\wedge (\forall x \in E. |g x| \leq p x)$

**proof** –

**have**  $\exists g. \text{linearform } E g \wedge (\forall x \in F. g x = f x)$

$\wedge (\forall x \in E. g x \leq p x)$

**proof** (*rule HahnBanach*)

**show**  $\forall x \in F. f x \leq p x$

**by** (*rule abs-ineq-iff* [*THEN iffD1*])

**qed**

**then obtain**  $g$  **where**  $*$  : *linearform*  $E g$   $\forall x \in F. g x = f x$

**and**  $\forall x \in E. g x \leq p x$  **by** *blast*

**have**  $\forall x \in E. |g x| \leq p x$

**proof** (*rule abs-ineq-iff* [*THEN iffD2*])

**show**  $E \trianglelefteq E$  ..

**qed**

**with**  $*$  **show** *?thesis* **by** *blast*

**qed**

### 12.3 The Hahn-Banach Theorem for normed spaces

Every continuous linear form  $f$  on a subspace  $F$  of a norm space  $E$ , can be extended to a continuous linear form  $g$  on  $E$  such that  $\|f\| = \|g\|$ .

**theorem** *norm-HahnBanach*:

**includes** *normed-vectorspace*  $E$  + *subspace*  $F E$  + *linearform*  $F f$  + *fn-norm* + *continuous*  $F \text{ norm } (\|\cdot\|) f$

**shows**  $\exists g. \text{linearform } E g$

$\wedge \text{continuous } E \text{ norm } g$

$\wedge (\forall x \in F. g x = f x)$

$\wedge \|g\|_E = \|f\|_F$

**proof** –

**have**  $E: \text{vectorspace } E$  .

**have**  $E\text{-norm}: \text{normed-vectorspace } E \text{ norm } ..$

**have**  $FE: F \trianglelefteq E$  .

**have**  $F: \text{vectorspace } F$  ..

**have**  $\text{linearform}: \text{linearform } F f$  .

**have**  $F\text{-norm}: \text{normed-vectorspace } F \text{ norm}$

**by** (*rule subspace-normed-vs* [*OF E-norm*])

**have**  $ge\text{-zero}: 0 \leq \|f\|_F$

**by** (*rule normed-vectorspace.fn-norm-ge-zero*  
[*OF F-norm continuous.intro, folded B-def fn-norm-def*])

We define a function  $p$  on  $E$  as follows:  $p x = \|f\| \cdot \|x\|$

**def**  $p \equiv \lambda x. \|f\|_F * \|x\|$

---

<sup>2</sup>This was shown in lemma *abs-ineq-iff* (see page 37).



$p$  is a seminorm on  $E$ :

```

have  $q$ : seminorm  $E$   $p$ 
proof
  fix  $x$   $y$   $a$  assume  $x$ :  $x \in E$  and  $y$ :  $y \in E$ 

```

$p$  is positive definite:

```

  have  $0 \leq \|f\|{-}F$  by (rule ge-zero)
  moreover from  $x$  have  $0 \leq \|x\|$  ..
  ultimately show  $0 \leq p\ x$ 
  by (simp add: p-def zero-le-mult-iff)

```

$p$  is absolutely homogenous:

```

show  $p\ (a \cdot x) = |a| * p\ x$ 
proof -
  have  $p\ (a \cdot x) = \|f\|{-}F * \|a \cdot x\|$  by (simp only: p-def)
  also from  $x$  have  $\|a \cdot x\| = |a| * \|x\|$  by (rule abs-homogenous)
  also have  $\|f\|{-}F * (|a| * \|x\|) = |a| * (\|f\|{-}F * \|x\|)$  by simp
  also have  $\dots = |a| * p\ x$  by (simp only: p-def)
  finally show ?thesis .
qed

```

Furthermore,  $p$  is subadditive:

```

show  $p\ (x + y) \leq p\ x + p\ y$ 
proof -
  have  $p\ (x + y) = \|f\|{-}F * \|x + y\|$  by (simp only: p-def)
  also have  $a$ :  $0 \leq \|f\|{-}F$  by (rule ge-zero)
  from  $x\ y$  have  $\|x + y\| \leq \|x\| + \|y\|$  ..
  with  $a$  have  $\|f\|{-}F * \|x + y\| \leq \|f\|{-}F * (\|x\| + \|y\|)$ 
    by (simp add: mult-left-mono)
  also have  $\dots = \|f\|{-}F * \|x\| + \|f\|{-}F * \|y\|$  by (simp only: right-distrib)
  also have  $\dots = p\ x + p\ y$  by (simp only: p-def)
  finally show ?thesis .
qed

```

$f$  is bounded by  $p$ .

```

have  $\forall x \in F. |f\ x| \leq p\ x$ 
proof
  fix  $x$  assume  $x \in F$ 
  show  $|f\ x| \leq p\ x$ 
    by (unfold p-def) (rule normed-vectorspace.fn-norm-le-cong
      [OF F-norm continuous.intro, folded B-def fn-norm-def])
qed

```

Using the fact that  $p$  is a seminorm and  $f$  is bounded by  $p$  we can apply the Hahn-Banach Theorem for real vector spaces. So  $f$  can be extended in a norm-preserving way to some function  $g$  on the whole vector space  $E$ .

```

with  $E$   $FE$  linearform  $q$  obtain  $g$  where
  linearformE: linearform  $E$   $g$ 
  and  $a$ :  $\forall x \in F. g\ x = f\ x$ 
  and  $b$ :  $\forall x \in E. |g\ x| \leq p\ x$ 
  by (rule abs-HahnBanach [elim-format]) iprover

```

We furthermore have to show that  $g$  is also continuous:

```

have  $g\text{-cont}$ : continuous E norm g using linearformE
proof
  fix  $x$  assume  $x \in E$ 
  with  $b$  show  $|g\ x| \leq \|f\|{-F} * \|x\|$ 
    by (simp only: p-def)
qed

```

To complete the proof, we show that  $\|g\| = \|f\|$ .

```

have  $\|g\|{-E} = \|f\|{-F}$ 
proof (rule order-antisym)

```

First we show  $\|g\| \leq \|f\|$ . The function norm  $\|g\|$  is defined as the smallest  $c \in \mathbb{R}$  such that

$$\forall x \in E. |g\ x| \leq c \cdot \|x\|$$

Furthermore holds

$$\forall x \in E. |g\ x| \leq \|f\| \cdot \|x\|$$

```

have  $\forall x \in E. |g\ x| \leq \|f\|{-F} * \|x\|$ 
proof
  fix  $x$  assume  $x \in E$ 
  with  $b$  show  $|g\ x| \leq \|f\|{-F} * \|x\|$ 
    by (simp only: p-def)
qed
from  $g\text{-cont}$  this ge-zero
show  $\|g\|{-E} \leq \|f\|{-F}$ 
  by (rule fn-norm-least [of g, folded B-def fn-norm-def])

```

The other direction is achieved by a similar argument.

```

show  $\|f\|{-F} \leq \|g\|{-E}$ 
proof (rule normed-vectorspace.fn-norm-least [OF F-norm, folded B-def fn-norm-def])
  show  $\forall x \in F. |f\ x| \leq \|g\|{-E} * \|x\|$ 
  proof
    fix  $x$  assume  $x: x \in F$ 
    from  $a$  have  $g\ x = f\ x$  ..
    hence  $|f\ x| = |g\ x|$  by (simp only:)
    also from  $g\text{-cont}$ 
    have  $\dots \leq \|g\|{-E} * \|x\|$ 
    proof (rule fn-norm-le-cong [of g, folded B-def fn-norm-def])
      from  $FE\ x$  show  $x \in E$  ..
    qed
    finally show  $|f\ x| \leq \|g\|{-E} * \|x\|$  .
  qed
show  $0 \leq \|g\|{-E}$ 
  using  $g\text{-cont}$ 
  by (rule fn-norm-ge-zero [of g, folded B-def fn-norm-def])
next
  show continuous F norm f by (rule continuous.intro)
qed
with linearformE a g-cont show ?thesis by blast

```

qed

end

## References

- [1] H. Heuser. *Funktionalanalysis: Theorie und Anwendung*. Teubner, 1986.
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